

12/02/2023

Till now we have studied about 1 D arrays.
Today we will be studying about the 2 D arrays.

2D Arrays

`int arr [] = { 1, 2, 5, 7 };`

arr → | 1 | 2 | 5 | 7 |

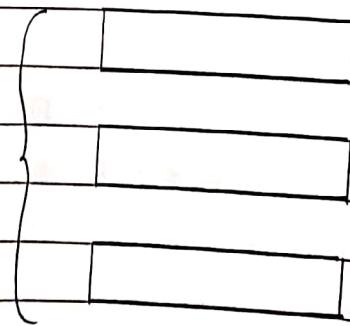
This is how 1D arrays are stored in the memory. But the question is what happens in case of 2D arrays. We can relate 2D arrays with Tic-tac-toe game which have 9 cells.

row 0			
row 1			
row 2			

col0 col1 col2

A grid is formed where there are multiple rows and multiple columns. The above grid can be formed by taking 3, 1D arrays of size 3 each.

3, 1D arrays



But this is a small grid, that's why we created 1D arrays but what if there are 1000 rows & it is not practical to create 1000 1D arrays. Hence 2D arrays make our job simpler.

```
↑ name           ↑ no. of columns
int arr [1000][1000];
↑               ↓ no. of rows
data type
```

Also we can mention different no. of rows & different number of columns.

Memory representation of 2D arrays

`int arr [2][2];` → Total elements = $2 \times 2 = 4$

`arr → [1 2 3 4] → 2D array`

Matrix was just used to visualize the 2D arrays but in memory, the 1D array representation is used.

	col0	col1	
row0	1	2	} Just way of visualizing 2D arrays
row1	3	4	

Accessing elements of 2D arrays

In the above 2D array, say we want to access 3. We can access it by arr [1][0]

`arr [0][0] → 1`

`arr [0][1] → 2`

`arr [1][1] → 4`

`arr [] []`
row index column
 index

Formulae for mapping in 1D array representation

is $C * i + j$

$C \rightarrow$ no. of columns

$i \rightarrow$ row-index

$j \rightarrow$ column-index

$\text{arr}[0][1] = 2$

Formulae $= 2 \times 0 + 1 = 0 + 1 = 1^{\text{st}}$ index in
1D array.

1	2	3	4
0	1	2	3

2 is present at 1st index in 1D array.

$\text{arr}[1][1] = 4$

Formulae $= 2 \times 1 + 1 = 2 + 1 = 3^{\text{rd}}$ index, 4 is
stored in 1D array.

Note $\rightarrow \text{arr}[0][1] \neq \text{arr}[1]$. It is just a way
that in memory 1D array is used but the
actual interface is the matrix only so
use $\text{arr}[i][j]$ to access the elements of
2D array.

Initialization of 2D array

`int arr[2][2] = {{1,2}, {3,4}};`

row-0 \rightarrow	1	2	
row-1 \rightarrow	3	4	

\uparrow \uparrow

col-0 col-1

Input and Output in 2 D Array

Input → `cin >> arr[i][j];`

Output → `cout << arr[i][j];`

We always have to play within 0 to $n-1$ index
where $n = \text{rows}$ and $\text{columns} = n$.

Note → `arr[i][j]`

No. of rows = n

No. of columns = m

$i \in [0, n-1]$ } Never go out this range
 $j \in [0, m-1]$

Accessing elements row-wise

```
int arr[3][3] = {{1,2,3}, {4,5,6}, {7,8,9}};
```

```
for (int i=0; i<3; i++) {
    for (int j=0; j<3; j++) {
        cout << arr[i][j] << " ";
    }
    cout << endl;
}
```

}

Output

1	2	3
4	5	6
7	8	9

arr

Accessing elements column-wise

```
for (int i=0; i<3; i++) {
```

Column

```

for (int j=0; j<3; j++) {
    cout << arr[j][i] << " ";
}
cout << endl;           (swapped in case of
                        column wise)

```

Output

1	4	7
2	5	8
3	6	9

Similarly we can take row-wise input & column-wise input.

Problem Solving

1) Row sum print

i/p →	1	2	3	
	4	5	6	
	7	8	9	

o/p → 6
15
24

Go to every row, calculate the sum & then print. We have to simply do row wise traversal.

