RD-Gen: Random DAG Generator Considering Multi-rate Applications for Reproducible Scheduling Evaluation

Abstract—Real-time systems include various requirements such as deadline and resource constraints. Such systems are becoming larger and more complex, and studies on performance analyses and efficient scheduling algorithms are becoming increasingly important. Directed acyclic graph (DAG) models. which can express task dependencies and parallelism, are used for such studies. Random DAG sets are used to demonstrate the effectiveness and objectivity of methods proposed for realtime systems. However, there is no random DAG generation tool available that can generate a DAG set that considers the latest multi-rate applications. Therefore, researchers need to generate random DAG sets on their own, leading to additional effort and reduced reliability and reproducibility. To solve this problem, we propose a random DAG generator considering multi-rate applications for reproducible scheduling evaluation (RD-Gen). RD-Gen also enables batch generation of random DAG sets with different parameters. Case studies are used to demonstrate that RD-Gen can manage various problem settings and DAG study requirements.

Index Terms—DAG, Random generation tool, Multi-rate applications

I. INTRODUCTION

Real-time systems, such as self-driving systems, need to successfully execute while meeting various requirements such as producing an output within a pre-determined time (i.e., meeting deadlines), having low power consumption and meeting resource constraints [1]–[3]. To fulfill these constraints, many studies have been conducted on task allocation and scheduling [4]–[6], as well as on analyzing the end-to-end latency and the response time of a system [7]–[9]. Systems are increasingly becoming larger and more complex, and many studies use models, such as directed acyclic graphs (DAGs), to represent the complex dependencies and parallelism of the tasks in these systems.

DAGs are used in many allocation, scheduling, and latency analysis studies [10]–[12] because they express the process flow from system input to output and can represent various types of information such as the dependencies between tasks, the task execution time, and the execution period. To evaluate the performance of proposed methods, it is important to compare them with existing methods using task sets. Accordingly, method evaluations using DAGs, randomly generated DAGs are used to ensure objectivity and demonstrate generality [2], [13], [14].

To aid in such evaluations, random DAG generation tools such as task graph for free (TGFF) [15] and () GGen [16], have been proposed and utilized in the latest publications [17]–[20].

These tools allow the user to parametrically specify the shape of a DAG (e.g., the number of nodes and the in-degree and out-degree per node) and the properties assigned to the tasks and edges (e.g., the execution time, the execution period, and the communication time). Furthermore, because these tools use a pseudo-random number generator, other researchers can easily reproduce a given DAG set by specifying the same options. However, TGFF and GGen have been proposed in 1999 and 2010, respectively, and cannot meet the requirements of multirate DAGs considering state-of-the-art real-time systems.

Because embedded systems in automobiles and avionics, as well as in self-driving systems, contain multiple tasks that operate over different periods (e.g., sensors [21], localization [14], and angle synchronization [22]), studies targeting multirate DAGs are becoming increasingly important [8], [23]. In studies of such multi-rate DAGs, not only the shape of the DAG but also the ratio of the execution time to the task execution period has a significant impact on the performance of the method (e.g., implicit deadlines [24], [25] and task utilization [26], [27]). However, TGFF cannot generate multirate DAGs, and GGen can only randomly set the period and the execution time of tasks. Therefore, most researchers who consider multi-rate DAGs have to implement their own random DAG sets [26]–[29]. It is laborious for researchers to prepare their own random DAG sets, and this practice further reduces the reliability and reproducibility of the evaluation results.

To solve these problems, this paper proposes a random DAG generator considering multi-rate applications for reproducible scheduling evaluation (RD-Gen). RD-Gen extends existing DAG generation methods and provides a flexible evaluation platform. Because RD-Gen uses a pseudo-random number generator, other researchers can reproduce the DAG sets used in an evaluation by specifying the same options.

Contributions: Our primary contributions are summarized as follows.

- RD-Gen provides flexible random DAG sets by adding parameters to existing DAG construction methods and a new chain-based construction method.
- RD-Gen automatically sets various properties, including complex properties calculated from multiple values.
- RD-Gen reduces implementation effort via the batch generation of random DAG sets.

The remainder of the paper is organized as follows. Section II describes a system model. Section III explains the design and implementation of RD-Gen. Section IV presents case

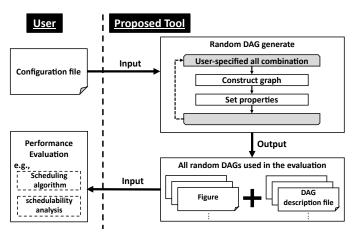


Fig. 1. System model.

TABLE I DAG NOTATIONS

Symbols	Descriptions
\overline{G}	DAG
V	Set of all nodes of G
V	Total number of nodes of G
E	Set of all edges of G
$ au_i$	<i>i</i> -th node
C_{i}	Worst-case execution time (WCET) of $ au_i$
$e_{i,j}$	Edge between $ au_i$ and $ au_j$
$comm_{i,j}$	Communication time of $e_{i,j}$
CCR	communication-to-computation ratio (CCR) of G
D	End-to-end deadline
$ au_i^{tm}$	<i>i</i> -th timer-driven node
ϕ_i	Offset of $ au_i^{tm}$
T_i	Execution period of τ_i^{tm}
d_i	Relative deadline of τ_i^{tm}
u_i	Utilization of $ au_i^{tm}$
\overline{U}	Total utilization of G
Γ_i	<i>i</i> -th chain
$ \Gamma $	Total number of chains of G
C_{Γ_i}	WCET of Γ_i
u_{Γ_i}	Utilization of Γ_i
T_{Γ_i}	Period of the head timer-driven node of Γ_i

studies. Section V compares RD-Gen with existing random DAG generation methods. Section VI discusses related work. Finally, Section VII presents the conclusions and future work.

II. SYSTEM MODEL

This section presents the system model, as shown in Fig. 1. Section II-A describes a basic single-rate DAG. Section II-B explains a multi-rate DAG. The DAG notation used in this paper is given in Table I.

A. Single-rate DAG

Single-rate DAGs are DAGs with either a single entry node or, all entry nodes entering at the same time; such DAGs can be used in real-time applications [30], cyber-physical systems [2], and cloud computing [3]. Here, the entry node represents

the input to the system (e.g., a sensor event or a command from the user), and the exit node represents the final output from the system.

