

(append  $\ell_1 \dots \ell_{n-1} \overset{\text{copied}}{\ell_n}$ ) shared in result

cost is  $O(|\ell_1| + |\ell_2| + \dots + |\ell_{n-1}|)$

(length  $\ell$ )  $O(|\ell|)$

(null?  $\ell$ )  $O(1)$

$(> (\text{length } \ell) \emptyset)$

## Prolog cost model

~~if X is a variable~~  
X

X = Y,

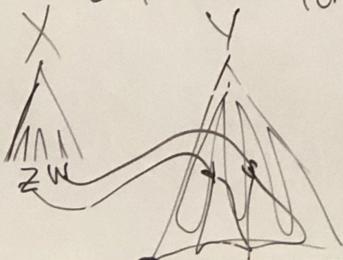
if X is still just a variable,  
if Y is still just a variable,

UNOC

$O(1)$

$O(1)$

if X & Y are both complicated structures



$O(\min(|X|, |Y|))$

$O(\max(|X|, |Y|))$

Scheme

(eq? A B)  
(eqv? A B)  
(equal? A B)

O(1)  
O(n)  
O( $\infty$ )

~~TAD BD TAD BD~~

~~TAD BD~~

### Array access cost

#### Powers of 2

$a[i]$

$\&a + (i - \text{lb}(a)) * \text{sizeof } a[0]$

`double a[1000: 2000];`

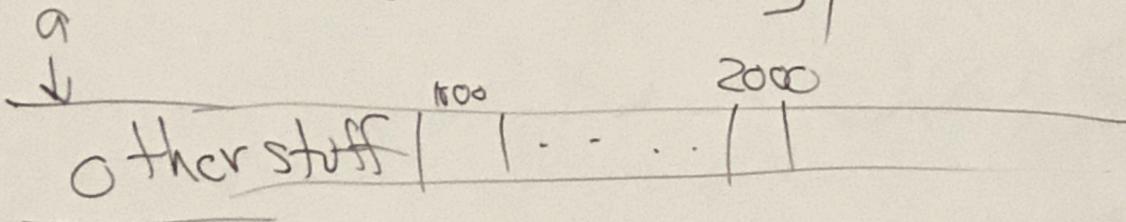
- if we use a power of 2, then we can get rid of multiply and replace with shift  
(better performance)

-

$a[i]$

$\&a + (i)$

double a [1000: 2000];

