

# Introduction to Mathematical Modeling

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## Introduction to Graph Theory

A **graph** is a mathematical structure used to model pairwise relations between objects. It is defined as:

$$G = (V, E)$$

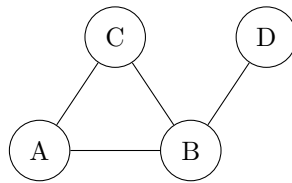
where:

- $V$  is a non-empty set of **vertices** (or **nodes**)
- $E$  is a set of **edges**, each connecting a pair of vertices

### Undirected Graph

In an **undirected graph**, edges have no direction. That is, an edge  $\{u, v\} \in E$  indicates a bi-directional relationship between vertices  $u$  and  $v$ . The edge set  $E$  is a set of **unordered** pairs of vertices.

**Example:**



Here:

$$V = \{A, B, C, D\}, \quad E = \{\{A, B\}, \{A, C\}, \{B, C\}, \{B, D\}\}$$

Degrees:

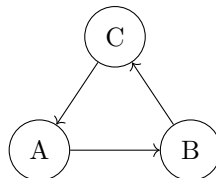
$$\deg(A) = 2, \quad \deg(B) = 3, \quad \deg(C) = 2, \quad \deg(D) = 1$$

Total degree sum:  $2 + 3 + 2 + 1 = 8 \Rightarrow \frac{8}{2} = 4$  edges

### Directed Graph

In a **directed graph** (or **digraph**), edges have a specific direction. An edge  $(u, v) \in E$  indicates a directed connection from vertex  $u$  to vertex  $v$ . Here, the edge set  $E$  is a set of **ordered** pairs.

**Example:**



Here:

$$V = \{A, B, C\}, \quad E = \{(A, B), (B, C), (C, A)\}$$

Degrees:

$$\deg^+(A) = 1, \quad \deg^-(A) = 1, \quad \deg^+(B) = 1, \quad \deg^-(B) = 1, \quad \deg^+(C) = 1, \quad \deg^-(C) = 1$$

## Vertex Degree

In an undirected graph, the **degree** of a vertex  $v$ , denoted  $\deg(v)$ , is the number of edges incident to  $v$ . If a loop exists at  $v$ , it contributes **2** to  $\deg(v)$ , since it connects to  $v$  twice.

In a directed graph:

- The **in-degree**  $\deg^-(v)$  is the number of edges coming *into*  $v$
- The **out-degree**  $\deg^+(v)$  is the number of edges going *out of*  $v$

## The Handshaking Lemma

**Statement:** In any undirected graph, the sum of the degrees of all vertices is equal to twice the number of edges.

$$\sum_{v \in V} \deg(v) = 2|E|$$

**Proof Sketch:** Each edge contributes exactly 1 to the degree of both of its endpoints. Therefore, every edge is counted twice in the sum of degrees.

**Corollary:** In every undirected graph, the number of vertices with an *odd* degree is even.

## Example (Handshaking Lemma):

Given a graph with vertex degrees:

$$\deg(A) = 1, \quad \deg(B) = 2, \quad \deg(C) = 3, \quad \deg(D) = 2, \quad \deg(E) = 2$$

Total degree sum:

$$1 + 2 + 3 + 2 + 2 = 10 \Rightarrow \frac{10}{2} = 5 \text{ edges}$$

Number of vertices with odd degree: A (1) and C (3): 2 vertices, Even

## Definition: Euler Circuit

An **Euler circuit** is a closed trail (a path that starts and ends at the same vertex) in a graph that visits every edge exactly once.

A graph has an Euler circuit **if and only if**:

- The graph is **connected**, and
- All vertices have **even degree**.

## Definition: Eulerizing a Graph

**Eulerizing** a graph means modifying it so that it contains an Euler circuit. This is done by:

- Adding edges **between existing vertices** (allowing duplicates),
- Until all vertices have **even degree**.

Note: Edges are added only between vertices that already have a path no new vertices are introduced.

## Key Rule:

The number of vertices with odd degree in any graph is always **even**. Therefore, Eulerizing involves pairing up odd-degree vertices and connecting them with duplicate edges.

## Step-by-Step Process to Eulerize a Graph

1. Identify all **odd-degree vertices**.
2. Find the **shortest path** between any pair of odd-degree vertices.
3. Add that path to the graph (duplicate the edges along the path).
4. Repeat steps 1-3 until all vertices have even degrees.

