1 Introduction

1.1 Protocol performances

G: Total load, S arrival rate of new packets.

1.1.1 Pure ALOHA

If you have data to send, send the data. If the message collides with another transmission, try resending later. On collision, sender waits random time before trying again.

$$P(k \text{ transm. in } 2Xs) = \frac{(2G)^k}{k!}e^{-2G}$$
$$S = G \cdot P(0) = Ge^{-2G}$$

1.1.2 Slotted ALOHA

Probability of k packets generated during a slot: P(k) = $\frac{G^k e^{-G}}{L!}$ Throughput: $P(1) = Ge^{-G}$

Goal: reduce the wastage of bandwidth due to packet collisions. Principle: sensing the channel before transmitting (never transmit when the channel is busy).

Non-persistent If channel is busy, directly run back off algo-

p-persistent If it is busy, they persist with sensing until the channel becomes idle. If it is idle:

- With probability p, the station transmits its packet
- With probability 1-p, the station waits for a random time and senses again

Performance of Unslotted nonpersistent CSMA : For a = t_{prop}/X , the normalized one-way propagation delay. S =

$$\frac{G^{-aG}}{G(1+2a)+e^{-aG}}$$

Performance of Slotted nonpersistent CSMA : $S = \frac{aG^{-aG}}{1 - aG}$

1.2 Exercises

Capacity of a link vs Transmission capacity (=total capacity of all the links). Wire: $C_t = \min\{C_1, C_2\}$ Wireless: $d/C_t = d/C_1 + d/C_2 \leftrightarrow C_t = (c_1c_2/c_1 + c_2)$ ALOHA: Aloha channel with infinite number of users gives

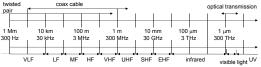
94% of idle slots. $P(0) = e^{-G} = 0.94 \rightarrow G = 0.062$

 $S = P(1) = Ge^{-G} \approx 5.8\%$

 $G < G_{peak} = 1$: channel underloaded.

Ration of busy slots occupied by collisions : $\frac{1-P(0)-P(1)}{1-P(0)} =$

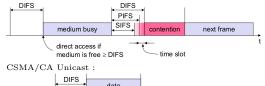
2 WLAN Engineering aspects

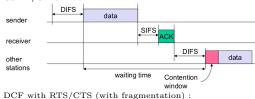


Frequency(f) and wave length(λ), $c = 3 \times 10^8 m/s$: $\lambda = c/f$

2.1 802.11

Physical layer : DSSS or FHSS, MAC Layer : best effort asynchronous data service, DCF CSMA/CA (mandatory), DCF with RTS/CTS or PCF (optional)





sender

receive other

MAC address format :

scenario	to DS	from DS	address 1	address 2	address 3	address 4
ad-hoc network	0	0	DA	SA	BSSID	-
infrastructure network, from AP	0	1	DA	BSSID	SA	-
infrastructure network, to AP	1	0	BSSID	SA	DA	-
infrastructure network, within DS	1	1	RA	TA	DA	SA

2.2 Exercises

 $A \rightarrow AP1$

Wireless LAN use polling between M workstations and a central access point. Channel at 25Mbps. Stations 100 m away from AP, polling messages 64 bytes long. Packet length: 1250 bytes. No more packet indicated with 64-byte message. Maximum arrival rate $\lambda_{max} = \rho_{max} * Br/P_{length} \rho_{max} =$ $M*N*T_{packet}$

 $\frac{Effective time}{Whole time}$ $\frac{P^{ac.ncc}}{M*(NT_{packet} + T_{poll} + T_{end} + 2t_{prop})}$ $t_{prop} = d/c \; T_{packet} = \frac{1250*8}{25*10^6}$ One station A sends a frame to another station B in a

different BSS in an IEEE 802.11 infrastructure network with DCF access method without RTS/CTS.