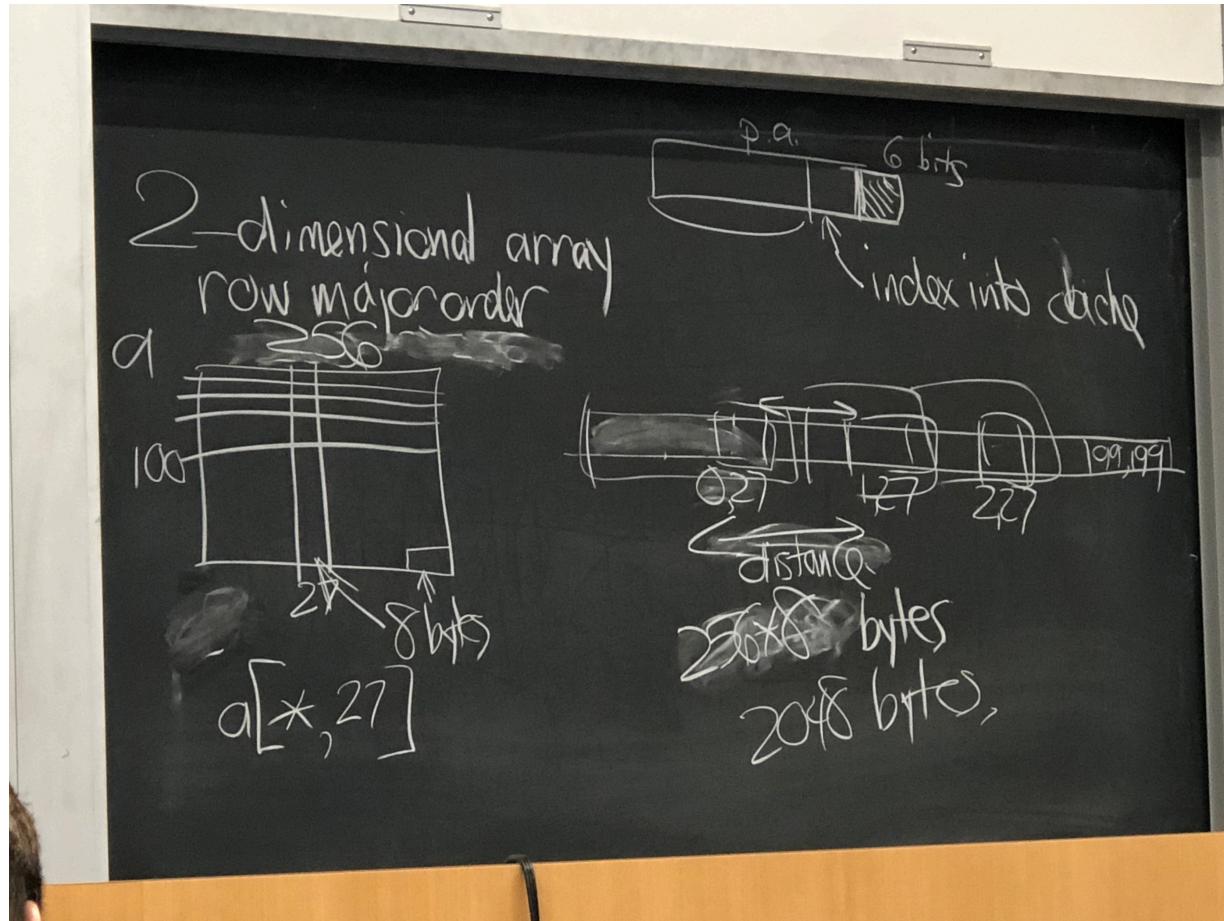


### Non-Powers of 2

Can actually be beneficial if we have 2D array

Want to access a column

But entries are  $200*8$  bytes away from each other, not power of 2 so access should not be power of 2



What if we increase shit to power of 2

$200 \rightarrow 256$

$256 \times 8$  bytes away from each other: fucks up our placement of memory into cache  
Kills spatial/temporal locality

## Semantics

What does a program mean? (as opposed to syntax, what does a program look like)

### static semantics

- how much you can figure out about a program while it's not running
- harder than syntax to find out
- **attribute grammars**

### dynamic semantics

- meaning when the program is running
- hardest to find out
- **operational semantics**
- **axiomatic semantics**