A DAG consists of a node set and an edge set, denoted G = (V, E). Nodes represent tasks in the system, and edges represent communication and dependencies between nodes or priority constraints. V is the set of all nodes, expressed as $V = \{\tau_1, ..., \tau_{|V|}\}$, where |V| is the total number of nodes. Each node has a worst-case execution time (WCET), and the WCET of τ_i is denoted as C_i . E is the set of all edges, where each edge $e_{i,j} \in E$ represents communication between τ_i and τ_i and a priority constraint. When $e_{i,j}$ exists in the DAG, τ_i cannot be executed until τ_i has completed its execution and the output of τ_i has arrived. If the communication time is given as an assumption, the communication time at $e_{i,j}$ is denoted as $comm_{i,j}$. The ratio of the sum of the communication times of all edges to the sum of the execution times of all nodes is called the communication-to-computation ratio (CCR) and is defined by Eq. (1).

$$CCR = \frac{\sum\limits_{e_{i,j} \in E} comm_{i,j}}{\sum\limits_{\tau_i \in V} C_i}$$
 (1)

An end-to-end deadline D is set at the exit node when it is necessary to guarantee the safety of hard real-time systems [31] or the cloud computing quality of service [32].

B. Multi-rate DAG

A multi-rate DAG is a DAG that contains multiple nodes that are triggered with different periods. Here, the definitions of nodes and edges in a multi-rate DAG are the same as those given in Section II-A. Multi-rate DAGs can be broadly classified into two categories: (i) DAGs in which all nodes are timer driven, such as in automotive systems [8], [14], and (ii) DAGs that combine a chain consisting of timer-driven nodes and a sequence of linked event-driven nodes, such as in self-driving systems [10], [33].

1) Multi-rate DAG consisting of only timer-driven nodes: Each timer-driven node in such multi-rate DAGs is denoted by τ_i^{tm} , and τ_i^{tm} is characterized by the tuple (ϕ_i, Ci, Ti, di) , where ϕ_i, T_i , and d_i represent the offset, execution period, and relative deadline, respectively. For DAGs that consider timer-driven nodes, every timer-driven node has a relative deadline of T_i time units indicating that every job of τ_i^{tm} has an absolute deadline at T_i time units after its release [25], [27]. Such a time constraint is called an implicit deadline.

The utilization of τ_i^{tm} is denoted as u_i , such that $u_i = C_i/T_i$. The total utilization U of a DAG consisting only of timer-driven nodes is defined by Eq. (2).

$$U = \sum_{\tau^{tm} \in V} u_i \tag{2}$$

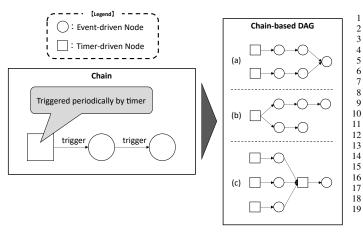


Fig. 2. Multi-rate DAGs consisting of multiple chains.

2) Chain-based multi-rate DAG: In the latest multi-rate applications, DAGs consisting of the multiple chains shown in Fig. 2 are considered. Each chain Γ_i is denoted as $\Gamma_i = \{\tau_i^{tm}, \tau_k, ..., \tau_{|\Gamma_i|}\}$, where $|\Gamma_i|$ is the number of nodes that compose Γ_i . The head τ_i^{tm} in the chain is triggered periodically, and subsequent event-driven nodes are triggered by their direct predecessors. This definition is the same as that used in existing studies [33], [34]. The WCET of Γ_i is denoted by C_{Γ_i} and defined in Eq. (3).

$$C_{\Gamma_i} = \sum_{\tau_i \in \Gamma} C_i \tag{3}$$

Because the chain is executed in a manner dependent on the period of the head timer-driven node, the utilization of the chain u_{Γ_i} is defined by Eq. (4).

$$u_{\Gamma_i} = \frac{C_{\Gamma_i}}{T_{\Gamma_i}} \tag{4}$$

Here, T_{Γ_i} is the period of the head timer-driven node τ_i^{tm} of the chain. The total utilization of the chain-based multi-rate DAG is defined by Eq. (5).

$$U = \sum_{\Gamma_i \in V} u_{\Gamma_i} \tag{5}$$

Chain-based multi-rate DAGs primarily exist in robot operating system (ROS)-based systems [10], [35]. In a typical ROS-based system, such as a self-driving system (e.g., Autoware [36]), different sensor data are processed and merged by multiple chains to output the final command. When modeling ROS-based systems as DAGs, it is necessary to consider DAGs in which multiple chains merge at the exit nodes ((a) in Fig. 2), where the chain branches ((b) in Fig. 2), and those in which multiple chains are vertically linked ((c) in Fig. 2). RD-Gen can generate all of these DAG types by using various parameters.

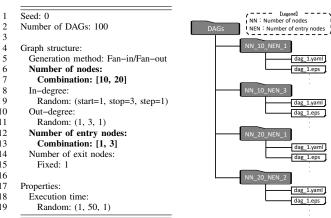


Fig. 3. Example of generating all combinations.

III. DESIGN AND IMPLEMENTATION

This section describes the design and implementation of RD-Gen. RD-Gen iterates over the construction of the graph and the set of properties for all parameter combinations described by the user in the configuration file as shown in the upper right of Fig. 1. Section III-A shows the common functionalities of RD-Gen, Section III-B describes the three graph construction methods of RD-Gen, and Section III-C explains how to set the properties for a graph.

A. Common functionality

RD-Gen receives as input a YAML file written by the user, and with a single command, RD-Gen generates an entire set of random DAGs with different parameter combinations. An example of an input parameter file in the YAML format is shown on the left side of Fig. 3. Because RD-Gen allows the user to specify seed values for pseudo-random number generation (line 1 in Fig. 3), the same YAML file can be used to easily reproduce random DAG sets generated by other researchers. *Number of DAGs* (line 2 in Fig. 3) indicates the number of DAGs randomly generated for each combination. The shape of the generated DAGs is determined by the parameters under *Graph structure*, and the properties of the nodes and edges are determined by the parameters under *Properties*.

For parameters that require a numeric input, users can specify values in the following three ways.

- Fixed: Only one value is specified, and this value is always used when a DAG is generated. This specification method can be used for DAGs in which end-to-end deadlines are considered and one exit node is always desired (line 15 in Fig. 3).
- 2) Random: When a DAG is generated, one value is selected at random from an input range. This is the most basic specification method that is used when there is no special intended value and when it is desirable to have a variety of values, as shown on lines 9, 11, and 19 in Fig. 3.

Generation methods	Parameters		Descriptions		
	Number of nodes		Number of nodes in a single DAG		
Fan-in/Fan-out	Number of entry nodes Number of exit nodes		Number of entry nodes in a single DAG		
G(n, p)			Number of exit nodes in a single DAG		
	Ensure weakly connected		When True is specified, the generated DAGs are always weakly connected		
Fan-in/Fan-out	In-degree		Number of edges input to one node		
Tun-in/Tun-oui	Out-degree		Number of edges output from one node		
G(n, p)	Probability of edge existence		Probability of an edge present between any nodes		
	Number of chains		Number of chains in a single chain-based DAG		
	Main sequence length		Length of the main sequence in a single chain		
	Number of sub sequences		Number of branches from the main sequence in a single chain		
	Vertically link chains	Number of entry nodes	Number of entry nodes in a single DAG		
Chain-based		Main sequence tail	When True is specified, the tail of the main sequence is randomly connected to the head of another chain		
		Sub sequence tail	When True is specified, the tail of the sub sequence is randomly connected to the head of another chain		
	Merge chains	Number of exit nodes	Number of exit nodes in a single DAG		
		Middle of chain	When True is specified, merge from a tail node of a random chain to a node other than the head and tail of another		
		Exit node	When True is specified, merge a head node of a random chain with the exit node.		

3) Combination: RD-Gen generates DAGs for all combinations of all lists of parameters for which Combination is specified. This specification method allows all DAG sets used in random evaluations in DAG studies to be generated in a single command execution. Therefore, RD-Gen can significantly reduce the time and effort required by researchers to generate DAGs. For example, when Combination of Number of nodes is [10, 20] and Combination of Number of entry nodes is [1, 3], RD-Gen generates 100 DAGs for all combinations, as shown on the right side of Fig. 3.

When specifying ranges for *Random* and *Combination* in RD-Gen, the following two intuitive descriptions are possible.

- 1) List format: The user specifies their choices in a list (array) format (such as [10, 20] on line 7 in Fig. 3).
- 2) *Tuple format*: The user specifies a *start* value, a *stop* value, and a *step* (such as (*start=1*, *stop=3*, *step=1*) on line 9 in Fig. 3). The tuple format is internally expanded into a list format of values that are sequentially added from the *start* value to the *stop* value at an interval of size *step* (e.g., (*start=1*, *stop=3*, *step=1*) expands to a list of [1, 2, 3]). Here, "*start=*", "*stop=*", and "*step=*" are optional (line 11 in Fig. 3).

RD-Gen can output DAG description files in the YAML, JSON, XML, and DOT formats, and DAG image files in the EPS, PDF, SVG, and PNG formats.

B. Graph construction

RD-Gen extends the Fan-in/Fan-out [15] and G(n, p) [16] methods widely used in the scheduling field by DAG researchers. In addition, RD-Gen provides a Chain-based method for generating state-of-the-art chain-based multi-rate DAGs as shown in Fig. 2. The parameters that can be specified in $Graph\ structure$ of RD-Gen are listed in Table II.

1) Fan-in/Fan-out method: Fan-in/Fan-out is a random DAG generation method proposed by Dick et al. that is provided by TGFF [15]. The Fan-in/Fan-out method can specify the in-degree (i.e., the number of edges to be input) and

TABLE III
HELPER FUNCTIONS IN THE ALGORITHMS

Symbols	Descriptions		
Random $(a, b) \mid a, b \in \mathbb{N}$	Function that returns a random natural		
$Random(a, b) \mid a, b \in \mathbb{N}$	number between a and b		
$In(au_i)$	Function that returns the in-degree of $ au_i$		
$\operatorname{Out}(au_i)$	Function that returns the out-degree of $ au_i$		
	Function that returns a linear sequence of nodes		
Sequence($\{\tau_i,, \tau_k\}$)	connected by edges consisting of input nodes.		
sequence $(\{\tau_i,, \tau_k\})$	For example, when the input is $\{\tau_1, \tau_2, \tau_3\}$,		
	it returns $\{\tau_1, e_{1,2}, \tau_2, e_{2,3}, \tau_3\}$.		

out-degree (i.e., the number of edges to be output) ranges for a single node; a DAG is then generated in which all nodes satisfy this condition. The *Fan-in/Fan-out* method extends the graph by randomly repeating the Fan-in and Fan-out phases. However, the original *Fan-in/Fan-out* method cannot completely satisfy the requirements of researchers using DAG.

The DAG generated by the *Fan-in/Fan-out* method has just one entry node, while in-vehicle and self-driving systems have multiple sensors [14], [21] and require a DAG with multiple entry nodes. Even though the latest version of TGFF allows multiple entry nodes, it may generate DAGs that are not weakly connected. It is also necessary to be able to specify the number of exit nodes because many studies consider DAGs with a single exit node [25], [32]. In addition, because scheduling and allocation methods using DAGs affect the performance depending on the number of nodes in a DAG, evaluations are performed with various changes in the number of nodes [2], [6]. However, TGFF does not allow a user to completely control the number of nodes to be generated in a DAG.

To meet these requirements, RD-Gen allows the specification of the number of nodes, entry nodes, and exit nodes in a single DAG and ensures that the DAG is weakly connected. The *Fan-in/Fan-out* method procedure provided by RD-Gen is shown in Algorithm 1, and the helper functions used in the algorithms are listed in Table III. First, the DAG is initialized with a specified number of entry nodes (line 1 in

Algorithm 1: Fan-in/Fan-out method in RD-Gen

```
Input: n \leftarrow Number of nodes
    nen \leftarrow Number of entry nodes
    nex \leftarrow Number\ of\ exit\ nodes
    id \leftarrow In\text{-}degree
    od \leftarrow Out\text{-}degree
    Output: A DAG that satisfies user-specified parameters
 1 Initialize G \leftarrow (V, E), with V \leftarrow \{\tau_1, ..., \tau_{nen}\} and E \leftarrow \{\emptyset\}
   while |V| \neq n - nex do
 2
 3
          if Random(0, 1) = 1 then
                 /* Fan-in phase
                S \leftarrow \text{Set of nodes in } V \text{ whose out-degree is less than or }
                  equal to od
                 r \leftarrow \text{Random}(1, id)
 5
                T \leftarrow \text{Randomly choose } r \text{ nodes from } S
 6
                V \leftarrow V \cup \{\tau_{|V|+1}\}
 7
                foreach \tau_i \in T do
 8
                      E \leftarrow E \cup \{e_{i,|V|}\}
                end
10
          end
11
          else
12
                 /* Fan-out phase
                \tau_i \leftarrow The node with the largest difference between od and its
13
                  out-degree in V
                diff \leftarrow od - Out(\tau_i)
14
                r \leftarrow \text{Random}(1, diff)
15
16
                foreach k \in |V|+1, ... |V|+m+1 do
                      V \leftarrow V \cup \{\tau_k\}
17
                      E \leftarrow E \cup \{e_{i,k}\}
18
19
                end
          end
20
21
          if |V| > n - nex then
                Initialize G \leftarrow (V, E), with V \leftarrow \{\tau_1, ..., \tau_{nen}\} and E \leftarrow
22
23
          end
24
   end
   G \leftarrow \text{add\_exit\_nodes}(G, nex)
25
   G \leftarrow \text{weakly\_connect}(G)
   return G
```

Algorithm 2: $add_exit_nodes(G, nex)$

```
Input: G: A DAG, nex: Number of exit nodes
Output: DAG with exit nodes added

1 NEX \leftarrow \{\tau_{|V|+1}, ..., \tau_{|V|+nex+1}\}
2 CEX \leftarrow Set of current exit nodes of G
3 while \exists n \in NEX, In(\tau_n) = 0, or \exists c \in CEX, out(\tau_c) = 0 do

4 \tau_i \leftarrow The node with the smallest out-degree in CEX
5 \tau_j \leftarrow The node with the smallest in-degree in NEX
6 E \leftarrow E \cup \{e_{i,j}\}
7 end
8 V \leftarrow V \cup NEX
9 return G
```

Algorithm 3: weakly connect(G)

```
Input: G: A DAG
  Output: A weakly connected DAG
  while G is not weakly connected do
       MC \leftarrow Weakly connected component with the maximum number
2
         of nodes in G
        RC \leftarrow Randomly selected a weakly connected component other
         than MC
       \tau_i \leftarrow \text{Randomly choose one of the exit nodes in } RC
4
          \leftarrow Randomly choose one node other than entry nodes in MC
5
          \leftarrow E \cup \{e_{i,j}\}
6
       Ĕ
7 end
8 return G
```

Algorithm 4: G(n, p) method in RD-Gen

```
Input: n \leftarrow Number of nodes
    p \leftarrow probability of edge existence
    nen \leftarrow Number of entry nodes
    nex \leftarrow Number\ of\ exit\ nodes
    Output: A DAG that satisfies user-specified parameters
 1 \quad n \leftarrow n - nen - nex
 2 Initialize G \leftarrow (V, E), with V \leftarrow \{\tau_1, ..., \tau_n\} and E \leftarrow \{\emptyset\}
    foreach i \in 1, ..., |V| do
 3
          foreach j \in 1, ..., |V| do
                if Random(0, 1) = 1 and i < j then
                    E \leftarrow E \cup \{e_{i,j}\}
 6
 7
                end
          end
 8
 9
   end
    /* Add entry nodes
\textit{10} \ NEN \leftarrow \{\tau_{|V|+1}, ..., \tau_{|V|+nen+1}\}
11 CEN \leftarrow Set of current entry nodes of G
12 while \exists n \in NEN, Out(\tau_n) = 0, or \exists c \in CEN, in(\tau_c) = 0 do
          \tau_i \leftarrow The node with the smallest in-degree in CEN
13
          \tau_j \leftarrow The node with the smallest out-degree in NEN E \leftarrow E \cup \{e_{i,j}\}
14
15
16 end
17 V \leftarrow V \cup NEN
18 G \leftarrow \text{add\_exit\_nodes}(G, nex)
19 G \leftarrow \text{weakly\_connect}(G)
20 return G
```

Algorithm 1). RD-Gen then iteratively expands the graph until the number of nodes fits within a user-specified value (line 2 in Algorithm 1). Next, the graph is randomly extended in the Fan-in phase or Fan-out phase (lines 3-20 in Algorithm 1). If the number of nodes exceeds the user-specified value, the DAG is initialized (lines 21-23 in Algorithm 1). After the loop ends, the *add_exit_nodes* function is called and nodes with no output edges are merged into a user-specified number of exit nodes with the minimum number of edges (Algorithm 2). Finally, the *weakly_connect* function is called to add edges until the DAG is weakly connected (Algorithm 3). With these extensions, the *Fan-in/Fan-out* method implemented by RD-Gen allows the user to fully control the number of nodes, entry nodes, and exit nodes in a single DAG and guarantees that the DAG is weakly connected.

2) G(n, p) method: G(n, p) is a well-known graph construction method proposed by Paul Erdős, and Alfréd Rényi et al. [16] that constructs a graph according to the number of nodes and the probability of an edge existing between any two nodes. Many of the random DAG sets used in evaluations in the latest studies using DAGs have been constructed based on the G(n, p) method [13], [28], [29], [37]. However, there are problems with the original G(n, p) method, such as the possibility of a cycle in a graph and the inability to guarantee that the graph is weakly connected.

Accordingly, in RD-Gen, the G(n, p) method is extended for direct use in the evaluation of DAG studies. The procedure for the G(n, p) method in RD-Gen is shown in Algorithm 4 First, nodes other than the entry and exit nodes are added to the graph (lines 1 and 2 in Algorithm 4). Then, for any two nodes, the addition of an edge is tried with a probability of 1 in 2. Here, if the index of the destination node is smaller than

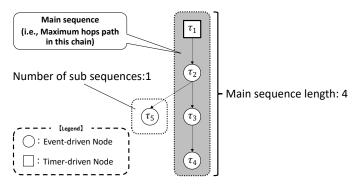


Fig. 4. Example of chain parameters.

Algorithm 5: Chain-based method in RD-Gen

```
Input: nc \leftarrow Number \ of \ chains
    m\hat{l} \leftarrow \textit{Main sequence length}
   ns \leftarrow \textit{Number of sub sequences}
    nen \leftarrow Number of entry nodes
    nex \leftarrow Number \ of \ exit \ nodes
   Output: A DAG that satisfies user-specified parameters
 1 Initialize G \leftarrow (V, E), with V \leftarrow \{\emptyset\} and E \leftarrow \{\emptyset\}
     /* Construct each chain
2 foreach i \in 1, ..., nc do
         main \leftarrow \text{Sequence}(\{\tau_{|V|+1}, ..., \tau_{|V|+ml+1}\})
3
4
         G \leftarrow main
5
         foreach j \in 1, ..., ns do
               r \leftarrow \text{Random}(|V|+1, |V|+ml)
6
                sl \leftarrow \text{Random}(1, ml - r)
               sub \leftarrow \mathsf{Sequence}(\{\tau_{|V|+1}, ..., \tau_{|V|+sl+1}\})
8
               E \leftarrow E \cup \{e_{r,|V|+1}\}
               G \leftarrow sub
10
11
         end
12 end
     /* Vertically link chains
   while Number of entry nodes of G \neq nen do
13
         \tau_i \leftarrow \text{Randomly choose one of the exit nodes in } V
14
             ← Randomly choose one node from the head of the chains
15
             \leftarrow E \cup \{e_{i,j}\}
16
17 end
    /* Merge chains
   while Number of exit nodes of G \neq nex do
18
         \tau_i \leftarrow \text{Randomly choose one of the exit nodes in } V
         	au_j \leftarrow 	ext{Randomly choose one node from other than head of the}
20
           chains
          E \leftarrow E \cup \{e_{i,j}\}
21
22 end
   return G
```

the index of the source node, the edge addition is canceled (line 5 in Algorithm 4). This condition results in no cycles being present in the graph [28], [38]. After the loop ends, a user-specified number of entry nodes and a node with an indegree of 0 are connected with the smallest edge. Finally, the add_exit_nodes and weakly_connect functions are called in the same manner as in the Fan-in/Fan-out method.

3) Chain-based method: The Chain-based method designed is a new random DAG generation method to meet the requirements of modern chain-based multi-rate DAGs. In the Chain-based method, multiple chains are combined to construct a single DAG. The parameters that determine the shape of a single chain are explained here using Fig. 4. The user can determine the shape of a given chain by specifying

Algorithm 6: Set utilization

```
Input: G \leftarrow \mathsf{DAG}
     pt \leftarrow Peridoic type
     U \leftarrow Total\ utilization
     p \leftarrow Period
     Output: A DAG that satisfies user-specified parameters
 1 if p\hat{t} = "All" then
           u[1,...,|V|] \leftarrow \text{UUniFast}(|V|,U) \text{ [39]}
           \text{for each } i \in {1,...,|V|} \text{ do}
 3
 4
                 T_i \leftarrow p
                  C_i \leftarrow u[i] \times T_i
 5
           end
 6
 7
    end
 s if pt = "Chain" then
           u[1,...,|\Gamma|] \leftarrow \text{UUniFast}(|\Gamma|,U)
 9
           foreach i \in 1, ..., |\Gamma| do
10
                 T_{\Gamma_i} \leftarrow p
11
12
                  C_{\Gamma_i} \leftarrow u[i] \times T_{\Gamma_i}
                 q \leftarrow \text{Number of nodes in } \Gamma_i
13
14
                  c[i,...,i+g] \leftarrow Grouping randomly into g with total equal to
                   C_{\Gamma_i}
                 foreach j \in i, ..., i + g do
15
16
                        C_j \leftarrow c[j]
                 end
17
18
           end
    end
19
20 return G
```

the main sequence length and the Number of subsequences. The main sequence is the path that has the largest number of hops from the head timer-driven node to the tail event-driven node in a single chain ($\{\tau_1, \tau_2, \tau_3, \tau_4\}$ in Fig. 4), and subsequences are straight lines of nodes branching off from the main sequence ($\{\tau_5\}$ in Fig. 4).

In the *Chain-based* method, two different ways of connecting chains can be specified. The first is to link chains vertically, i.e., the event-driven node at the end of each chain is randomly connected to the timer-driven node at the head of another chain (e.g., (c) in Fig. 2). The second method is to merge multiple chains into one specific node. The *Chain-based* method allows nodes other than the head node of the chain to be specified as integration nodes, and the event-driven node at the tail of each chain is merged into a randomly selected integration node (e.g., (a) in Fig. 2).

The procedure for the *Chain-based* method is shown in Algorithm 5. First, chains are constructed for the *Number of chains* specified by the user (lines 2-12 in Algorithm 5). In the construction of each chain, the main sequence is generated (lines 3 and 4 in Algorithm 5). Subsequences branch from random nodes other than the tail of the main sequence and are coordinated to not exceed the length of the main sequence (lines 6-10 in Algorithm 5). After all the chains are constructed, the chains are linked vertically until the number of entry nodes in the DAG is equal to the user specification (lines 13-17 in Algorithm 5). Finally, the chain is randomly merged until the number of exit nodes in the DAG matches the specified number (lines 18-22 in Algorithm 5). Consequently, the *Chain-based* method can flexibly construct DAGs consisting of multiple chains.

TABLE IV ${\rm ALL\ PARAMETERS\ OF\ } Properties\ {\rm IN\ } {\rm RD\text{-}GEN}$

Parameters			Descriptions			
Execution time			Execution time of nodes			
Communication time			Communication time of edges			
CCR			CCR of a single DAG			
End-to-end deadline	Ratio of	deadline to critical path	Ratio of the end-to-end deadline to the critical path of DAG			
	Periodic	type	Specify a group of nodes to be timer driven			
	Period		Period of timer-driven nodes			
	Entry no	ode period	Period of entry timer-driven nodes			
Multi-rate	Exit noo	le period	Period of exit timer-driven nodes			
	Offset		Offset of timer-driven nodes			
	Total utilization		Total utilization of a single DAG			
	Maximum utilization		Maximum utilization of one timer-driven node			
Additional properties	Node properties		User-defined numeric parameters with any name for nodes			
Additional properties	Edge properties		User-defined numeric parameters with any name for edges			
	DAG	YAML				
		JSON	Outputs DAG description files in the format specified by True			
		XML	Outputs Divo description mes in the format specified by True			
		DOT				
Output formats	Figure	Draw legend	When True is specified, a legend is drawn on the output DAG figure			
		PNG				
		SVG	Outputs DAG figures in the format specified by True			
		EPS				
		PDF				

C. Property setting

RD-Gen automatically sets the typical properties that characterize the nodes and edges of a DAG to meet the user requirements. All of the properties that can be set automatically in RD-Gen are listed in Table IV. RD-Gen can generate (i) single-rate DAGs, (ii) multi-rate DAGs consisting of only timer-driven nodes, and (iii) chain-based multi-rate DAGs depending on the *Periodic type* parameter. If the *Periodic type* is not specified, DAGs of type (i) are generated; if the *Periodic type* is "All" DAGs of type (ii) are generated; and if the *Periodic type* is "Chain" DAGs of type (iii) are generated.

In RD-Gen, all properties described in Section II can be specified using "Fixed", "Random" or "Combination". Even though existing random DAG generation tools such as TGFF and GGen provide the functionality to randomly assign properties, they do not allow the user to control the values calculated by multiple properties such as CCR and the total utilization. Because CCR and the total utilization have a significant impact on the performance of the scheduling algorithms and the analysis methods, evaluations are performed by varying these values [2], [13], [38], [40]. Therefore, RD-Gen also supports the specification of such complex property values.

An example of property setting based on the total utilization in RD-Gen is shown in Algorithm 6. RD-Gen uses the UUniFast method [39] to set the utilization to each node in a uniform distribution (lines 2 and 9 in Algorithm 6). If the *Periodic type* is "All", the utilization for each node is determined based on the *Total utilization* specified by the user, and the period and execution time are set to satisfy this utilization (lines 1-7 in Algorithm 6). Here, when *Maximum utilization* is specified, RD-Gen allocates utilization to each node not to exceed *Maximum utilization*. If the *Periodic type* is "Chain", the utilization of each chain is determined based on the *Total utilization*, and the period of each chain and the sum

of the execution time are calculated accordingly (lines 8-12 in Algorithm 6). The total execution time is randomly divided by the number of nodes in the chain to set the execution time for each node (lines 13-17 in Algorithm 6). In this way, RD-Gen randomly and automatically sets properties according to complex parameters specified by the user.

RD-Gen provides other parameters that allow for setting the ratio of the end-to-end deadlines to the critical path length and the period of entry and exit nodes. In addition, users can define unique parameters for simple properties that randomly assign numerical values to nodes or edges.

IV. CASE STUDY

This section illustrates that RD-Gen can generate the random DAG sets used in the evaluation of existing studies based on DAGs. Case studies of a single-rate DAG, a multi-rate DAG consisting of only timer-driven nodes, and a chain-based multi-rate DAG are shown, respectively.

A. Case study 1

Since CCR changes the nature of DAGs and affects the performance of scheduling algorithms, existing studies of single-rate DAGs have used random DAG sets with different CCR values in their evaluations. An example of the generation in RD-Gen of a random DAG set with varying CCR, number of nodes, execution time, and communication time using the *Fan-in/Fan-out* method as used in existing studies [41]–[43] is shown in Fig. 5. Because the *Combination* is specified for the *Number of nodes* (lines 6-7 in Fig. 5) and *CCR* (lines 21-22 in Fig. 5), RD-Gen generates 100 random DAGs each for all combinations of these parameters (i.e., $\{10, 20, 30, ..., 1, 000\} \times \{0.1, 0.2, 0.5, 1.0, 2.0, 5.0, 10.0\}$). For each DAG, after graph construction, execution time is randomly assigned to each node in the range of 1 to 30 (lines 19-20 in Fig. 5). Then, the total communication time is calculated

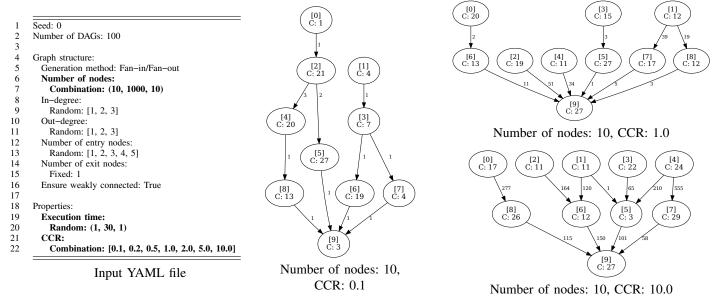


Fig. 5. Example of the generation a random DAG set with varying CCR, number of nodes, execution time, and communication time using the Fan-in/Fan-out method.

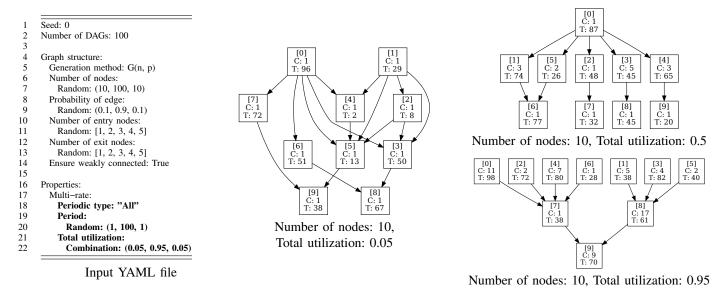


Fig. 6. Example of generating multi-rate DAGs consisting of only timer-driven nodes of various total utilization.

from the CCR and the sum of execution time based on Eq. (1), and the total communication time is randomly distributed to each edge. Thus, the user can generate all DAG sets used in the evaluation with a single command, without adjusting the number of nodes and CCR.

