

Sunday, February 27, 2022

## Assignment 4

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CSCI 411

1.

- a. Given  $A = [5, 8, 8, 3, 4, 1, 7, -3, 2, 9, 12]$ ,

A sequence is called Bitonic if it is first increasing, then decreasing.

Hence, the longest Bitonic sequence of A should be:

$[8, 3, 1, -3, 2, 9, 1]$

- b. A sequence is called Bitonic if it is first increasing, then decreasing.

For example,

$S = [1, 11, 2, 10, 4, 5, 2, 1]$

Skip  $\implies 11, 5$

$[ \text{---} \rightarrow i \text{---} \rightarrow d ]$

$LBS = [1, 2, 10, 4, 2, 1]$

A sequence, sorted in increasing order is considered Bitonic with the decreasing part as empty. Similarly, decreasing order sequence is considered Bitonic with the increasing part as empty. Using two subsequences arrays, the Longest Increasing Subsequence(LIS) and the Longest Decreasing Subsequence(LDS). The LIS will hold the length of increasing subsequence ending with  $\text{array}[i]$ . The LDS will store the length of decreasing subsequence starting from  $\text{array}[i]$ . Using these two arrays, we can get the length of longest bitonic subsequence.

c.

LBS(a):

inc\_sub\_seq.size = a.size

For i = 1..... a.size -1:

For j = 0.....i - 1:

If  $a[i] > a[j]$  and  $\text{inc\_sub\_seq}[i] < \text{inc\_sub\_sum}[j] + 1$ :

$\text{inc\_sub\_sum}[i] = \text{inc\_sub\_sum}[j] + 1$

dec\_sub\_seq.size = a.size

For i = a.size - 2.....0:

For j = a.size - 1.....i + 1:

If  $a[i] > a[j]$  and  $\text{dec\_sub\_seq}[i] < \text{dec\_sub\_sum}[j] + 1$ :

$\text{dec\_sub\_seq}[i] = \text{dec\_sub\_seq}[j] + 1$

$M = \text{inc\_sub\_seq}[0] + \text{dec\_sub\_seq}[0] - 1$

For i = 1..... a.size:

If  $\text{inc\_sub\_seq}[i] + \text{dec\_sub\_seq}[j] - 1 > M$ :

$M = \text{inc\_sub\_seq}[i] + \text{dec\_sub\_seq}[i] - 1$

Return M

d. Time complexity is  $O(n^2)$ , where n is the size of A because we are looping twice for the calculation of LIS and LDS.

2.

a. "Exponential" and "Polynomial"

0	1	2	3	4	5	6	7	8	9	10	11
1	1	2	2	3	4	5	6	7	8	9	10
2	2	2	3	2	3	4	5	6	7	8	9
3	3	3	3	3	3	4	5	6	7	8	8
4	4	4	4	4	4	4	5	6	7	8	10
5	5	5	5	5	4	5	4	5	6	7	8
6	6	6	6	5	5	5	5	5	6	7	8

```

7  7  7  7 6 6 6 6 6 7 8
8  8  8  8 7 7 7 7 7 6 7 8
9  9  9  9 8 8 8 8 8 7 6 7
10 10 10 10 9 9 9 9 9 8 7 6

```

6 is the answer!

b. Given A = "cat" and B = "cut",  
ins = 3, sub = 2, del = 6

replacing 'a' with 'u'  
cost = 2

Given A = "sunday" and B = "saturday",  
ins = 3, sub = 2, del = 6

Convert 'un' with 'atur'  
Replace 'n' with 'r'  
Insert 't'  
Insert 'a'  
cost = 2 + 3 + 3 = 8

c.

editDistance(A, B, ins, del, sub):

```

buffer[A.size][B.size]
buffer[0][0].score = 0
buffer[0][0].act = "match"

```

```

For i=1.....A.size:
    buffer[0][i].act = "delete"
    buffer[0][i].score = buffer[i-1][0].score + del

```

```

For i=1.....B.size:
    buffer[i][0].act = "insert"
    buffer[i][0].score = buffer[0][i-1].score + ins

```

```

For i=1.....A.size:
    For j=1.....A.size:

```

```
    If (a[i-1] != b[j-1]):  
        buffer[i][j].act = min_cost()  
        buffer[i][j].score = min_cost()
```

```
Result->score = buffer[A.size][B.size].score
```

```
While A.size > 0 || B.size > 0:
```

```
    If "insert":
```

```
        Replace with '_'
```

```
        A.size--
```

```
    Else If "delete":
```

```
        Replace with '_'
```

```
        B.size--
```

```
    Else:
```

```
        Replace with '_'
```

```
        A.size--
```

```
        B.size--
```

```
Result->a = modified_a
```

```
Result->b = modified_b
```

```
Return Result
```

d. Time complexity of this algorithm is  $O(m \times n)$  where  $m$  and  $n$  are the sizes of string  $A$  and  $B$  respectively.