From Туре Dur Α1 A2 0 $T_d + SIFS + T_A$ BSS1 Data Free space loss, ideal isotropic antenna: $AP1 \rightarrow A$

To DS	From DS	Type	Duration	Addr. 1		
0	0	ACK	0	A		
$\mathrm{AP1} o \mathrm{AP1}: 1$, 1 , Data , $T_d + S + T_A$, $\mathrm{AP1}$, B , $\mathrm{AP1}$						

 $AP2 \rightarrow AP1 : 0, 0, ACK, 0, AP1$ $AP2 \rightarrow B: 0, 1, Data, T_d + S + T_A, B, BSS2, A$ $B \rightarrow AP2 : 0, 0, ACK, 0, BSS2$

3 Bianchi model

1

 π , probability of transmission, p, probability of collision, $b_{i,k}$ stationary probability of state i, k: $p = 1 - (1 - \pi)^{N-1}$ $\pi = \sum_{i=0}^{m} b_{i,0} = \frac{b_{0,0}}{1-p} = \frac{2(1-2p)}{(1-2p)(W_{min}+1) + pW_{min}(1-(2p)^m)}$ $= \frac{i=0}{1+W_{\min}+pW_{\min}\sum_{k=0}^{m-1}(2p)^k}$ $\left|\begin{array}{l} b_{i,k} = \frac{CW_i-k}{CW_i} \cdot \left\{ \begin{array}{l} (1-p)\sum_{j=0}^{m}b_{j,0} & i=0\\ p\cdot b_{i-1,0} & 0< i< m\\ p\cdot (b_{m-1,0}+b_{m,0}) & i=m \end{array} \right.$

3.1 Saturation throughput

$$\begin{split} \tau &= \frac{E[\text{Payload Transmitted by user i in a slot time}]}{E[\text{Duration of slot time}]} \\ &= \frac{P_{\text{S}}P_{\text{tr}}L}{P_{\text{S}}P_{\text{tr}}T_{\text{S}} + P_{\text{tr}}(1-P_{\text{S}})T_{\text{C}} + (1-P_{\text{tr}})T_{\text{id}}}, \\ P_{\text{S}} &= \frac{N\pi(1-\pi)^{N-1}}{1-(1-\pi)^{N}}, \\ P_{\text{tr}} &= 1-(1-\pi)^{N}, \\ T_{\text{S}} &= t_{\text{header}} + t_{\text{payload}} + \text{SIFS} + t_{\text{ACK}} + \text{DIFS} + 2\sigma, \\ T_{\text{C}} &= t_{\text{header}} + t_{\text{payload}} + \text{SIFS} + \sigma \end{split}$$

3.2 DOMINO Cheating detection

Cheating Method	Detection Test		
Frame scrambling	Number of retransmissions		
Oversized NAV1	Comparison of the declared and		
	actual NAV values		
Transmission before	Comparison of the idle time after		
DIFS	the last ACK with DIFS		
Backoff manipulation	Actual Backoff/ Consecutive		
_	Backoff		
Frame scrambling	Periodic dummy frame injection		
with MAC forging			

4 Antennas & Propagation

Free space propagation, received power: $P_R = P_T \frac{A_R}{4\pi d^2} \eta_R$ with $\eta_{\rm R}$ an efficiency parameter, $A_{\rm R}$ the receiving antenna

Focusing capability, depends on size in wavelength λ : $G_{\rm T} = 4\pi \eta_{\rm T} A_{\rm T}/\lambda^2$

Directional emitter, received power: $P_{\rm R} = P_{\rm T} G_{\rm T} \frac{A_{\rm R}}{4\pi d^2} \eta_{\rm R}$ Free space received power: $P_{\rm R} = P_{\rm T} G_{\rm T} G_{\rm R} (\frac{\lambda}{4\pi d})^2$ Loss: $L = \frac{P_T}{P_R} = \frac{(4\pi d)^2}{G_R G_T \lambda^2}$

Parabola: $G = \frac{7A}{12}$

Mobnet Decibels : $B = 10 \log(\frac{P}{P_0})$

Propagation modes Ground Wave: $f \leq 2 \text{ Mhz}$, Sky Wave, Line of Sight: f > 30 Mhz

