



**The University of Lahore**  
*Faculty of Information Technology*

**Assignment Cover Letter**  
*(Individual Work)*

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SAP ID	70082385	Title of Assignment	03
Course Code	CS 11303	Due Date	26/12/2022
Course Name	Operating Systems	Submission Date	26/12/2022
Section	T		

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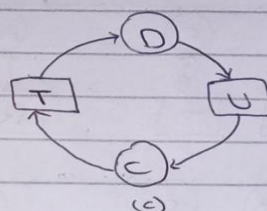
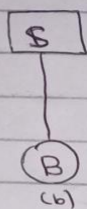
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Signature	
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## Question-1

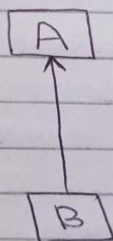
Question 1a



So, this process can complete until it gets the required resource and it can't release the already present resources before complete, so the system is in a deadlock.

This graph also helps us to make plan in order to prevent deadlock, we just have to prevent occurring of a loop in the graph.

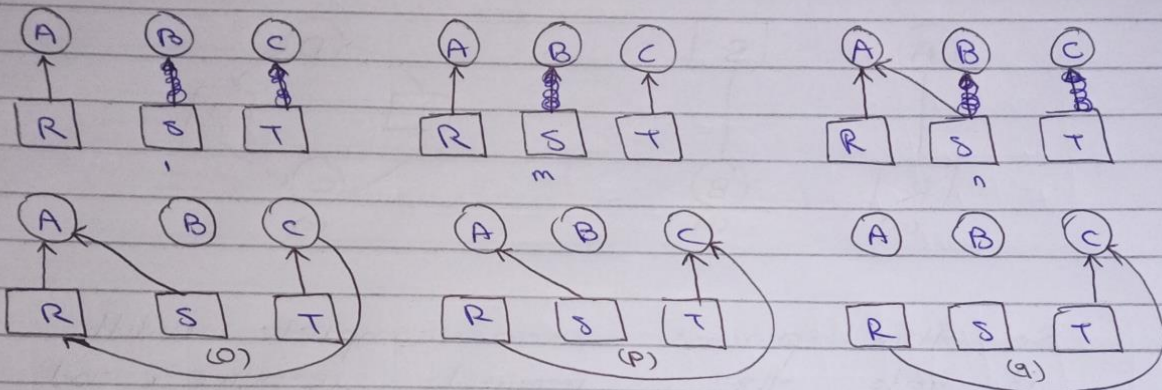
The resource graph can be illegal if one process is requesting for another process instead of a resource.



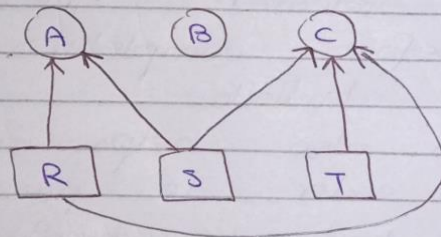
①

## Question-2

Question 2:



Step (o):



As we look C. request S instead of R but deadlock won't occur because no cycle will be formed when an arrow is added from S to C.

As we added another arrow from R to C. In here deadlock won't occur because no cycle will be formed. ~~when an arrow is added~~  
C can access S after A releases it.



### Question-3

Question 3a

Available<sub>execute</sub>  $\begin{Bmatrix} 2 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{Bmatrix}$

$P3 \Rightarrow \begin{Bmatrix} 2 & 2 & 2 & 0 \\ 1 & 2 & 1 & 0 \end{Bmatrix}$

$P2 \Rightarrow \begin{Bmatrix} 4 & 2 & 2 & 1 \\ 2 & 2 & 1 & 0 \end{Bmatrix}$

$P1 \Rightarrow \begin{Bmatrix} 4 & 2 & 2 & 1 \end{Bmatrix}$

Request

$\begin{Bmatrix} 2 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ -2 & 1 & 0 & 0 \end{Bmatrix}$

