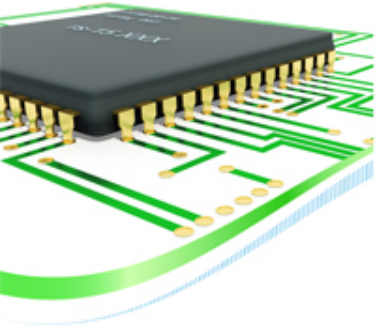


**Intermediate
Representations**



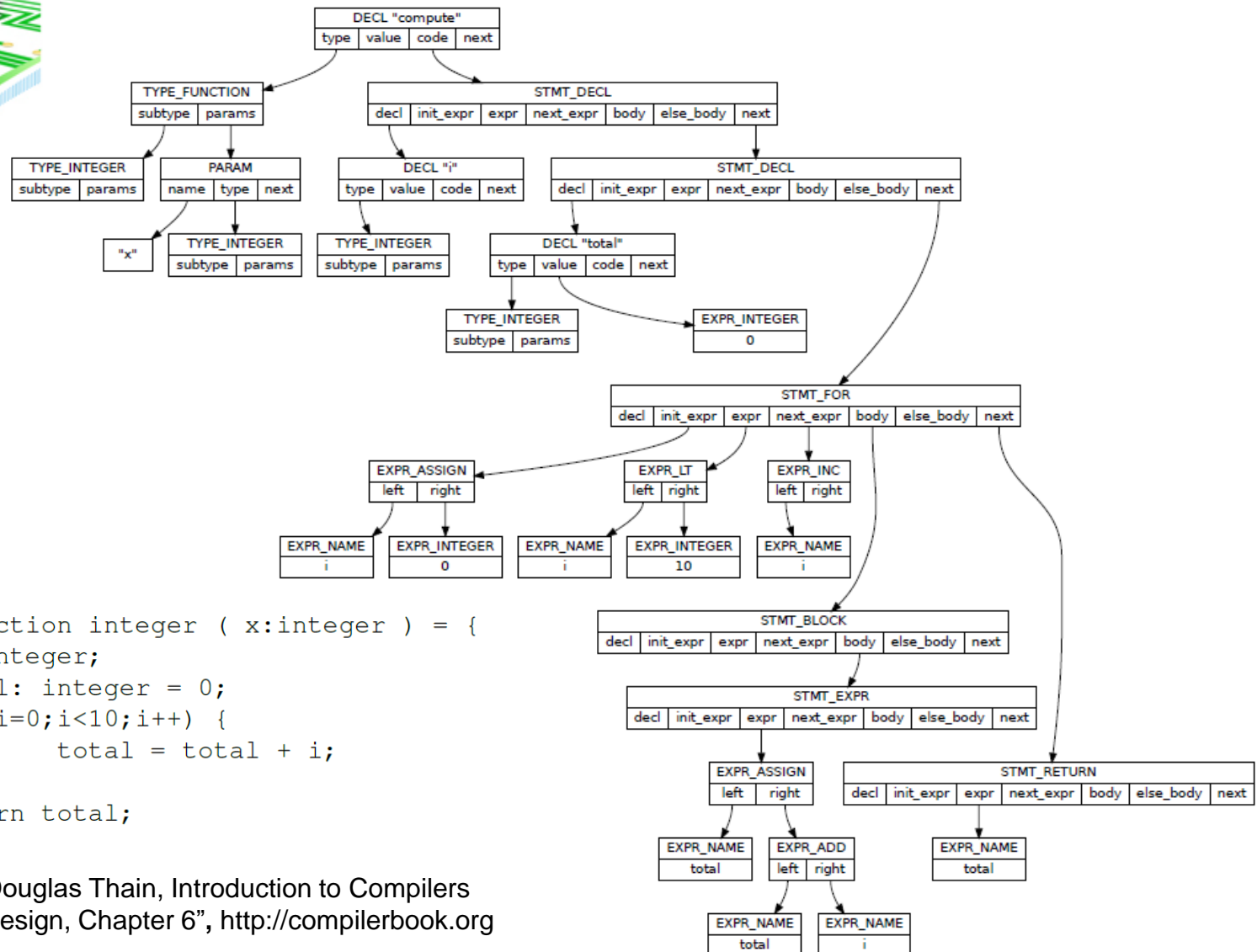
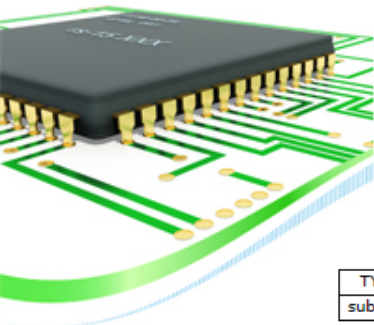
Definition

An intermediate representation is a composite derivative structure that supports generation of target code.

Extends the meaning,
regards the target architectures.

