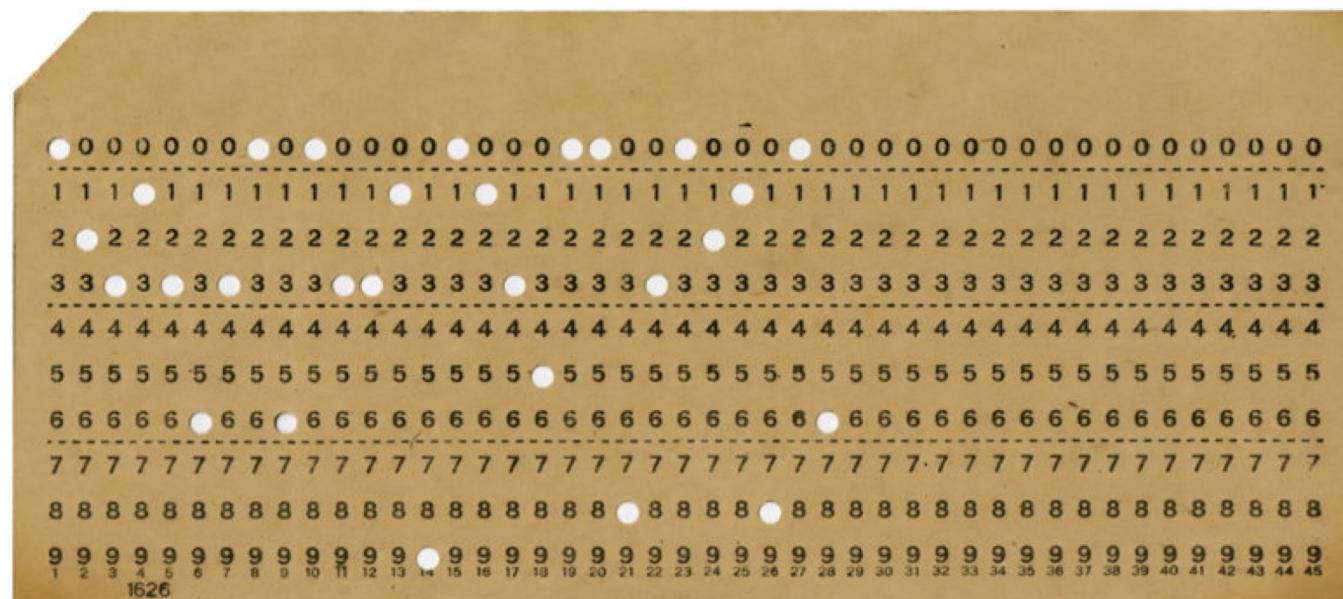


# Computer Architecture, Fall 2019

# *Course Introduction, Number Representation*



# Agenda

- Course Overview
- Number Representation
  - Number Bases
  - Signed Representations
  - Overflow
  - Sign Extension

# My Computers

Late 1980's



- CPU: Zilog Z80 3.58 MHz
- ROM: 32KB  
(BIOS 16KB + MSX Basic v1.0  
16KB)
- RAM: 8KB

Fast forward ~30 years:

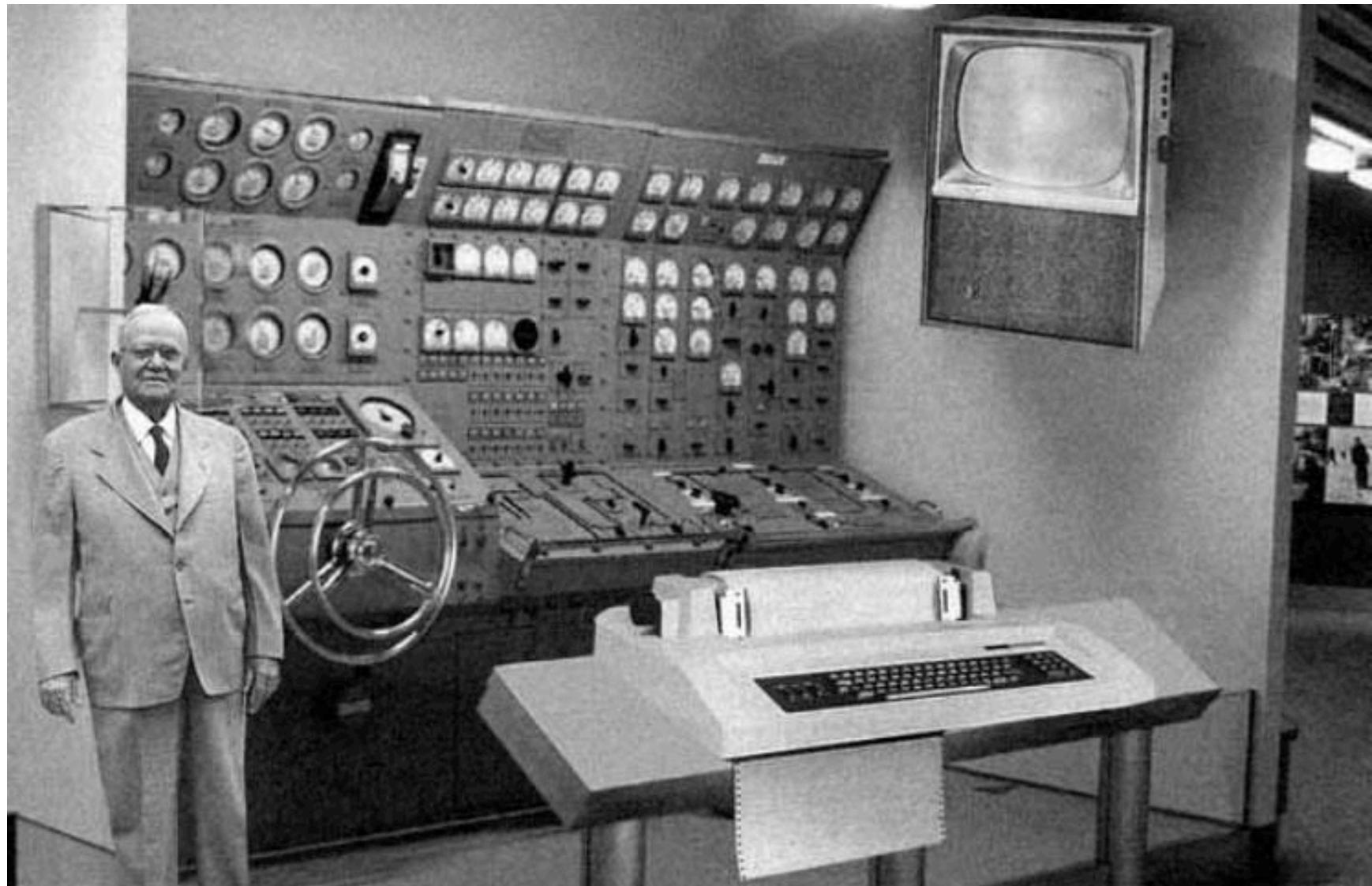


Present 2018



- CPU: Intel Xeon 40 CPUs 2.4 GHz
- ROM: SSD 8TB
- RAM: 256 GB
- 10 NVIDIA GTX Titan XP GPUs
- 10 Gbps network

# Computer Architecture Course in Old School



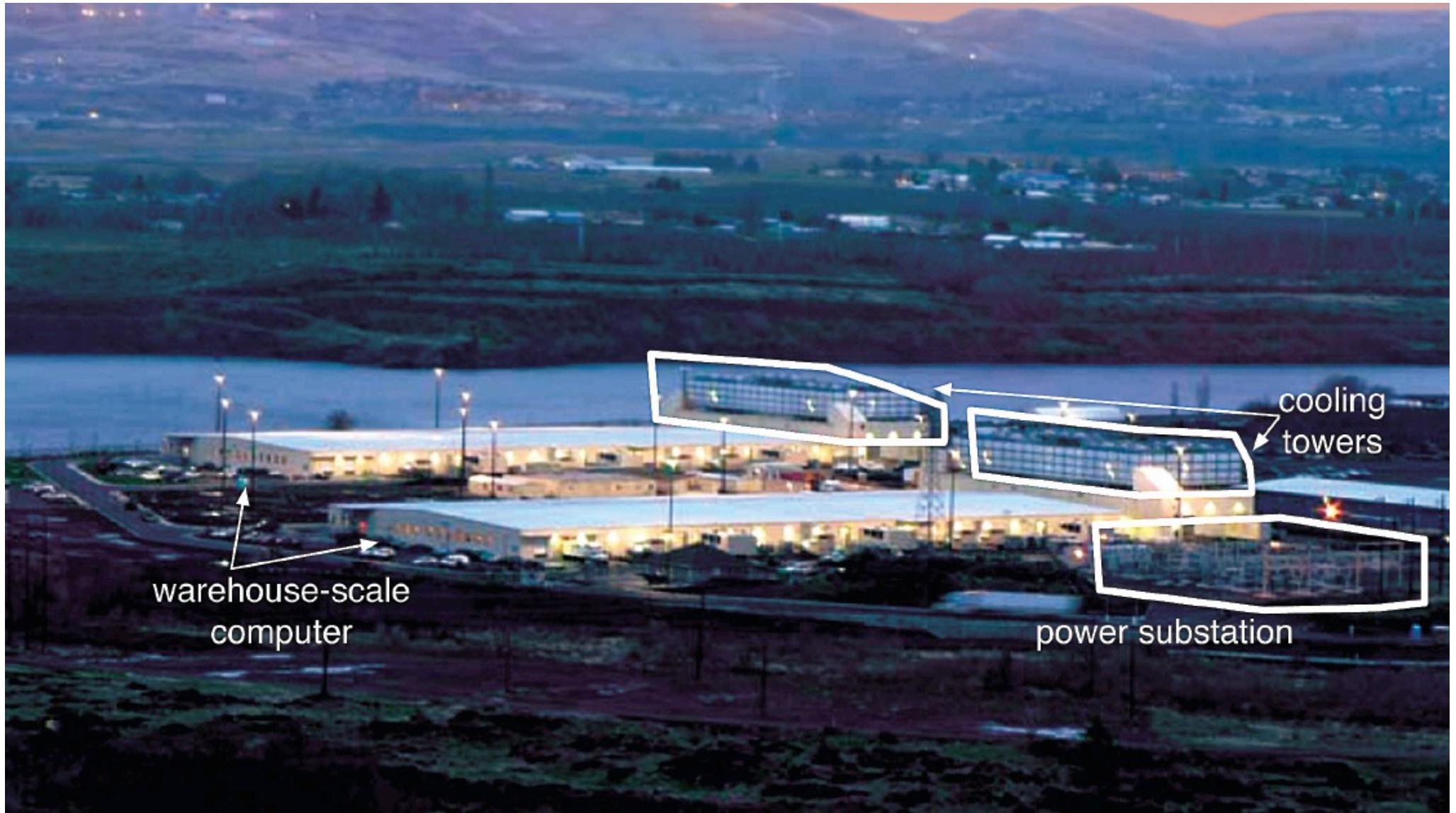
# Computer Architecture Course in New School (1/3)

