# CS 5153/5053 Network Security, Spring 2023 Project 2: Buffer Overflow Attack Report

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## Link to Source Code <a href="https://github.com/austinc3030/buffer">https://github.com/austinc3030/buffer</a> m11809075

## Host Environment Used

Operating System: macOS Ventura 13.0.1

Hardware: Apple M1 Mac Mini (ARM Architecture)

Hypervisor: UTM

## **Additional Notes**

Due to the current state of virtualization on M1 Macs (ARM based), the pre-built virtualbox image provided by SEED Labs was not used. Instead, the instructions to build the virtual machine for the lab provided by SEED Labs were followed and used for the project.

### Virtual Machine Environment Used

## Operating System: Ubuntu Desktop 22.03.2

```
hetworksecvm% uname –a
Linux networksecvm 5.15.0–60–generic #66–Ubuntu SMP Fri Jan 20 14:29:49 UTC 2023 x86_64 x86_64 x86_64 GNU/Linux
```

#### Shell: zsh v5.8.1

```
networksecvm% ps –p $$
    PID TTY
                     TIME CMD
                 00:00:04 zsh
   1283 tty1
hetworksecvm% zsh ––version
zsh 5.8.1 (x86_64–ubuntu–linux–gnu)
hetworksecvm%
```

#### gcc:

```
letworksecvm% gcc --version
(cc --version
(cc (Ubuntu 11.3.0-1ubuntu1~22.04) 11.3.0
opyright (C) 2021 Free Software Foundation, Inc.
his is free software; see the source for copying conditions. There is NO
warranty; not even for MERCHANTABILITY or FITNESS FOR A PARTICULAR PURPOSE.
```

#### Caveats:

1. Due to differences between 32/64-bit assembly, the 'gcc-multilib' package is required.

```
retworksecvm% sudo apt-get install gcc-multilib
leading package lists... Done
leading package lists... Done
leading state information... Done
libsG-dev-x32 libsG-i386 libsG-x32 libx32asan6 libx32atomic1 libx32gcc-s1 lib32gomp1 libx32gomp1 libx32tdc++6 lib32ubsan1 libcG-dev-i386
libsG-dev-x32 libsG-i386 libsG-x32 libx32asan6 libx32atomic1 libx32gcc-s1-dev libx32gcc-s1 libx32gomp1 libx32itm1 libx32quadmath0 libx32stdc++6 libx32ubsan1
libsG-dev-x32 libsG-i386 libsG-x32 libx32asan6 libx32atomic1 libx32gcc-s1 libx32gomp1 libx32itm1 libx32quadmath0 libx32stdc++6 libx32ubsan1
libsG-dev-x32 libsG-i386 libsG-x32 libx32asan6 libx32atomic1 libx32gcc-s1-dev libx32gcc-s1 libx32gomp1 libx32itm1 libx32quadmath0 libx32stdc++6 libx32ubsan1
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libsG-dev-x32 libsG-i386 libsG-x32 libx32asan6 libx32atomic1 libx32gcc-s1 libx32gcc-s1 libx32gomp1 libx32itm1 libx32quadmath0 libx32stdc++6 libx32ubsan1
lead to get 21.8 MB of archives.

When the this operation, 82.7 MB of additional disk space will be used.

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```

- a. Further, when compiling c code in this environment, the gcc flag `-m32` is required to instruct gcc to compile for a 32-bit architecture.
- 2. The gdb plugin Python Exploit Development Assistance for GDB (PEDA) is installed on this virtual machine for ease of use in gdb

```
networksecvm% git clone https://github.com/longld/peda.git
Cloning into 'peda'...
remote: Enumerating objects: 382, done.
remote: Counting objects: 100% (9/9), done.
remote: Compressing objects: 100% (7/7), done.
remote: Total 382 (delta 2), reused 8 (delta 2), pack–reused 373
Receiving objects: 100% (382/382), 290.84 KiB | 1.37 MiB/s, done.
Resolving deltas: 100% (231/231), done.
hetworksecvm% ls
buffer_m11809075 peda
networksecvm% echo "source /home/user/peda/peda.py" >> /home/user/.gdbinit
networksecym%
```

## How do you perform the attack in your VM

1. Disable address space randomization

```
networksecvm% sudo sysctl –w kernel.randomize_va_space=0
kernel.randomize_va_space = 0
hetworksecvm% _
```

- 2. Compile *stack.c* with the flags:
  - a. `-m32` instructs gcc to compile for 32-bit architectures
  - b. `-z execstack` make the stack executable
  - c. `-fno-stack-protector` disabled Stack-Guard Protection Scheme

```
networksecvm% gcc –m32 –z execstack –fno–stack–protector –o stack stack.c
networksecvm% ls
badfile stack stack.c stack_debug
networksecvm% _
```

**Note:** it is assumed here that the steps in the sections below, "How do you find the value of ebp" and "How do you decide the content of badfile" have been completed previously

3. Make the *stack* program a root-owned Set-UID program