4.0.1 Line of sight equations

Horizon distance d[km] in **kilometers**, antenna height h[m]and refraction adjustment factor K = 4/3:

Optical LOS : $d = 3.57\sqrt{h}$

Effective LOS : $d = 3.57\sqrt{Kh}$

Max LOS distance for two antennas

$$3.57(\sqrt{Kh_1} + \sqrt{Kh_2})$$

A431 Free Space Loss

 $\frac{P_t}{P_r} = \frac{(4\pi d)^2}{\lambda^2} = \frac{(4\pi f d)^2}{c^2}$

Free space loss equation can be recast:

$$L_{DB} = 10\log\frac{P_t}{P_r} = 20\log(f) + 20\log(d) - 147.56dB$$

Free space loss accounting for gain of other antennas

$$\frac{P_t}{P_r} = \frac{(4\pi d)^2}{G_r G_t \lambda^2} = \frac{(cd)^2}{f^2 A_r A_t}$$

 $G_t = gain of transmitting antenna$

 A_r = effective area of receiving antenna

Categories of noise : Thermal Noise, Intermodulation Noise, Cross-talk, Impulse Noise.

Thermal Noise $N_0 = kT \quad (W/Hz)$

For signal power S, bitrate R, $k = 1.3806 \cdot 10^{-23} JK^{-1}$ the Boltzmann constant and T the temperature: $\frac{E_b}{N_0} = \frac{S/R}{N_0} =$ $\frac{S}{kTR}$

4.2 Forward Error Correction (FEC)

Redundancy in packets to allow limited error correction at the receiver: used in 802.11a (Convolutional), HSDPA (Turbo Codes) and 802.11n (LDPC).

5 Cellular Networks

For a trunk of N channels, an offered load $A = \lambda E[X]$, X the call duration, Y the call arrival per $\sec \sim \text{Poisson}(\lambda)$ and ρ the traffic carried by each channel:

> $P_{\text{Blocking}} = P(\text{Drop a call because busy line})$ $= \frac{A^N}{N! \sum_{i=0}^N (\frac{A^i}{i!})}$ $\rho = \frac{(1 - P_{\text{blocking}})A}{N}$

Cellular efficiency $E = \frac{Conversations}{collection}$

Area: $A = 1.5R^2\sqrt{3}$

Distance btw. adjacent cells: $d = \sqrt{3}R$

5.1 Co-channel interference

Co-channel reuse ratio : $Q = \frac{D}{R} = \sqrt{3N}$ with D the distance to the nearest co-channel cell, R the radius of a cell and N the cluster size.

Signal to Interference ratio (SIR) : $SIR = \frac{S}{I} = \frac{S}{\sum_{i=1}^{i} I_i}$. With

S the desired signal power, I_i the interference power from the ith interfering co-channel base-station, i₀ the number of co-channel interfering cells

Signal to Interference plus Noise ratio (SINR) : SINR = $\frac{S}{I+No}$

Average received power $P_r: P_r = P_0(\frac{d}{d_0})^{-\alpha}$ or

 $P_r(dBm) = P_0(dBm) - 10\alpha \log(\frac{d}{d_0})$ with P_0 the power received from a small distance d_0 from the transmitter and α the path loss exponent.

SIR in the corner of a cell : $\frac{S}{I} = \frac{R^{-\alpha}}{\sum^{i_0} D_{\alpha} D_{\alpha}^{-\alpha}}$

First interfering layer approximation : $\frac{S}{I} = \frac{(\frac{D}{R})^{\alpha}}{i_0} = \frac{(\sqrt{3N})^{\alpha}}{i_0}$ eg. $= (\frac{D}{R})^2 \frac{1}{2} \text{ for two first layer interferers (cell divided into 3 sectors with directional antennas.)}$

5.2 Capacity of a cellular network

For B_{t} the total allocated spectrum and B_{c} the channel

$$m = \frac{B_t}{B_c \frac{Q^2}{3}} = \frac{B_t}{B_c \left(\frac{6}{3\frac{\alpha}{2}} \left(\frac{S}{T}\right)_{\min}\right)^{\frac{2}{\alpha}}} = \lfloor \frac{C}{N} \rfloor$$

For a cluster size N, $N = (i + j)^2 - ij$ for i, j = 0, 1, 2, ...and number of channels C.

