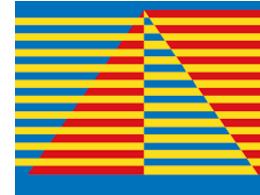


Introducing Fast and Secure Deterministic Stash Free

Write Only Oblivious RAMs for Demand Paging in Keystone

Workshop on Computer Architecture Research with RISC-V
(CARRV'21 @ISCA '21)



Mriganka Chakravarty
IIT Kanpur



Biswabandan Panda
IIT Bombay

Introducing Fast and Secure Deterministic
Stash Free

Write Only Oblivious RAMs for
Demand Paging in **Keystone**

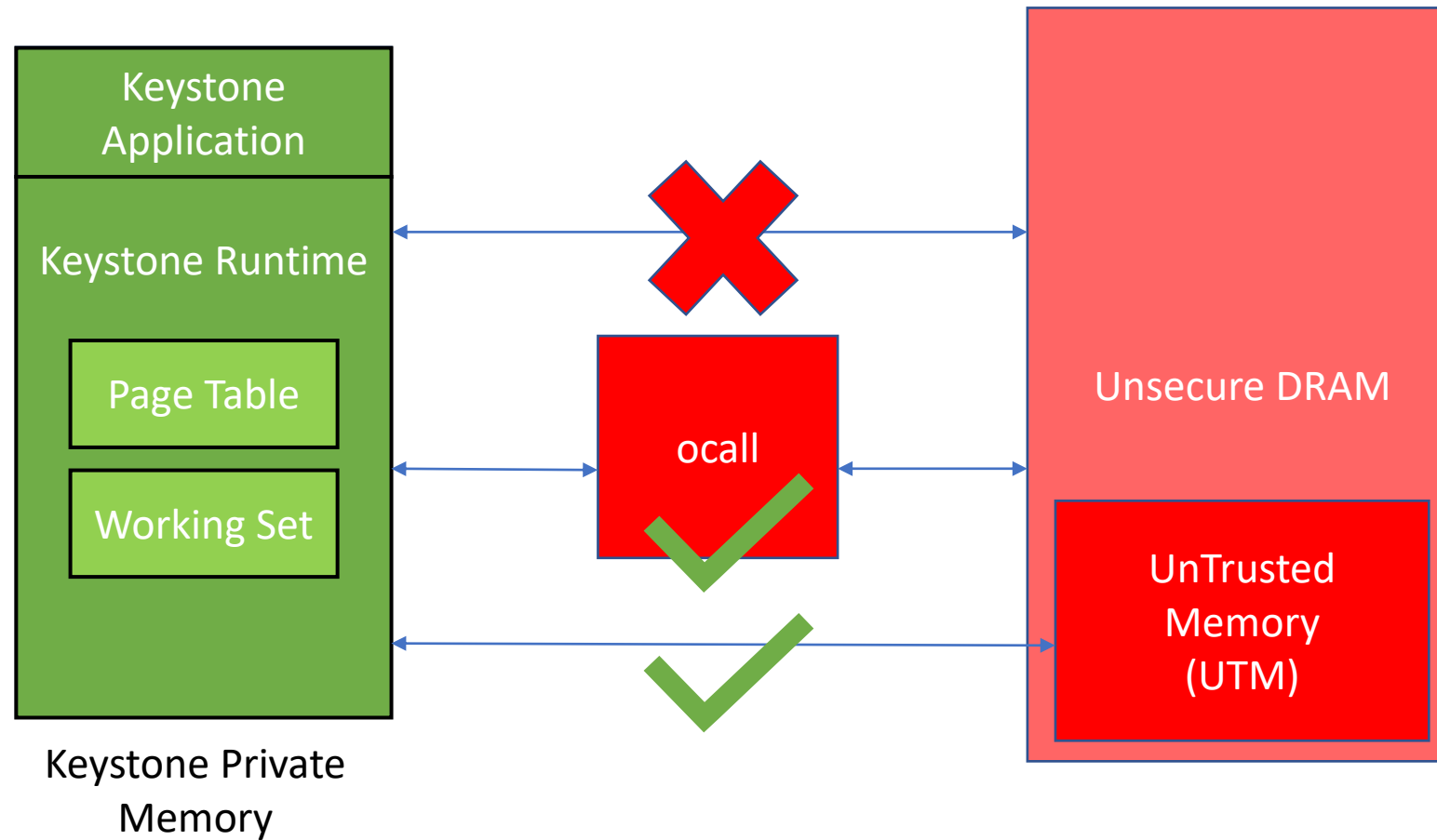
Keystone

A trusted execution environment based on RISC-V architecture.

Isolate memory into secure **Keystone private** and **Unsecure non-Keystone** memory.

Allows application to run securely in presence of privileged adversary.

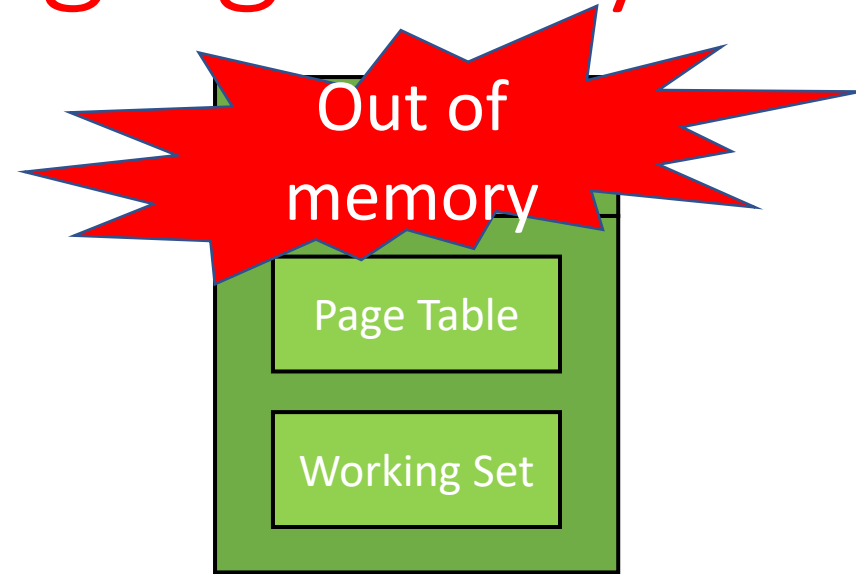
Components of Keystone



Introducing Fast and Secure Deterministic
Stash Free

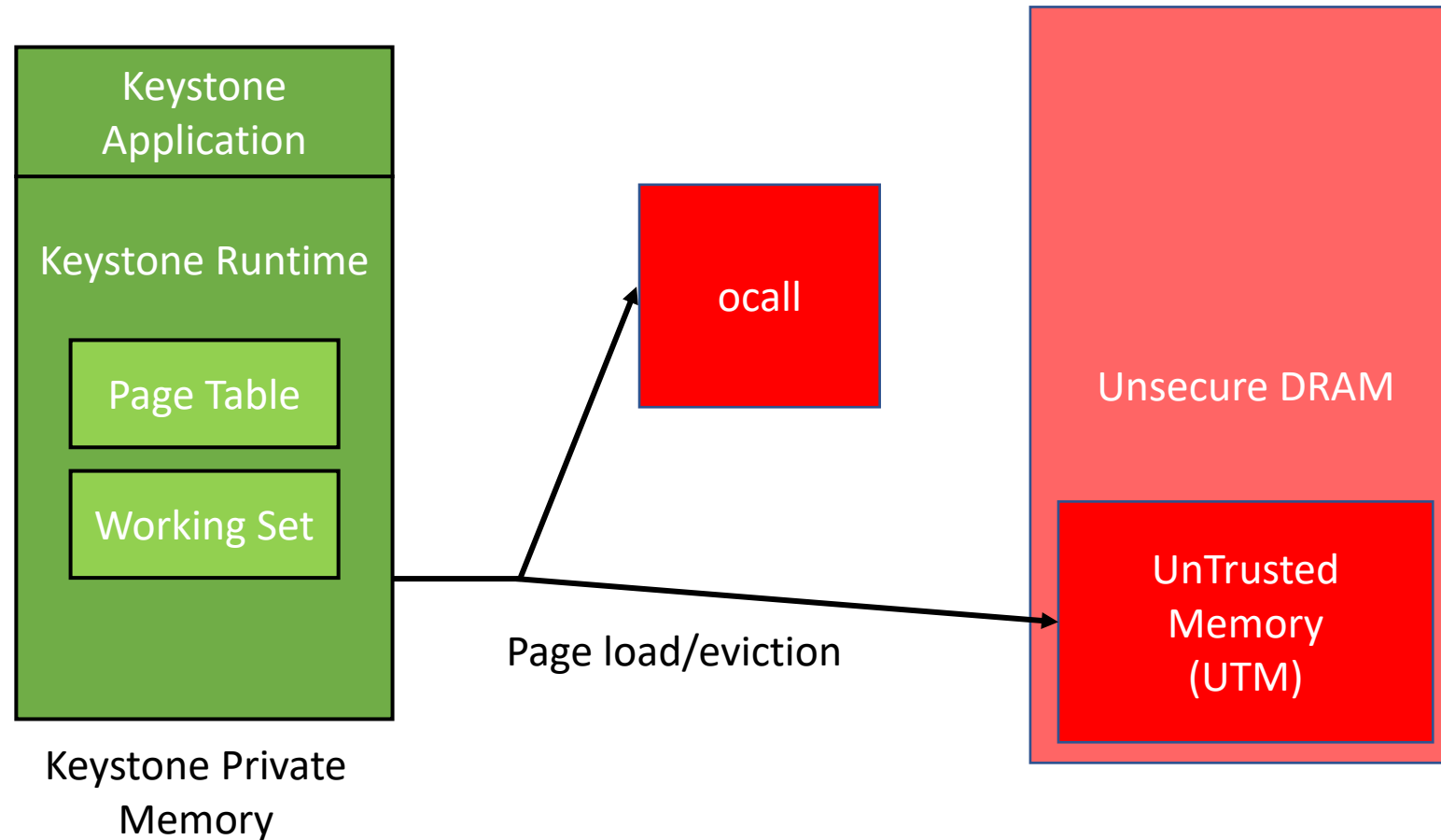
Write Only Oblivious RAMs for
Demand Paging in Keystone

Demand paging in Keystone



- To run large applications, we need demand paging.
- We utilize the unsecure non-Keystone memory as backing store.

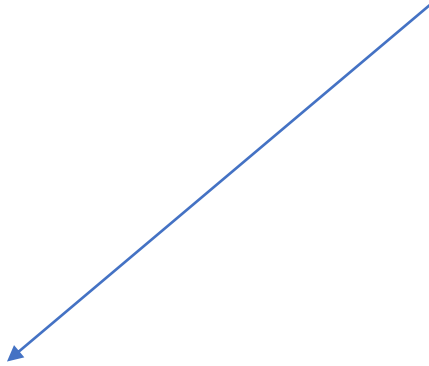
Demand paging in Keystone



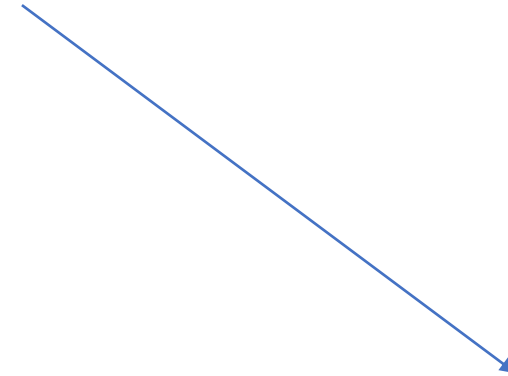
What makes life difficult?

- OS oversees the unsecure DRAM region.
- As we shall see, demand paging leaks access patterns.

Leakage due to demand paging



Page loads are visible



Page evictions are visible

Leakage due to Page Load

```
Page *p = (Page*)malloc(2*sizeof(Page));  
.  
.  
.  
int secret = input();  
  
if (secret == 1)  
    write(p + 0, 0);  
  
else  
    write(p + 1, 0);
```



Working Set

Leakage due to Page Load

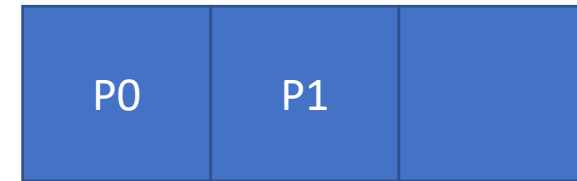
```
Page *p = (Page*)malloc(2*sizeof(Page));
```

```
·  
·  
·
```

```
int secret = input();
```

```
if (secret == 1)  
    write(p + 0, 0);
```

```
else  
    write(p + 1, 0);
```



Working Set

Leakage due to Page Load

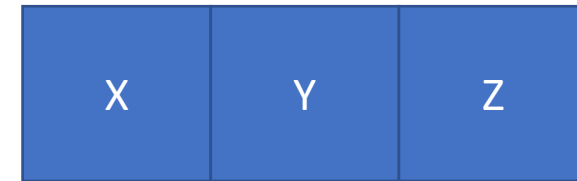
```
Page *p = (Page*)malloc(2*sizeof(Page));
```

```
•  
•  
•
```

```
int secret = input();
```

```
if (secret == 1)  
    write(p + 0, 0);
```

```
else  
    write(p + 1, 0);
```



Working Set

Leakage due to Page Load

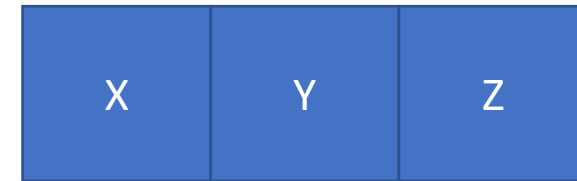
```
Page *p = (Page*)malloc(2*sizeof(Page));
```

```
·  
·  
·
```

```
int secret = input();
```

```
if (secret == 1)  
    write(p + 0, 0);
```

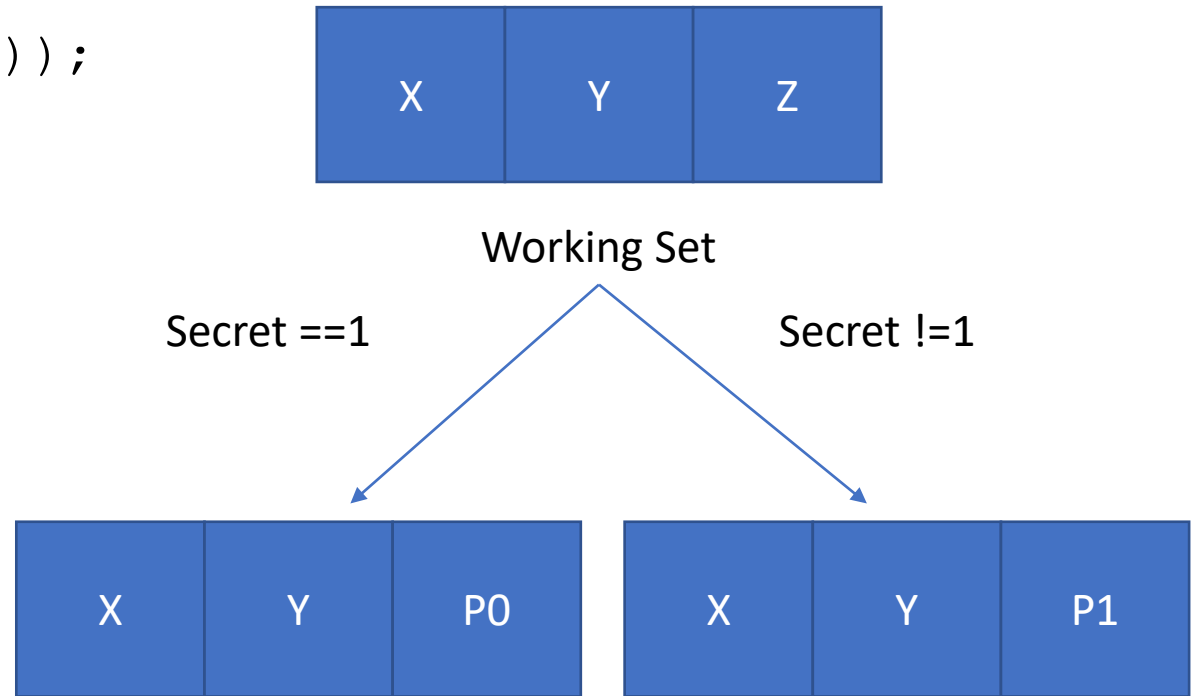
```
else  
    write(p + 1, 0);
```



Working Set

Leakage due to Page Load

```
Page *p = (Page*)malloc(2*sizeof(Page));  
.  
.  
.  
int secret = input();  
  
if (secret == 1)  
    write(p + 0, 0);  
  
else  
    write(p + 1, 0);
```



Checkmate!

Leakage due to page evictions

```
Page *P = (Page *)malloc(4*sizeof(Page));
```

```
.  
.   
.
```

```
int secret = input();
```

```
if (secret == 1){  
    write(P[1],0);  
    write(P[2],0);  
}
```

```
else{  
    write(P[2],0);  
    write(P[1],0);  
}
```

```
write(P[0],0);
```

```
write(P[3],0);
```

More Recent →



Working set
(LRU replacement policy)

Leakage due to page evictions

```
Page *P = (Page *)malloc(4*sizeof(Page));
```

```
.  
. .  
. .
```

```
int secret = input();
```

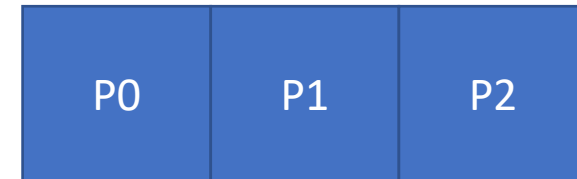
```
if (secret == 1){  
    write(P[1], 0);  
    write(P[2], 0);  
}
```

```
else{  
    write(P[2], 0);  
    write(P[1], 0);  
}
```

```
write(P[0], 0);
```

```
write(P[3], 0);
```

More Recent →



Leakage due to page evictions

```
Page *P = (Page *)malloc(4*sizeof(Page));
```

```
•
```

```
•
```

```
•
```

```
int secret = input();
```

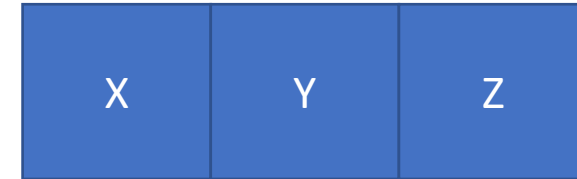
```
if (secret == 1){  
    write(P[1], 0);  
    write(P[2], 0);  
}
```

```
else{  
    write(P[2], 0);  
    write(P[1], 0);  
}
```

```
write(P[0], 0);
```

```
write(P[3], 0);
```

More Recent →



Leakage due to page evictions

```
Page *P = (Page *)malloc(4*sizeof(Page));
```

```
.  
.   
.
```

```
int secret = input();
```

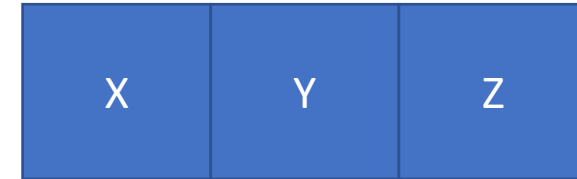
```
if (secret == 1){  
    write(P[1],0);  
    write(P[2],0);  
}
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```
else{  
    write(P[2],0);  
    write(P[1],0);  
}
```

```
write(P[0],0);
```

```
write(P[3],0);
```

More Recent →



Leakage due to page evictions

```
Page *P = (Page *)malloc(4*sizeof(Page));
```

```
.  
.   
.
```

```
int secret = input();
```

```
if (secret == 1){
```

```
    write(P[1],0);
```

```
    write(P[2],0);
```

```
}
```

```
else{
```

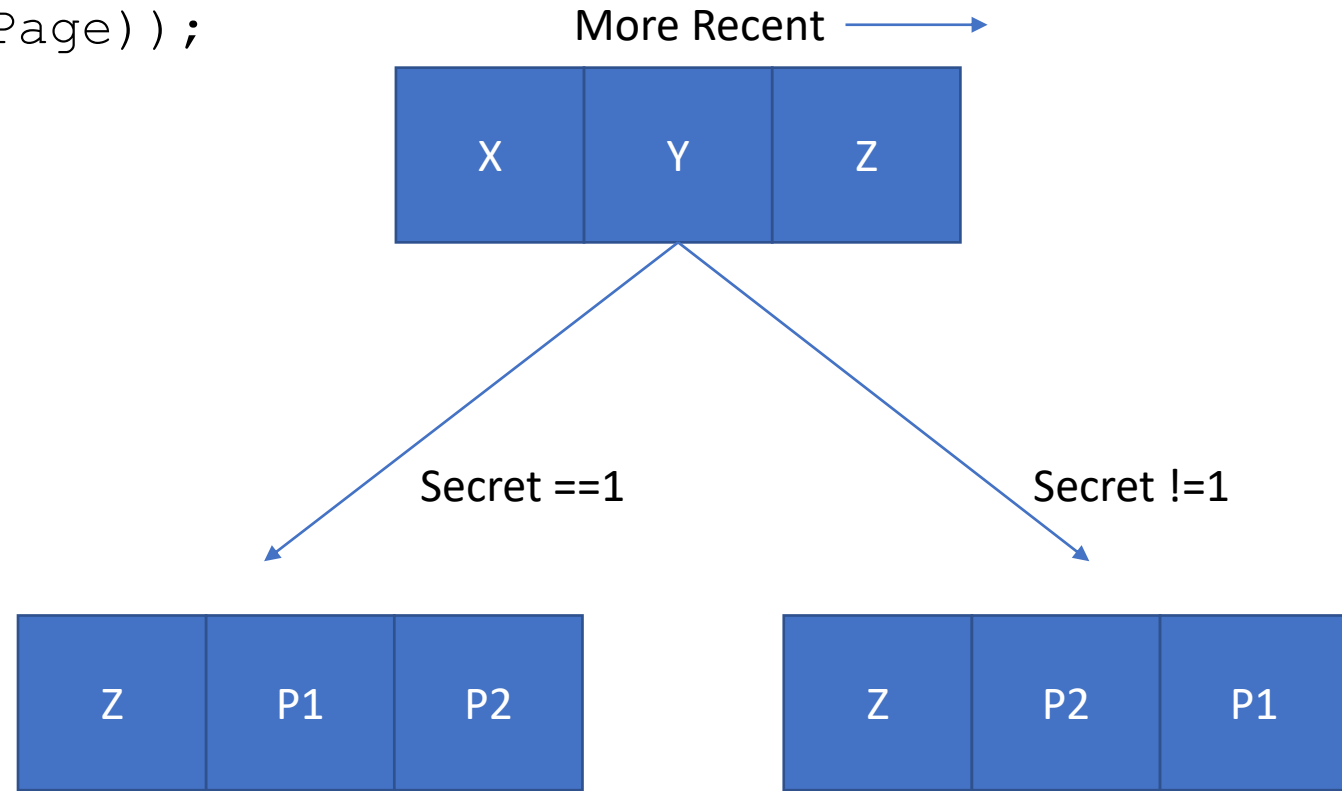
```
    write(P[2],0);
```

```
    write(P[1],0);
```

```
}
```

```
write(P[0],0);
```

```
write(P[3],0);
```



Leakage due to page evictions

```
Page *P = (Page *)malloc(4*sizeof(Page));
```

```
.  
.   
. 
```

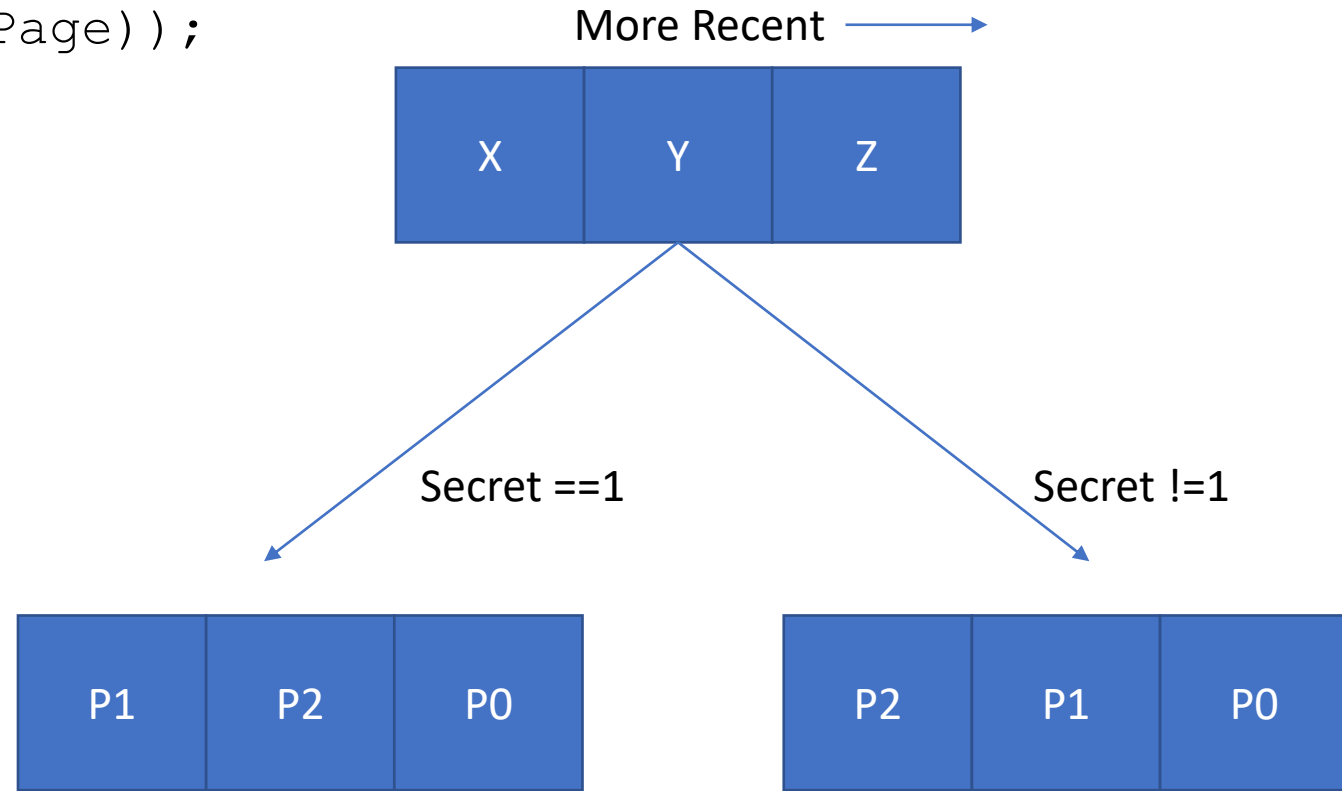
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if (secret == 1){  
    write(P[1],0);  
    write(P[2],0);  
}
```

```
else{  
    write(P[2],0);  
    write(P[1],0);  
}
```

```
write(P[0],0);
```

```
write(P[3],0);
```



Leakage due to page evictions

```
Page *P = (Page *)malloc(4*sizeof(Page));
```

```
.  
.   
. 
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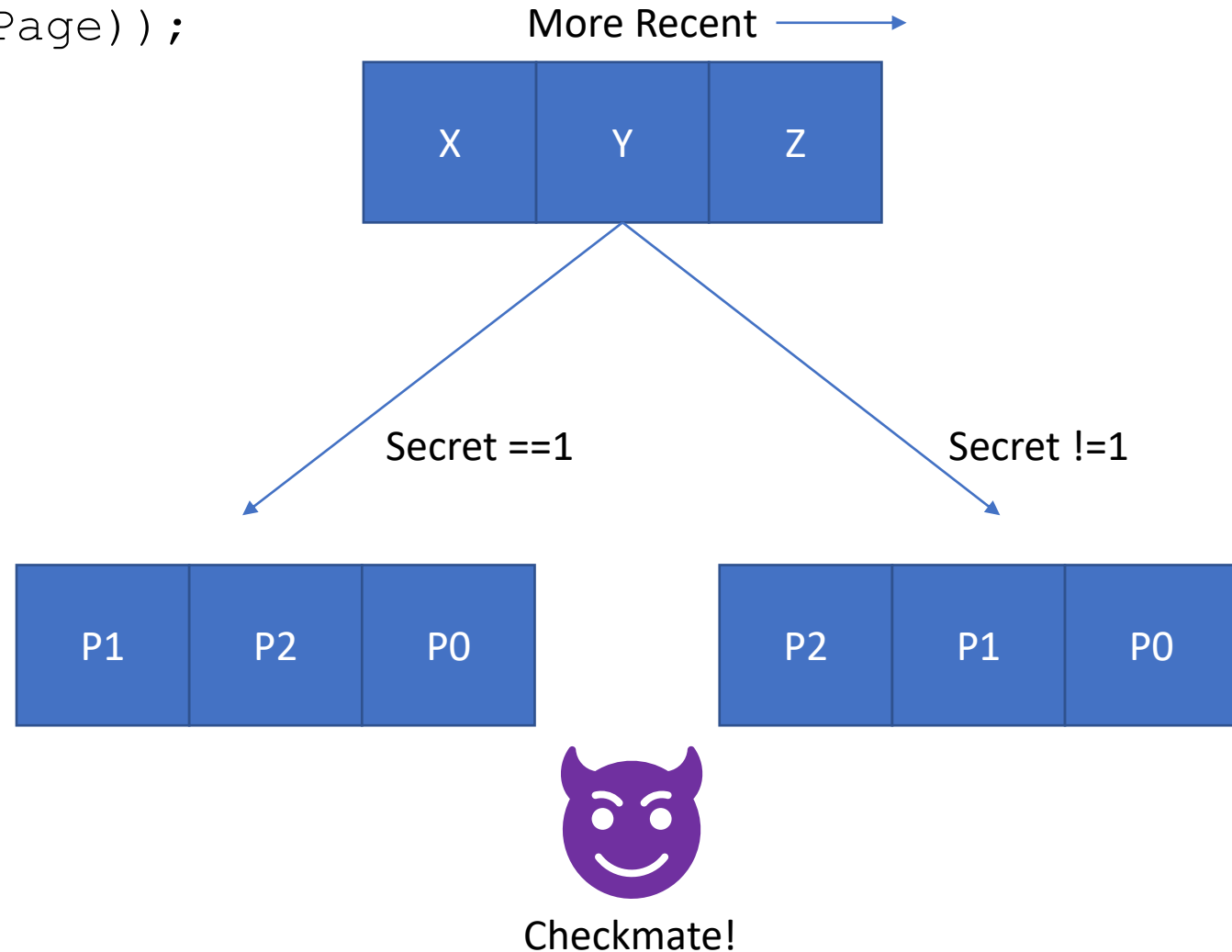
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    write(P[1],0);  
    write(P[2],0);  
}
```

```
else{  
    write(P[2],0);  
    write(P[1],0);  
}
```

```
write(P[0],0);
```

```
write(P[3],0);
```



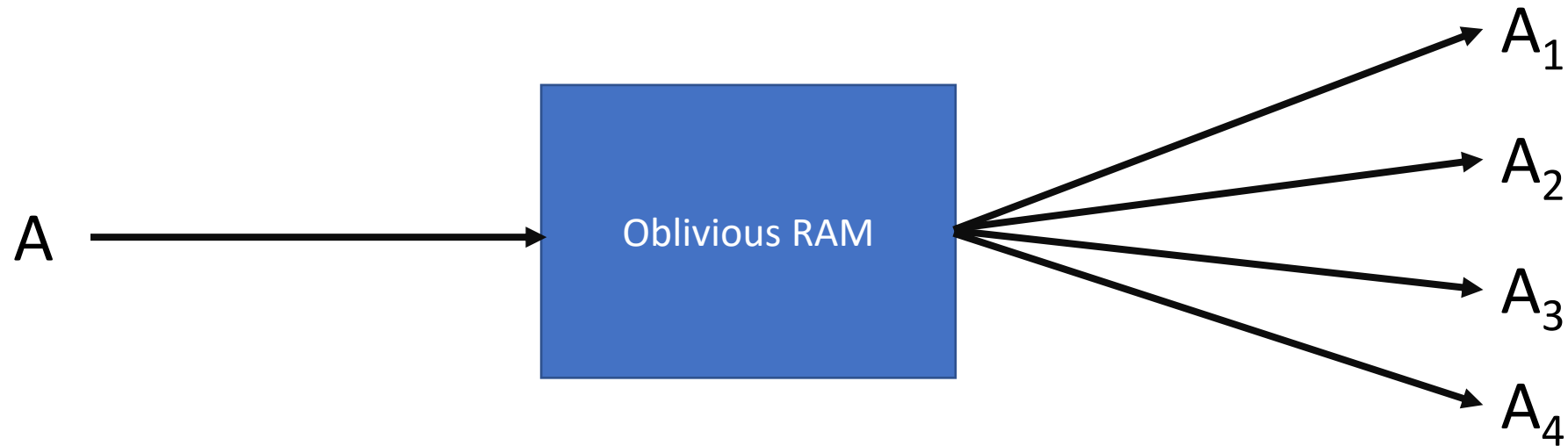
Solution?

Use Oblivious RAM for loading and evicting pages.

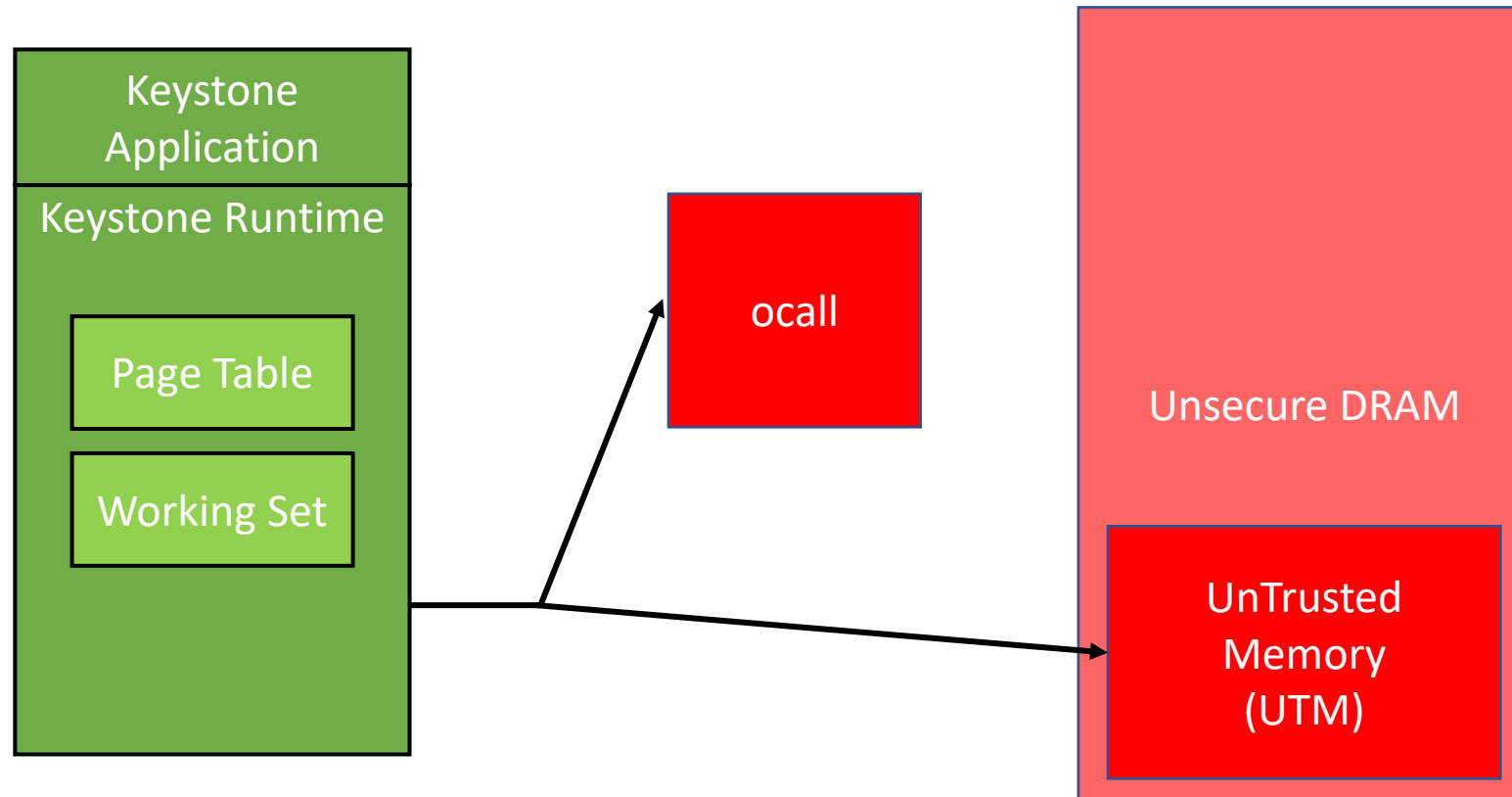
Introducing Fast and Secure Deterministic
Stash Free

Write Only **Oblivious RAMs** for
Demand Paging in Keystone

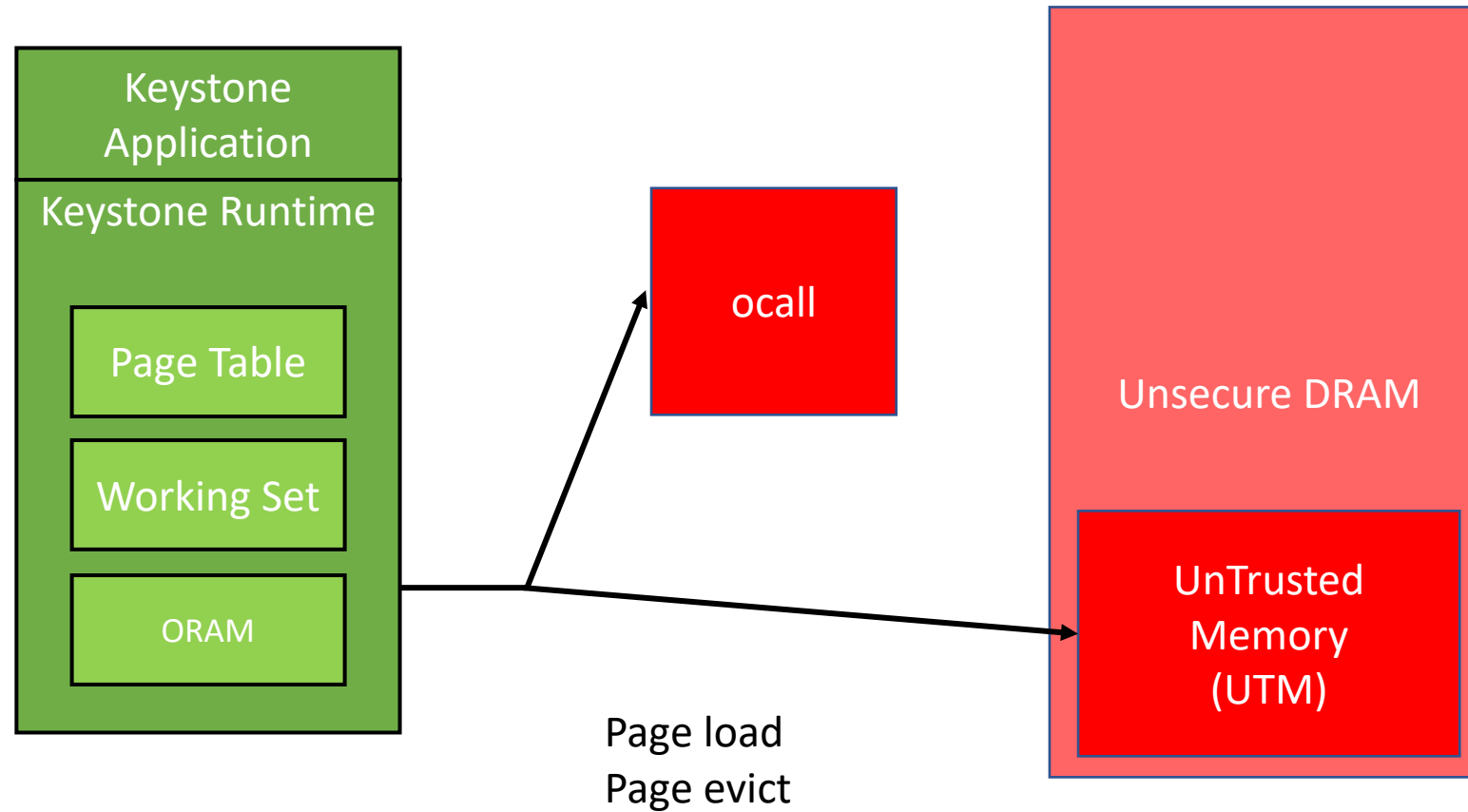
What is Oblivious RAM?



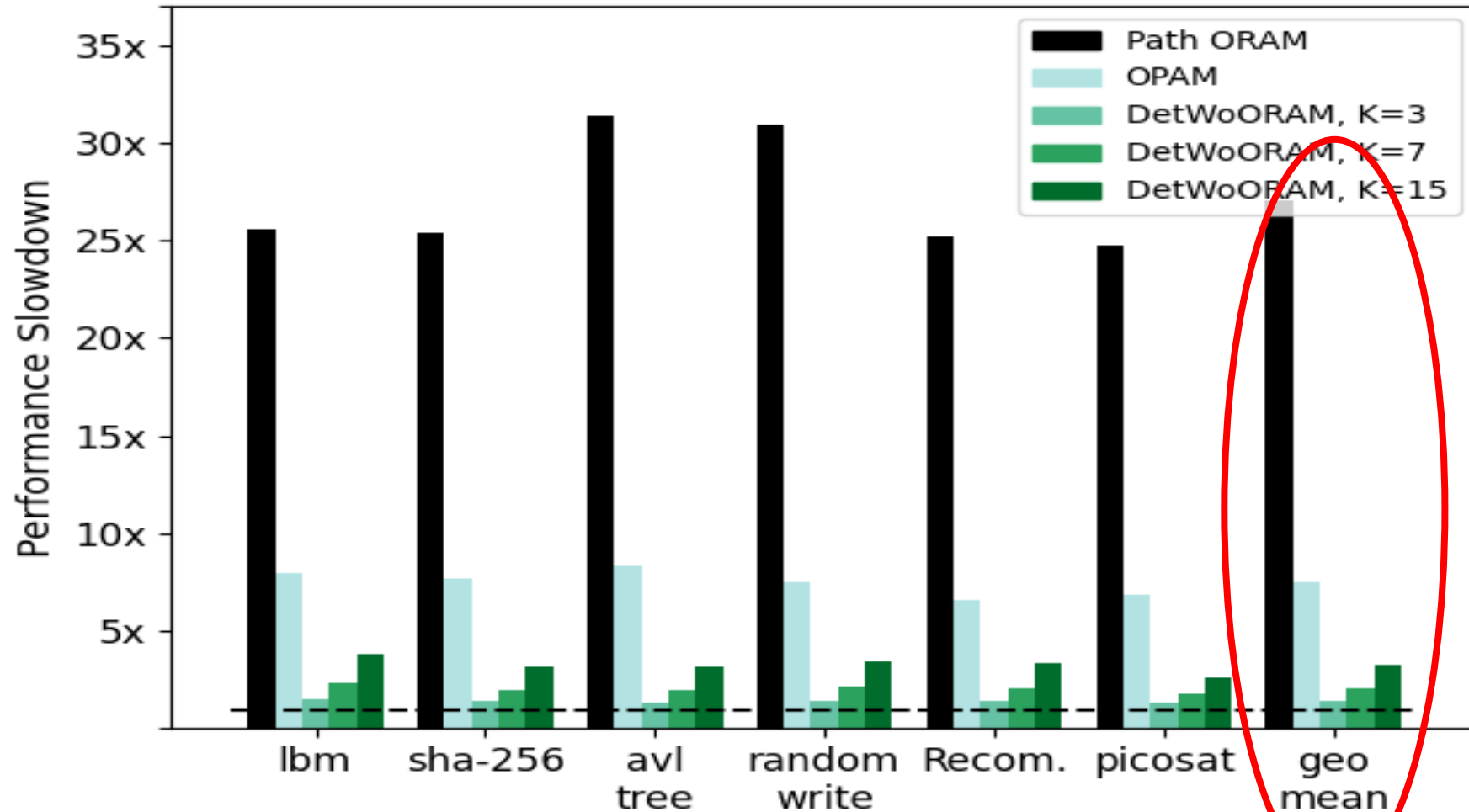
Oblivious Demand paging in Keystone



Oblivious Demand paging in Keystone



Motivation



Path ORAM
and
OPAM are
very slow.

DetWoORAM
is very fast.

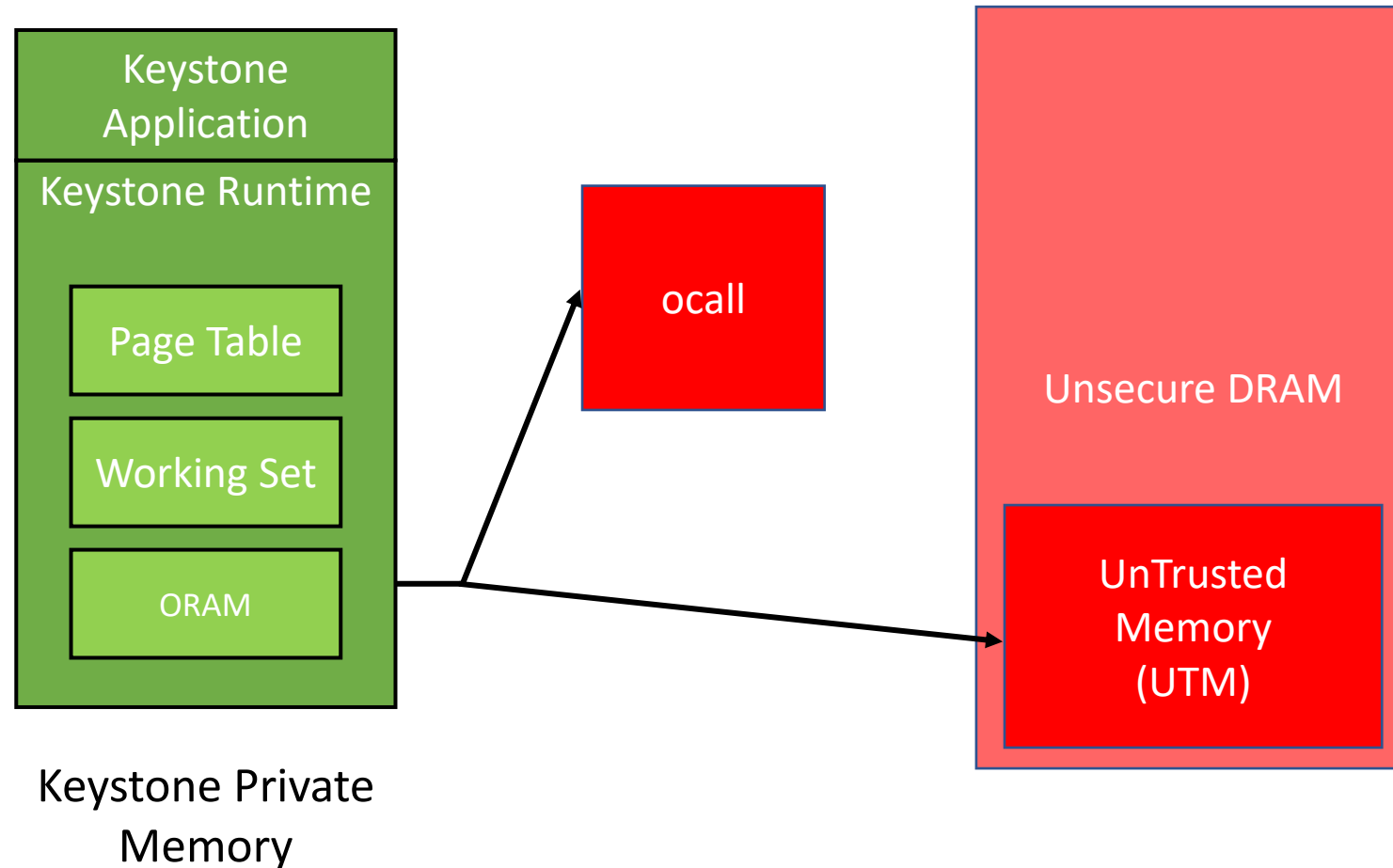
Introducing Fast and Secure Deterministic
Stash Free

Write Only Oblivious RAMs for
Demand Paging in Keystone

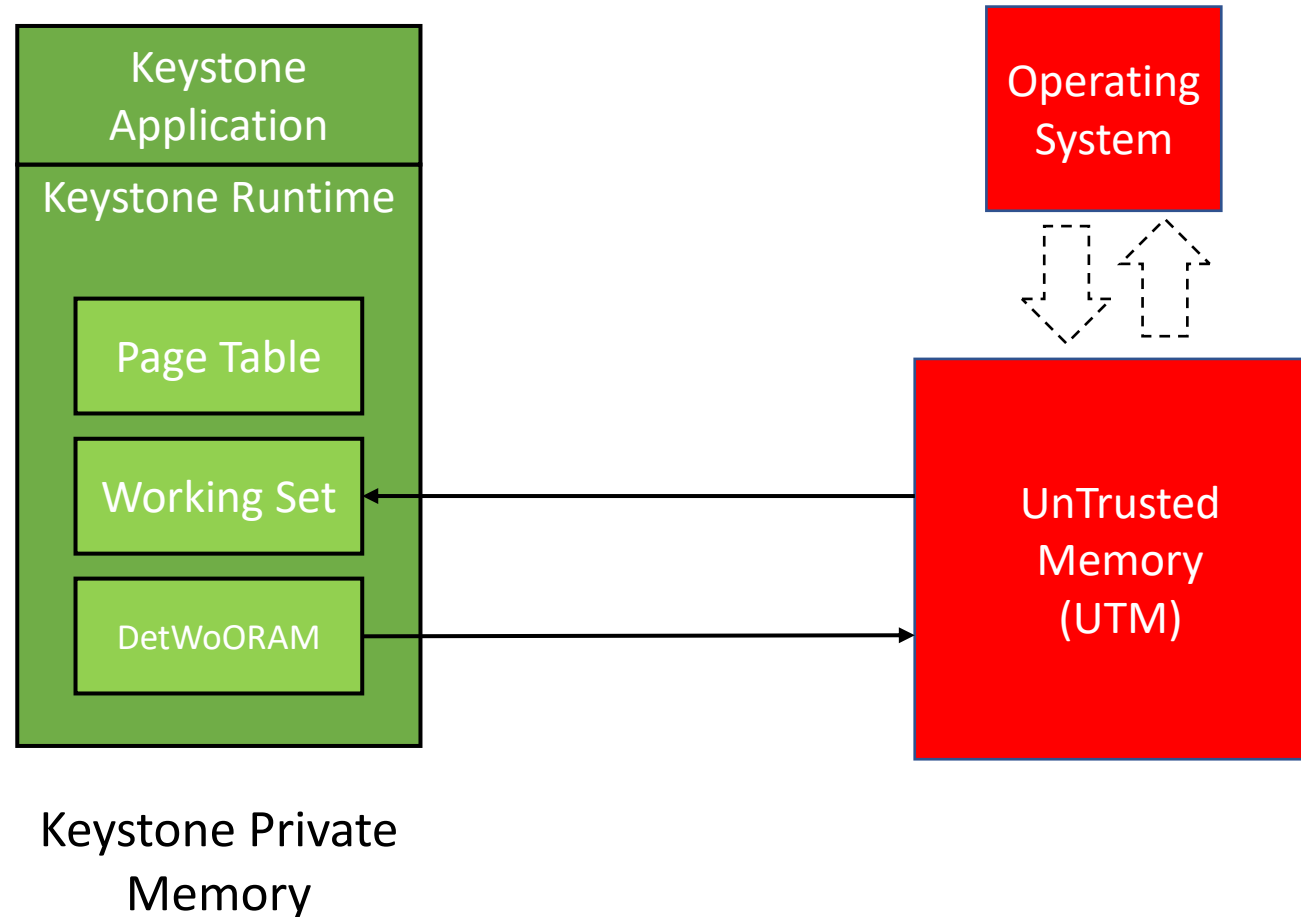
Write only ORAM

- WoORAM assumes that the reads are oblivious.
 - No one can see where a read is happening.
- But paging involves reading(loading) and writing(evicting).
 - So how do we assume that the reads are invisible?

Write Only ORAM for demand paging



Write Only ORAM for demand paging



Challenges

- The entire application should be able to fit into the UTM.
- We expect **future work** in increasing the UTM size.

Summary

- So far, we discussed how we can use DetWoORAM for demand paging.



Can we make
something good out
of this work?



Summary

- So far, we discussed how we can use DetWoORAM for demand paging.

That's not enough work



Can we make something good out of this work?



Summary

- So far, we discussed how we can use DetWoORAM for demand paging.

We need to
optimize further.



Can we make
something good out
of this work?



Summary

- So far, we discussed how we can use DetWoORAM for demand paging.

We need to
optimize further.



But... DetWoORAM is
algorithmically
optimal. 😊



Summary

- So far, we discussed how we can use DetWoDRAM for demand paging.

Can we go beyond
the algorithm? 🤔



But... DetWoDRAM is
algorithmically
optimal. 😊



Summary

- So far, we discussed how we can use DetWoORAM for demand paging.

Can we go beyond
the algorithm? 🤔



Got it! 😎





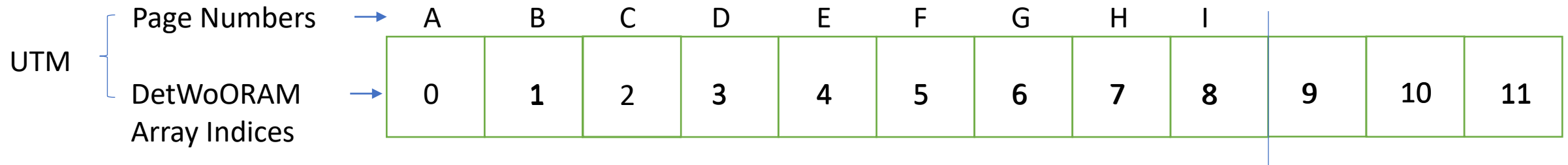
Lets shorten it as DetWoORAM

Introducing Fast and Secure Deterministic Stash Free Write Only Oblivious RAMs for Demand Paging in Keystone

Our Contribution

- We introduce enhancements to DetWoORAM, namely:
 - Eager DetWoORAM
 - Parallel DetWoORAM

Working of DetWoORAM*

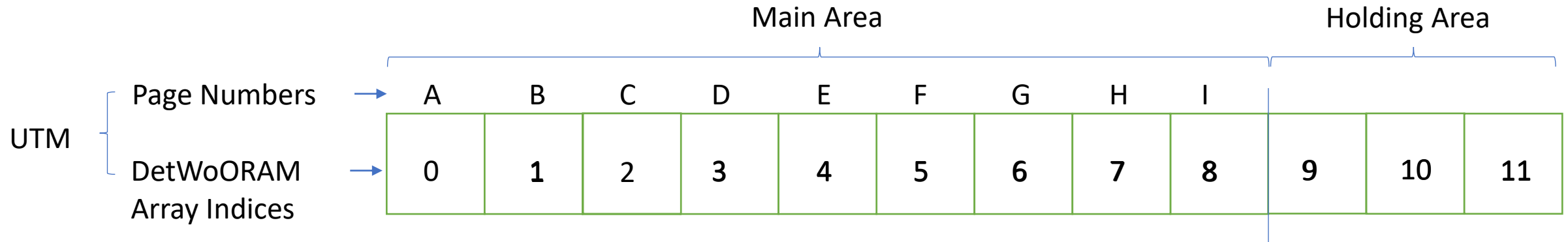


Position Map

A	0
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*

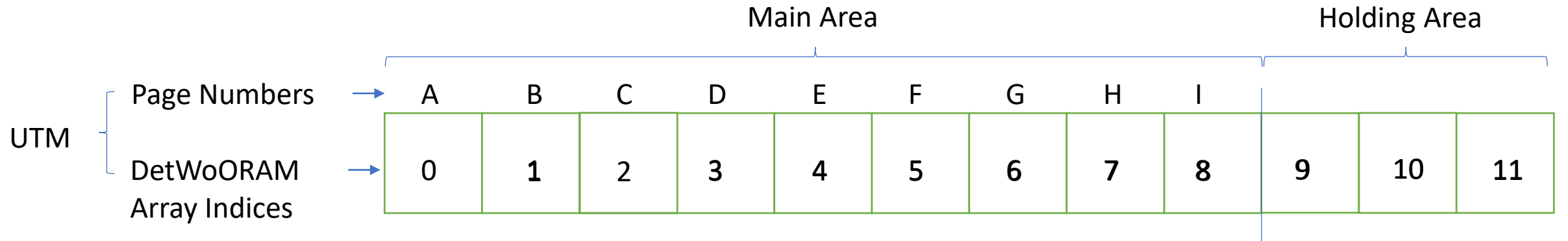


Position Map

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B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



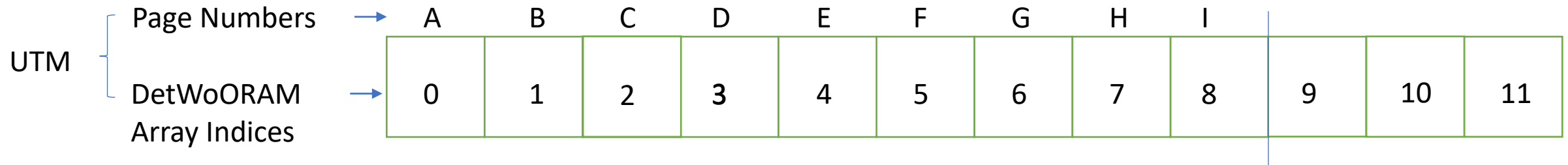
Position Map

A	0
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

$$K = \frac{\text{Main Area}}{\text{Holding Area}}$$

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Position Map

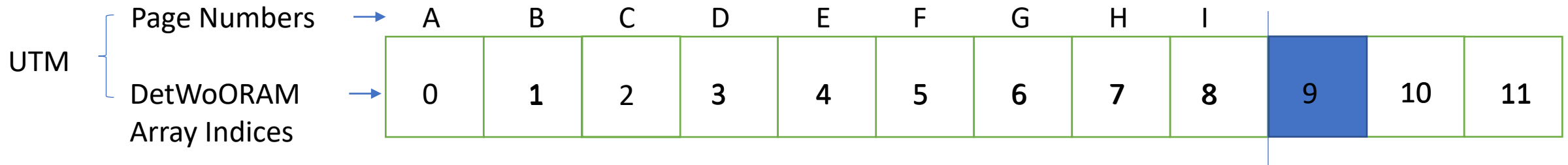
A	0
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

Operations

1. Write A

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

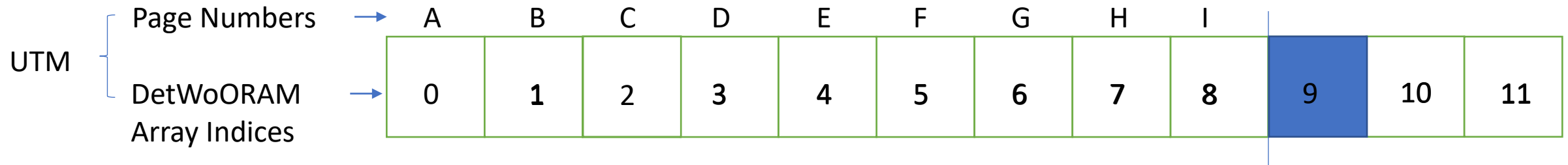
1. Write A

Position Map

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F	5
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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

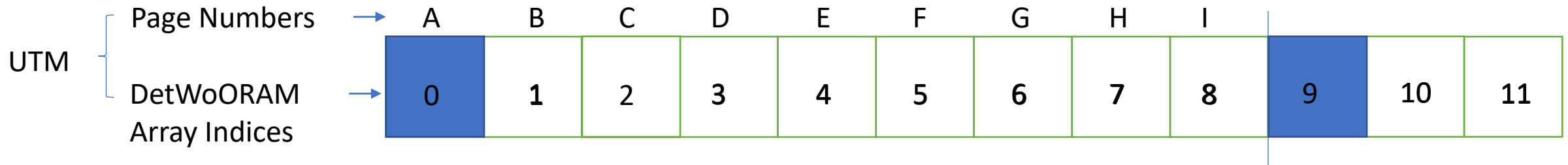
1. Write A

Position Map

A	9
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

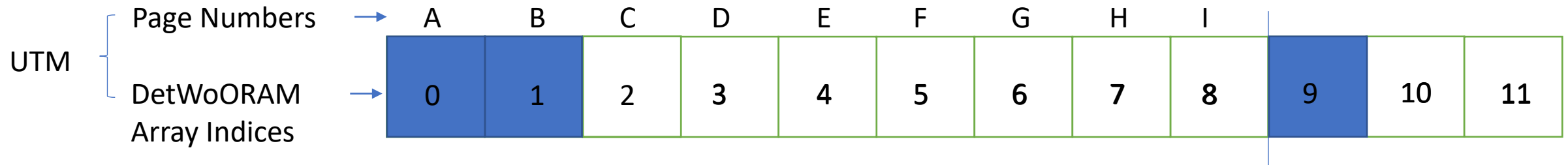
1. Write A

Position Map

A	0
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

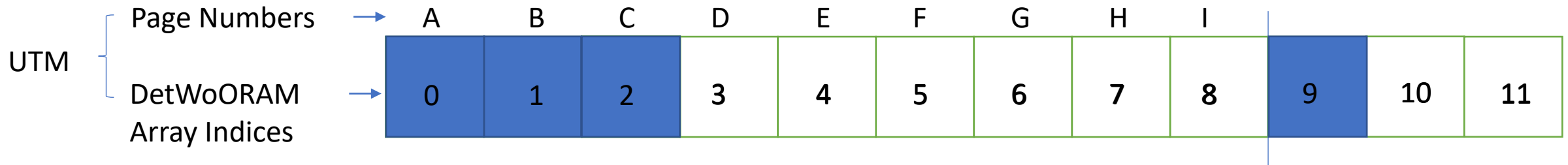
1. Write A

Position Map

A	0
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

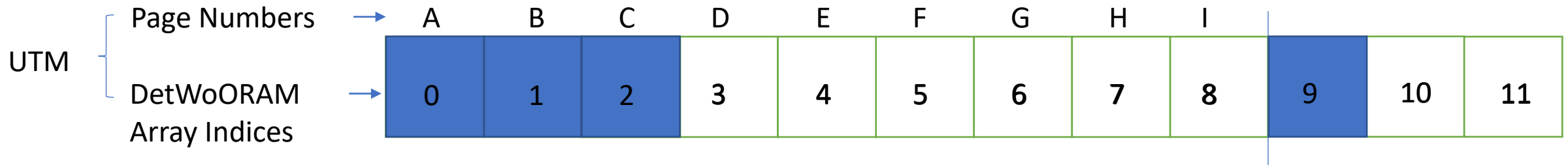
1. Write A

Position Map

A	0
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

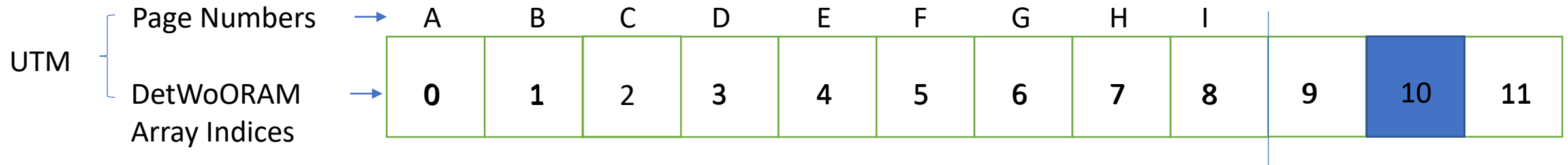
1. Write A
2. Write B

Position Map

A	0
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

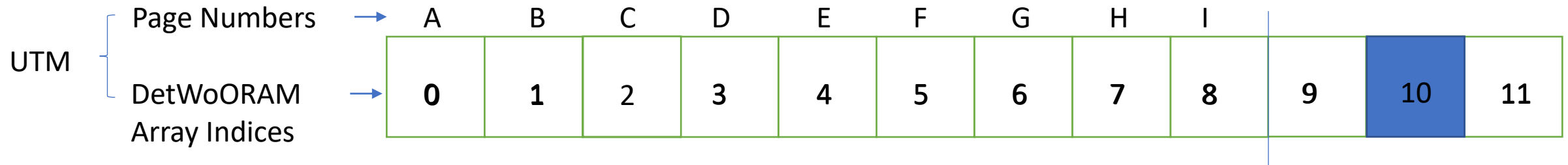
1. Write A
2. Write B

Position Map

A	0
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

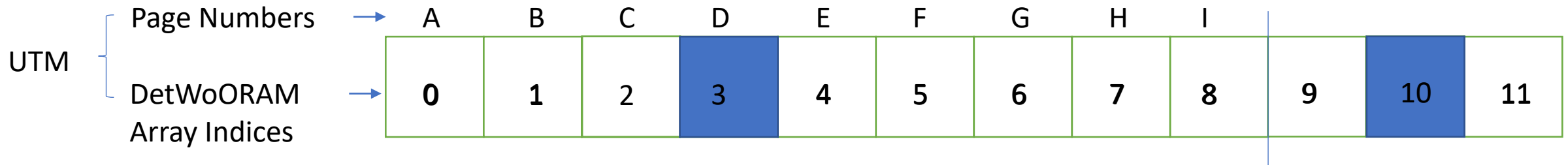
1. Write A
2. Write B

Position Map

A	0
B	10
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Position Map

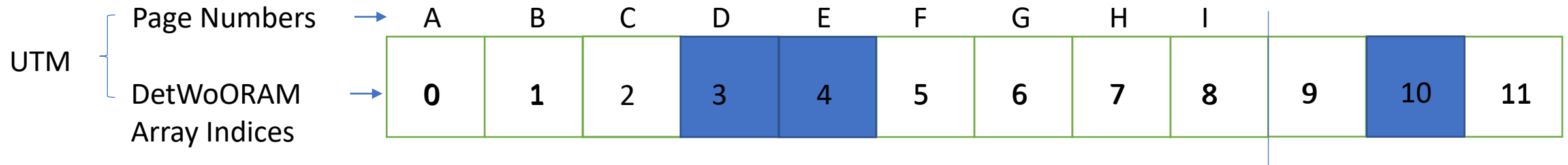
A	0
B	10
C	2
D	3
E	4
F	5
G	6
H	7
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Operations

1. Write A
2. Write B

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

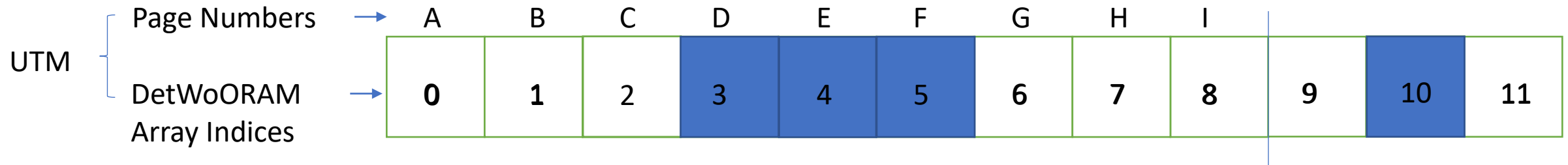
1. Write A
2. Write B

Position Map

A	0
B	10
C	2
D	3
E	4
F	5
G	6
H	7
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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

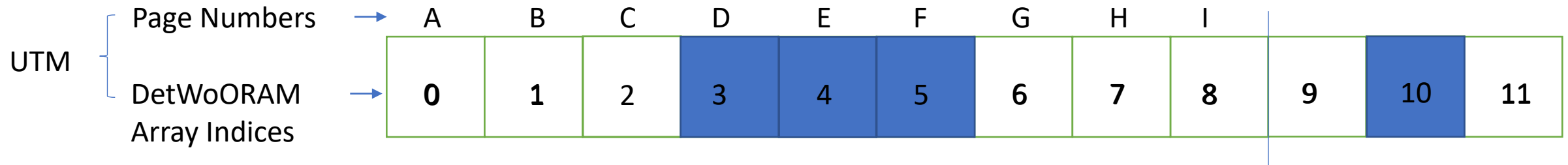
1. Write A
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Position Map

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B	10
C	2
D	3
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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

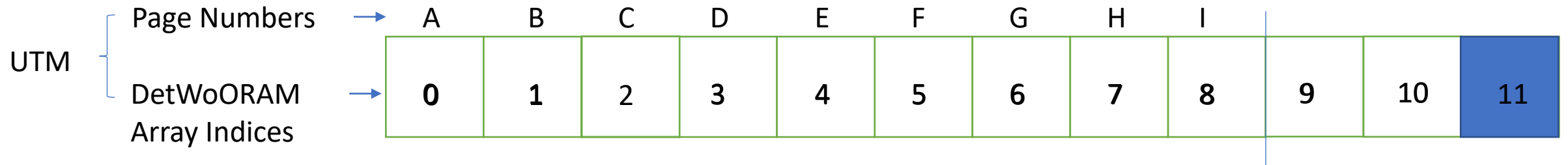
1. Write A
2. Write B
3. Write C

Position Map

A	0
B	10
C	2
D	3
E	4
F	5
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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Position Map

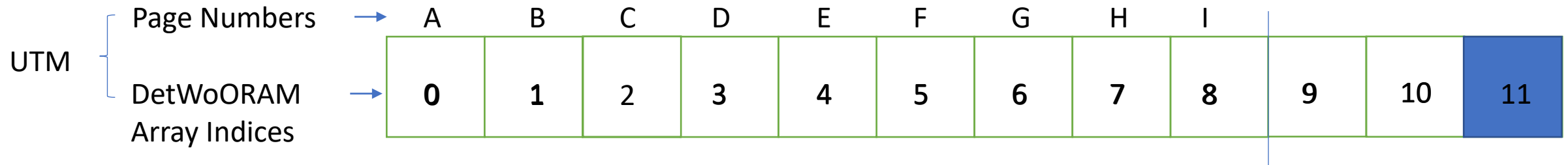
A	0
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E	4
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Operations

1. Write A
2. Write B
3. Write C

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Operations

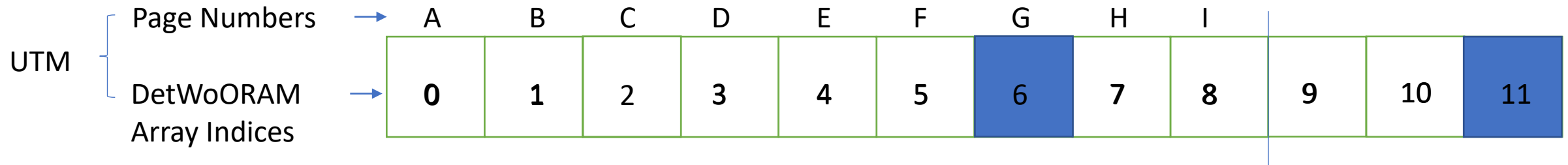
1. Write A
2. Write B
3. Write C

Position Map

A	0
B	10
C	11
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Position Map

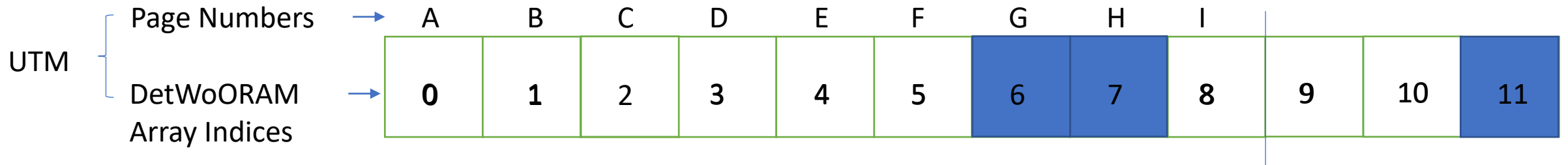
A	0
B	10
C	11
D	3
E	4
F	5
G	6
H	7
I	8

Operations

1. Write A
2. Write B
3. Write C

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



