Complexity

Name	Complexity
Selection sort	All case: $O(n^2)$
Bubble Sort	All case: $O(n^2)$
Shell sort	Worst: $O(n^{3/2})$
Early-termination Bubble sort	Best: $O(n)$ Worst: $O(n^2)$ Avg: $O(n^2)$
Insertion sort	Best: $O(n)$ Worst: $O(n^2)$ Avg: $O(n^2)$
Merge sort	All case: $O(nlogn)$
Quick sort	Best: $O(nlogn)$ Worst: $O(n^2)$ Avg: $O(1.3999nlogn)$ Avg case 39% faster than merge sort
Heap sort	Build + Deque $O(nlogn) + O(n) = O(nlogn)$
DFS, BFS	Adj Matrix: $\Theta(V^2)$ Adj List: $\Theta(V + E)$
Binary search	Best: $O(1)$ Worst: $O(\log(n))$ If the array is sorted or has $O(1)$ access to any positon

BST Search	Best: $O(1)$ Average: $O(\log n)$ Worst: $O(n)$ If it's a stick, that's $h=n-1$
Search in sorted list	Merge sort + binary search $O(nlogn) + O(logn) = O(nlogn)$ Worth to presort when: $k > = rac{nlogn}{n/2 - logn}$
AVL trees	Height: $h <= 1.4404log(n+2) - 1.3277$ Search and Insert: $O(logn)$ Delete: $O(logn)$ Rebalance (rotations): $O(logn)$ Disadvantages: Rotations are frequent and implementation is complex
2-3 trees	Height: $log_3(n+1)-1 <= h <= log(n+1)-1$ SEARCH, INSERT, DELETE: $O(logn)$ Rebalancing: $< O(logn)$
Heaps	Height: $h = logn$ Repair: $O(logn)$ C(n) = O(n) Building: $O(nlogn)$
Distribution sorting	Worst case: $O(n)$ if $n>n_{max}$ Space: $O(n)+O(n_{max})$
	List - Insert: $O(1)$ - Delete: $O(n)$ - Search: $O(n)$ Balance Tree - Insert: $O(\log n)$

Hash tables	<pre>- Delete: O(logn) - Search: O(logn) Array - Insert: O(logn) - Delete: O(logn) - Search: O(logn) Open Hashing - Insert: O(1) - Delete: propotional to list's length - Search: propotional to list's length - Average: O(1)</pre>
Prim's algorithm	Min-heap + Adjacency list $ -\ O(E log V) $ Fibonacci heap + adjacency list $ -\ O(E + V log V) $
Kruskal's algorithm	If there are edges: $ - O(E log E) $ If there is no edge: $ - O(E log V) $
Dijkstra's Algorithm	Min-heap: - $\Theta(E log V)$
Huffman's Code	Min-heap: $-\ \Theta(nlogn)$ Sorted by probabilities: $-\ \Theta(n)$
Edit distance	Worse and Average: - Construction: $\Theta(nm)$ where n and m is length of 2 strings. - Backtrace complexity: $\Theta(m+n)$

Knapsnack problem Build table: $\Theta(nW)$ Backtrace complexity: $\Theta(n+W)$ Complexity: $\Theta(n^3)$ Space: $\Theta(n^2)$ DFS & BFS: $\Theta(n^2+nm)$ for n, $\Theta(n+m)$ for 1 If not many: $\Theta(n^2)$

• Stack:

- Add and remove at front.
 - [a] add b => [b,a]. remove b => [a]
- All log, ln has the same speed
- (n-1)! < n!
- When rotating for AVL trees. If one element changes its position to the right. We say R(that element). So draw first and then write rotation direction.