5.2.1 CDMA Capacity: single cell case

For the bitrate R, available bandwidth W, noise spectral density N_0 , thermal noise η , received user signal (at base station) S, we have a possible number N of users:

$$N = 1 + \frac{W/R}{E_b/N_0} - (\frac{\eta}{S})$$

With a duty cycle δ (Discontinuous transmission mode: takes advantage of intermittent nature of speech):

$$N = 1 + \frac{1}{\delta} \frac{W/R}{E_b/N_0} - (\frac{\eta}{S})$$

And if we have m sectors, the effective capacity becomes mN. 5.2.2 CDMA multiple cells

Frequency reuse factor on the uplink $f = \frac{N_0}{N_0 + \sum_i U_i N_{ai}}$ where $N_0 = \text{total}$ interference power received from N-1 in-cell users, $U_i = \text{number of users in the}_i^{\text{th}}$ adjacent cell and N_{ai} = average interference power from a user located in the ith

Average received power from users in adjacent cell $N_{a\,i}$ = $\sum_{i} N_{ij}/U_{i}$ where N_{ij} = power received at the base station of interest from the j^{th} user in the i^{th} cell

5.3 Ad-hoc Netowrks

Upper Bound for the Throughput If we have n identical randomly located nodes each capable of transmitting W bits/s. Then the throughput $\lambda(n)$ obtainable by each node for a randomly chosen destination is $\lambda(n) = \Theta\left(\frac{W}{\sqrt{n \log n}}\right)$

Routing proactive: DSDV, OLSR. reactive: AODV, DSR DSR: Route discovery only when source S attempts to send a packet to destination D, by flooding Route Requests (RREQ). Route maintenance by allowing S to detect when a link is broken with a Route Error message RERR, S try other route in its cache, otherwise route disc. AODV: Similar to DSR but maintains routing tables at the nodes (smaller header). AODV ages the routes and maintains a hop count.

6 Mobile Network Laver

ACO Authenticated Cipher Offset

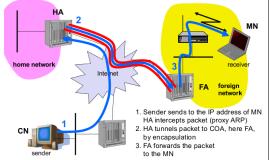
AIFS Arbitrary Inter-Frame Space

AMF Authentication and Key management

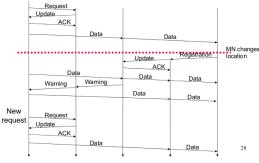
Mobile Network Layer : Transparency, Compatibility, Security, Efficiency, Scalability,

Network Layers Top-down: Application, Transport, (HIP layer), Network, Data-link, Physical.

Mobile IP :Issues : Security(Authentication to FA is problematic), Firewalls, QoS. IPSec can provide CIA by adding layer btwn IP and TCP/UDP. Mobile IPv6 : no FA, COA alwys co-loc, IPsec, bidirectional tunnel HA<->COA.

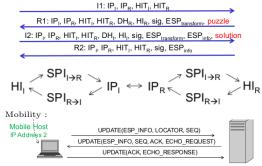


Agent Discovery: -MN discovers its location, HA and FA periodically send advertisement messages or via DHCP -MN learns a COA. Registration : -MN securely signals the COA to the HA (via the FA) Tunneling: -HA encapsulates IP packets from CN and sends them to the COA -FA (or MN) decapsulates these packets and sends them to the MN (Route optimization: HA provides the CN with the current location of MN (FA).CN sends tunneled traffic directly to FA. Picture of route and FA handover optim :



HIP New layer btw IP and transport, integrate security, mobility and multi-homing, decouple name and locator role of IP HI = public key, HIT=h(HI), DH : Diffie-Hellman key

material, sig signature generated with private key of $HI_{I/R}$ Initiator (I) Responder (R)



7 TCP

7.1 Standard

Tahoe Basic TCP. Three duplicate ACK's provoke fast retransmit (resend 1st missing packet), set ssthresh to cwnd/2, cwnd to 1 and provoke slow start.

