

Sunday, April 24, 2022

Assignment 7

CSCI 411

1.

- a. $(0+1)(00+11+01+10)^*$
- b. $\epsilon + (01+1)(0+1)^*$
- c. $(0+1)^*(01+10+11) + \epsilon + 0 + 1$
- d. $((10+1)^* + (01+1)^*)(00 + \epsilon)((10+1)^* + (01+1)^*)$

2.

- Remove states that are not reachable from starting state

	0	1
->A	B	F
B	D	G
*D	G	E
E	H	E
*F	A	G
G	D	G
H	D	G

- Keep on separating the final and non final states ill we get all states output is lies in a common class for both 0 and 1.

	0	1
A	B	F
B, G, H	D	G
D	G	E
E	H, B	E
*F	A	G

	A	B	C	D	E	F	G	
B	X							
C	X	X						
D	X	X						
E	X	X		X				
F	X	X	X	X	X			
G	X		X	X	X	X		
H	X		X	X	X	X		
I	X	X		X	X	X	X	X

EQ State: {B, G}, {B, H}, {C, D}, {C, E}, {C, I}, {G, H}

3.

* I - Intermediate state, F - Final state

If the pairs forming with 0 and 1 as input are symmetric(i.e. I, I and F, F), then M1 and M2 are equivalent.

	0	1
A, D	B, E (I, I)	A, D (F, F)
B, F	C, G (F, F)	A, D (F, F)
C, G	C, F (F, F)	A, G (F, F)
B, E	C, F (F, F)	A, G (F, F)

We can see that all the pairs formed in the transition table are symmetric. Hence, M1 and M2 are equivalent.

4.

* I - Intermediate state, F - Final state

If the pairs forming with 0 and 1 as input are symmetric(i.e. I, I and F, F), then M1 and M2 are equivalent.

	0	1
A, C	A, D (F, F)	B, E (I, I)
B, E	A, C (F, F)	B, E (I, I)
A, D	A, D (F, F)	B, E (I, I)

We can see that all the pairs formed in the transition table are symmetric. Hence, M1 and M2 are equivalent.

5.

a. let L be regular length and P be pumping length.

w belongs to L , where $|w| \leq P$

$w = xy$

conditions: $|y| < P$, (y is composed only of zeros)

$|xy| \geq p$, (x and y lies within p)

$\forall i < \infty, xy^i$ belongs to L (if $i=0$, $xy^0 = x = 0^{(p-k)}$)

this contradicts L . Hence, L is not a regular language.

b.

c. let L be regular length and P be pumping length.

let $w^2((w^r)^2)$ belongs to L , where $|w| > 1$ and $|ww| \leq P$ and w belongs to L

now let $P \geq 4$

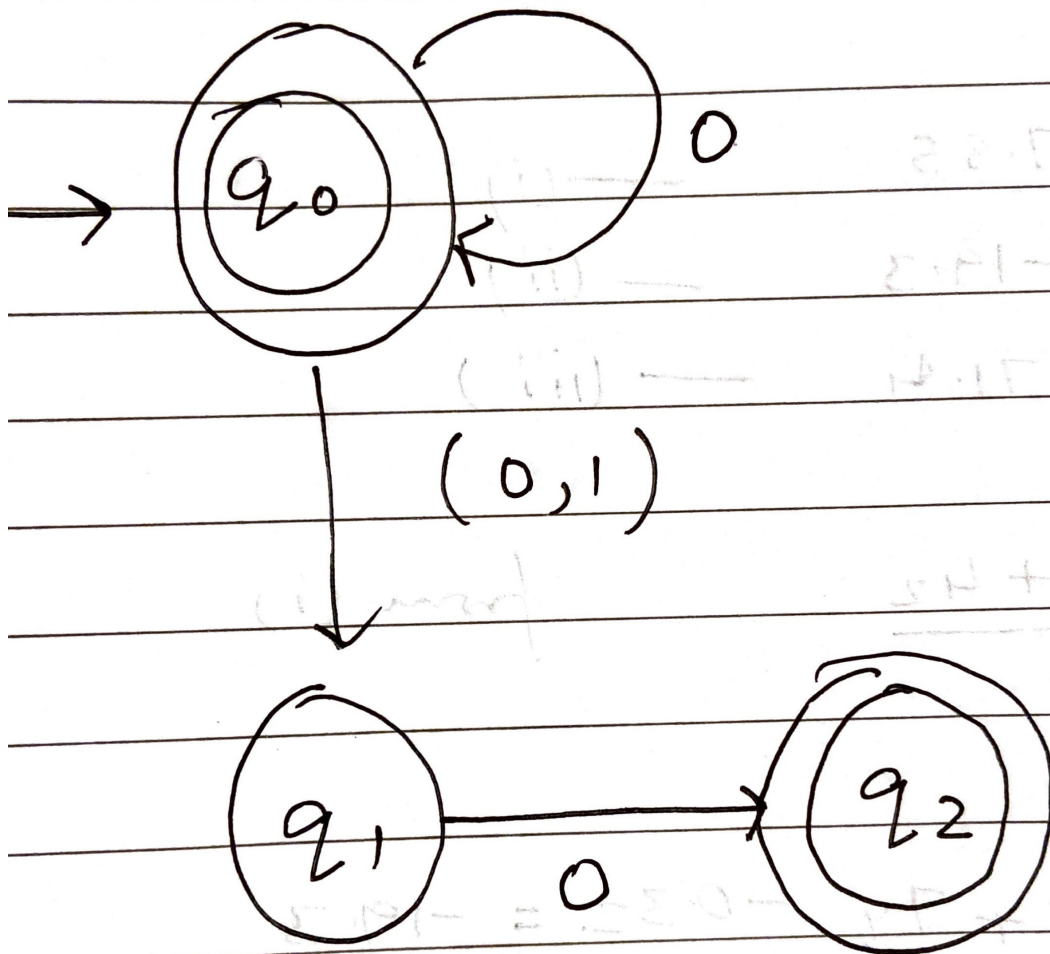
and xyz where $x = y = z$, and $z = (w^r)^2$ which is does not

$$\{ww^R | w \in \{0, 1\}^*\}.$$

Hence, $\{ww^R | w \in \{0, 1\}^*\}$ is not a regular language.

6.

- a. Initial state of finite automata will be final because empty string is included.



Binary number will end with a 0 for even numbers.

b. This is not a regular language because there is comparison.

c.

This is not a regular language because there is comparison.

For a regular language L , there exists an integer value ' n ', such that for all ' w ' belongs to L with $|w| \geq n$.

$w = xyz$, $|xy| \leq n$, $|y| > 0$ or $|y| \geq 1$ and for all $i \geq 0$: $xy^i z$ belongs to L

Let L be regular length and P be pumping length.

Let n be the states of DFA for which and example string $s = 0^n 1 0^n 1$

Now, $s = uvw$, if u and v were empty and then we could pump the string, but according to the pumping lemma Condition, v must be all 0's.

Thus, $uv^i w$ is not in language L because contains 1 which indicates that s does not work.

Hence, L is not regular.

- d. This is a regular language.

NFA:

