

Reading: Kleinberg & Tardos Ch. 3.4

Additional resource: CLRS Ch. 22

Icebreaker

Suppose you are trying to organize a first-year orientation event, and your goal is to get students who don't already know each other to meet. (Students that already know each other will probably start clustering...)

Can you send the students into two different rooms so that each contains students who don't know anyone else there?



Thanks to Jessica Sidman for posing this problem!

Icebreaker

Construct (undirected) graph

- Vertex for each student
- Edge between two students if they know each other



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Working together

Suppose you and your friend have a research project to study the poetry of Emily Dickinson. You've been given a list of books about her poems upon which to base the project. However, neither one of you wants to read analysis of the same poem more than once...



Is it possible to divide the books between you so that no one reads about the same poem more than once?

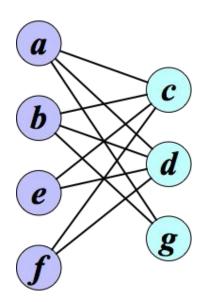
Working together

Construct (undirected) graph

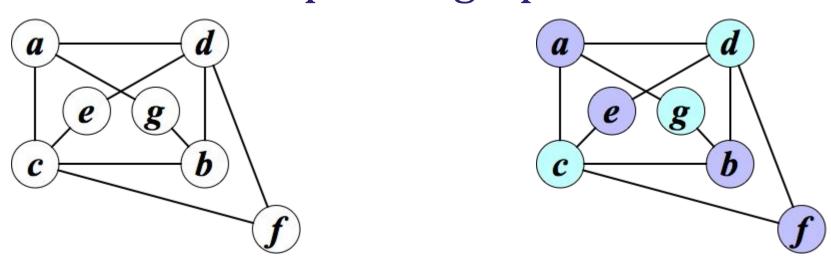
- Vertex for each book
- Edge between two books if analyze same poem



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Bipartite graphs



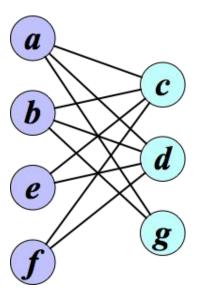
Bipartite graphs

- Bipartite graph G = (V, E)
- *V* can be partitioned into disjoint vertex sets: *X* and *Y*
- Edges: $E \subseteq X \times Y$
 - One endpoint in each set

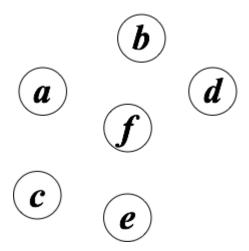
Example:

$$X = \{a, b, e, f\}, Y = \{c, d, g\}$$

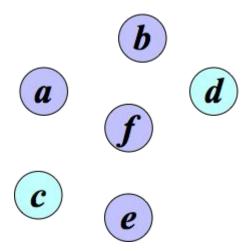
 $E = \{ac, ad, ag, bc, bd, bg, ec, ed, fc, fd\}$

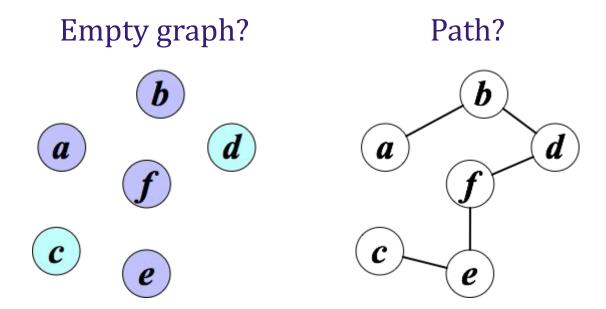


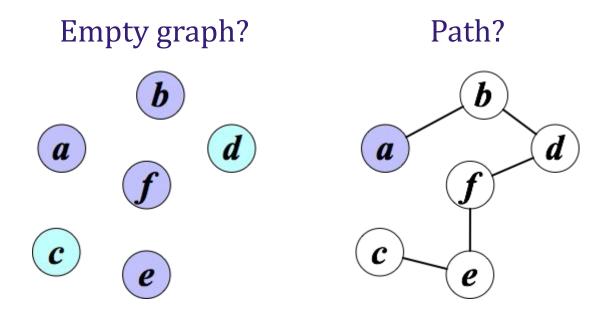
Empty graph?