Abstract Syntax Tree (AST)

Semantic Representation as Graphical IR



```

compute: function integer ( x:integer ) = {
    i: integer;
    total: integer = 0;
    for(i=0;i<10;i++) {
        total = total + i;
    }
    return total;
}
  
```

Example from "Douglas Thain, Introduction to Compilers and Language Design, Chapter 6", <http://compilerbook.org>

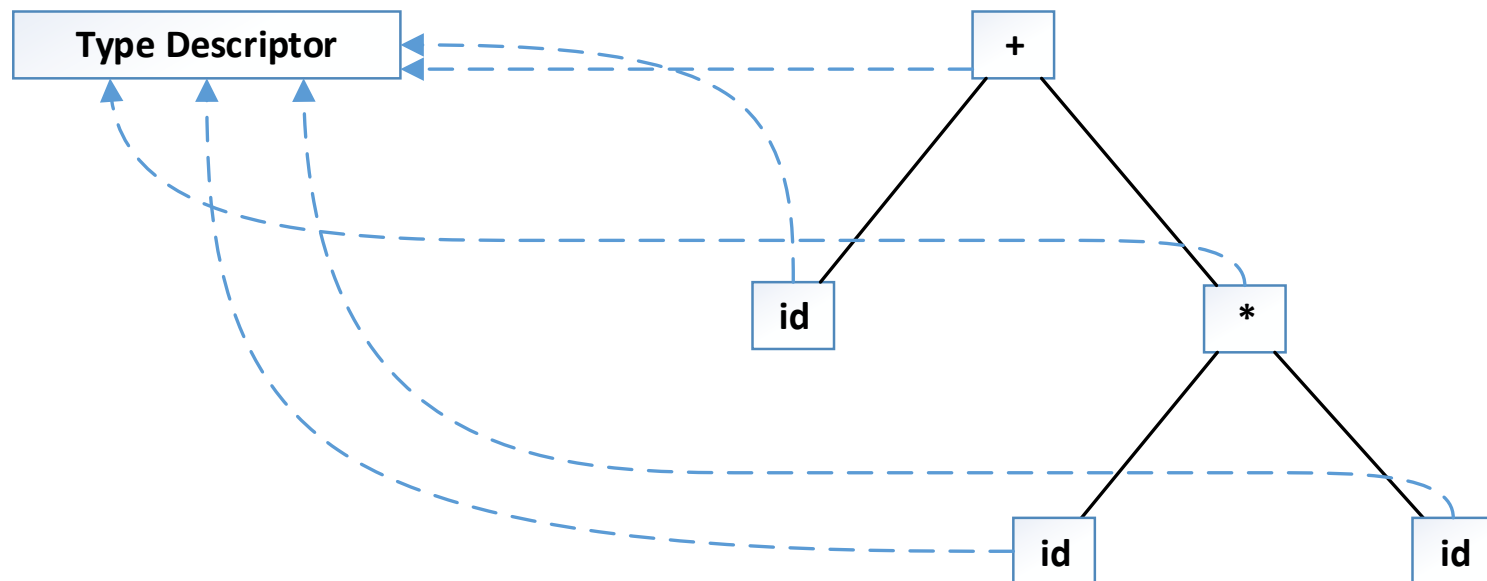
Abstract Syntax Tree (AST)

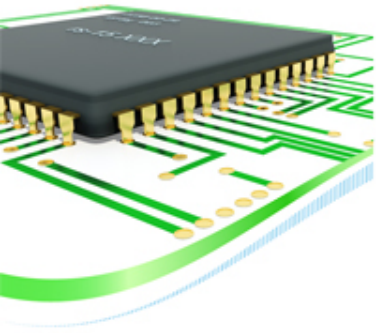
Semantic Representation as Graphical IR

Syntax is hierarchical.

Meaning is not!

- **Expressions**
- **Structure – Member Relationships**
- **Flow Control Statements**
- **Other Declarations**
- ...

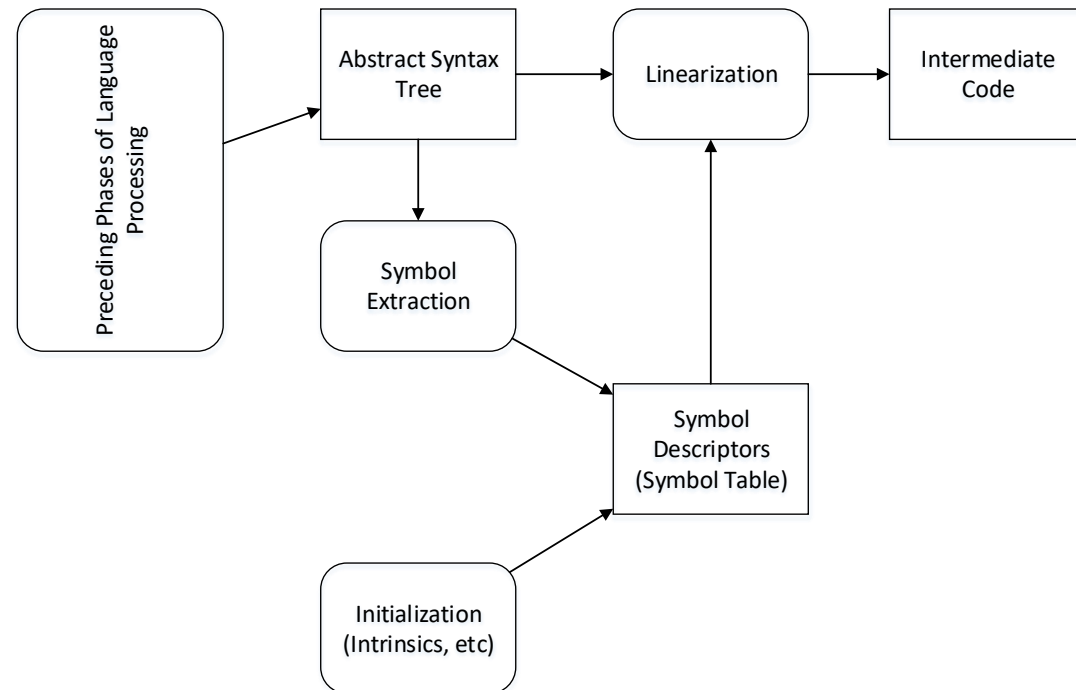




IR Generation

Linear Representations

Transforming into linear structure



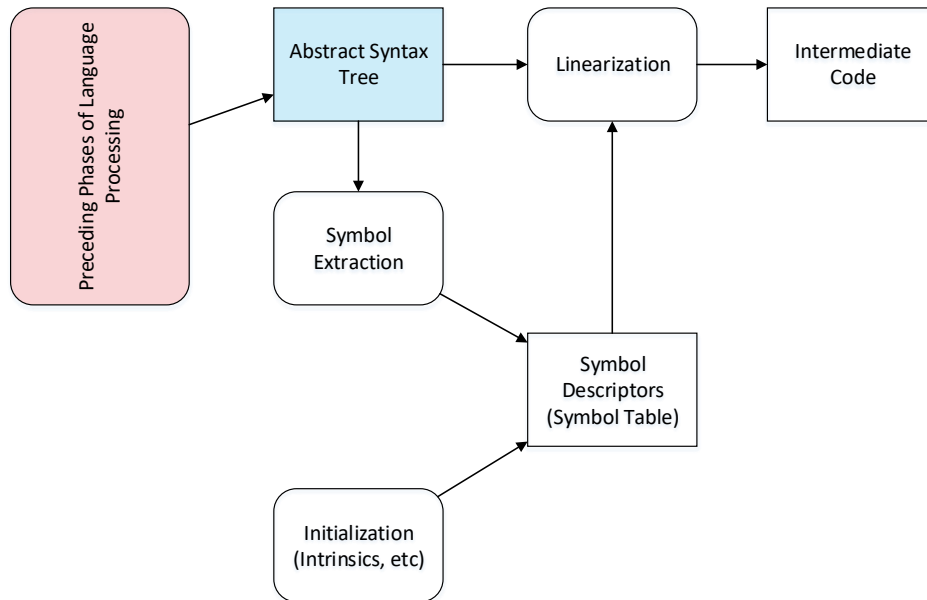
Just a scheme for transformation.



IR Generation

Graphical Representations

Building up AST



Semantic analysis phase

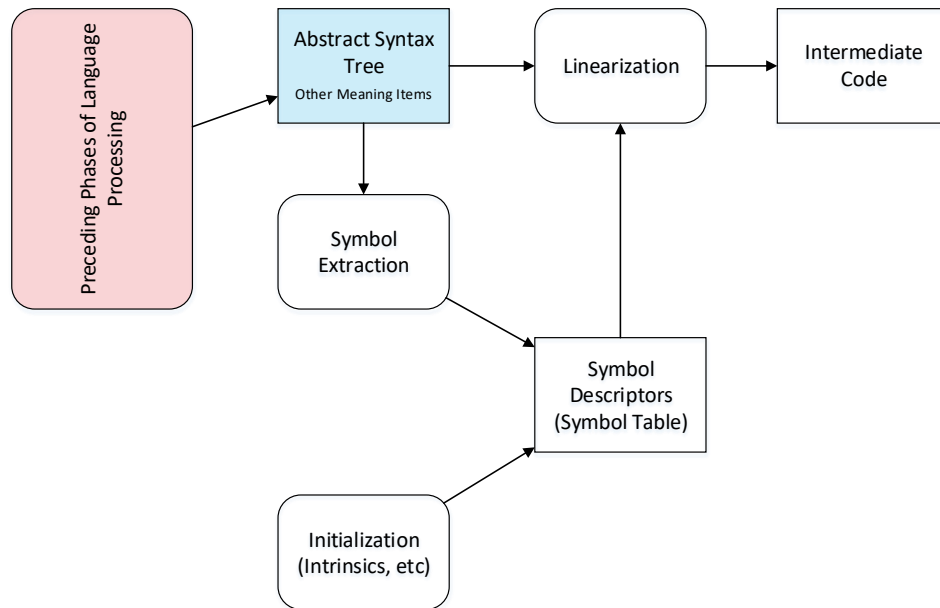
- **Inherited Attributes (Top-Down)**
- **Synthesized Attributes (Bottom-Up)**
- **Opportune Moments in Parsing!**