Position Map

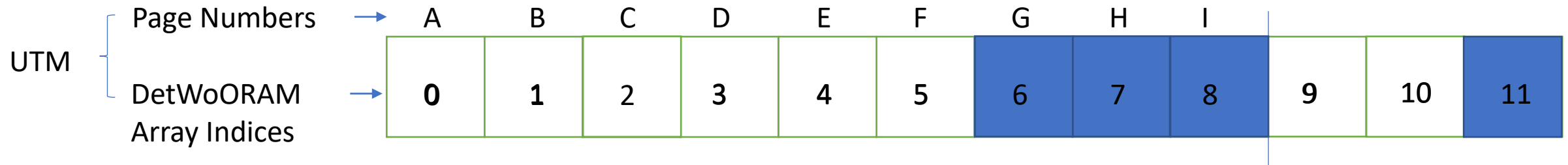
A	0
B	10
C	11
D	3
E	4
F	5
G	6
H	7
I	8

Operations

1. Write A
2. Write B
3. Write C

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of DetWoORAM*



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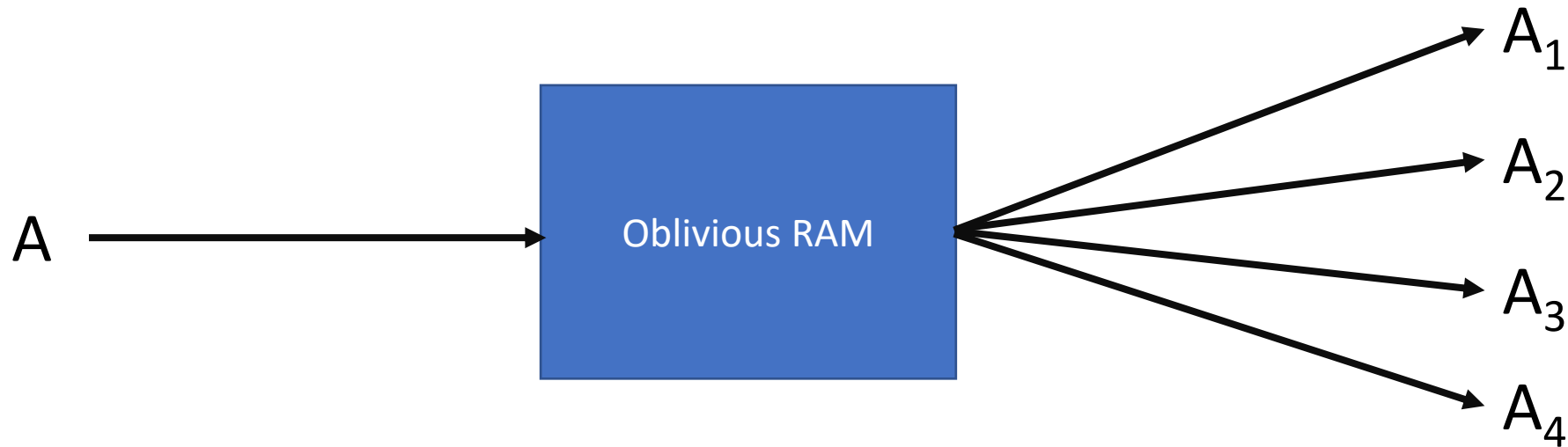
*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Enhancements

1) Eager DetWoORAM (EDetWoORAM)

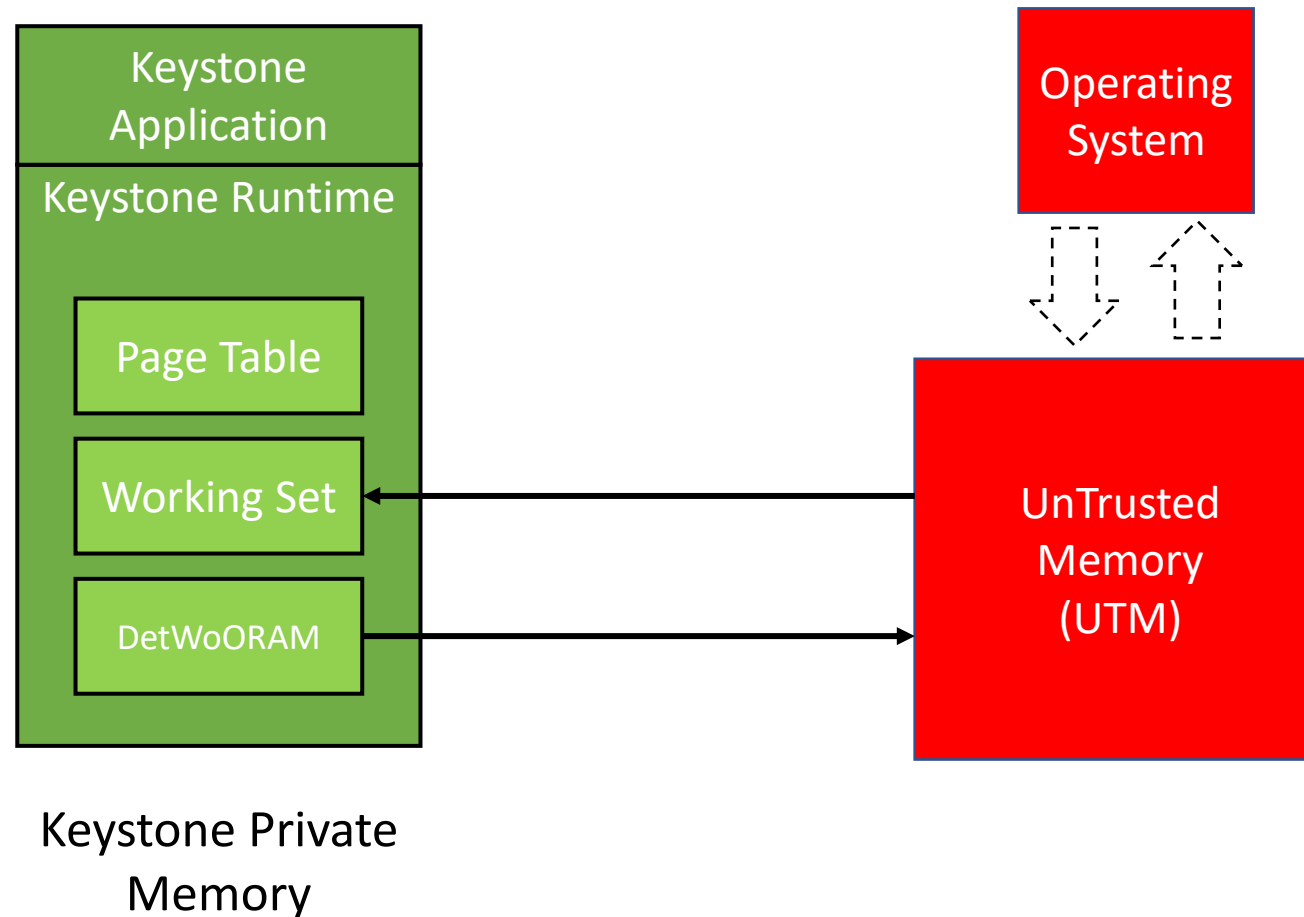
Observation

DetWoORAM is “deterministic”.

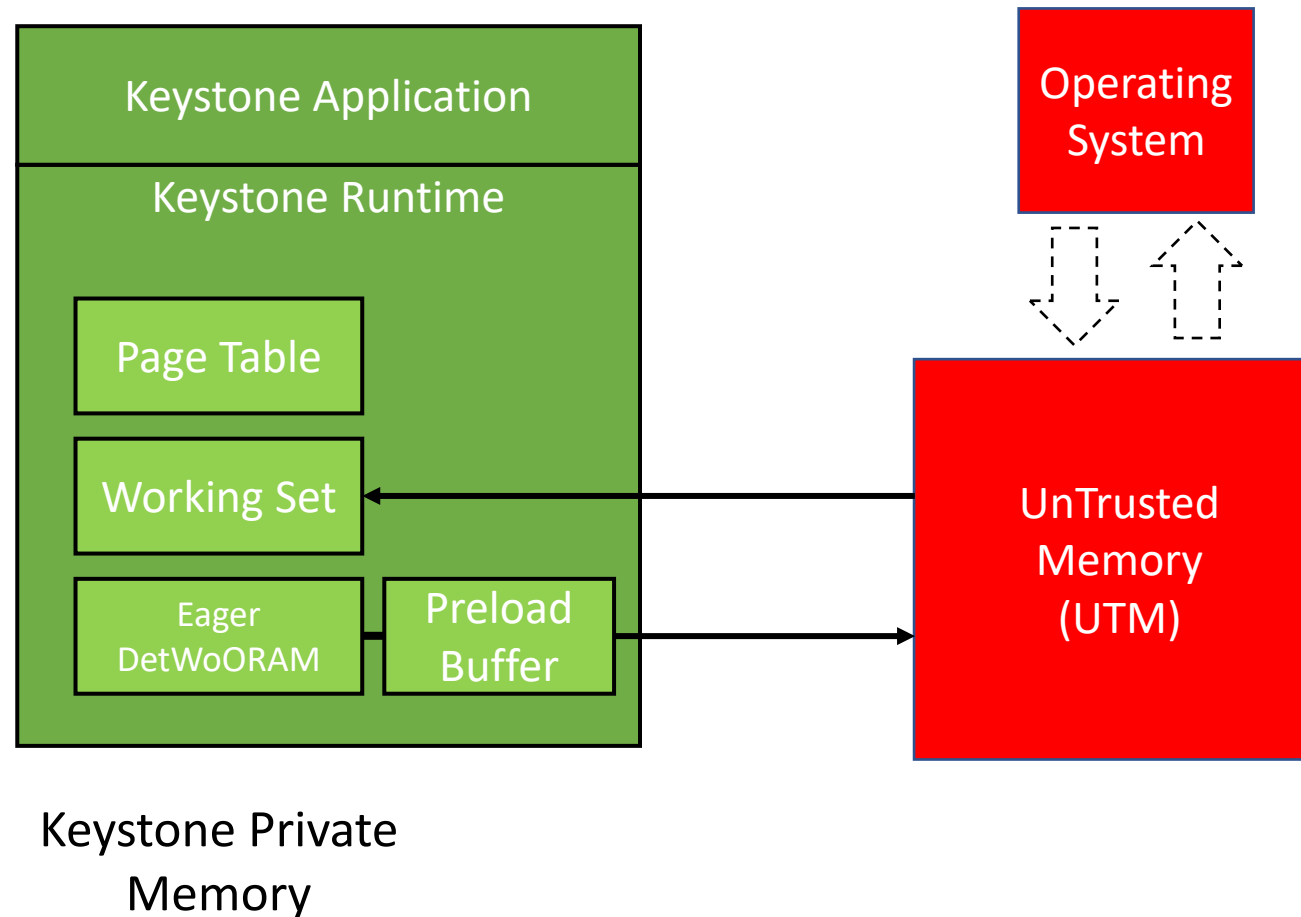


Irrespective of A ,
 A_1, \dots, A_4 will always be the same

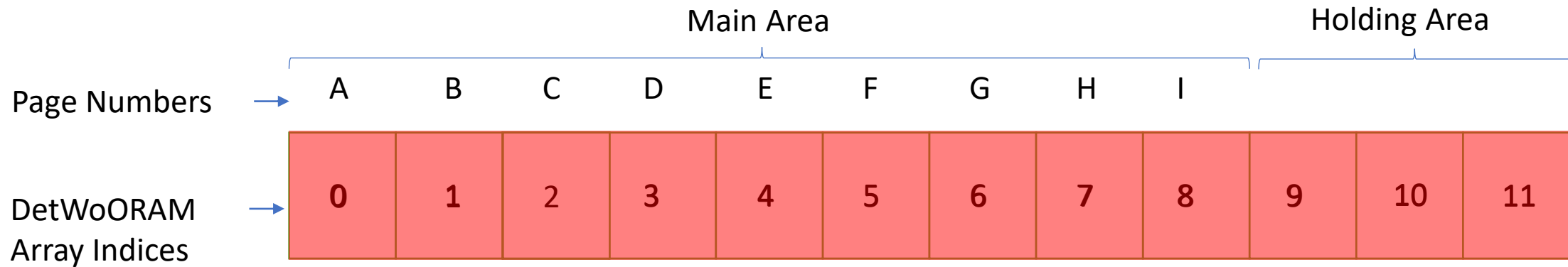
DetWoORAM for demand paging



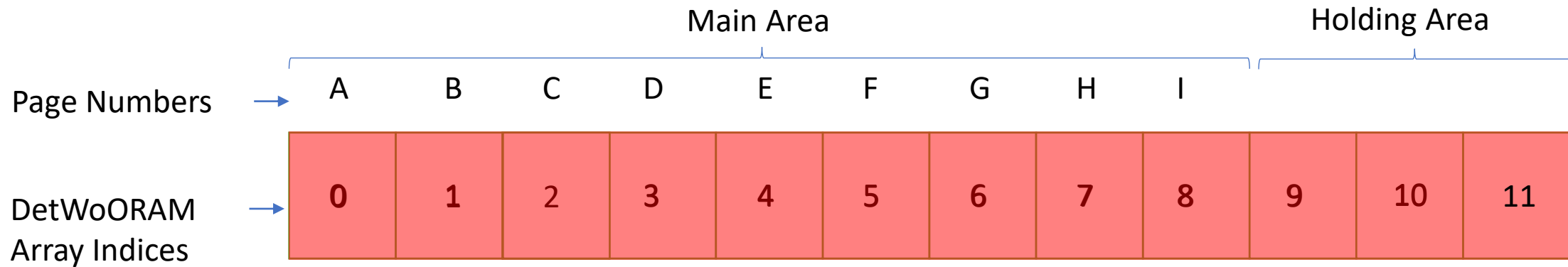
EDetWoORAM for demand paging



Working of EDetWoORAM

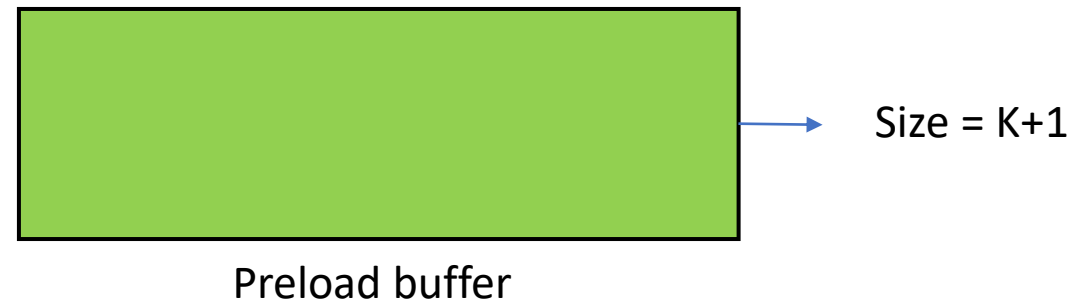


Working of EDetWoORAM

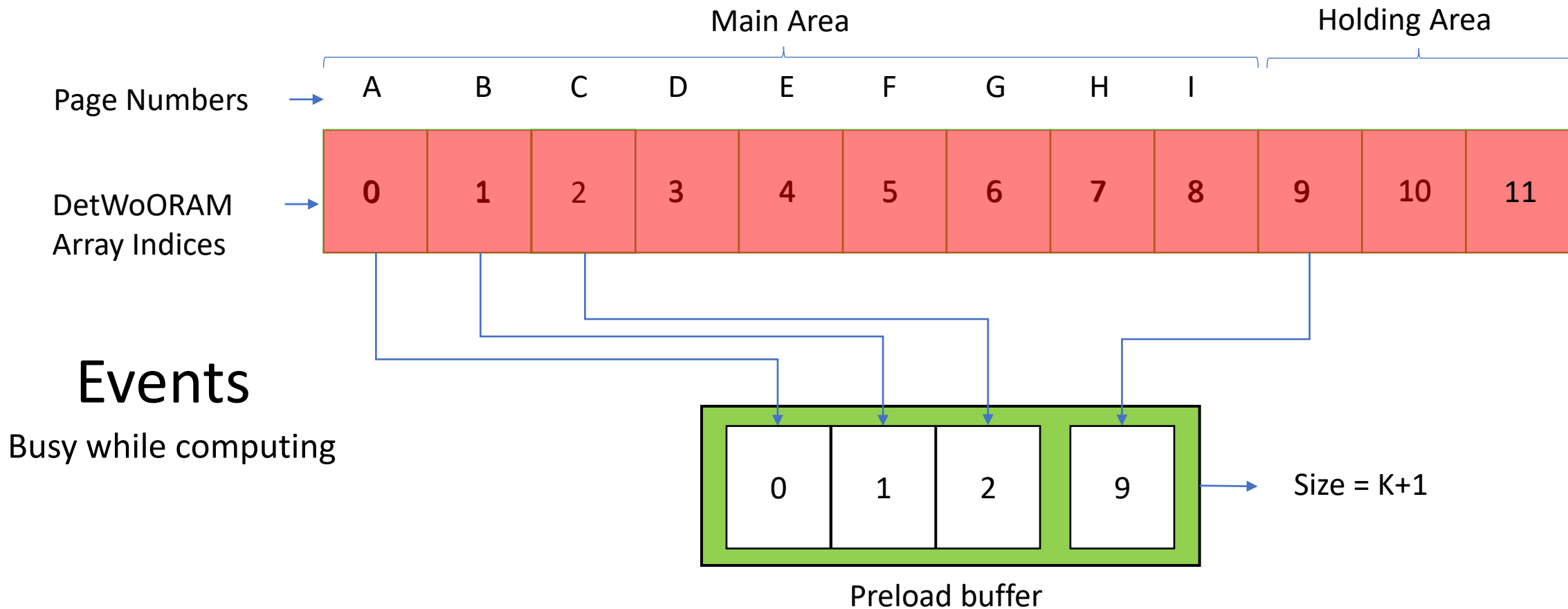


Events

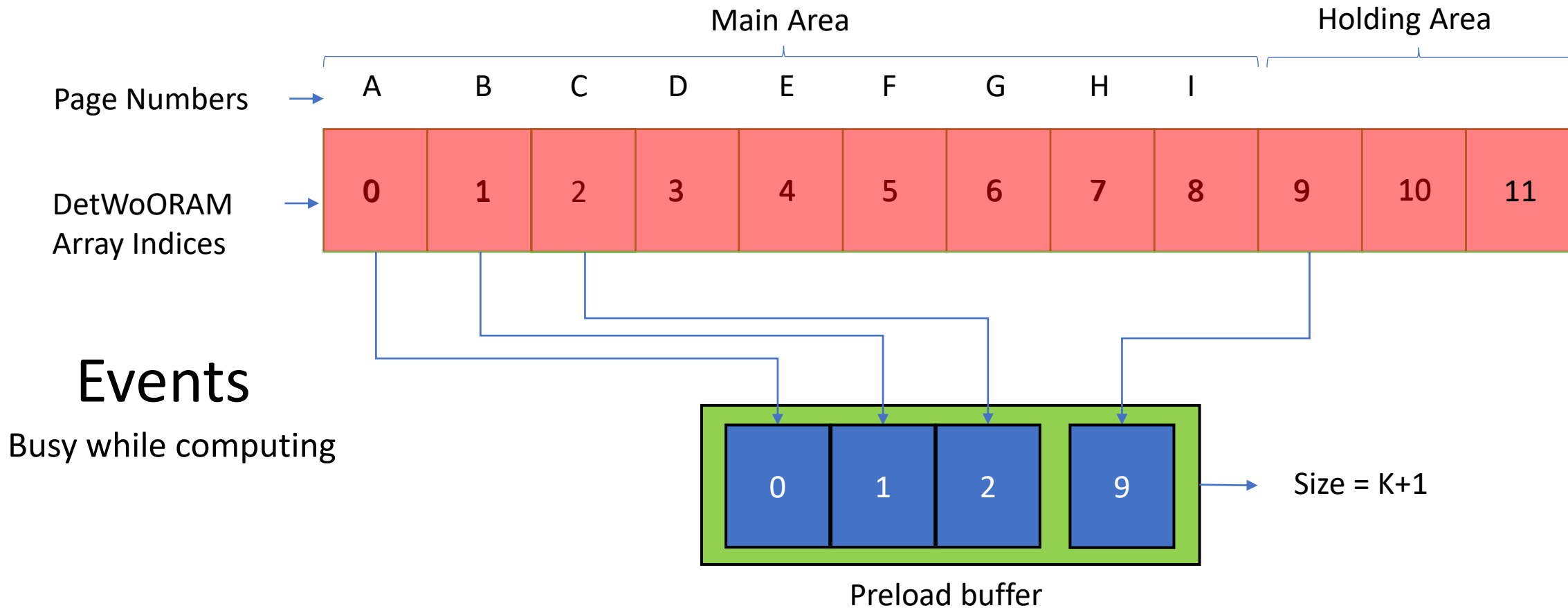
Busy while computing



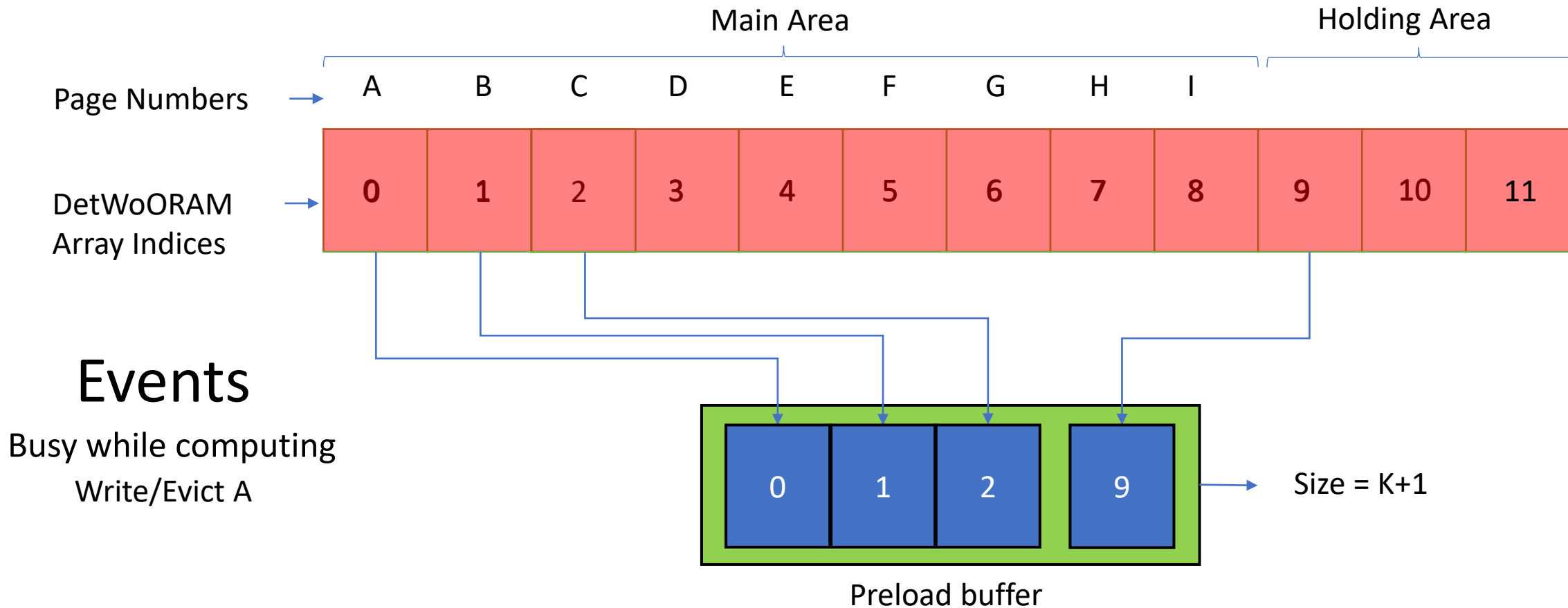
Working of EDetWoORAM



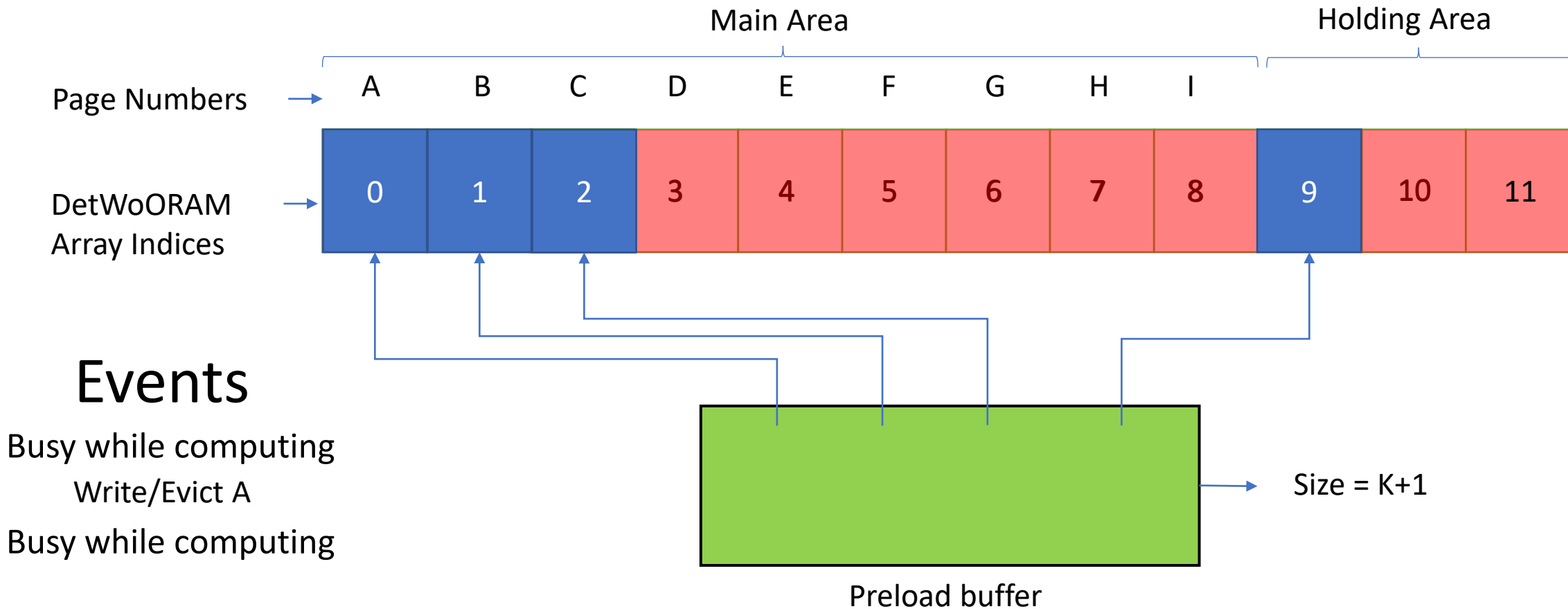
Working of EDetWoORAM



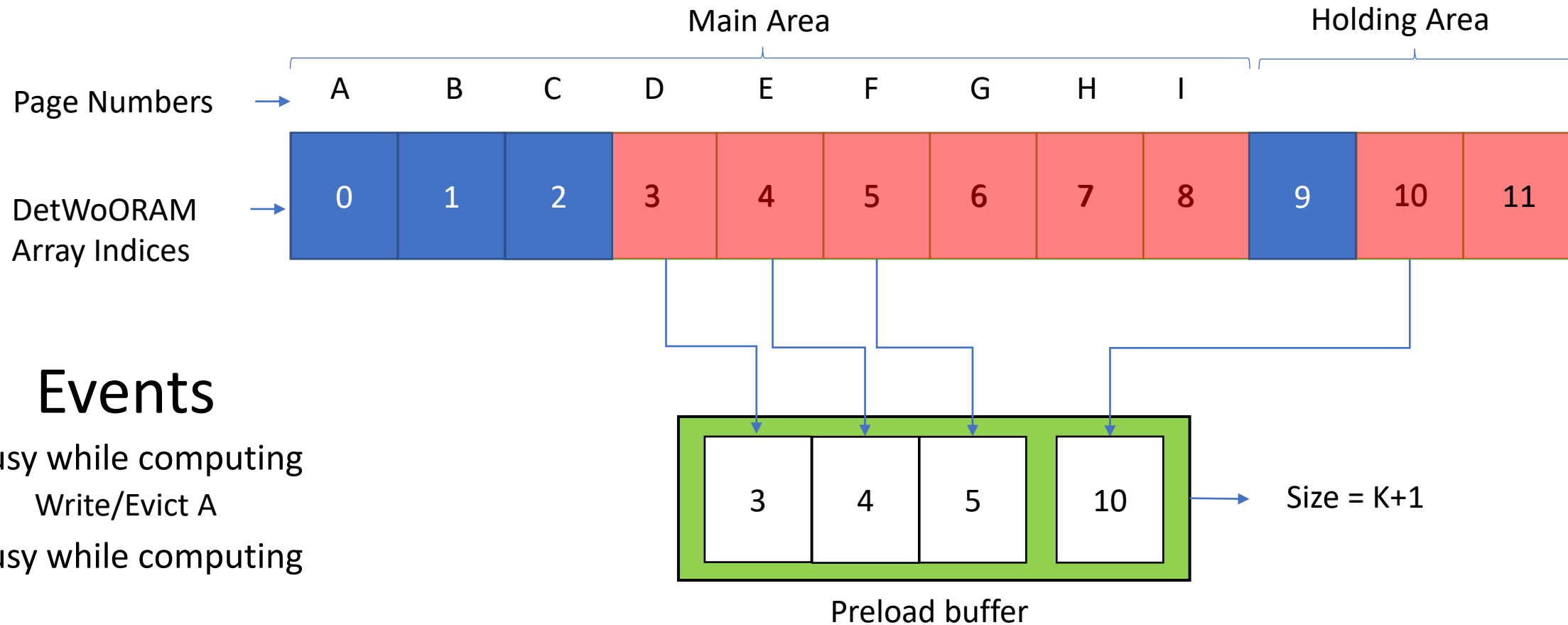
Working of EDetWoORAM



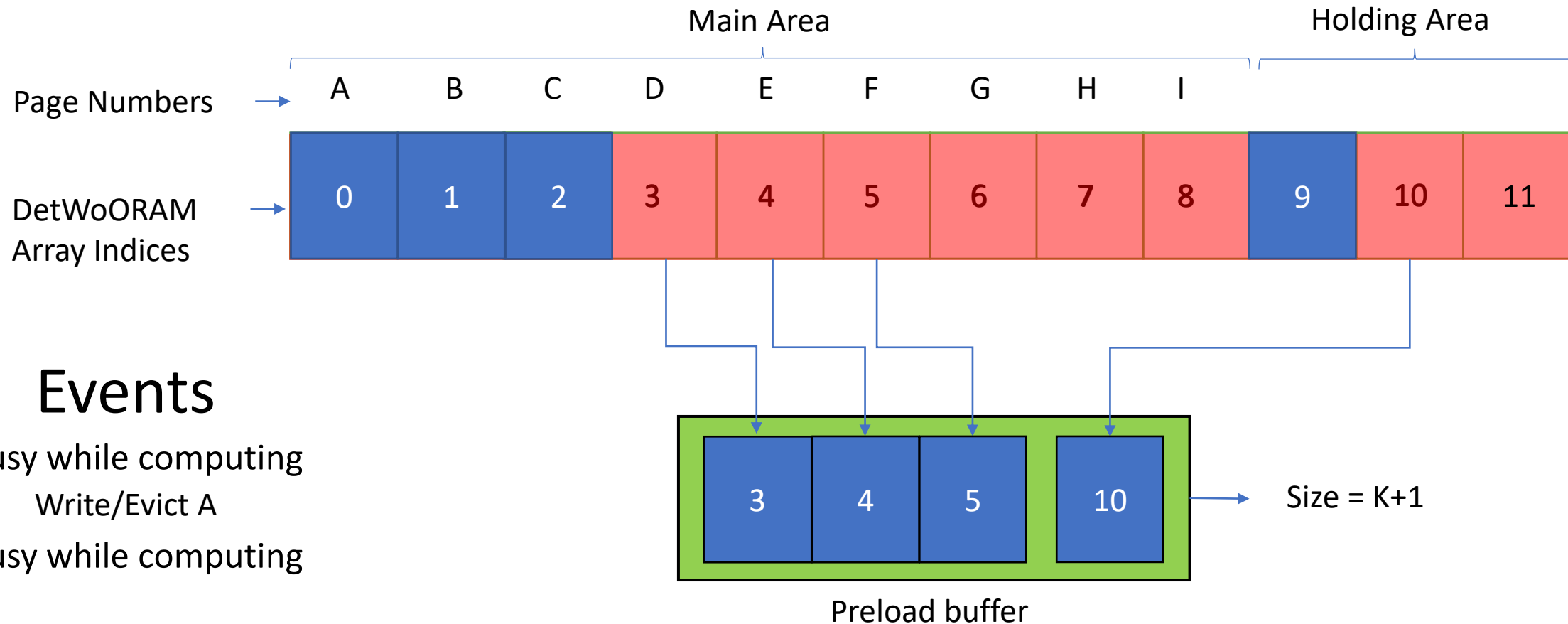
Working of EDetWoORAM



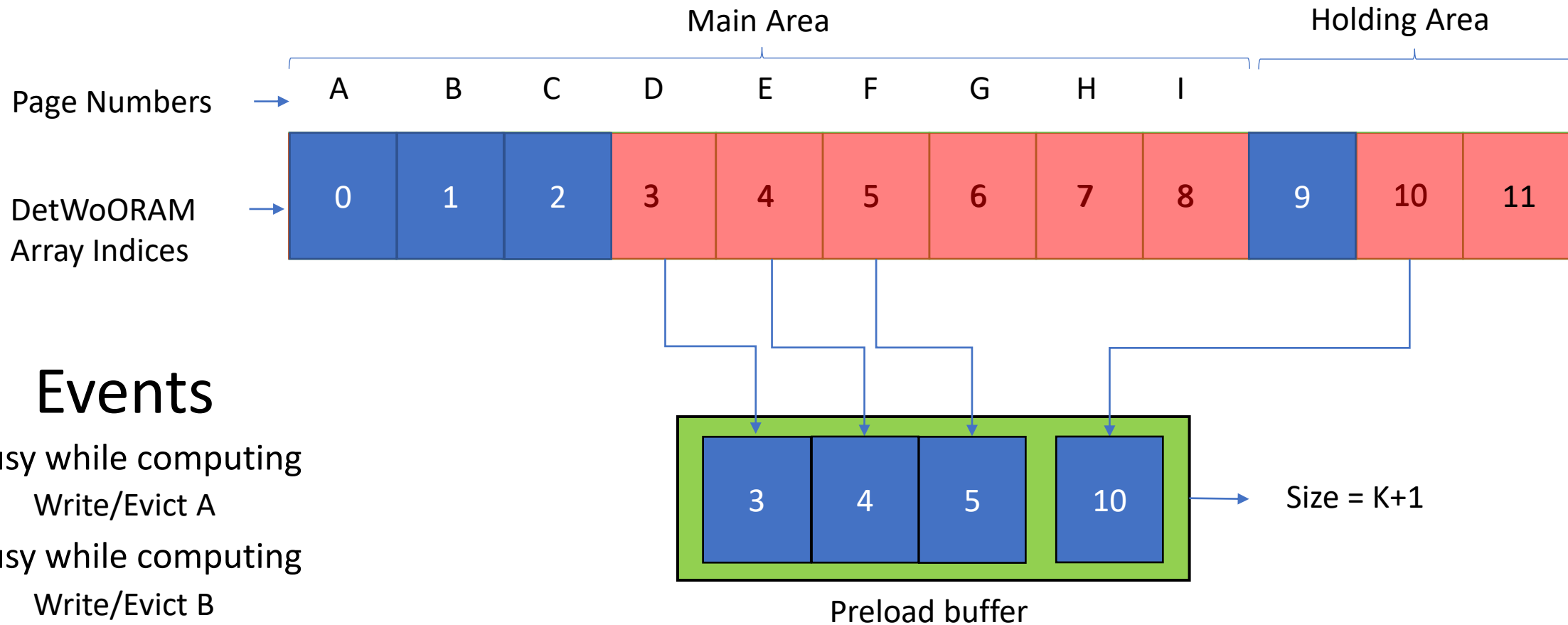
Working of EDetWoORAM



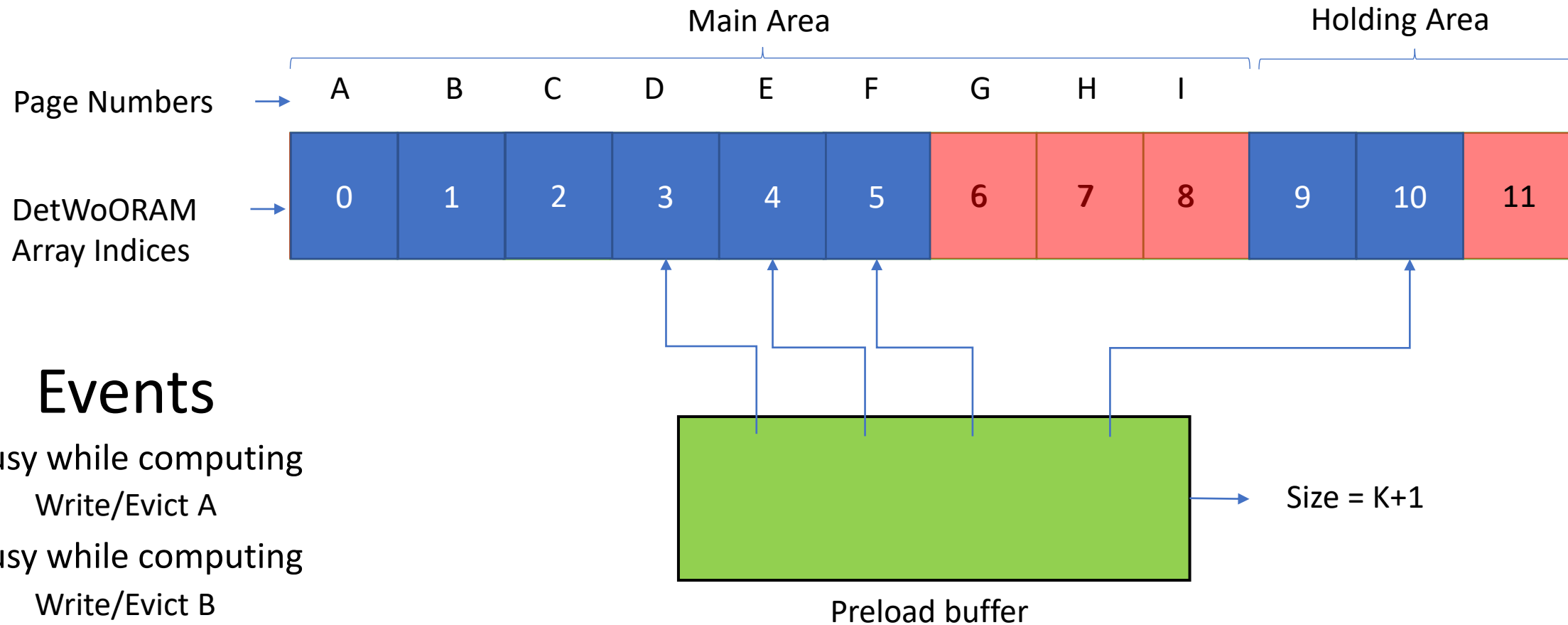
Working of EDetWoORAM



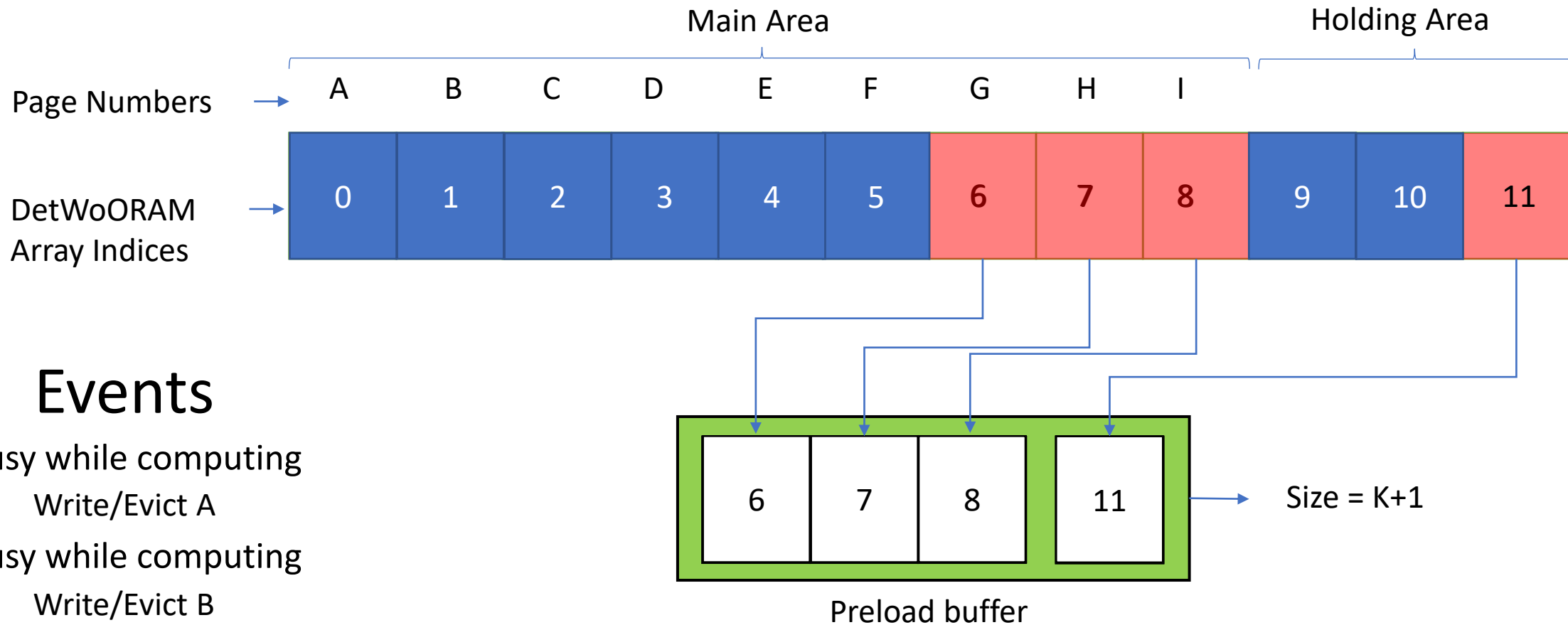
Working of EDetWoORAM



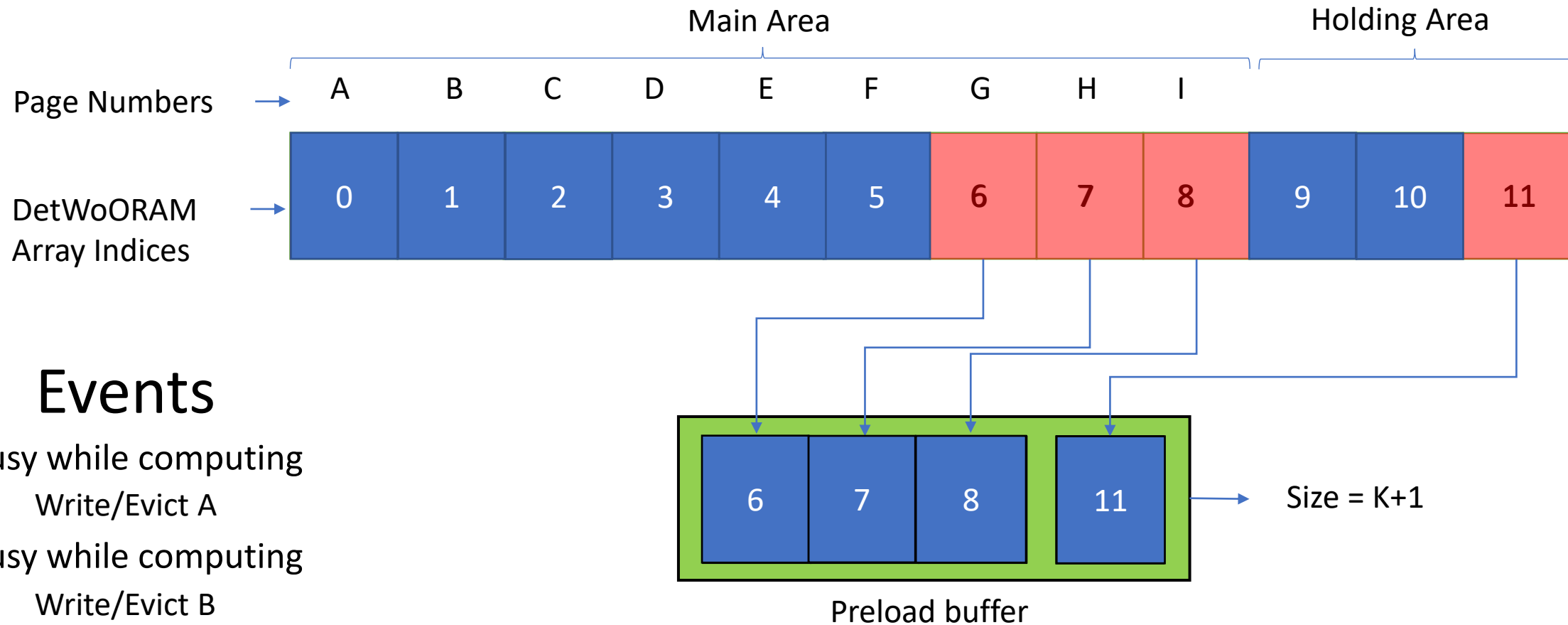
Working of EDetWoORAM



Working of EDetWoORAM



Working of EDetWoORAM



Simulation of EDetWoORAM

$X \leftarrow$ DetWoORAM latency that's supposed to be hidden by EDetWoORAM.

$Y \leftarrow$ Time between consecutive page fault.

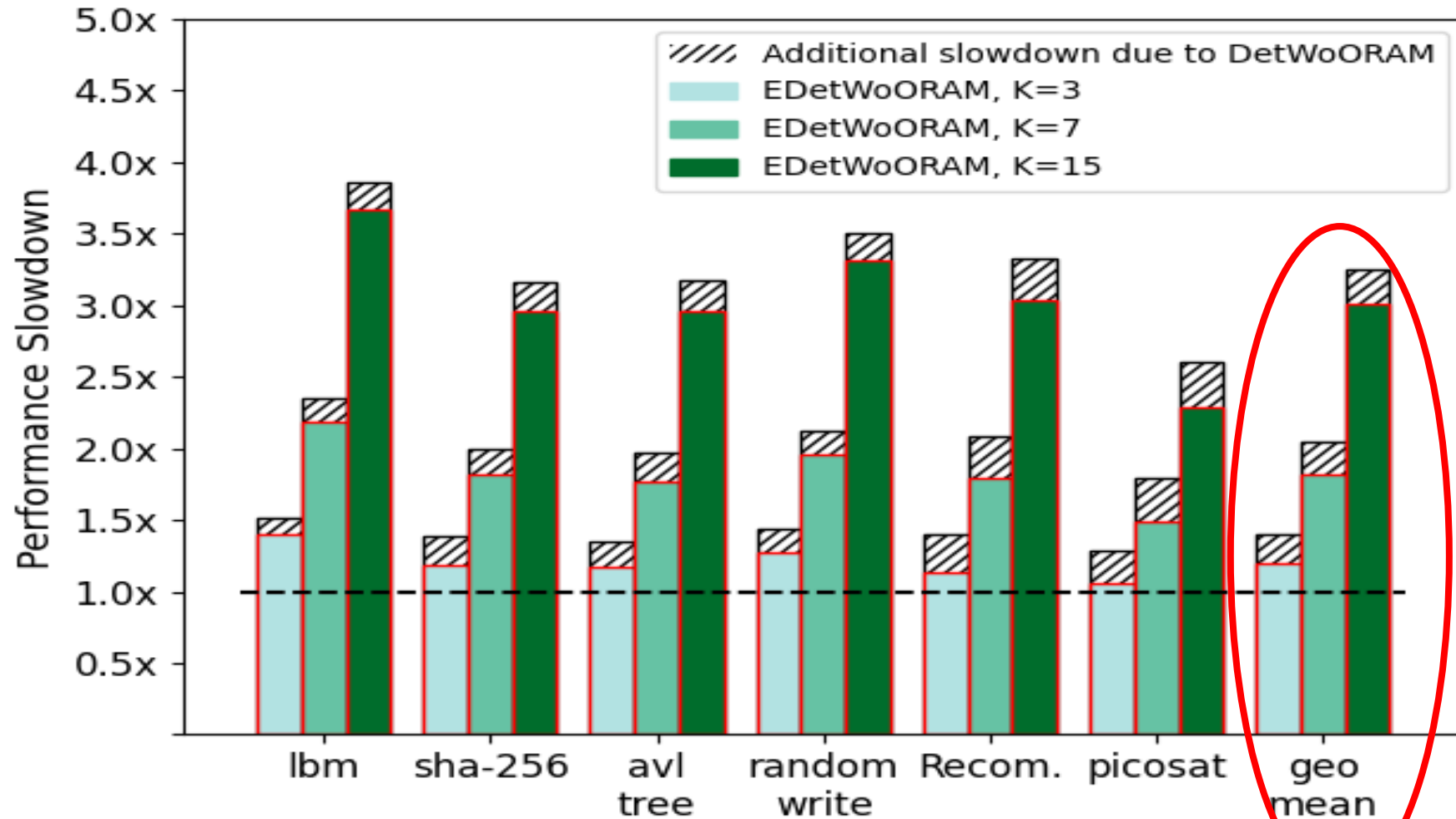
if $X < Y$:

 discount X from total execution time.

else:

 discount Y from total execution time.

Speedup with EDetWoORAM



Marginal, but considerable reduction in slowdown.

This is all I would have had to present....But...