Personal  
Mobile  
Devices

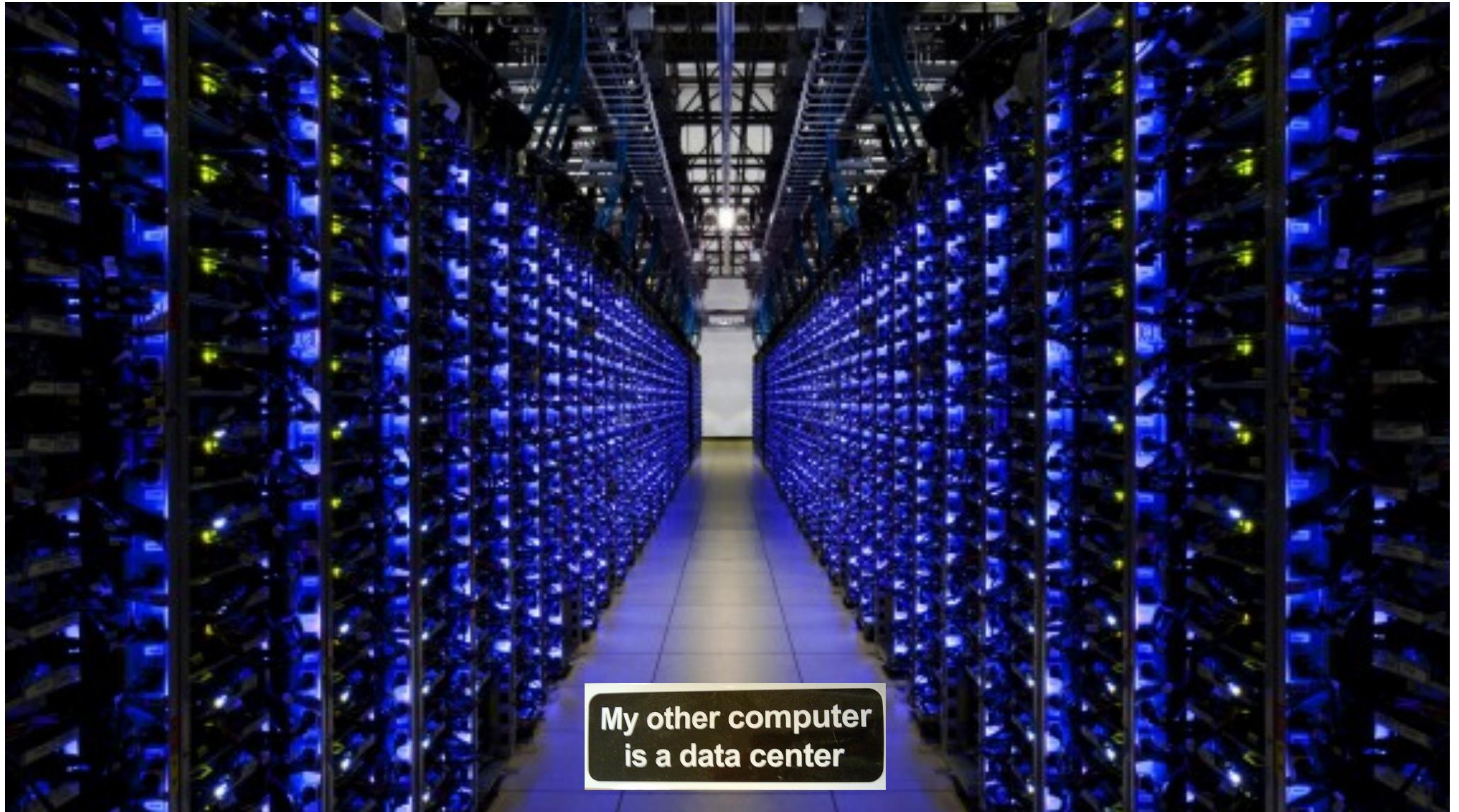


Network  
Edge  
Devices

# Computer Architecture Course in New School (2/3)

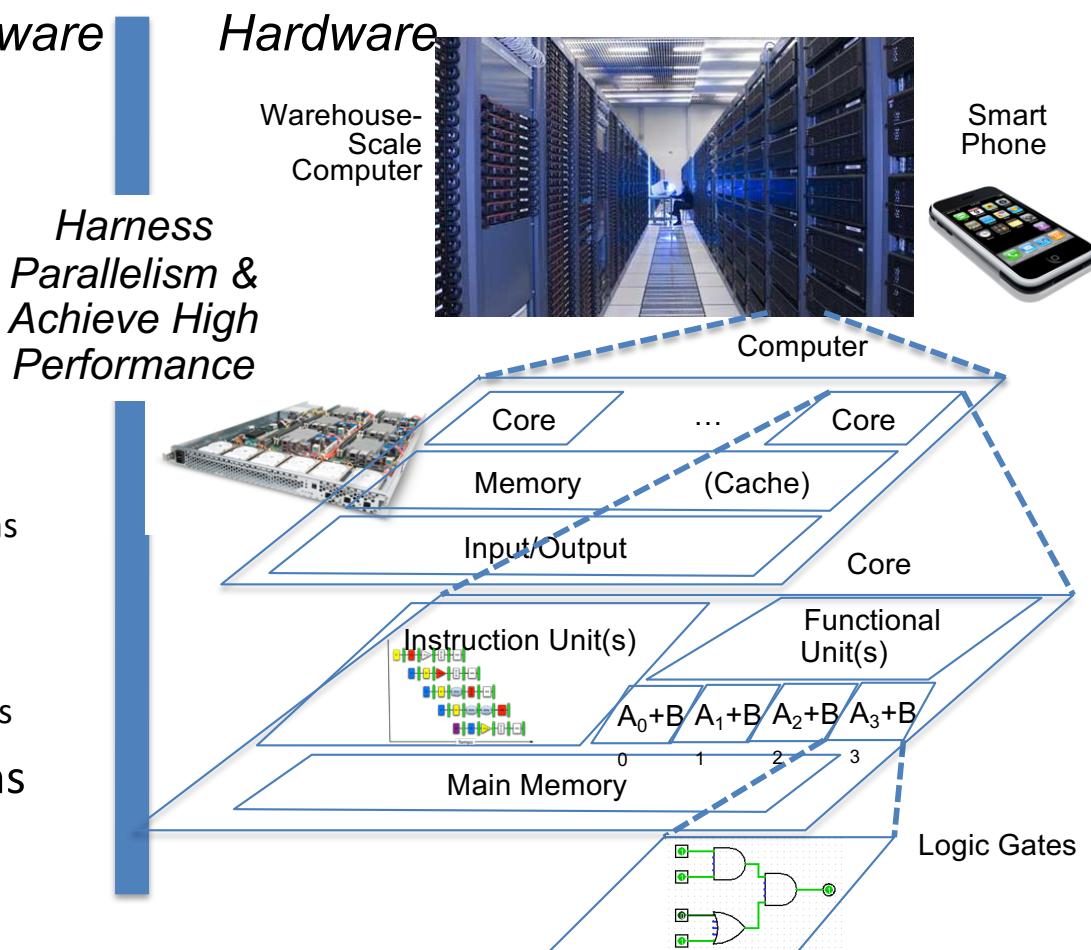


# Computer Architecture Course in New School (3/3)



# New-School Machine Structures

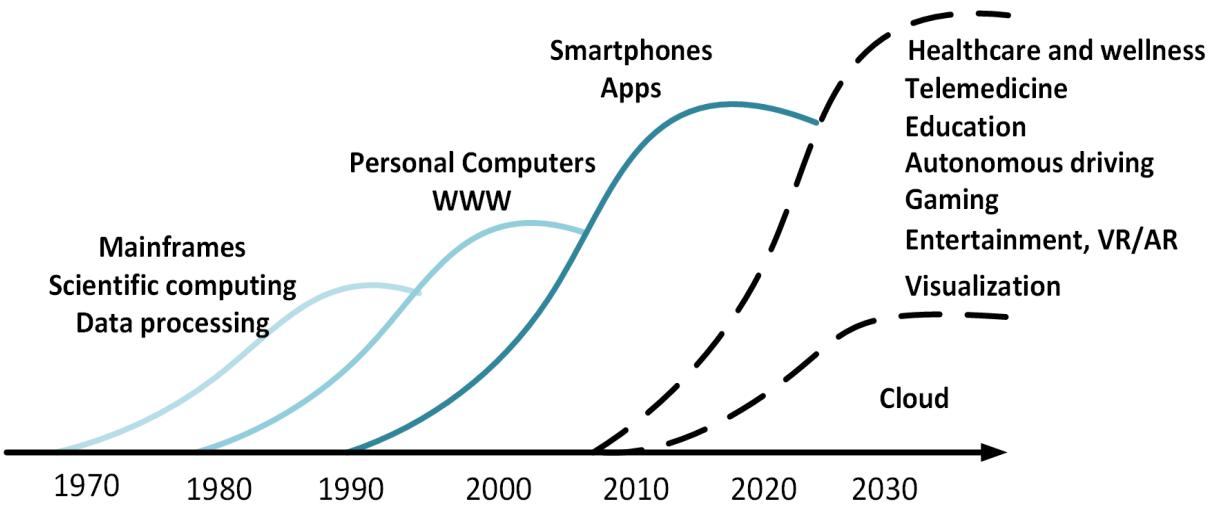
- Parallel Requests Software  
Assigned to computer  
e.g., Search “cats”
- Parallel Threads  
Assigned to core  
e.g., Lookup, Ads
- Parallel Instructions  
 $>1$  instruction @ one time  
e.g., 5 pipelined instructions
- Parallel Data  
 $>1$  data item @ one time  
e.g., Add of 4 pairs of words
- Hardware descriptions  
All gates functioning in parallel at same time