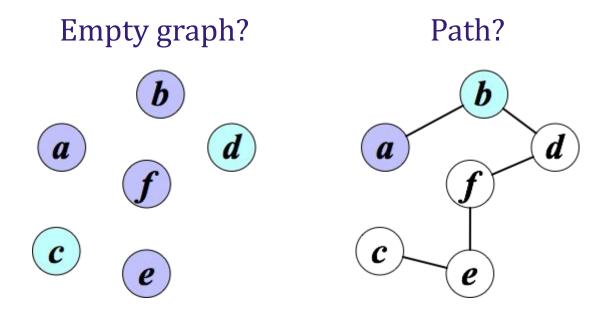


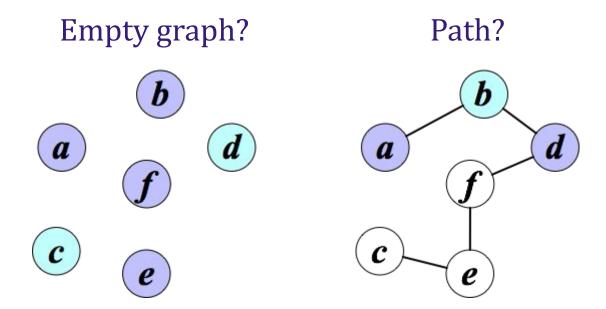
Empty graph?

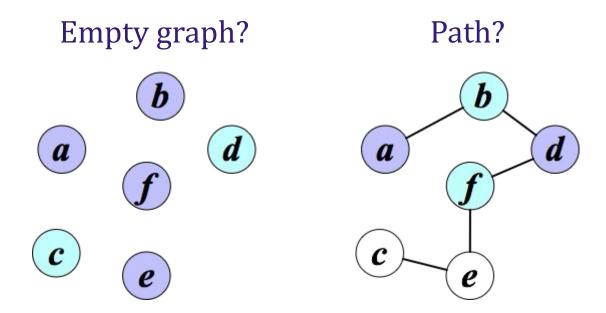


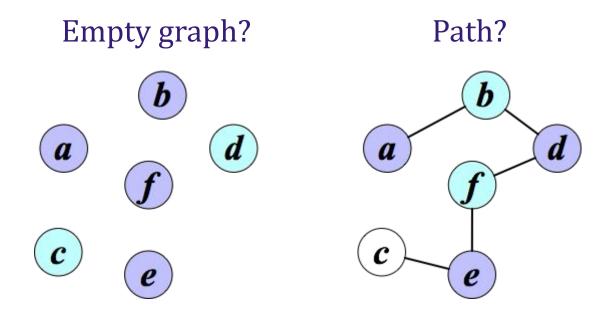


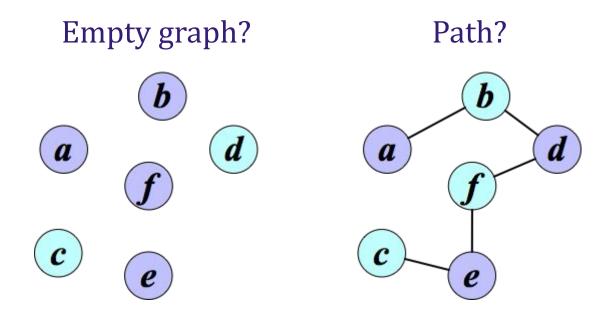


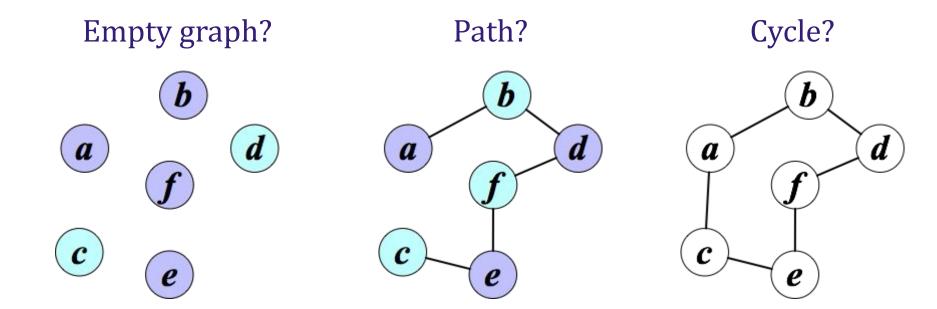


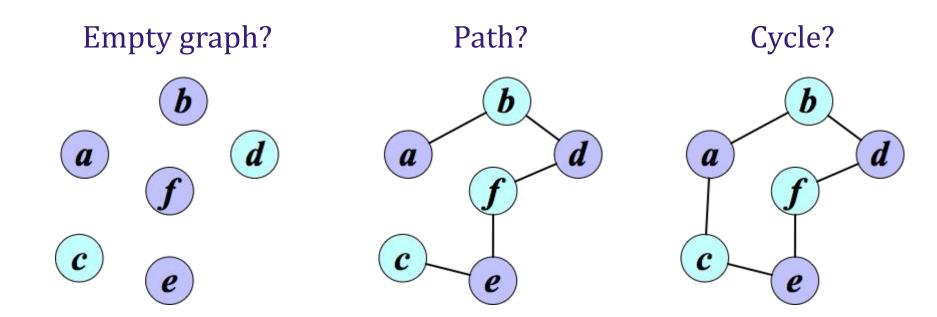








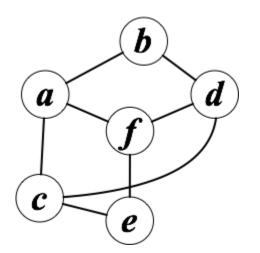


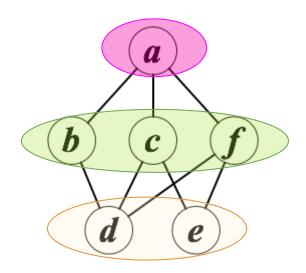


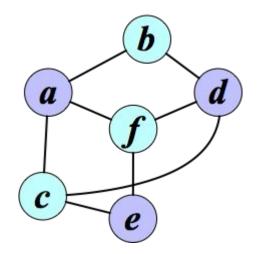
Can we generalize this to an algorithm?

What about...

- BFS from a vertex
- *X*: even layers
- *Y*: odd layers



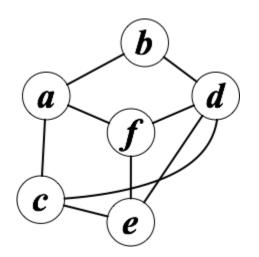


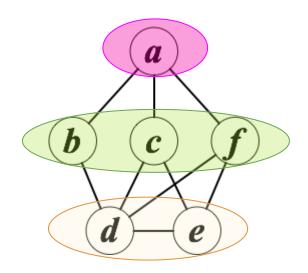


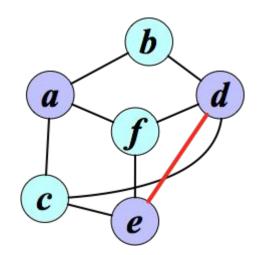
Uh-oh!

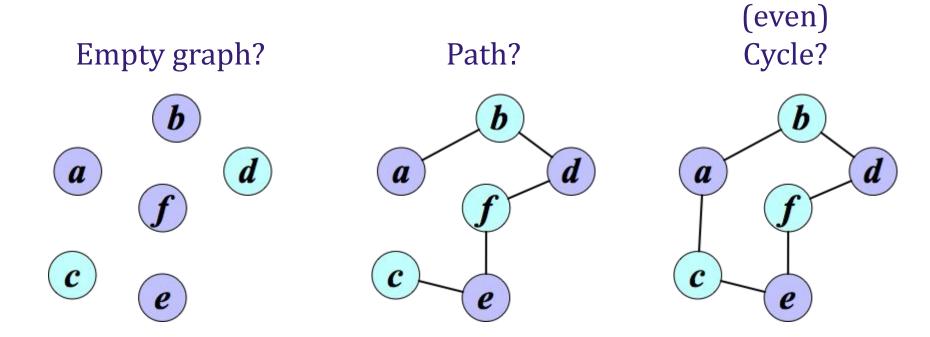
What about...

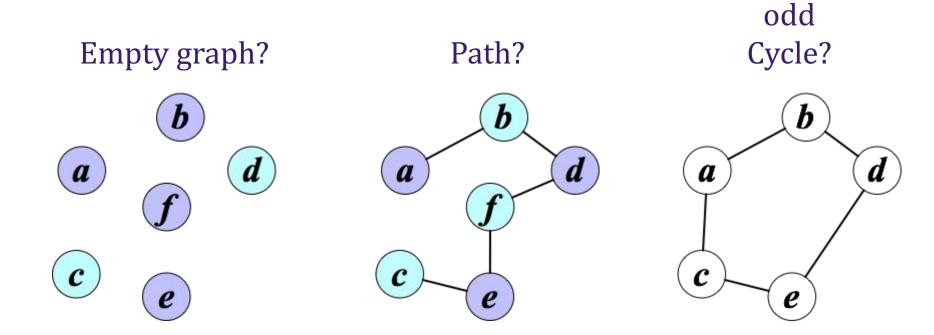
- BFS from a vertex
- *X*: even layers
- *Y*: odd layers

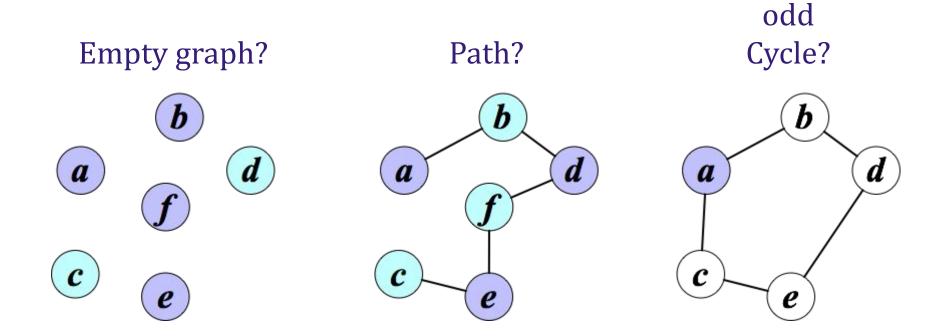


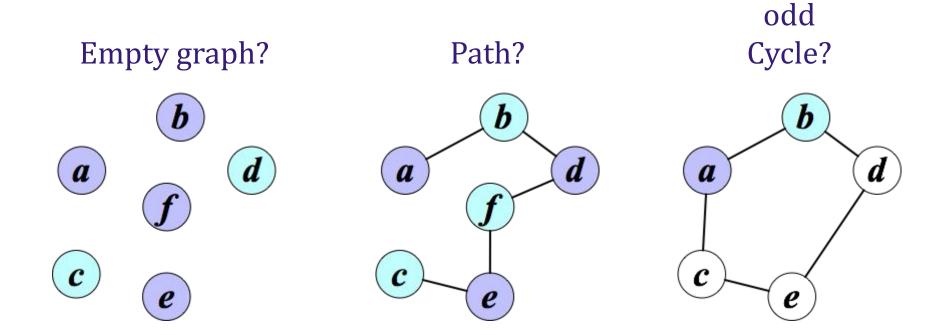


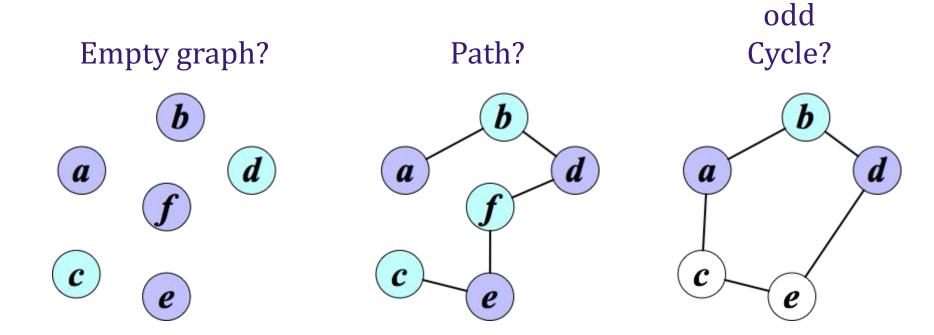


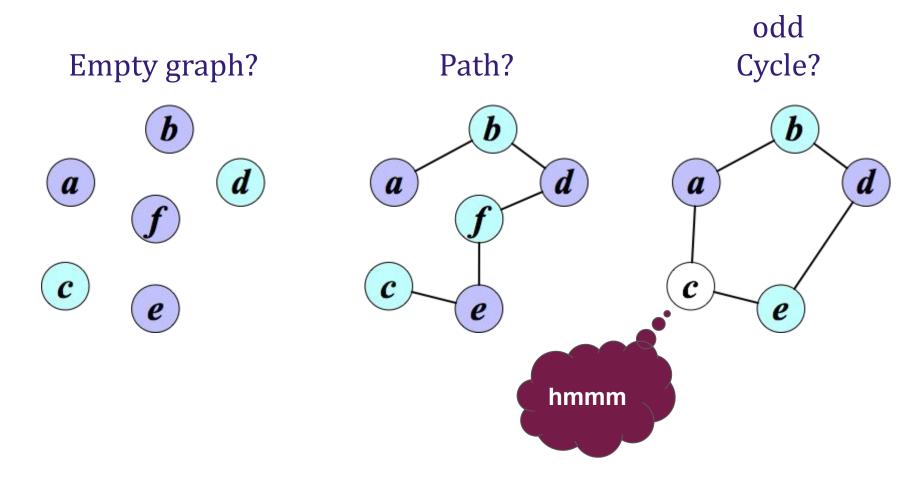












Bipartite graphs & odd cycles

Claim: Bipartite graphs => no odd cycles

Proof: Via contrapositive: *odd cycle => not bipartite*

Suppose, for a contradiction, *G* is bipartite with its vertices partitioned into sets *X* and *Y*.

By assumption, *G* has an odd cycle $P = v_1, v_2, ..., v_k$, where *k* is an even number.

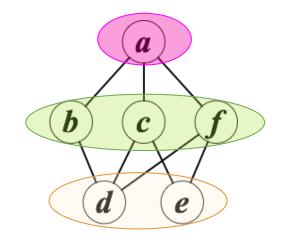
Since there is an edge between each consecutive pair of vertices, the vertices must alternate between *X* and *Y*.

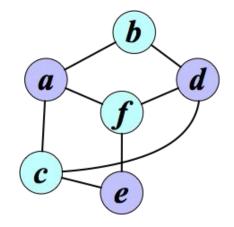
WLOG assume $v_1 \in X$; then $v_i \in X$ and $v_{i+1} \in Y$ for all odd i = 1,3,..., k-1. In particular, $v_k \in Y$. But, since P is a cycle, $v_1 = v_k$, contradicting the supposition that X and Y are disjoint.

So what happened during our proposed process?

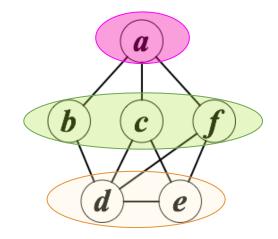
Remember... non-tree BFS edge endpoints are:

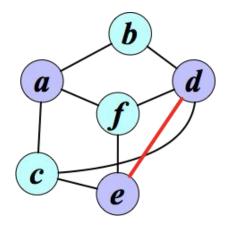
- Same layer, or
- One layer apart





same layer seems problematic?





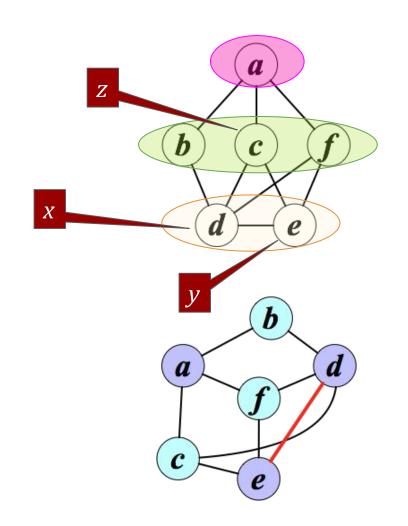
BFS non-tree edge: same layer => odd cycle

Proof:

Let T be the tree output by BFS and let xy be an edge with both x and y in layer L_i .

Consider z, the least common ancestor of x and y, in layer L_i

The cycle z-xy-z is of length (i-j) + 1 + (i-j) = 2*(i-j) + 1, so the graph contains an odd cycle. Therefore, it is not bipartite.

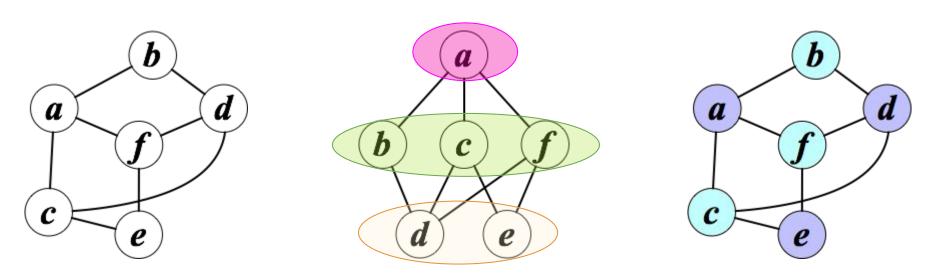


We have an algorithm!

isBipartite:

construct BFS tree from any vertex for each layer

if exists non-tree edge, return FALSE return TRUE with (*X*: even layers, *Y*: odd layers)



CS 312 - Algorithms - Mount Holyoke College

Correctness

Claim: *G* is bipartite if and only if **isBipartite** returns TRUE

Proof:

(=>) Show: *G* is bipartite => **isBipartite** returns TRUE Via contrapositive: **isBipartite** returns FALSE => *G* is not bipartite. **isBipartite** returns FALSE => non-tree edge in same layer of BFS tree => odd cycle => *G* is not bipartite

(<=) Show: isBipartite returns TRUE => G is bipartite
Tree edges alternate X and Y.
BFS non-tree edges between adjacent layers or in same layer.
 Adjacent layers respect bipartite
 no edges in same layer <=> isBipartite returns TRUE
G is bipartite.

Running time

isBipartite:

```
construct BFS tree from any vertex for each layer if exists non-tree edge, return FALSE return TRUE with (X: even layers, Y: odd layers)
```

- BFS: O(m + n)
- Good data structure \rightarrow non-tree edge check if same layer: O(m)
 - array to label vertex with layer #
- Total: O(m + n)