```
networksecvm% sudo chown root stack
[sudo] password for user:
networksecvm% sudo chmod 4755 stack
networksecvm% ls –la stack
–rwsr–xr–x 1 root user 15088 Feb 28 01:35 stack
networksecvm%
```

4. Run exploit.py to generate badfile

5. Run stack

```
networksecvm% ./stack
# whoami
root
#_
```

## How do you find the value of ebp

- 1. Compile stack.c with the flags:
  - a. `-m32` instructs gcc to compile for 32-bit architectures
  - b. `-g` debug flag
  - c. `-z execstack` make the stack executable
  - d. `-fno-stack-protector` disabled Stack-Guard Protection Scheme

```
networksecvm% gcc -m32 -g -z execstack -fno-stack-protector -o stack_debug stack.c
networksecvm% ls
stack.c stack_debug
networksecvm% _
```

2. Create a temporary, blank file badfile (failure to do so leads to a segmentation fault)

```
networksecvm% touch badfile
networksecvm% ls
badfile stack.c stack_debug
networksecvm%
```

3. Run stack\_debug with gdb

4. Set a breakpoint on the bof function

```
gdb-peda$ b bof
Breakpoint 1 at 0x11f1: file stack.c, line 21.
vdb-peda$
```

5. Continue execution of the program

```
EAX: 0x56558fc8 --> 0x3ed0
EBX: 0x56558fc8 --> 0x3ed0
ECX: 0x5655a238 --> 0x0
EDX: 0x0
ESI: Oxffffdbd4 --> Oxffffdd15 ("/home/user/buffer_m11809075/src/stack_debug")
EDI: 0xf7ffcb80 --> 0x0
EBP: 0xffffd858 --> 0xffffdb08 --> 0xf7ffd020 --> 0xf7ffda40 --> 0x56555000 --> 0x464c
ESP: 0xffffd7c0 --> 0x36 ('6')
EIP: 0x565561f1 (<bof+20>: sub esp,0x8)
EFLAGS: 0x212 (carry parity ADJUST zero sign trap INTERRUPT direction overflow)
   0x565561e1 <bof+4>: sub
                                  esp,0x94
   0x565561e7 <bof+10>: call
   0x565561ec <bof+15>: add
                                  eax,0x2ddc
=> 0x565561f1 <bof+20>: sub
                                  esp.0x8
   0x565561f4 <bof+23>: push
                                  DWORD PTR [ebp+0x8]
   0x565561f7 <bof+26>: lea
                                  edx, [ebp-0x8a]
   0x565561fd <bof+32>: push
                                  edx
   0x565561fe <bof+33>: mov
                                  ebx,eax
0000| 0xffffd7c0 --> 0x36 ('6')
0004| 0xffffd7c4 --> 0x56557008 --> 0x61620072 ('r')
0008| 0xffffd7c8 --> 0xf7e04d9d (<_IO_doallocbuf+13>: add
0012| 0xffffd7cc --> 0x1000
                                                                      ebx,0x1a7263)
0016| 0xffffd7d0 --> 0x5655a1a0 --> 0xfbad2498
0020| 0xffffd7d4 --> 0x205
0024| 0xffffd7d8 --> 0xffffd8f7 --> 0x0
0028| 0xffffd7dc --> 0xf7e02d19 (add
                                            esp,0x10)
Legend: code, data, rodata, value
Breakpoint 1, bof (str=0xffffd8f7 "") at stack.c:21
             strcpy(buffer, str);
 (db-peda$
```

6. Obtain the value of ebp

```
gdb-peda$ p $ebp
$1 = (void *) 0xffffd858
```

ebp = 0xffffd858

## How do you decide the content of badfile

- 1. Complete the steps in the section above, "How do you find the value of ebp" as the *ebp* value found will be used in the calculations for the construction of the badfile
- 2. In addition to the value of ebp, the address where buffer starts is required

```
gdb-peda$ p &buffer
$1 = (char (*)[130]) 0xffffd7ce
buffer starts at 0xffffd7ce
```

3. The return address ebp is after the starting address of buffer, therefore, we can overflow the buffer such that the content at that return address will be an address we want to return to. To determine where this address should be stored, we must determine the difference between the return address (ebp) and the start address of buffer

```
gdb-peda$ p/d 0xffffd858 - 0xffffd7ce
$3 = 138
gdb-peda$
```

- a. This means that 138 bytes after the start of *buffer*, the next byte will be the beginning of the return address we want to alter. Thus, our offset will be 142, meaning the address we want to return will be stored 142 bytes into the content of badfile
- 4. Since we are placing our shellcode at the end of the badfile, and we do not want to return exactly where the program would normally return, we need to add some padding to the address we store so that the return lands at least in the NOP bytes of our payload. This can be an arbitrary number greater than 4 such that this number lies within the content of the badfile. I chose 96 such that:
  - a. ebp+offset+96 = 0xFFFFD858 + 0x60 = 0xFFFFD8B8
- 5. The address determined in step 5a needs to be stored at the offset in the badfile

# Whether your attack is successful

1. When running *stack* with the generated *badfile* alongside, the exploit works correctly. We can verify we have a root shell by running `whoami` which returns *root* 

```
networksecvm% ./stack
# whoami
root
#
```