B. Case study 2

Random DAG sets with different total utilization are used in the evaluation of most studies considering multi-rate DAGs. An example of generating a multi-rate DAG with a random period and execution time for each node based on the total utilization of the DAG, as used in existing studies [13], [23], [24], is shown in Fig. 6. Here, timer-driven nodes are drawn as squares. Because *Combination* is specified for *Total utilization*

(lines 21-22 in Fig. 6), RD-Gen generates 100 DAGs each with total utilization from 5% to 95% in 5% increments. For each DAG, the utilization of each node is determined by the UUniFast method to satisfy the total utilization, and the execution time is calculated from the utilization and a randomly selected period from a uniform distribution ranging from 1 to 100.

C. Case study 3

In the evaluation of chain-based multi-rate DAGs, a random DAG set is used with varying total utilization for the entire system and each chain. An example of generating a random DAG set, as used in existing studies [10], [33], in which the utilization of each chain is determined from the total utilization

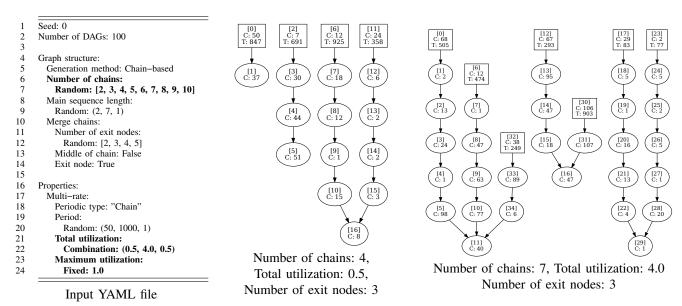


Fig. 7. Example of generating chain-based multi-rate DAGs of various total utilization.

TABLE V PROCESSES THAT MUST BE IMPLEMENTED BY THE USER AND LINES OF CODE

Symbols	Processes	Lines of code
(a)	Loop with different parameter values	3
(b)	Separate DAGs generated for each parameter into directories	5
(c)	Load .tgff file as DAG	40
(d)	Construct a DAG using G(n, p) method	10
(e)	Construct a DAG using chain-based method	80
(f)	Merge the exit nodes to a specific number	10
(g)	Add edges to be weakly connected	15
(h)	Set the properties randomly within a specific range	10
(i)	Set the execution time for each node and the communication time for each edge based on the CCR	25
(j)	Set the utilization for a multi-rate DAG consisting of only timer-driven nodes (lines 1-7 in Algorithm 6)	30
(k)	Set the utilization for a chain-based multi-rate DAG (lines 8-19 in Algorithm 6)	50

of the entire system, and the chain period and execution time of each node is randomly assigned to satisfy the utilization of the chain is shown in Fig. 7. RD-Gen randomly sets the utilization of each chain using the UUniFast method to achieve a determined total utilization. Here, since *Maximum utilization* is specified as 1.0 (line 24 in Fig. 7), RD-Gen never generates a chain with a utilization greater than 1.0. The period of each chain is randomly set in the range of 50 to 1,000, and the execution time of each node is calculated based on the utilization and period (lines 10-18 in Algorithm 6).

V. EVALUATION

This section shows that compared to existing random DAG generation tools, RD-Gen can generate a random DAG set with fewer lines of code and no unique user implementation. Table V shows the processes that must be implemented by the user and the number of lines of code. Here, the process in Table V is written in Python and shell scripts, and the Python NetworkX library is used for processes related to DAGs.

TABLE VI EFFORT REQUIRED TO GENERATE RANDOM DAG SETS FOR Case Study 1 [41]–[43]

	Lines of description for tool options	Lines of code
TGFF [15]	6	$(a) \times 2 + (b) + (c) + (f) + (g) + (i)$ = 101
GGen [16]	3	$(a) \times 2 + (b) + (g) + (i)$ = 51
RD-Gen	20	0

The amount of description required to generate a random DAG set for the case studies in Section IV-A is shown in Table VI. In this case, while TGFF and GGen require the user to implement more than 50 lines, RD-Gen can generate all random DAG sets using only the tool's functionality. Since TGFF and GGen do not have the functionality to generate DAGs with different parameter settings at once, the user must execute the generation command many times while changing the parameters, using shell scripts or other means. In contrast, RD-Gen generates DAGs with all parameter combinations by specifying Combination as a parameter and automatically divides them into directories. Although TGFF uses the Fanin/Fan-out method to construct DAGs, the .tgff files output by TGFF are in a proprietary format and require users to implement a reading process. Furthermore, since TGFF and GGen do not guarantee that the graph is weakly connected, users may unintentionally use a non weakly connected DAG for evaluation. TGFF and GGen cannot specify parameters such as CCR that are calculated from multiple properties. Therefore, the CCR must be adjusted for the DAG set after it is generated. In RD-Gen, users can automate all of the above processes by inputting the YAML file as shown on the left side of Fig. 5.

The effort required to generate a random DAG set for

TABLE VII
EFFORT REQUIRED TO GENERATE RANDOM DAG SETS
FOR Case Study 2 [13], [23], [24]

	Lines of description for tool options	Lines of code
TGFF [15]	-	(a) + (b) + (d) + (g) + (j) = 63
GGen [16]	2	(a) + (b) + (g) + (j) = 53
RD-Gen	20	0

TABLE VIII
EFFORT REQUIRED TO GENERATE RANDOM DAG SETS
FOR Case Study 3 [10], [33]

	Lines of description for tool options	Lines of code
TGFF [15]	-	(a) + (b) + (d) + (k) = 138
GGen [16]	-	(a) + (b) + (d) + (k) = 138
RD-Gen	22	0

the case studies in Section IV-B is shown in Table VII. TGFF does not support the G(n, p) method and multi-rate DAG. GGen allows properties to be randomly set to nodes or edges within a user-specified range after graph construction. However, GGen does not allow users to specify the total utilization used in the evaluation of most studies that consider multi-rate DAGs. Therefore, the user has to go through the trouble of implementing such as lines 1-7 in Algorithm 6. In contrast, RD-Gen uses the UUniFast method to randomly and automatically set the period and execution time for each node to meet the specified total utilization.

The amount of description to generate a chain-based random DAG set, as in the case study in Section IV-C, is shown in Table VIII. Chain-based multi-rate DAGs cannot be generated by TGFF and GGen because they are considered in state-of-the-art self-driving systems research. RD-Gen can generate a batch of random chain-based DAG sets with different total utilization without any user implementation by inputting a YAML file as shown in Fig. 7.