Allocation

$\begin{Bmatrix} 0 & 0 & 1 & 0 \\ 2 & 0 & 0 & 1 \\ 0 & 1 & 2 & 0 \end{Bmatrix}$

Allocation

$\begin{Bmatrix} 0 & 0 & 1 & 0 \\ 2 & 0 & 0 & 1 \\ 2 & 2 & 2 & 0 \end{Bmatrix} \Rightarrow \begin{Bmatrix} 0 & 0 & 1 & 0 \\ 2 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{Bmatrix} \Rightarrow \begin{Bmatrix} 0 & 0 & 1 & 0 \\ 3 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \end{Bmatrix}$

$\begin{Bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{Bmatrix} \Rightarrow \begin{Bmatrix} 2 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{Bmatrix}$

$\begin{Bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{Bmatrix}$

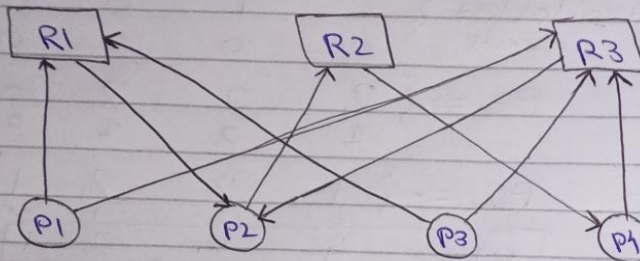
Deadlock not occurs and available are

Available  $= \{4, 2, 2, 1\}$

(8)

#### Question-4

Question 4:



Q3

In this graph P1 request to R1 and R3. Where R1 assign to P2 and P2 request to R2 and R2 assign to P4.

P2 request R1 & R3 and R3 assign to P2. P4 request to ~~R2~~ R3.

~~P3~~ ~~P4~~

In this graph P3 and P4 are waiting for R3 where R3 assign to P2.

b3

As we know if there is circle between processes and request there may be or may not deadlock occurs.

In this graph there will circle between processes and request.



# Question-5

Question 5:

	Has	May
A	1	6
B	1	5
C	2	4
D	4	7

Free: 2

=>

A	1	6
B	1	5
C	<del>4</del>	<del>4</del>
D	4	7

Free: 0

A	1	6
B	1	5
C	<del>4</del>	<del>4</del>
D	4	7

Free: 4

=>

A	1	6
B	<del>5</del>	<del>5</del>
C	<del>4</del>	<del>4</del>
D	4	7

Free: 0

A	1	6
B	<del>5</del>	<del>5</del>
C	<del>4</del>	<del>4</del>
D	4	7

Free: 5

A	<del>6</del>	<del>6</del>
B	<del>5</del>	<del>5</del>
C	<del>4</del>	<del>4</del>
D	4	7

Free: 0

A	<del>6</del>	<del>6</del>
B	<del>5</del>	<del>5</del>
C	<del>4</del>	<del>4</del>
D	4	7

Free: 6

When C complete we have 4 unit then B we have 5 unit and A we have 6 unit free.

A	<del>6</del>	<del>6</del>
B	<del>5</del>	<del>5</del>
C	<del>4</del>	<del>4</del>
D	<del>7</del>	<del>7</del>

Free: 10

So we know if D ask for one more unit. This will lead to safe state. And sequence is C → B → A → D.

# Question-6

Question 6:

	Allocation	Maximum	Available	(Max-alloc) F. Need
A	1 0 2 1 1	1 1 2 1 2	0 0 x 1 1	0 1 0 0 1
B	2 0 1 1 0	2 2 2 1 0	0 0 1 1 1	0 2 1 0 0
C	1 1 0 1 0	2 1 3 1 0		1 0 3 0 0
D	1 1 1 1 0	1 1 2 2 1		0 0 1 1 1

Only D will released its allocated resource.  
Let  $x = 1$

$$\begin{array}{rcl}
 & 00111 & \text{allow with D} \\
 D = & \underline{11110} & \\
 & 11221 & \\
 A = & \underline{10211} & \text{allow with A} \\
 & 21432 & \\
 B = & \underline{20110} & \text{allow with B} \\
 & 41542 & \\
 C = & \underline{11010} & \text{allow with C} \\
 & 52552 & 
 \end{array}$$

$$x = 1$$

It is the smallest value.

The safe sequence is

D, A, B, C