IR Generation

Linear Representations

Building up AST



Build

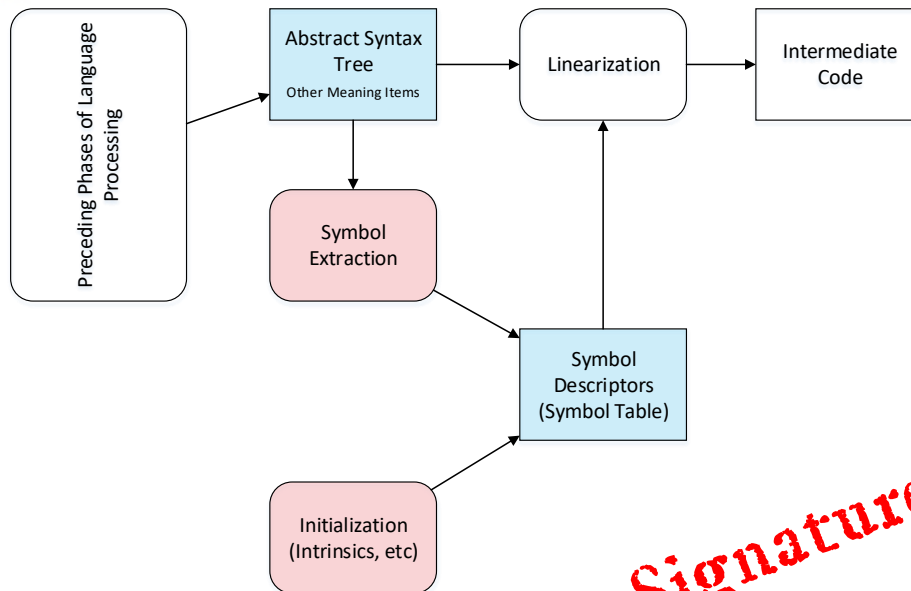
- **Expression Trees**
- **Control Flow Structures**
- **Symbolic Representation Items**
(namespaces, classes, types, formal parameters, methods, procedures, variables, ...)