CAR-RV 2021 deadline extended!

TIME = MONEY



TIME

=

PERFORMANCE



Enhancements

2) Parallel DetWoORAM (PDetWoORAM)

Effect of ``K`` on performance

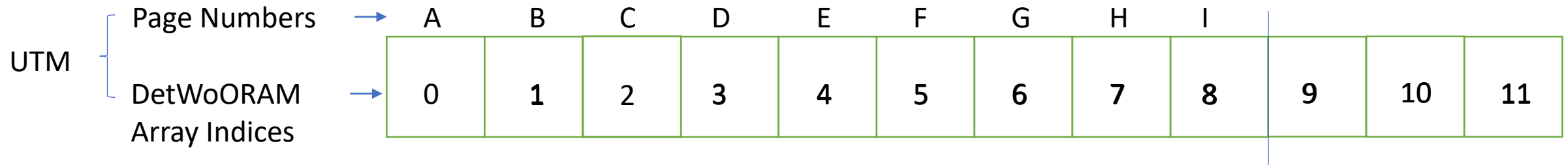
High value of K \Rightarrow More space in Main Area.

Good

\Rightarrow More writes in Main Area.

Bad

Working of PDetWoORAM*

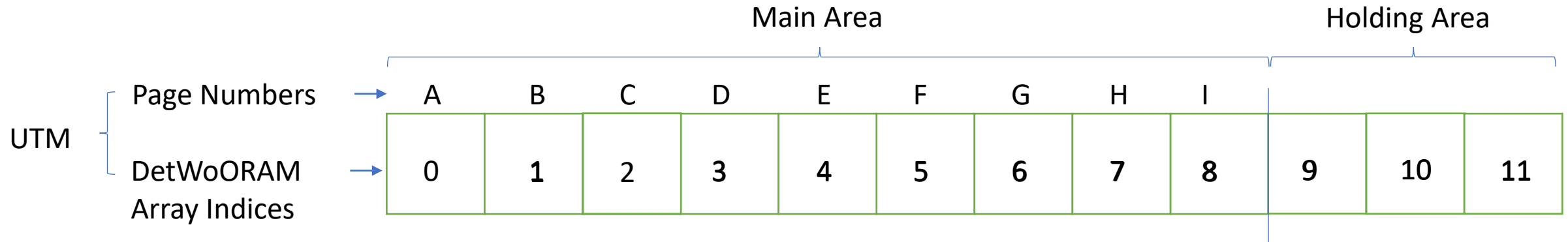


Position Map

A	0
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of PDetWoORAM*

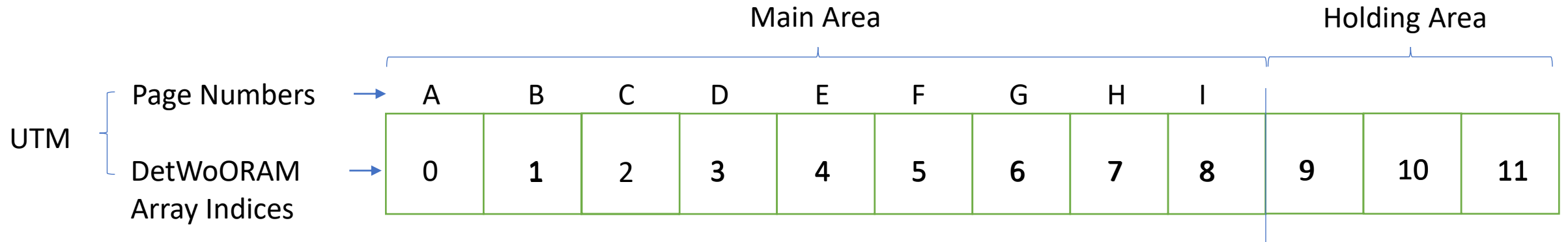


Position Map

A	0
B	1
C	2
D	3
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F	5
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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of PDetWoORAM*



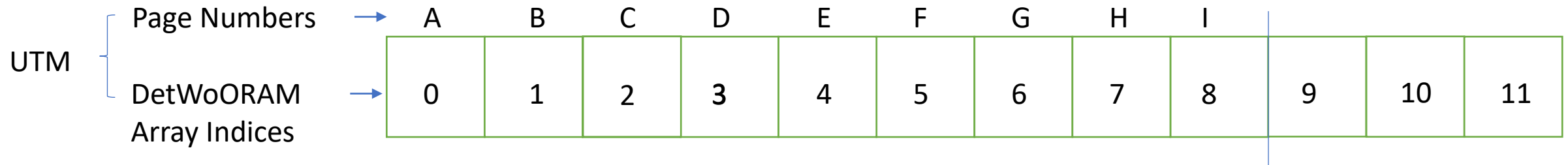
Position Map

A	0
B	1
C	2
D	3
E	4
F	5
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H	7
I	8

$$K = \frac{\text{Main Area}}{\text{Holding Area}}$$

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of PDetWoORAM*



Position Map

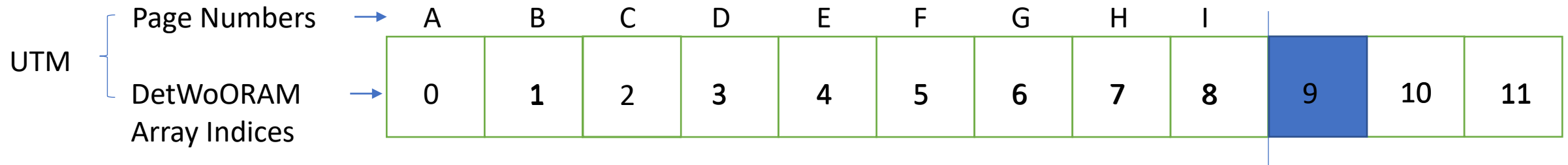
A	0
B	1
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D	3
E	4
F	5
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I	8

Operations

1. Write A

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of PDetWoORAM*



Operations

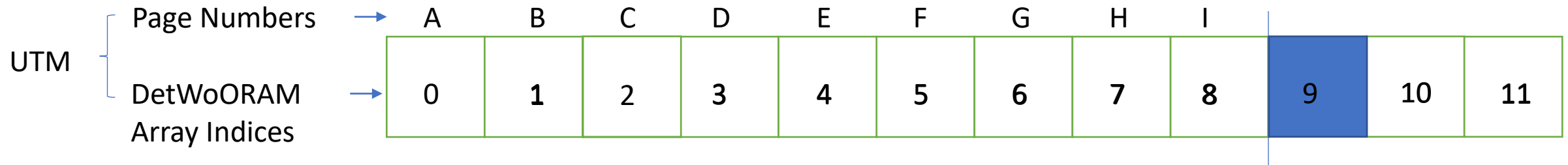
1. Write A

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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of PDetWoORAM*



Operations

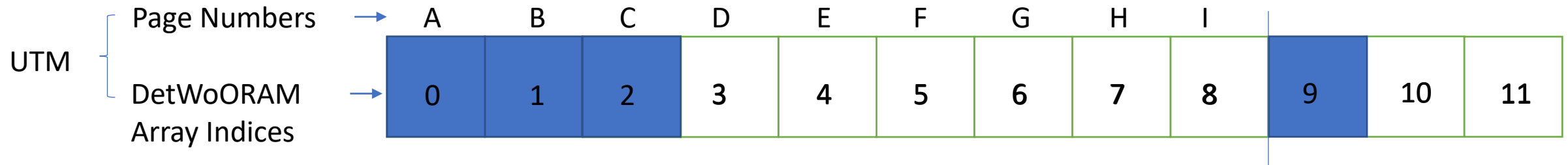
1. Write A

Position Map

A	9
B	1
C	2
D	3
E	4
F	5
G	6
H	7
I	8

*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of PDetWoORAM*



Position Map

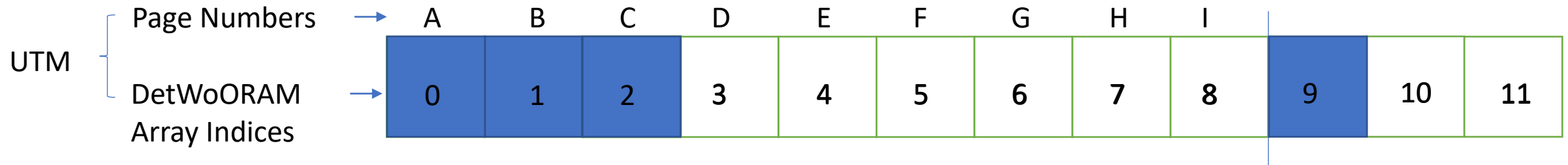
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Operations

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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of PDetWoORAM*



Operations

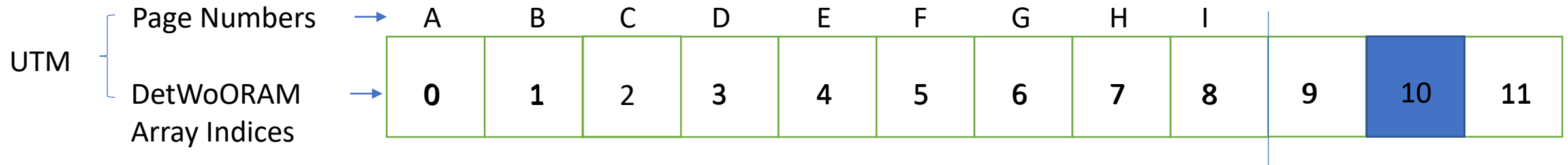
