

Final Exam CHEAT SHEET

<p>comparing</p> <ul style="list-style-type: none"> - move ← + move → o stay <p>Tree Traversal</p> <p>pre: P L R (read/write) in: L P R (print) Post: L R P (destroy)</p>	<p>comparator Object specific order</p> <p>comparable natural order implementation</p>	<p>ArrayList · O(n) add</p>	<p>x2 LL</p> <ul style="list-style-type: none"> · Remove end ↳ go bktwds. Set pointers O (z+n) 												
			Insert (B/w) GetMin (peek/pop)												
		<table border="0" style="width: 100%;"> <tr> <td style="width: 33%;">unsorted vector</td> <td style="width: 33%;">O(1)</td> <td style="width: 33%;">O(n)</td> </tr> <tr> <td>sorted vector</td> <td>O(log(n+n)) / O(n)</td> <td>O(1)</td> </tr> <tr> <td>Heap</td> <td>O(log n)</td> <td>O(1) / O(log n)</td> </tr> <tr> <td>Bal. BST</td> <td>O(log n)</td> <td>O(1) / O(log n)</td> </tr> </table>	unsorted vector	O(1)	O(n)	sorted vector	O(log(n+n)) / O(n)	O(1)	Heap	O(log n)	O(1) / O(log n)	Bal. BST	O(log n)	O(1) / O(log n)	
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			<p>[Huff Node] impl. comparator value, weight, left, right return myWeight - other.myWeight</p> <p>[readHeader](in) → HN if (in.rB(1) == 0) HN.left = readHeader(in) HN.right = same↑ r n HN(o, o, left, right) e r n HN(in.rB(1), o)</p>												
<p>PQ</p> <ul style="list-style-type: none"> Search O(N) Insert O(1) / O(log N) Delete O(N) FindMin O(1) <p>Trie O(w) for all TreeSet O(log N) for all</p> <p>HashSet</p>	<p>[Huff Decoding] (S arch, S[] dict) → S Map <S → S> put (dict[""], ... character.toS(c char)(i+fs))) out = newSB int start = 0 end = 1 while (end < curr.l+1){ S str = curr.subS(start, end) if (map.containsKey(str)) out.append (map.get(str)) start = end end = start + 1 else end++ } } out.toS</p> <p>[parseBody](HN curr, in, out) HN.root = curr w(T) int i = in.rB(1) if (i == -1) break if (i == 1) curr = curr.right else curr = curr.left ← if 0 if (isLeaf(curr)) if curr.value == PSEUDO_EOF return else { out.wB(s, curr.value) curr = root throw n HuffException }</p>	<p>[makeTree](in → HuffNode) int[] f = count (in) //arr char f PQ<HN> pq for (i=0 → ALPH_SIZE){ if (f[i] != 0) pq.add (n HN(i, f[i])) } pq.add (n HN(PSEUDO_EOF, 0)) w (pq.size > 1) HN a = pq.poll HN b = pq.poll HN c = n HN(-1, a.wt + b.wt, a, b) pq.add (c) } pq.poll * array = n S[ALPH_SIZE+1] [extract](+N curr, S path) if (!isLeaf(curr)) array[curr.value] = path return extract (curr.left, path + 0) extract (curr.right, path + 1)</p> <p>[writeHeader](HN curr, out) if (!isLeaf(curr)) out.writeBits (1, 1) out.writeBits (9, curr.value) return out.writeBits (1, 0)</p> <p>[compress](in, out) HN huff = makeTree (in) array = n S[ALPH_SIZE + 1] extract (huff, "") out.writeBits (B_P_I, HUFF #) writeHeader (huff, out) writeBody (in, out)</p>	<p>[decompress](in, out) if (in.rB(B_P_I) != HUFF #) throw n HuffException HN curr = readHeader (in) parseBody (current, in, out)</p> <p>[Map Iterate] for (Map.Entry<S, I> entry : M.entrySet())</p> <p>[Scanner Tricks] w (in.hasNext ()) S str = in.nextLine ().replace (":", ",") S[] words = str.split (", ") in.close()</p>												
			* Make sure to init maps + check contains												