IR Generation

Augmentation of AST, DAG

Multi-pass AST



Initial Setup

- Register Intrinsic Symbols

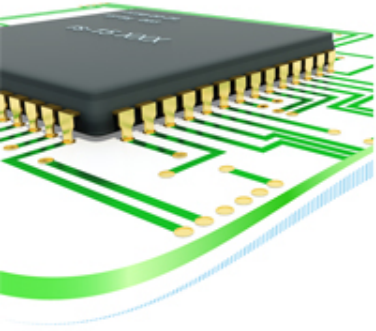
Symbolic Representation Items

- Populate the Symbol Entries
- Detect Duplications and Conflicts

Calculation

- Calculate Symbol Address
- Type Data

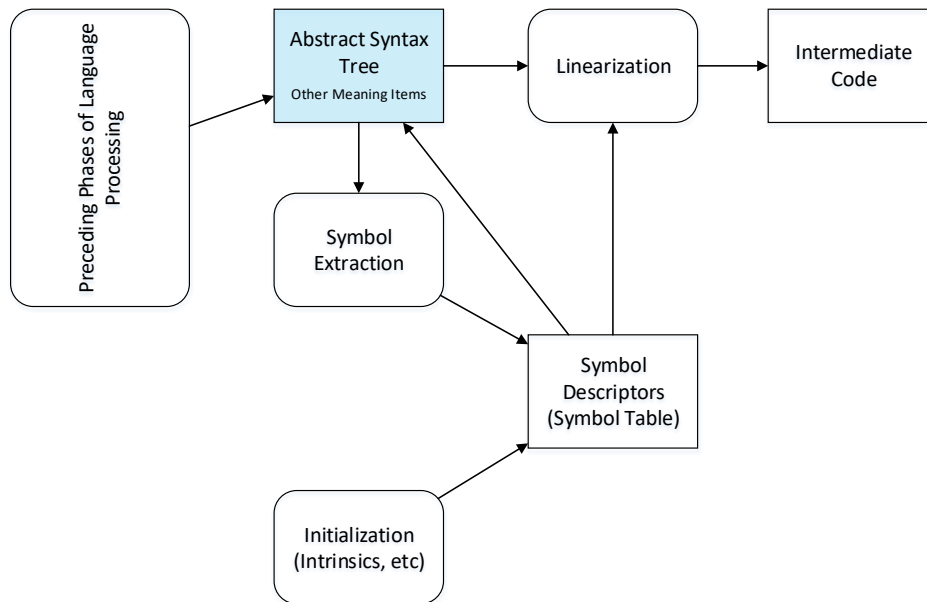
Meta-types & Type Signatures



IR Generation

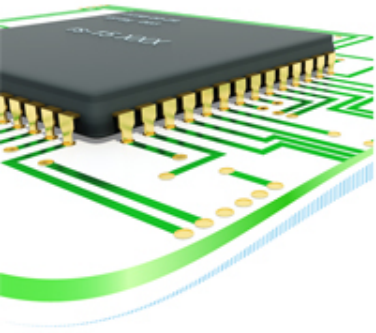
Augmentation and Transformation

Multi-pass AST



Expression Trees

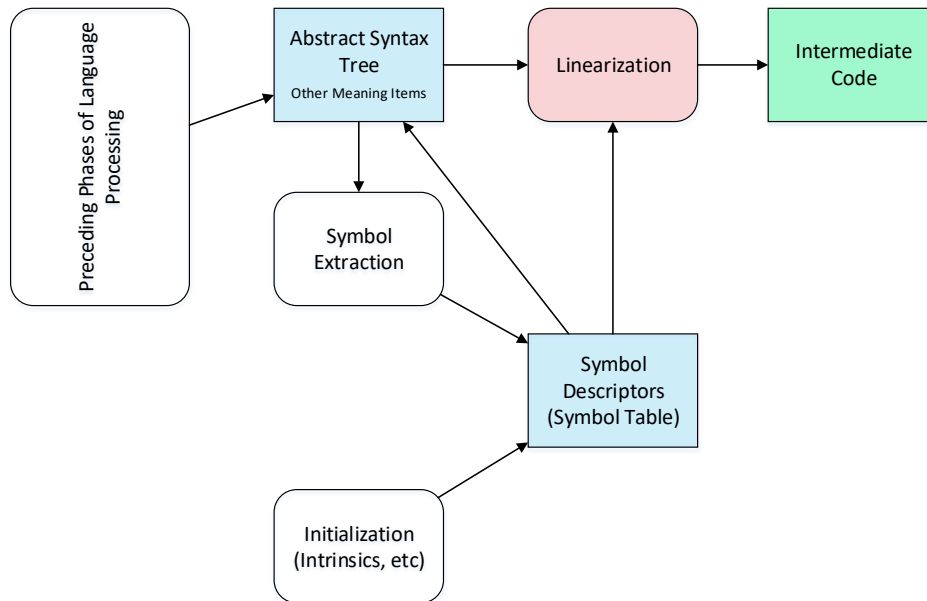
- **Apply Transformation (DAG)**
- **Inject Cast / Pseudo Operators**
- **Evaluate Node Types**
- **Detect errors** (Unresolved symbols, operator applicability, etc.)
- **Apply Constant Folding**
- **Calculate Node Addresses**
- **Calculate Stack Loads**



IR Generation

Linearization

Multi-pass AST

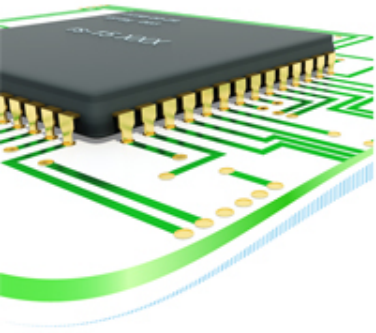


Control Flow Structure

- Apply Code Templates to Emit IC Instructions

Expression Trees

- Traverse and Emit IC Instructions



IR Generation

Intermediate Code

Purpose

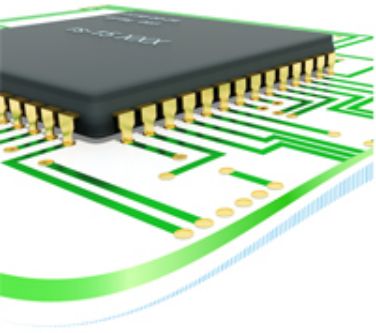
- **Architecture Independence (!)**
- **Standard / Improved Target Code Quality**
- **Reducing Cost of Overall Implementation**

Basic Types

- **Stack Machine Code**
- **Three Address Code**

Examples

- **Gimple / LLVM / JVM**

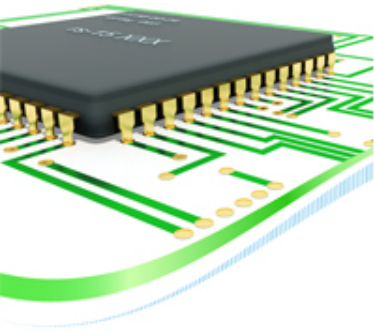


IC Generation

Example: Gimple

```
float f( int a, int b, float x ) {  
    float y = a*x*x + b*x + 100;  
    return y;  
}
```

```
f (int a, int b, float x)  
{  
    float D.1597D.1597;  
    float D.1598D.1598;  
    float D.1599D.1599;  
    float D.1600D.1600;  
    float D.1601D.1601;  
    float D.1602D.1602;  
    float D.1603D.1603;  
    float y;  
  
    D.1597D.1597 = (float) a;  
    D.1598D.1598 = D.1597D.1597 * x;  
    D.1599D.1599 = D.1598D.1598 * x;  
    D.1600D.1600 = (float) b;  
    D.1601D.1601 = D.1600D.1600 * x;  
    D.1602D.1602 = D.1599D.1599 + D.1601D.1601;  
    y = D.1602D.1602 + 1.0e+2;  
    D.1603D.1603 = y;  
    return D.1603D.1603;  
}
```



