

# MATHEMATICAL PROOF

LECTURE 1 JAN 7

1A

## Mathematical Proofs (Question)

- In mathematics, a proof is
- A. Similar to a scientific experiment. *Hypothesis* *Data to support*
  - B. A preponderance of evidence for the truth of a statement.
  - C. A plausible argument that a statement is true.
  - D. An evaluation of a boolean-valued expression. *{could be}*
  - E. None of the above.

Plausibility isn't enough for a statement to be true

## Writing Traditional Proofs (Question)

- To learn how to write traditional proofs, you should
- A. Read proofs given in mathematics articles and textbooks.
  - B. Write proofs of facts you already know.
  - C. Translate formal proofs into traditional proofs.
  - D. Do exercises that ask for a proof.

Mathematical Proof? argument showing that a conclusion follows from a set of premises

Theorem Statement with proof

Conjecture could be true, no proof to support +

- Proofs are unique to math, they're means of establishing truth.  $\square$ .

- ↳ communicate ideas, vehicles
- ↳ certify a result is correct
- ↳ Discovery through proofs
- ↳ 'univerality' true no matter what

Other fields establish truth through other means  
i.e experiments

## Styles of Mathematical Proof

→ description, describe the deduction

→ Prescription, Prescribe how to produce the proof; reasoning involved.

→ two-col, left statements (intermediate), right reasons (justifications)

→ computation

→ construction

→ geometrically, rules of geometry (visual proof)

→ 'classical', non constructive reasoning, proof it exists w/o product

→ 'constructive' uses only constructive reasoning must have product

## Traditional Proof Style:

- natural language (expression, ambiguous notation)
- may contain gaps (reader fills these)
- key ideas are understandable
- low-level details are computed or left for reader  
good for communication
  - ↳ organization
  - ↳ discovery / beauty
- found in txt books, research papers

## Formal Proof Style:

- clear rules, constrained
- described derivation
  - or
  - prescribed to create derivation
- Software can be implemented
- details must be correct
- not effective at communication but likely its correct

## Proving a Conjunction (Question)

Which is not a valid way to prove  $A \wedge B$ ?

- A. Prove  $A$  and  $B$  separately.
- B. Prove  $A$  and  $B$  together.
- C. Prove  $A$  and then prove  $B$  assuming  $A$ .
- D. Prove  $A$  assuming  $B$  and prove  $B$  assuming  $A$ .

C. is okay since we use  $A$  which is true to prove  $B$

D. is invalid since you're basing your proof on an unproven assumption.

# METHODS OF PROOF FOR PROPOSITIONAL FORMULA

implication  $A \Rightarrow B$  to prove. We assume  $A$  to be true and use that to prove  $B$

negation  $\neg A$  Assume  $A$ , derive a contradiction

conjunction  $A \wedge B$  Prove separately or together, Prove  $A$ , Assuming  $A$  Prove  $B$

disjunction  $A \vee B$  Prove one or the other, Assume  $\neg A$  then prove  $B$ , vice versa

iff  $A \Leftrightarrow B$  Prove  $A \Rightarrow B$ ,  $B \Rightarrow A$

# METHOD OF PROOF FOR QUANTIFIED FORMULAS

Universal stmt  $\forall x \in S \cdot A$

- Assume  $x \in S$ , Prove  $A$
- $\exists x \in S \cdot \neg A$ , derive a contradiction
- if  $S$  is finite, subs each member and prove  $A$  holds for each
- Induction

Existential stmt  $\exists x \in S \cdot A$

- for some ' $a$ '  $\in S$ ,  $A[x := a]$
- Assume  $\forall x \in S \cdot \neg A$ , derive a contradiction

# THEOREM TERMINOLOGY

theorem statement result of proof, fundamental importance

axiom theorem with assumed truth

lemma theorem to help prove other theorems

corollary theorem of fundamental importance

proposition theorem that's obviously true

# PROOF TERMINOLOGY

iff logical equivalence  $a \Leftrightarrow b$

'brute force'

verify each case individually

obvious no thinking

'symmetric argument'

structure preserving transformation from another argument

clearly result is verified w/ little effort

'well-defined'

definition is fully given

trivial extremely simple

'TFAE'

list of  $\equiv$  statements

straight-forward every step is obvious

Without loss of generality

'WLOG', Prove special case is ok

similar same structure/technique

'QED'

$\square$ ,  $\blacksquare$  end of proof