# Six Great Ideas in Computer Architecture

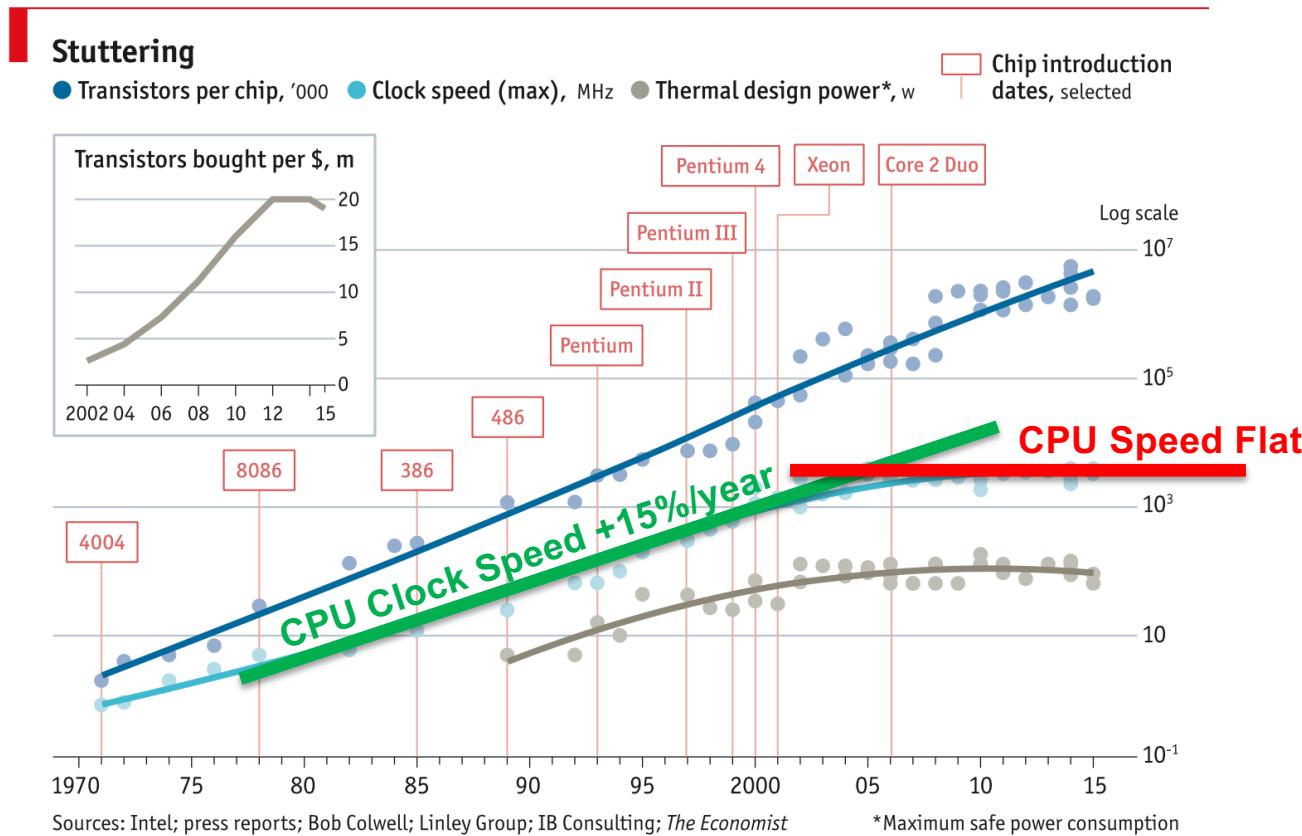
- 1) Abstraction
- 2) Technology Trends
- 3) Principle of Locality/Memory Hierarchy
- 4) Parallelism
- 5) Performance Measurement & Improvement
- 6) Dependability via Redundancy

# Why is Architecture Exciting Today?



- Number of deployed devices continues growing, but no single killer app
  - Diversification of needs, architectures

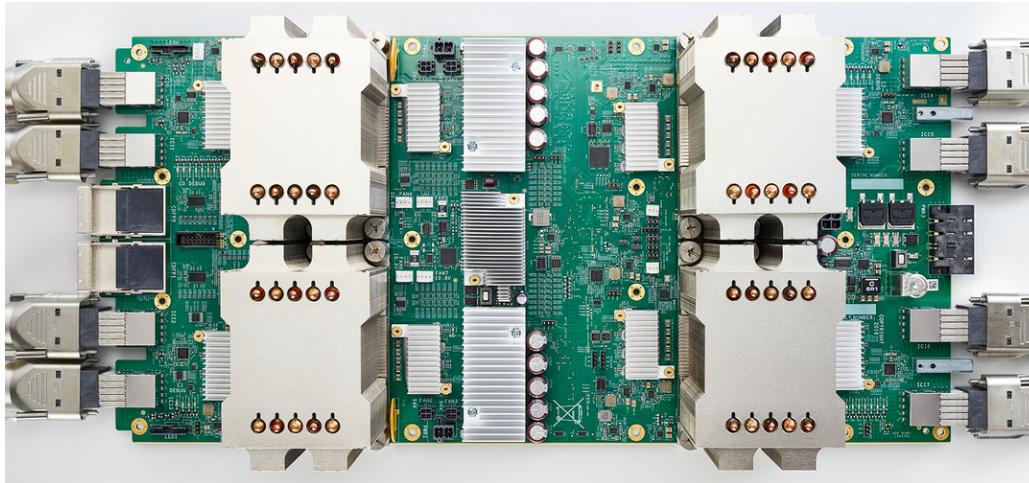
# Why is Architecture Exciting Today?