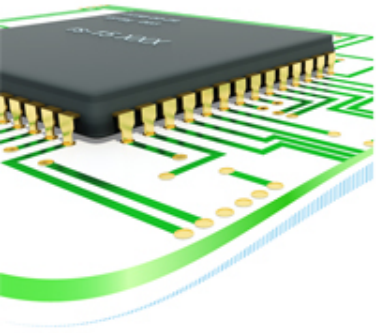
```
float f( int a, int b, float x ) {  
    float y = a*x*x + b*x + 100;  
    return y;  
}
```

IC Generation

Example: LLVM

```
define float @f(i32 %a, i32 %b, float %x) #0 {  
    %1 = alloca i32, align 4  
    %2 = alloca i32, align 4  
    %3 = alloca float, align 4  
    %y = alloca float, align 4  
    store i32 %a, i32* %1, align 4  
    store i32 %b, i32* %2, align 4  
    store float %x, float* %3, align 4  
    %4 = load i32* %1, align 4  
    %5 = sitofp i32 %4 to float  
    %6 = load float* %3, align 4  
    %7 = fmul float %5, %6  
    %8 = load float* %3, align 4  
    %9 = fmul float %7, %8  
    %10 = load i32* %2, align 4  
    %11 = sitofp i32 %10 to float  
    %12 = load float* %3, align 4  
    %13 = fmul float %11, %12  
    %14 = fadd float %9, %13  
    %15 = fadd float %14, 1.000000e+02  
    store float %15, float* %y, align 4  
    %16 = load float* %y, align 4  
    ret float %16  
}
```

Example from “Douglas Thain, Introduction to Compilers and Language Design, Chapter 8”, <http://compilerbook.org>

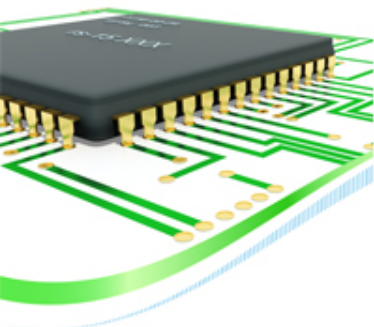


IC Generation

Example: JVM

```
float f( int a, int b, float x ) {  
    float y = a*x*x + b*x + 100;  
    return y;  
}
```

```
0: iload    1  
1: i2f  
2: fload    3  
4: fmul  
5: fload    3  
7: fmul  
8: iload    2  
9: i2f  
10: fload    3  
12: fmul  
13: fadd  
14: ldc      #2  
16: fadd  
17: fstore   4  
19: fload    4  
21: freturn
```



IC Generation

Example: Proprietary

```
float f(int32 a, int32 b, float x)
begin
    float y=a*x*x + b*x + 100;

    return y;
end
```

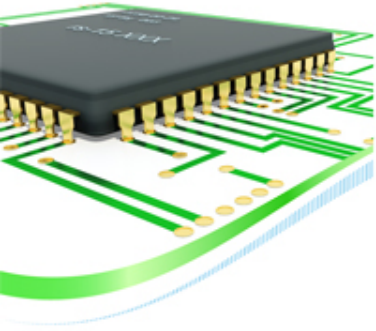
```
0 ssr 1[1 0x00000001]
1 ssr 29[52 0x00000034]
2 psh mwp base pointer offset 0
3 pmw base pointer offset -5
4 pmw base pointer offset -3
5 cvt_F_i -2 regs
6 mul_F
7 pmw base pointer offset -3
8 mul_F
9 pmw base pointer offset -4
10 pmw base pointer offset -3
11 cvt_F_i -2 regs
12 mul_F
13 add_F
14 psh int8 100 0x64
15 cvt_F_t -1 regs
16 add_F
17 ssr 4[0 0x00000000]
18 ssr 3[1 0x00000001]
19 ssr 29[89 0x00000059]
20 pmw base pointer offset 0
21 ssr 5[1 0x00000001]
22 ssr 2[1 0x00000001]
23 rtf 3
24 hlt
```

```
; Prologue f S0n(F0ni0ni0nF0n)
; Debug expression prologue y=a*x*x + b*x + 100;
```

```
; flat
```

```
; Debug expression prologue y;
```

```
; Epilogue f S0n(F0ni0ni0nF0n)
```



IC Generation

Multi-pass IR Traversal

Some Basics

- Multi-Pass IR Node Traversal
- Virtual instructions and types
- Control Flow Code Templates
- Type Calculation
- Address Calculation
- Optimizations
- More...