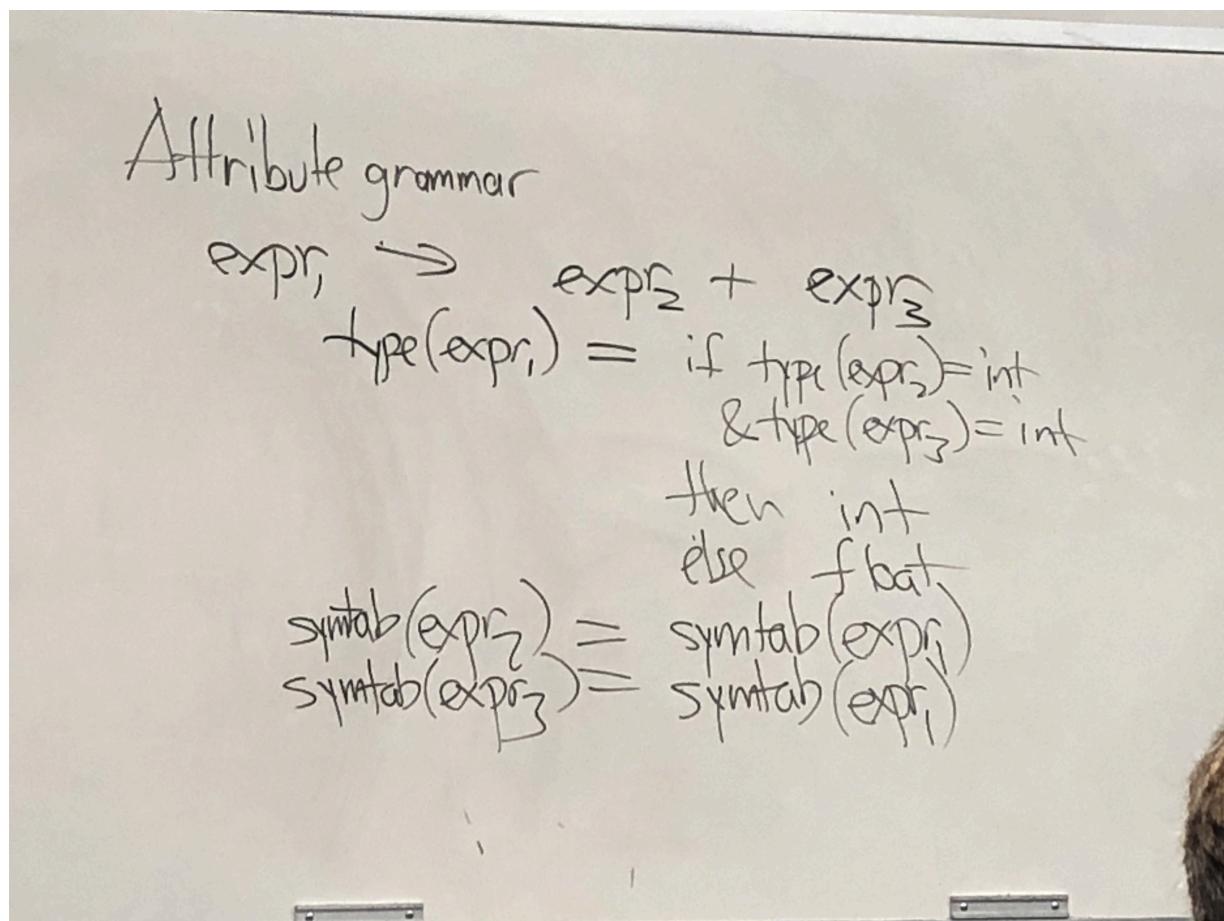
- denotational semantics

### Static Semantics:

#### Attribute Grammar

$\text{expr} \rightarrow \text{expr} + \text{expr}$

- we know is ambiguous: but don't care right now because that's a syntax problem



Attribute grammar is harder to find out, but once you have it you can write a compiler and

- specify type checking
- scope checking

### Dynamic Semantics :

#### Operational Semantics

- uses abstract imperative interpreter

#### Axiomatic Semantics

- uses axioms and rules of inference to define stuff

#### Denotational semantics

- maps programs to functions (mathematical)
- hardest one

Types of languages we've studies (language paradigms)

- logic
- functional
- imperative

People use specific semantics to talk about each paradigm of language: i.e.  
functional uses denotational semantics usually

The image shows handwritten notes on a whiteboard. At the top, there is a horizontal line with two metal brackets on either side. Below the line, there are three semantic rules:

$$\frac{\langle E_1, C \rangle \rightarrow v_1, \langle E_2, C \rangle \rightarrow v_2}{\langle E_1 + E_2, C \rangle \rightarrow v_1 + v_2}$$
$$\langle k, C \rangle \rightarrow k, \quad k \text{ a constant}$$
$$\langle v, C \rangle \rightarrow \text{lookup}(v, C), \quad v \text{ a variable}$$

$$\frac{\frac{\frac{\langle E_1, C \rangle \rightarrow v_1, \langle E_2, \text{bind}(x, v_1) :: C \rangle \rightarrow v_2}{\langle \text{let } x = E_1 \text{ in } E_2, C \rangle \rightarrow v_2}}{\langle \text{fun } x \rightarrow E, C \rangle \rightarrow \underline{\lambda(x, E, C)}}$$

$$\frac{\frac{\langle E, C \rangle \rightarrow \underline{\lambda(x_f, E_f, C)} \quad \langle E_2, C \rangle \rightarrow v_2, \langle E_f, \text{bind}(x_f, v_2) :: C_f \rangle \rightarrow v_3}{\langle E_1, E_2, C \rangle \rightarrow v_3}}{\langle E_1, E_2, C \rangle \rightarrow v_3}$$

the fact that the bottom line is  $c_f$  at the end rather than  $c$ , we know that we have static scoping (why??)

Semantics can show how to figure out if there's static scoping or dynamic scoping, more information, etc etc.