# Old Conventional Wisdom

- Faster, cheaper, lower-power general-purpose computers each year
- In glory days, 1%/week performance improvement!
- Dumb to compete by designing parallel or specialized computers
- By time you've finished design, next generation of general-purpose will beat you

# New Conventional Wisdom



Google TPU2  
Specialized Engine for neural  
network training  
Deployed in cloud  
45 TFLOPS/chip



Serious heatsinks!

# Agenda

- Course Overview
- Number Representation
  - Number Bases
  - Signed Representations
  - Overflow
  - Sign Extension

# Number Representation



0	1	2	3	4	5	6	7	8	9
	•	••	•••	••••		•	••	•••	••••
10	11	12	13	14	15	16	17	18	19
=====	=====	=====	=====	=====	=====	•	••	•••	••••

### Example:

$$28 = (1 \times 20) + 8 = \underline{\dots}$$

$$433 = (1 \times 400) + (1 \times 20) + 13 =$$

# MAYANS

7 1	𠂇 11	𠂇 21	𠂇 31	𠂇 41	𠂇 51
7 2	𠂇 12	𠂇 22	𠂇 32	𠂇 42	𠂇 52
7 3	𠂇 13	𠂇 23	𠂇 33	𠂇 43	𠂇 53
7 4	𠂇 14	𠂇 24	𠂇 34	𠂇 44	𠂇 54
7 5	𠂇 15	𠂇 25	𠂇 35	𠂇 45	𠂇 55
7 6	𠂇 16	𠂇 26	𠂇 36	𠂇 46	𠂇 56
7 7	𠂇 17	𠂇 27	𠂇 37	𠂇 47	𠂇 57
7 8	𠂇 18	𠂇 28	𠂇 38	𠂇 48	𠂇 58
7 9	𠂇 19	𠂇 29	𠂇 39	𠂇 49	𠂇 59
7 10	𠂇 20	𠂇 30	𠂇 40	𠂇 50	

# BABYLONIANS



# Number Representation

- Numbers are an abstract concept!
  - Recognize that these are *arbitrary* symbols
- Inside a computer, everything stored as a sequence of 0's and 1's (bits)
  - Even this is an abstraction!
- How do we represent numbers in this format?
  - Let's start with integers

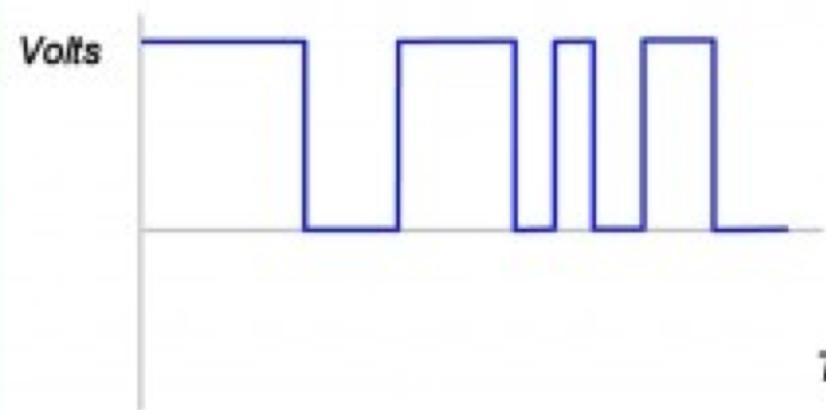
# Computers: Base 2



**Analog signal**



**Digital Signal**



# Number Bases

- **Key terminology:** digit ( $d$ ) and base ( $B$ )
- In base  $B$ , each digit is one of  $B$  possible symbols
- Value of  $i$ -th digit is  $d \times B^i$  where  $i$  starts at 0 and increases from right to left
  - $n$  digit number  $d_{n-1}d_{n-2} \dots d_1d_0$
  - value =  $d_{n-1} \times B^{n-1} + d_{n-2} \times B^{n-2} + \dots + d_1 \times B^1 + d_0 \times B^0$
- Notation: Base is indicated using either a prefix or a subscript

# Commonly Used Number Bases

- **Decimal** (base 10)
  - Symbols: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9
  - Notation:  $9472_{\text{ten}} = 9472$
- **Binary** (base 2)
  - Symbols: 0, 1
  - Notation:  $101011_{\text{two}} = 0b101011$
- **Hexadecimal** (base 16)
  - Symbols: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F
  - Notation:  $2A5D_{\text{hex}} = 0x2A5D$

Decimal	Binary	Hex
0	0000	0
1	0001	1
2	0010	2
3	0011	3
4	0100	4
5	0101	5
6	0110	6
7	0111	7
8	1000	8
9	1001	9
10	1010	A
11	1011	B
12	1100	C
13	1101	D
14	1110	E
15	1111	F

# Number Base Examples

- Examples:

3 2 1 0

$$9472_{\text{ten}} = 9000 + 400 + 70 + 2$$

$$9 \times 1000 + 4 \times 100 + 7 \times 10 + 2 \times 1$$

$$9 \times 10^3 + 4 \times 10^2 + 7 \times 10^1 + 2 \times 10^0$$

$$9472_{\text{ten}} = 2 \times 16^3 + 5 \times 16^2 + 0 \times 16^1 + 0 \times 16^0$$

$$= 2500_{\text{hex}}$$

$$0xA15 = 0b\ 1010\ 0001\ 0101$$

# Bits Can Represent Anything

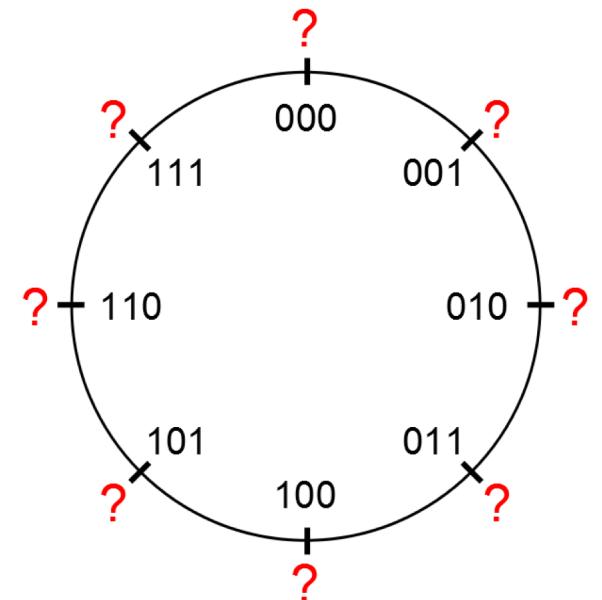
- *n* digits in base  $B$  can represent at most  $B^n$  things!
  - Each of the  $n$  digits is one of  $B$  possible symbols
  - Have more things? Add more digits!
- Example: Logical values (1 bit) – 0 is False, 1 is True
- Example: Characters
  - 26 letters require 5 bits ( $2^5 = 32 > 26$ )
- Example: Students in this class (6 bits)
- For convenience, can group into *nibbles* (4 bits) and *bytes* (8 bits)

# So What Number Is It Anyway?

- Here we are using binary bit patterns to **represent** numbers
  - Strictly speaking they are called *numerals* and have no meaning until you **interpret** them
  - Is CAB a word (taxi) or a number ( $3243_{\text{ten}}$ )?
  - Is 0x999999 a number or a color (RGB)?
- Keep in mind that the same bit pattern will mean different things depending on how you choose to interpret it

# Numbers in a Computer

- Numbers really have  $\infty$  digits, but hardware can only store a finite number of them (fixed)
  - Usually ignore *leading zeros*
- “Circle” of binary numerals:  
(3-bit example)



# Unsigned Integers

Represent only non-negative (unsigned) integers:

$$0000_{\text{two}} = 0_{\text{ten}}$$

$$0001_{\text{two}} = 1_{\text{ten}}$$

...

$$0111_{\text{two}} = 7_{\text{ten}}$$

$$1000_{\text{two}} = 8_{\text{ten}}$$

...

$$1110_{\text{two}} = 14_{\text{ten}}$$

$$1111_{\text{two}} = 15_{\text{ten}}$$



**Zero?**

$$0\dots 0_{\text{two}} = 0_{\text{ten}}$$

**Most neg number?**

$$0\dots 0_{\text{two}} = 0_{\text{ten}}$$

**Most pos number?**

$$1\dots 1_{\text{two}} = (2^n - 1)_{\text{ten}}$$

**Increment?**

# Signed Integers

- $n$  bits can represent  $2^n$  different things
  - Ideally, want the range evenly split between positive and negative
- How do we want to encode zero?
- Can we encode them in such a way that we can use the same hardware regardless of whether the numbers are signed or unsigned?
  - e.g. incrementing pos/neg numbers

# Sign and Magnitude

“first” bit gives sign, rest treated as unsigned (magnitude):

$$000_{\text{two}} = +0_{\text{ten}}$$

$$001_{\text{two}} = +1_{\text{ten}}$$

$$010_{\text{two}} = +2_{\text{ten}}$$

$$011_{\text{two}} = +3_{\text{ten}}$$

---

$$100_{\text{two}} = -0_{\text{ten}}$$

$$101_{\text{two}} = -1_{\text{ten}}$$

$$110_{\text{two}} = -2_{\text{ten}}$$

$$111_{\text{two}} = -3_{\text{ten}}$$

**Zero?**

$$0\dots0_{\text{two}} \text{ and } 10\dots0_{\text{two}} = \pm 0_{\text{ten}}$$

**TWO ZEROS?!**

**Most pos number?**

$$01\dots1_{\text{two}} = (2^{(n-1)} - 1)_{\text{ten}}$$

**Most neg number?**

$$1\dots1_{\text{two}} = -(2^{(n-1)} - 1)_{\text{ten}}$$

**Increment?**

# Biased Notation

Like unsigned, but “shifted” so zero is (roughly) in the middle: (value = “unsigned value” - bias)