Some Advanced Processing

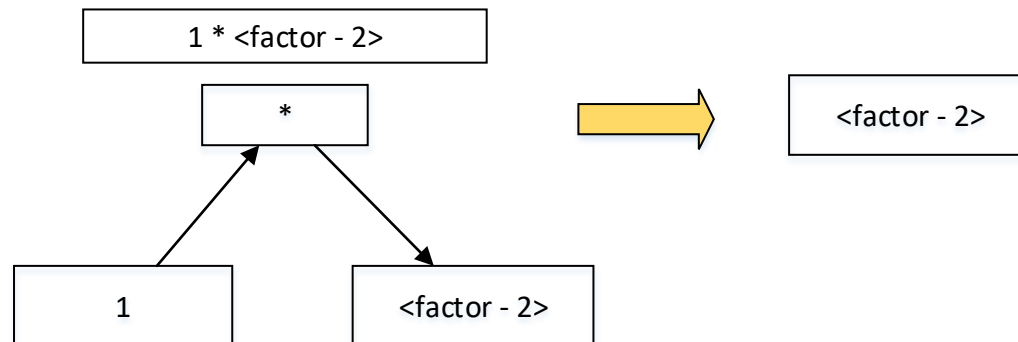
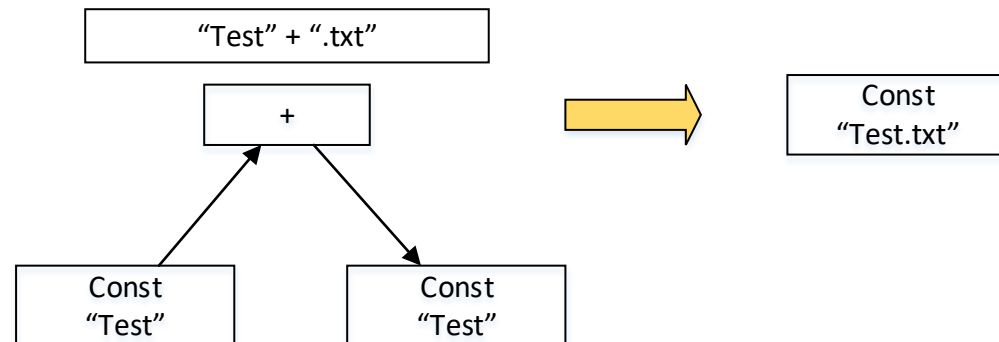
- Type System Operations
- Structured Exception Handling
- Closures
- Concurrency Related Patterns
- More ...

Simple Optimizations on IR

Constant Folding

Arithmetic Logic Operation Properties

- Constant operands
- Identity elements

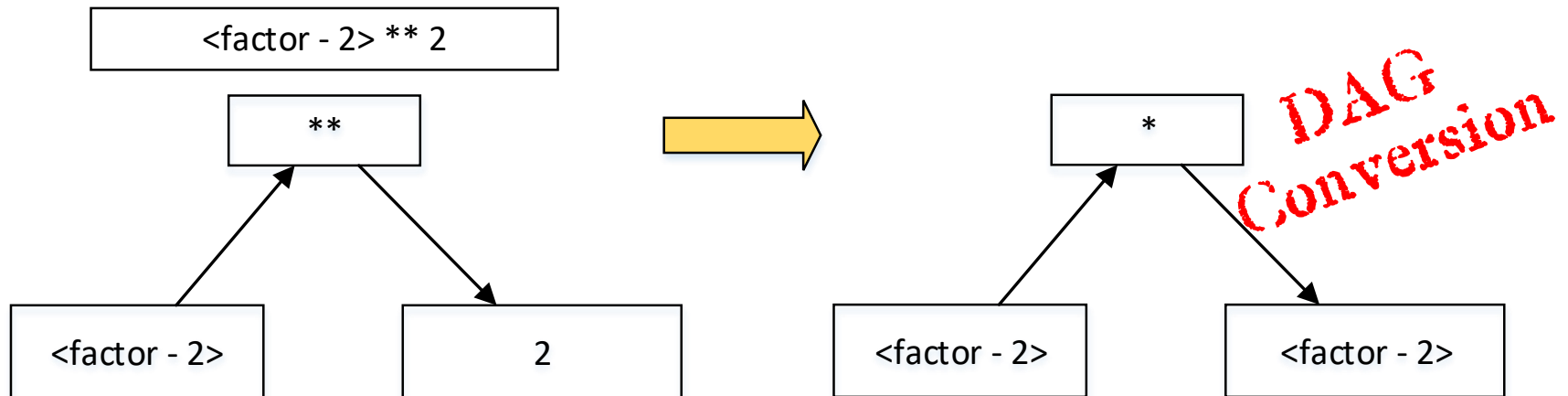


Simple Optimizations on IR

Strength Reduction

Arithmetic Logic Operation Properties

- Exponentials with constants
- Multiplications with constant
- More ...