Reno Three duplicate ACK's provoke fast retransmit, ssthresh to cwnd/2, cwnd to ssthresh + 3 and enter fast recov-

Fast Recovery Increase cwnd by 1 segment for every received duplicate ACK. (Warning, unlogical: When new ACK is received, cwnd = ssthresh and enter congestion avoidance). If a timeout occurs, set cwnd to 1 and enter slow start.

New Reno Fast Recovery More intelligent fast recovery where you remember the last received ACK.

7.2 Mobile

Indirect TCP (I-TCP) Connection split at FA. Standard TCP on the wire line, wireless optimized TCP on the wifi side: shorter timeout, faster retransmission. Loss of end-to-end semantics, security issues.

Mobile TCP (M-TCP) Split connection at FA. Monitor packets, if a disconnect is detected, report receiver window = 0: sender will go into persist mode and doesn't timeout or modify his congestion window. Preserves end-to-end semantics. Disadv.: wifi losses propagate to the wire network, link-errors pkt loss must be resent by sender, security issues. Summary: only handles mobility errors, no transmission errors.

Snooping-TCP TCP-aware link layer: Split connections, FA buffers and retransmits segments, does not ACK buffered packets (preserves end-to-end semantics).

Transaction oriented TCP (T-TCP) TCP phases: connection setup, data transmission, connection release. T-TCP combines these steps and only 2-3 packets are needed for short messages. Efficient for single packet transactions, but requires TCP modifications on all hosts.

8 Security

Security Requirements : Confidentiality, Authenticity, Replay Detection, Integrity, Access Control, Jamming Protection.

GSM Shared secret and challenge responses, one-way au-

3GPP (Improvements from GSM) Two-way authentication, avoid fake base station, cipher keys and auth data is now encrypted, integrity. Privacy/Anonymity not completely protected however.

9 Privacy

Privacy Related Notions Anonymity, untraceability, unlinkability, unobservability, pseudonymity

Best to worst against information leakage: GPS: no third-party, determined 'alone'. Cell-ID: requires the operator database that is relatively protected (they won't easily mine you). Wireless: requires one or several third-party owned databases that can track you, and it is relatively precise due to short radio range.

9.1 Privacy Metrics

Entropy-Based Anonymity A the anonymity set, p_x the probability for an external observer that the action was performed

$$\sum_{\forall x \in A} p_x \log(p_x)$$

Entropy-Based Unlinkability I_1, I_2 , sets of elements to be related, p_r , the probability two elements are related for an external observator:

$$\sum_{\forall R \subseteq I_1 \times I_2} p_r \log(p_r)$$

This cheat-sheet is an update by Aubry Cholleton of the amazing work of Julien Perrochet, Christopher Chiche and Tobias Schlatter, GitHub:https://github.com/aubry/mobnet2012

AODV Ad Hoc On-demand Distance-Vector AP Access Point AP Access Point ATIM Ad-hoc Traffic Indication Map AUTN Authentication Token AV Authentication Vector AV Authentication Vector
BSS Basic Service Set Identifier
BSS Basic Service Set
CARA Collision Avoidance and Resolution Multiple Access

CA Collision Avoidance CCA Clear Channel Assessment CDMA Code Division Multiple Access CH Correspondant Host CN Correspondant Node

COA Care-Of Address CRC packet received CoRreCtly CSMA/CD CSMA with Collision Detection CSMA Carrier Sense Multiple Access

CTS Clear To Send CW Contention Window DAMA Demand-Assigned Multiple Access DA Destination Address DBPSK Differential Binary Phase Shift

Keying DCF Distributed Coordination Function

DQPSK Differential Quadrature Phase DSDV Destination Sequenced Distance Vector
DSRC Dedicated Short Range Communica-DSR Dynamic Source Routing DSSS Direct Sequence Spread Spectrum DS Differentiated Service DS Distribution System DTIM Delivery Traffic Indication Map DIIM Delivery Traine Indicated and DoS Denial of Service EAP-ILS TLS over EAP EAPOL EAP Over LAN EAP Extensible Authentication Protocol EDCA Enhanced Distributed Channel Ac-EHF Extra High Frequency