**PRO**  $000_{\text{two}} = -3_{\text{ten}}$

$001_{\text{two}} = -2_{\text{ten}}$

$010_{\text{two}} = -1_{\text{ten}}$

**CON**  $011_{\text{two}} = 0_{\text{ten}}$  Conventional Bias:  $(2^{n-1}-1)$

$100_{\text{two}} = +1_{\text{ten}}$

$101_{\text{two}} = +2_{\text{ten}}$

$110_{\text{two}} = +3_{\text{ten}}$

$111_{\text{two}} = +4_{\text{ten}}$

**PRO**

**Zero?**

$01\dots1_{\text{two}} = 0_{\text{ten}}$

**Most neg number?**

$0\dots0_{\text{two}} = -(2^{n-1}-1)_{\text{ten}}$

**Most pos number?**

$1\dots1_{\text{two}} = 2^{(n-1)}_{\text{ten}}$

**Increment?**

**PRO** just like unsigned (sorta)

# One's Complement

New negation procedure – complement the bits:

$$000_{\text{two}} = +0_{\text{ten}}$$

$$001_{\text{two}} = +1_{\text{ten}}$$

$$010_{\text{two}} = +2_{\text{ten}}$$

$$011_{\text{two}} = +3_{\text{ten}}$$

---

$$100_{\text{two}} = -3_{\text{ten}}$$

$$101_{\text{two}} = -2_{\text{ten}}$$

$$110_{\text{two}} = -1_{\text{ten}}$$

$$111_{\text{two}} = -0_{\text{ten}}$$

**Zero?**

$$0\dots 0_{\text{two}} \text{ and } 1\dots 1_{\text{two}} = \pm 0_{\text{ten}}$$

**TWO ZEROS AGAIN.....**

**Most neg number?**

$$10\dots 0_{\text{two}} = -(2^{(n-1)} - 1)_{\text{ten}}$$

**Most pos number?**

$$01\dots 1_{\text{two}} = (2^{(n-1)} - 1)_{\text{ten}}$$

**Increment?**

# Two's Complement

Like One's Complement, but “shift” negative #s by 1:

$$000_{\text{two}} = +0_{\text{ten}}$$

$$001_{\text{two}} = +1_{\text{ten}}$$

$$010_{\text{two}} = +2_{\text{ten}}$$

$$011_{\text{two}} = +3_{\text{ten}}$$

$$100_{\text{two}} = -4_{\text{ten}}$$

$$101_{\text{two}} = -3_{\text{ten}}$$

$$110_{\text{two}} = -2_{\text{ten}}$$

$$111_{\text{two}} = -1_{\text{ten}}$$

**Zero?**

$$0\dots0_{\text{two}} = 0_{\text{ten}}$$

**Most neg number?**

$$10\dots0_{\text{two}} = -2^{(n-1)}_{\text{ten}}$$

**Most pos number?**

$$01\dots1_{\text{two}} = (2^{(n-1)} - 1)_{\text{ten}}$$

**Increment?**

# Two's Complement Summary

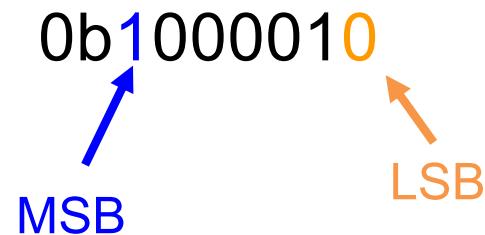
- Used by all modern hardware
- Roughly evenly split between positive and negative
  - One more negative # because positive side has 0
- Can still use MSB as sign bit
- To negate: Flip the bits and add one
  - Example:  $+7 = 0b\ 0000\ 0111$ ,  $-7 = 0b\ 1111\ 1001$

# Two's Complement Review

- Suppose we had 5 bits. What integers can be represented in two's complement?
  - (A) -31 to +31 ← need 6 bits
  - (B) -15 to +15 ← one's complement
  - (C) 0 to +31 ← unsigned
  - (D) -16 to +15 ← two's complement**
  - (E) -32 to +31 ← need 6 bits

# Numbers in a Computer

- Numbers really have  $\infty$  digits, but hardware can only store a finite number of them (fixed)
  - Usually ignore *leading zeros*
  - Leftmost is *most significant bit* (MSB)
  - Rightmost is *least significant bit* (LSB)



# Overflow

- **Overflow** is when the result of an arithmetic operation can't be represented by the hardware bits
  - i.e. the result is mathematically incorrect
- Examples:
  - Unsigned:  $0b1\dots1 + 1_{10} = 0b0\dots0 = 0?$
  - Two's:  $0b01\dots1 + 1_{10} = 0b10\dots0 = -2^{(n-1)}_{ten}?$



# Sign Extension

- Want to represent the same number using more bits than before
  - Easy for positive #'s (add leading 0's), more complicated for negative #'s
  - Sign and magnitude: add 0's *after* the sign bit
  - One's complement: copy MSB
  - Two's complement: copy MSB
- Example:
  - Sign and magnitude:  $0b\ 11 = 0b\ 1001$
  - One's/Two's complement:  $0b\ 11 = 0b\ 1111$

# Summary

- Number Representation: How to represent positive and negative integers using binary
  - Unsigned: Interpret numeral in base 2
  - Signed: Two's Complement
  - Biased: Subtract bias
  - Sign extension must preserve signed number