The evaluation results demonstrated that RD-Gen can generate the random DAG sets used in the evaluation with only an intuitive YAML file description, regardless of the single rate multi rate. Therefore, RD-Gen is a flexible evaluation platform that can be adapted to various requirements and provides reliability and reproducibility for the latest DAG studies.

VI. RELATED WORK

This section describes existing random DAG generation tools and existing studies using random DAGs and compares them with RD-Gen. Table IX shows a comparison of RD-Gen with existing methods.

A. Random DAG generation tools

Random DAG generation tools provide reliability and reproducibility for evaluations of scheduling and latency analysis studies. TGFF [15] is the first tool proposed for this purpose

TABLE IX RD-GEN VS EXISTING METHODS

	RMD	RPU	RDT	RTE	OCG
DAGEN [44]			✓		
MRTG [45]			✓		
TGFF [15]			✓		
GGen [16]	✓		✓		
Kordon et al. [8]	✓				
Yang et al. [27]	✓	✓			
Gunzel et al. [23]	✓	✓			
Ueter et al. [24]	✓	✓			
Dong et al. [29]	✓	✓			
He et al. [13]	✓	✓			
Voronov et al. [28]	✓	✓			
Verucchi et al. [14]	✓	✓			
Klaus et al. [12]	✓	✓			
Tang et al. [33]	✓	✓		✓	
Choi et al. [10]	✓	✓		✓	
RD-Gen	✓	✓	✓	✓	✓

RMD: Random generation of multi-rate DAGs

RPU: Random property settings based on total utilization

RDT: Random DAG generation tool

RTE: Random generation of chain-based multi-rate DAGs OCG: One-command batch generation of random DAGs

and has been used to evaluate recent studies [46]–[48]. TGFF determines the shape of a DAG primarily by specifying either or both the maximum and minimum input degree and maximum (first) output degree for a single node (the *Fan-in/Fan-out* method). TGFF can quickly generate many DAGs, and the task set can be easily reproduced by other researchers by inputting the same parameters. However, TGFF was released in 1998 and has many problems, such as its output format (.tgff), which is difficult to handle, and it cannot generate the multi-rate DAGs.

GGen [16] is a unified implementation of classical task graph generation methods used in the scheduling domain. GGen allows the user to add properties such as the execution period and the communication time to nodes and edges after generating DAGs using a user-specified generation method. However, GGen does not allow constraints to be specified between different properties, such as implicit deadlines (i.e., the execution time of a node must not exceed its period). Therefore, users with such requirements must adjust these values themselves.

Other random DAG generation tools such as DAGEN [44] and MRTG [45] have also been proposed. DAGEN generates random workflow applications by specifying the node load balancing, edge connection probability, and workflow shape. MRTG is a random DAG generation tool with a module-based implementation for user extensibility. However, these tools are not capable of generating multi-rate DAGs. Conversely, RD-Gen can flexibly generate multi-rate DAGs of various types. In addition, RD-Gen can automatically set properties calculated by multiple values such as CCR and the total utilization.

B. Unique implementation of random DAG generation

This section presents existing studies in which the authors generated random DAG sets to evaluate their own implemented algorithms and settings. Because there are no random DAG generation tools that can generate multi-rate DAGs by specifying the total utilization, as described in Section VI-A, most studies that consider multi-rate DAGs include unique implementations.

In real-time systems such as in-vehicle systems and self-driving systems, multi-rate DAGs consisting of only timer-driven nodes are considered. Many real-time system researchers use the G(n, p) method to construct DAGs and generate random task sets with different numbers of nodes, different total utilizations, and different periods. [13], [28], [29]. While proprietary algorithms are used to construct graphs in some cases, the approach is similar in that the utilization is determined using the UUniFast method, and the WCET value of the nodes are assigned based on the utilization and the periods [23], [24], [27].

In a multi-rate DAG, as considered in automotive systems, a period is randomly assigned to each node based on the period observed in the automotive application. Verucchi et al. [14] randomly extended the automotive benchmark proposed by BOSCH in the 2015 WATERS Challenge [49] to analyze the performance of their proposed method. Verucchi et al. randomly set the task period from the values [1, 5, 10, 20, 50, 100, 200, 1, 000] in milliseconds, as found in automotive applications, for the DAG utilization. Klaus et al. [12] set the utilization, the period, and the number of nodes for each DAG node. Verucchi et al. and Klaus et al. create random task sets based on node chains consisting of only timer-driven nodes. RD-Gen can also generate chains consisting of only timer-driven nodes using the Chain-based method and by specifying the *Period type* as "All". Using the Python NetworkX library, Kordon et al. [8] randomly set the period, the number of edges per task, the release time, and the number of entry nodes for their DAG sets. RD-Gen can automatically set all combinations of the total utilization, periods, and the number of nodes as described above.

Studies of chain-based multi-rate DAGs, such as the latest ROS-based systems, have also been evaluated using random task sets. Tang et al. [33] allocated a value of the utilization to each chain based on the total utilization of the entire system and the number of chains and then assigned the utilization to the execution units in the chain. Choi et al. [10] similarly assigned a utilization value to each chain using the UUniFast method based on the total system-wide utilization. Multi-rate DAGs based on such chains can be generated flexibly with the *Chain-based* method in RD-Gen. In addition, because RD-Gen allows the utilization to be automatically set based on the chain, researchers do not need to implement this functionality on their own.

VII. CONCLUSIONS

In this paper, we proposed a random DAG generator considering multi-rate applications for reproducible scheduling evaluation called RD-Gen that can generate both single-rate DAGs and state-of-the-art multi-rate DAGs. RD-Gen extended the existing random graph generation methods for DAGs and provided a new chain-based method. RD-Gen could

automatically set complex properties such as CCR and the total utilization. Moreover, RD-Gen could support researchers by providing functions such as the batch generation of a given number of DAG sets for different parameters. Case studies indicated that RD-Gen could meet the requirements of DAG studies in a variety of problem settings. Therefore, RD-Gen can provide reliability and easy reproducibility for DAG studies. In future work, we plan to extend RD-Gen to cover additional graph generation methods and more complex properties.

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