DECT Digital Enhanced Cordless Telecom-

munications
DHCP Dynamic Host Configuration Proto-

DH Diffie-Hellman DNS Domain Name System

ESP_{info} Contains SPI

FA Foreign Agent

EPC Electronic Product Code ESP Encapsulating Security Payload $ESP_{transform}$ Supported crypto suites ESS Extended Service Set FAMA Floor Acquisition Multiple Access

GMK Group Master Key cation HA Home Agent HF High Frequency HIP Host Identity Protocol HIT Host Identity Tag IFS Inter Frame Spacing IHL Internet Header Length IKE Internet Key Exchange LDPC Low Density Parity Check

FDD Frequency Division Duplex LEAP Light EAP FDMA Frequency Division Multiple Access LFSR Linear Feedback Shift Register FEC Forward Error Correction FHSS Frequency Hopping Spread Spec-LF Low Frequency LTE Long Term Evolution FQDN Fully Qualified Domain Name MACA-BI MACA By Invitation MACA Multiple Access with Collision GFSK Gaussian Frequency Shift Keying Avoidance (RTS-CTS(+ACK)) MAC Message Authentication Code GPRS General Packet Radio Service GSM Global System for Mobile Communi-MAHO Mobile Assisted Handover MAP Mobility Anchor Point MD Mobile Device MF Medium Frequency HCCA HCF Controlled Channel Access MH Mobile Host MIB Management Information Base MIC Message Integrity Code MN Mobile Node MSC Mobile service Switching Center HI Host Identifier HMIP Hierarchical Mobile IP HSPDA High Speed Downlink Packet Ac-MTSO Mobile Telecommunications Switching Office NAASS Normalized Average Anonymity ICMP Internet Control Message Protocol NAT Set Size

NAT Network Address Translation
NAV Net Allocation Vector

OFDMA Orthogonal Frequency-Division
Multiple Access IMSI International Mobile Subscriber OLSR Optimized Link- State Routing OTP One-Time Password PCF Point Coordination Function PEAP Protected EAP PEP Performances Enhancing Proxies ISI InterSymbol Interference KISS Keep It Simple and Stupid

PIN Personal Identification Number PLCP Physical Layer Convergence Protocol PMD Physical Medium Dependent PMK Pairwise Master Kev PN Pseudo-random Noise PSTN Public Switched Telephone Network PTK Pairwise Transient Kev QoS Quality of Service RADIUS, Remote, Authentication Dial-In User Service
RA Receiver Address
RERR Route ERRor
RFID Radio Frequency Identification RREP Route REPly RREQ Route REQuests RSN Robust Security Network RTCP Real Time Control Protocol RTM Retransmission TimeOut RTP Real Time Protocol RTS Request To Send RVS Rendez-Vous Server RWND Receiver Window SACK Selective ACKnowledgment SA Security Association SA Source Address SDMA Space Division Multiple Access SHF Super High Frequency SIFS Short Inter Frame Spacing SIM Subscriber Identity Module

SIP Session Initiation Protocol SPI Security Parameter Index SFI security Parameter Index STResh Slow Start Threshold STA STAtion STA Station TA Transmitter Address TCP Transmission Control Protocol TDD Time Division Duplex TDMA Time Division Multiple Access TIM Traffic Indication Map TKIP Temporal Key Integrity Protocol TLS Transport Layer Security TMSI Temorary Mobile Subscriber Iden-TOS Type Of Service TSF Timing Synchronisation Function TTL Time To Live UHF Ultra High Frequency UMTS Universal Mobile Telecommunica-tions System UV Ultraviolet Light VANET Vehicular Ad-hoc NETwork VHF Very High Frequency VLF Very Low Frequency

WAP Wireless Access Point WEP Wired Equivalent Privacy

WPA WiFi Protected Access

WLAN Wireless Local Area Network WMN Wireless Mesh Network WPAN Wireless Personal Area Network