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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of PDetWoORAM*



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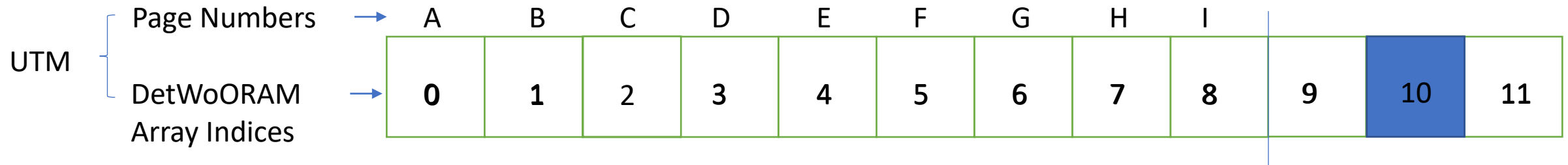
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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

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Operations

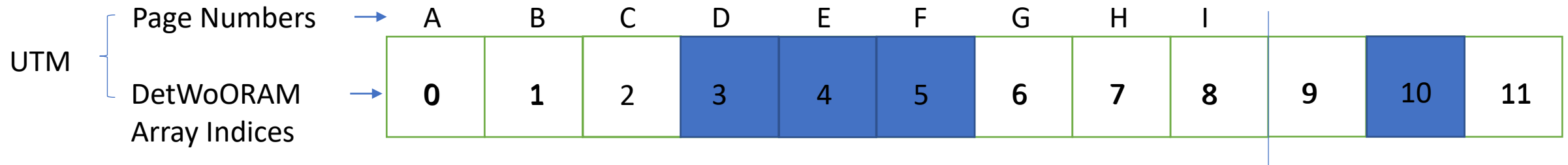
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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Working of PDetWoORAM*



Operations

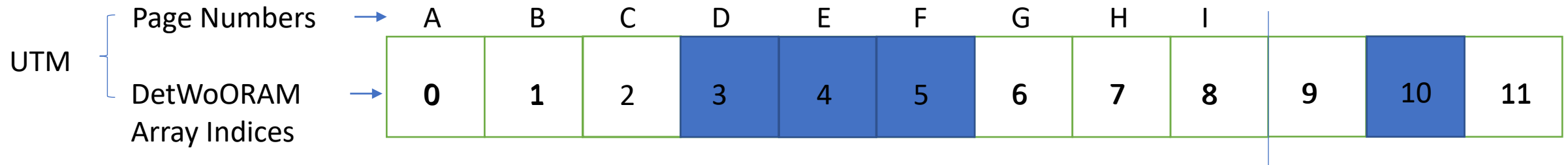
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Position Map

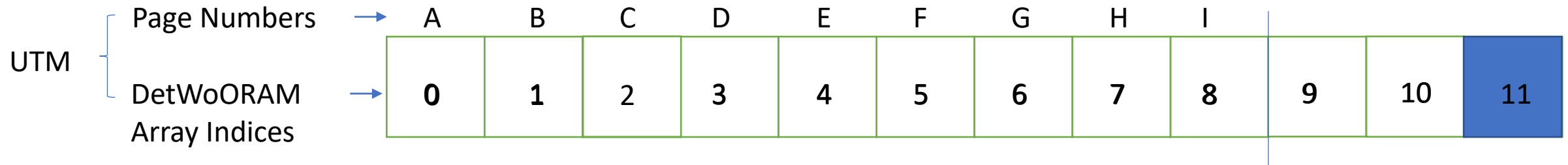
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C	2
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Position Map

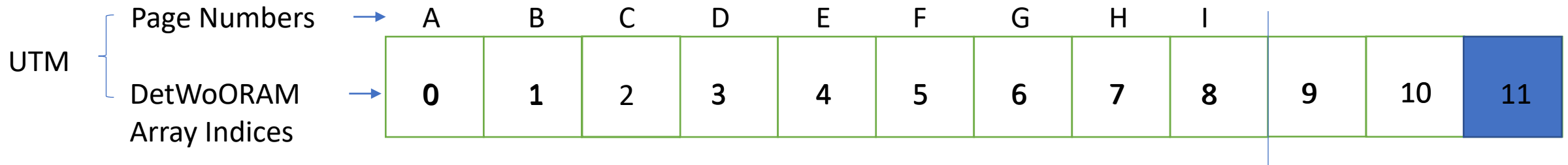
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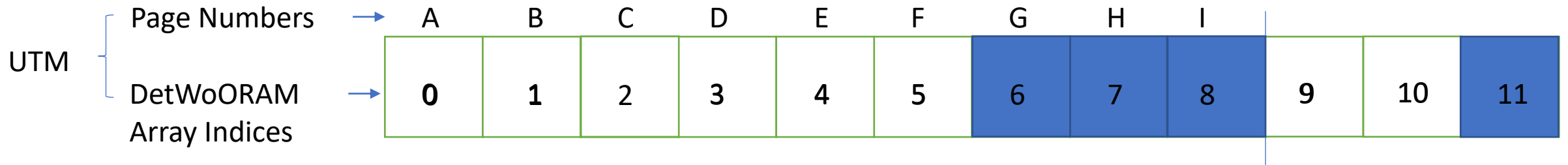
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*Deterministic, Stash-Free Write-Only ORAM by Roche et al, 2017

Simulation of PDetWoORAM

Divide the time spent in fault handling with DetWoORAM by c .

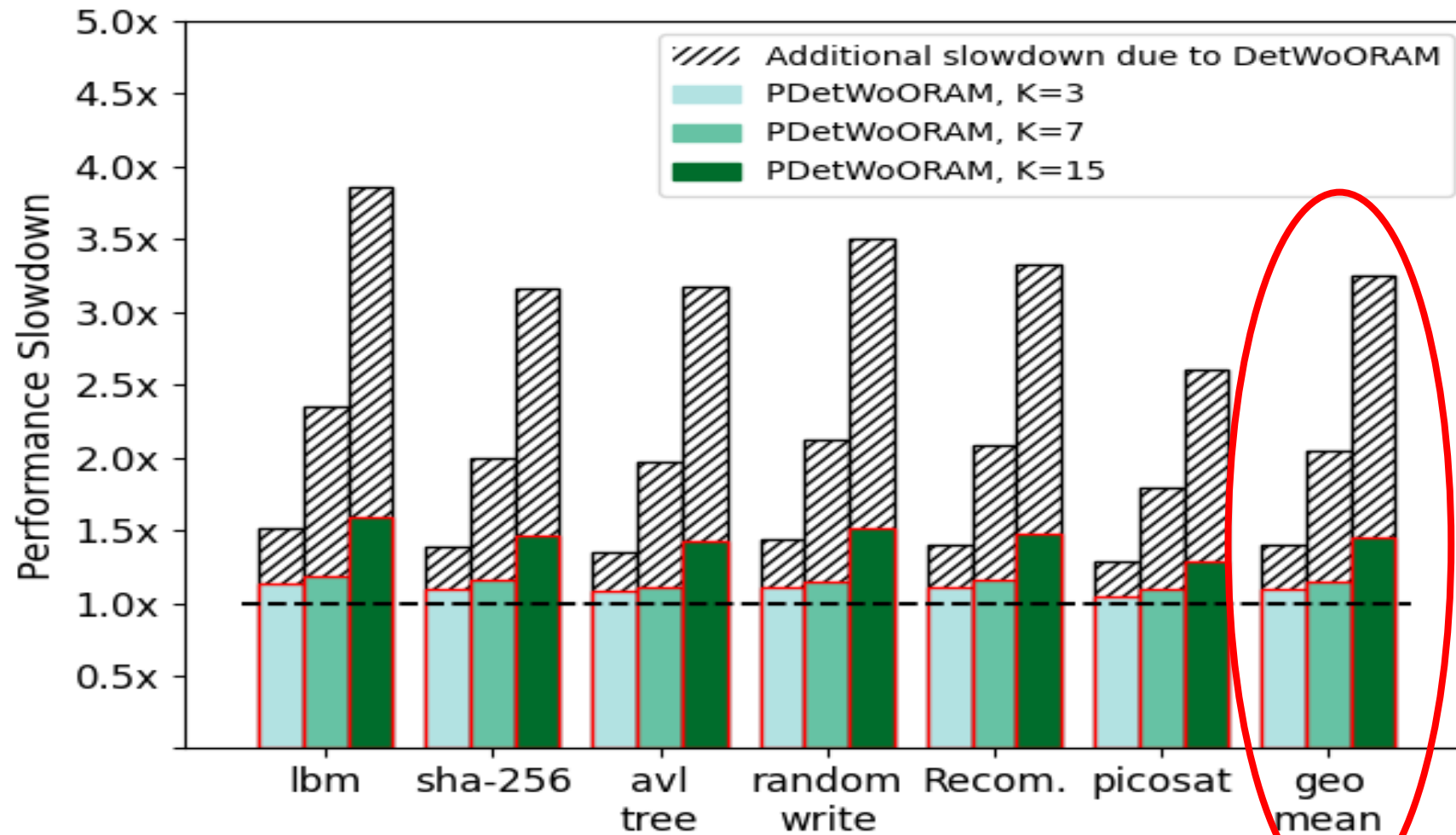
c is the PDetWoORAM speedup on a native implementation.

$$T' = T - t + \frac{t}{c}$$

The diagram illustrates the components of the equation $T' = T - t + \frac{t}{c}$ for the simulation of PDetWoORAM. Blue arrows point from the following labels to their respective terms in the equation:

- Simulated time** points to T' .
- Actual running time in Keystone** points to T .
- Time spent in DetWoORAM** points to t .
- PDetWoORAM Speedup from a real implementation** points to c .

Slowdown with PDetWoORAM



Substantial improvement in slowdown.

Almost optimal!

Summary

- We discussed how Write Only ORAMS are useful in Keystone.
- We introduced EDetWoORAM, that provides marginal speedup.
- We also introduced a PDetWoORAM that provides aggressive speedup.

THANK YOU