Code

```

Void printRowSum ( int arr [ ] [ 3 ] , int rows
                    int cols ) {
```

```

for (int i=0; i<rows; i++) {
    int sum = 0; → After each row, sum = 0
    for (int j=0; j<cols; j++) {
        sum = sum + arr[i][j];
    }
    cout << "Row " << i << "sum is "
    << sum << endl;
}

```

Note → While passing 2D arrays in function, we need to mention the bound of column. It is a rule. Formulae of $C * i + j$ is the reason behind this.

2) Column sum print

i/p → Same array as that of Ques-1

O/p → 12
15
18

We just have to do column-wise traversal & print the sum just like we did in row wise sum print.

In the above code just change

$sum = sum + arr[i][j]$ to

$sum = sum + arr[j][i]$

3) Linear search in 2D array

i/p → Same array as that of ques-1

element = 3

O/p → True

We can do either row-wise traversal or column-wise traversal & compare each element with the element that is given to us as in i/p. If element is matched return true else return false if not found even after traversing full array.

Code

bool

find key (int arr[][3], int r, int c, int k)

for (int i=0; i<r; i++) {

 for (int j=0; j<c; j++) {

 if (arr[i][j] == k) {

 return true;

}

} return false;

}

4) Maximum and minimum in 2D array.

i/p → Same array as that of ques-1

o/p →

Maximum = 9

Minimum = 1

INT_MAX

Initialize $\text{maxi} = \text{INT_MIN}$ & $\text{mini} = \text{INT_MAX}$ and compare maxi & mini with each element.

✓ $\text{maxi} < \text{arr}[i][j] \rightarrow \text{update maxi}$

✓ $\text{mini} > \text{arr}[i][j] \rightarrow \text{update mini}$

After

the maxi & mini

traversing full array, return

we can leave only first index empty

Codeit is necessary to give bound to
2d array ,while declaring in fuxn

```
int maximum (int arr [][3], int r, int c) {
    int maxi = INT_MIN;
    for (int i=0; i<r; i++) {
        for (int j=0; j<c; j++) {
            if (arr[i][j] > maxi) {
                maxi = arr[i][j];
            }
        }
    }
    return maxi;
}
```

```
int minimum (int arr [][3], int r, int c) {
    int mini = INT_MAX;
    for (int i=0; i<r; i++) {
        for (int j=0; j<c; j++) {
            if (arr[i][j] < mini) {
                mini = arr[i][j];
            }
        }
    }
    return mini;
}
```

5) Transpose of matrix

i/p same array as that of ques-1

o/p →

1	4	7
2	5	8
3	6	9

We just have to change the rows into the column. We need to perform the swap on each cell.

swap (arr[i][j], arr[j][i]);

But the inner for loop won't run fully as some elements will be swapped twice & would reach to the original position and hence the matrix would be same.

Code

```
void transpose (int arr [ ] [ 3 ], int r , int c ) {
```

```
    for (int i = 0 ; i < r ; i + +) {
```

```
        for (int j = 0 ; j < i ; j + +) {
```

```
            swap (arr [i] [j] , arr [j] [i] )
```

3

3

Better way to use 2D arrays

1 D array → vector <int> arr;

2 D array → We will be using vector of vector concept

3 elements

vector 1				→ Vector of vector having
vector 2				3 cells
vector 3				
vector 4				
vector 5				
vector 6				total elements = $6 \times 3 = 18$

vector < vector < int >> v

`↳ data type` name of vector

Each block will contain vector of integers.

How to mention no. of rows & no. of columns?

`vector<vector<int>> arr(3, vector<int>(5, 0))`

no. of rows = 3 ↳ no. of initialised with 0
 columns = 5

Inserting data in the 2D vector /vector of vector.

```
vector <vector <int>> arr;
```

In arr, vectors will be pushed.

```
vector <int> a {1, 2, 3};
```

arr.push_back(a);

No. of rows in vector of vector

`arr.size()` will tell the no. of rows

in vector of vector

No. of columns in vector of vectors.

`arr[i].size()` where i is row-index

will tell the no. of columns in each row.

Note →

```
vector<vector<int>> arr (rows ,
```

```
vector<int> (col, 0);
```

Inner vector will be

initialized to 0.