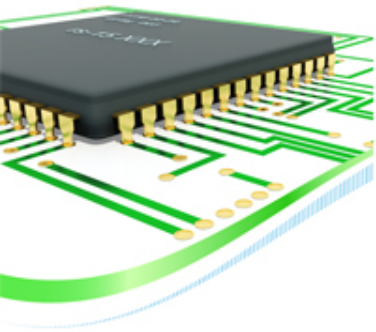


These properties may have impact on instruction selection, too.



Simple Optimizations on IR

Invariant Code Motion – Code Hoisting



```
let k=0,  
    m=0,  
    x=random(),  
    y=random();  
  
while (m<1000)  
{  
    delete(num2str(x*y)+".txt");  
    m=m+1;  
}
```

```
00004 setv      1 type:[]  
00005 pop      1 type:[]  
00006 call     22 type:[number]  
00007 setv     2 type:[]  
00008 pop      1 type:[]  
00009 call     22 type:[number]  
00010 setv     3 type:[]  
00011 pop      1 type:[]  
00012 id       1 type:[number]  
00013 const    0 type:[number] 1000  
00014 lt       0 type:[number]  
00015 pop      1 type:[]  
00016 jf       31 type:[]  
00017 id       2 type:[number]  
00018 id       3 type:[number]  
00019 mul      0 type:[number]  
00020 call     1 type:[string]  
00021 const    0 type:[string] ".txt"  
00022 add      0 type:[string]  
00023 call     18 type:[string]  
00024 pop      1 type:[]  
00025 id       1 type:[number]  
00026 const    0 type:[number] 1  
00027 add      0 type:[number]  
00028 asn      1 type:[number]  
00029 pop      1 type:[]  
00030 jmp      12 type:[]
```

Move

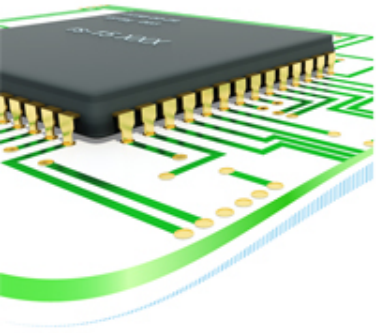
Example from the TQL assignment





Simple Optimizations on IR

Dead Code Elimination



```
let m=0,  
    x=0,  
    y=random();
```

```
x=2*m*y/(m+2);  
y=y/3;  
x=7/2;
```

00000	const	0	type:[number]	0
00001	setv	0	type:[]	
00002	pop	1	type:[]	
00003	const	0	type:[number]	0
00004	setv	1	type:[]	
00005	pop	1	type:[]	
00006	call	22	type:[number]	
00007	setv	2	type:[]	
00008	pop	1	type:[]	
00009	const	0	type:[number]	2
00010	id	0	type:[number]	
00011	mul	0	type:[number]	
00012	id	2	type:[number]	
00013	mul	0	type:[number]	
00014	id	0	type:[number]	
00015	const	0	type:[number]	2
00016	add	0	type:[number]	
00017	div	0	type:[number]	
00018	asn	1	type:[number]	
00019	pop	1	type:[]	
00020	id	2	type:[number]	
00021	const	0	type:[number]	3
00022	/	0	type:[number]	
00023	asn	2	type:[number]	
00024	pop	1	type:[]	
00025	const	0	type:[number]	3.5
00026	asn	1	type:[number]	
00027	pop	1	type:[]	

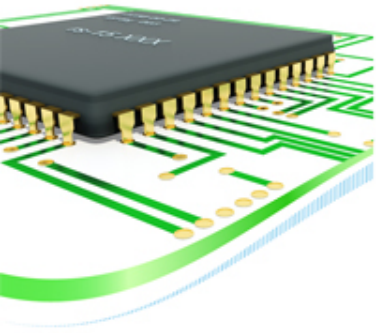
**Not
Effective**

Example from the TQL assignment



Simple Optimizations on IR

Peephole Optimization



```
let m=0,  
    x=0,  
    y=random();  
  
x=2*m*y/(m+2);  
y=x+1;
```

00000	const	0	type:[number]	0
00001	setv	0	type:[]	
00002	pop	1	type:[]	
00003	const	0	type:[number]	0
00004	setv	1	type:[]	
00005	pop	1	type:[]	
00006	call	22	type:[number]	
00007	setv	2	type:[]	
00008	pop	1	type:[]	
00009	const	0	type:[number]	2
00010	id	0	type:[number]	
00011	mul	0	type:[number]	
00012	id	2	type:[number]	
00013	mul	0	type:[number]	
00014	id	0	type:[number]	
00015	const	0	type:[number]	2
00016	add	0	type:[number]	
00017	/	0	type:[number]	
00018	asn	1	type:[number]	
00019	pop	1	type:[]	
00020	id	1	type:[number]	
00021	const	0	type:[number]	1
00022	add	0	type:[number]	
00023	asn	2	type:[number]	
00024	pop	1	type:[]	
00025	ret	0	type:[]	

**Not
Effective**

Example from the TQL assignment