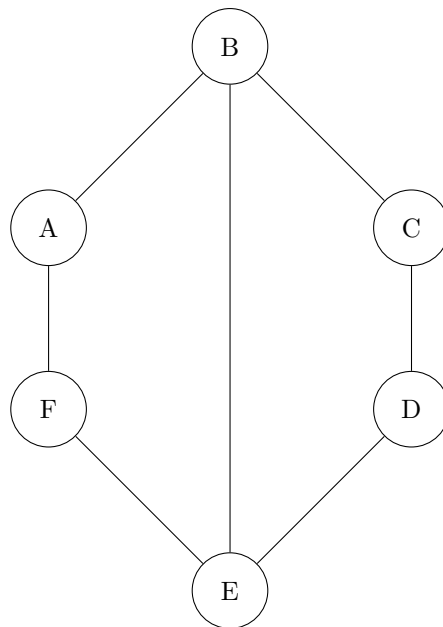
## Example Graph

**Vertices:** A, B, C, D, E, F

**Edges:**

A-B, B-C, C-D, D-E, E-F, F-A, B-E

### Initial Graph



## Step 1: Identify Odd-Degree Vertices

- $\deg(A) = 2$
- $\deg(B) = 3 \rightarrow \text{odd}$
- $\deg(C) = 2$

- $\deg(D) = 2$
- $\deg(E) = 3 \rightarrow \text{odd}$
- $\deg(F) = 2$

Odd-degree vertices: B, E

## Step 2: Find Shortest Path Between Odd-Degree Vertices

The shortest path between B and E is the direct edge **B–E** (already exists).

## Step 3: Add the Path to the Graph

Duplicate the edge B–E:

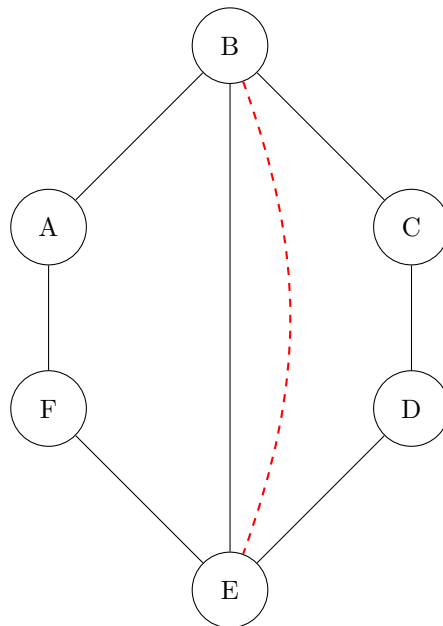
- $\deg(B): 3 \rightarrow 4$
- $\deg(E): 3 \rightarrow 4$

## Step 4: Confirm All Degrees Are Even

- All vertices now have even degree.
- Graph remains connected.

**Conclusion:** The graph is now **Eulerized** and contains an Euler circuit.

## Updated Graph with Duplicated Edge (Dashed in Red)



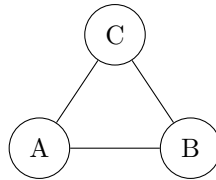
## 5. Adjacency Matrix

The **adjacency matrix** of a graph with  $n$  vertices is an  $n \times n$  matrix  $A = [a_{ij}]$ , where:

$$a_{ij} = \begin{cases} 1, & \text{if there is an edge from vertex } i \text{ to vertex } j \\ 0, & \text{otherwise} \end{cases}$$

### 5.1 Undirected Graph Example

Consider this undirected graph with vertices  $A$ ,  $B$ , and  $C$ :



Label vertices:  $A = 1$ ,  $B = 2$ ,  $C = 3$

$$A = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix}$$

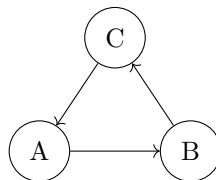
**Note:**

- The matrix is **symmetric** since the graph is undirected.
- Diagonal entries are 0 (no loops).
- Row sums (or column sums) give vertex degrees.

$$\deg(A) = 2, \quad \deg(B) = 2, \quad \deg(C) = 2.$$

### 5.2 Directed Graph Example

Now consider a directed graph:



Label:  $A = 1$ ,  $B = 2$ ,  $C = 3$

$$A = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

**Note:**

- The matrix is **not symmetric**.
- Row  $i$  shows edges **from** vertex  $i$  (outgoing)

- Column  $j$  shows edges **to** vertex  $j$  (incoming)

Out-degrees:  $\deg^+(A) = 1, \deg^+(B) = 1, \deg^+(C) = 1$

In-degrees:  $\deg^-(A) = 1, \deg^-(B) = 1, \deg^-(C) = 1$

### 5.3 Properties of Adjacency Matrices

- For an undirected graph:
  - $A$  is symmetric
  - The sum of all elements in  $A$  is  $2|E|$
- For a directed graph:
  - The sum of row  $i = \deg^+(v_i)$
  - The sum of column  $j = \deg^-(v_j)$
- For both types:
  - Diagonal entries  $a_{ii}$  count loops
  - You can compute the number of paths of length  $k$  using powers of  $A$ :  $(A^k)_{ij}$  = number of walks from  $i$  to  $j$  of length  $k$

## 1. Can a Matrix Entry Be 2?

Yes! This happens if two points are joined by more than one edge.

### Example: Simple Graph

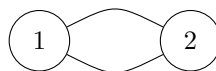
$$A = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$



Here, the 1 means there is one edge between vertices 1 and 2.

### Example: Multigraph

$$A = \begin{bmatrix} 0 & 2 \\ 2 & 0 \end{bmatrix}$$



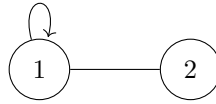
Here, the 2 means there are two edges between vertices 1 and 2.

## What If There Is a Loop?

A loop is when a vertex connects to itself. In the matrix, loops appear on the diagonal (the entries  $A_{ii}$ ).

**Example: Graph with a Loop**

$$A = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$



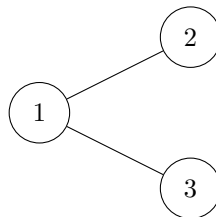
- The 1 in the top-left means vertex 1 has a loop.
- The 1 in position (1,2) means there is an edge from 1 to 2.

**Can the Matrix Be Used to Count Connections?**

Yes! The number of edges connected to a vertex is called its degree.

**Undirected Graph Example**

$$A = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \end{bmatrix}$$



- Row 1 has two 1s  $\Rightarrow$  vertex 1 has degree 2.
- Row 2 has one 1  $\Rightarrow$  vertex 2 has degree 1.
- Row 3 has one 1  $\Rightarrow$  vertex 3 has degree 1.

**Worksheet: For each graph, write the adjacency matrix.**

**Figure A**

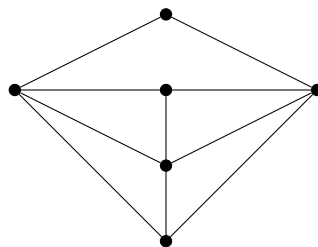


Figure B

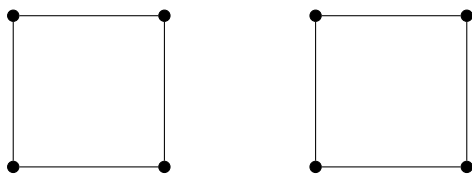


Figure C

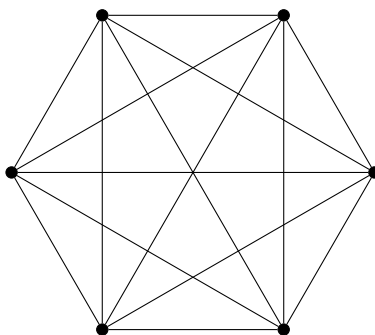


Figure D

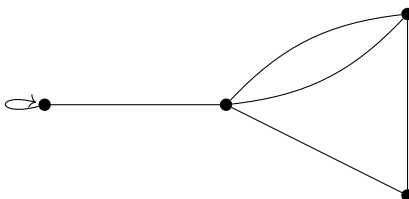
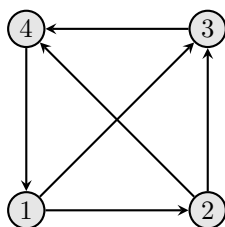


Figure E



For each adjacency matrix, draw the graph.



## Matrix A

$$A = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$$

## Matrix B

$$B = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

## Matrix C

$$C = \begin{bmatrix} 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 \end{bmatrix}$$

## Matrix D

$$D = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 \end{bmatrix}$$

## Matrix E

$$E = \begin{bmatrix} 0 & 1 & 3 & 0 \\ 1 & 0 & 2 & 2 \\ 3 & 2 & 0 & 2 \\ 0 & 2 & 2 & 0 \end{bmatrix}$$