[Erdos Number] HM vert $\langle S \rightarrow \text{Cell}(S) \rangle$ HM erdos $\langle S, I \rangle$	[Insertion] · insert unsort \rightarrow sorted list	[Quick Sort] · Used for primitive [] · choose pivot \rightarrow split vector · very fast when fine-tuned	[MergeSort] · Used for Objects · Recursive (split) · good for // (work on diff. sections)
[addToMap] $\langle S[] o, HM M \rangle$ $Set(S) t = n HS(S)(A.asList(a))$ for $(S s : t)$ if $(M.containsKey(s))$ $M.get(s).addAll(t)$ $M.get(s).remove(s)$ else ~ same ↑ (Map: name \rightarrow collab)	[Selection Sort] · Div. into sort/unsort regions · Move max unsort to sorted · Heap Sort is improved version	Array $O(1)$ access, $O(N)$ everything else Stack $O(N)$ access/search, $O(1)$ insert/delete LLx1 $O(N)$ access/search, $O(1)$ insert/delete (unsorted) LLx2 Same as above HashTable $O(1)$ best, $O(N)$ worst BST $O(\log(n))$ balanced, $O(N)$ unbalanced Binary Heap $O(1)$ find max, $O(\log N)$ insert/delete/poll	
[getAdjList] $\langle S[] pub \rangle \rightarrow HM(S, \text{Cell}(S))$ HM $\langle S, \text{Cell}(S) \rangle M$ for $(i = 0 \rightarrow pubs.l)$ SC names = pubs[i].split(" ") addToMap(names, M) return M		Adj. List $O(V+E)$ storage, $O(1)$ add V/E worst space QuickSort $O(N \log N)$ best/avg, $O(N^2)$ worst, $O(\log N)$ MergeSort $O(N \log N)$ all, $O(N)$ worst space HeapSort $O(N \log N)$ all, $O(1)$ worst space InsertionSort $O(N)$ best, $O(N^2)$ avg/worst, $O(1)$ worst space SelectionSort $O(N^2)$ all, $O(1)$ worst space	
[bfs] $\langle S \text{ start} \rangle \rightarrow TS(S)$ erdos = $n HM(S, I)$ erdos.put(start, 0) // base TS(S) visited = $n TS(S)$ Q(S) qu = $n LL(S)$ visited.add(start) qu.add(start) while (qu.size > 0) S v = qu.remove for (S adj : vert.get(v)) if (!visited.contains(adj)) visited.add(adj) qu.add(adj) erdos.put(adj, erdos.get(v)+1)		[ith] $\langle LN \text{ list}, \text{int } i \rangle \rightarrow LN$ for $(k = 0 \rightarrow i \& \& list \neq \emptyset)$ list = list.next r list	
[calc Number] $\langle S[] pub \rangle$ vert = getAdjList(pub) Set(S) noΣ = vert.keySet TS(S) myV = bfs("ERDOS") noΣ.removeAll(myV) TS(S) out = $n TS(S)$ for $(S s : myV)$ S = S + " " + myΣ.get(s) out.add(s) TS(S) ref ref.addAll(out) ref.addAll(noΣ) return ref.toA($n S[\text{out.size}]$)	[splice] $\langle LN \text{ list}, \text{int } begin, \text{int } end \rangle$ +1 if (list = \emptyset) r \emptyset +2 LN splice = ith(list, begin-1) if (splice = \emptyset) r splice LN result = splice.next LN resultTail = ith(result, end - begin - 1) if (resultTail $\neq \emptyset$) splice.next = resultTail.next resultTail.next = \emptyset else // end = l-1 splice.next = \emptyset r result		

Comparisons

ArrayList vs. ArrayList<String>

new String[n] new ArrayList<String>()

.length .size()
[k] .get(k)

Arrays.sort(un) Collections.sort(un)

Comparable vs. Comparator

compareTo method
• "natural" order
• ⚡ if this < 02
• o if this == 02
• ⚡ if this > 02

compare method
• mult. sorting methods
• 2 arguments

Path Exists
// check out of bounds

```
if (isGoal(row, col)) r T
if (isVisited(row, col)) r F
markVisited(row, col)
r (!isBlocked(row, col, row1, col))
&& pathExists(row1, col)) || ...
```

Exam Two Cheat Sheet

by Brian Lin

Forbes' Pro Stacks

Reverse Sort → Arrays.sort(a, Collections.reverseOrder())
Collections.reverse(arrlist);

TreeSet<String>((s1, s2) → -s1.compareTo(s2));

Big O

N=100 → t = 5 ms
N=500 → t = ? using nlogn
S=k (100 log₂ 100)
? = k (500 log₂ 500)

Recurrence Relations

T(N) = T(?) + O(?)
Previous recursive step Other operations
Balanced
Unbalanced
$$\begin{aligned} \text{curr} &= \text{root} \\ \text{for } (i=0 \rightarrow w, l) & \\ \text{if } (\text{curr}.children.contains(\text{curr}.children[i])) & \\ \text{w}.charAt(i)) & \\ \text{curr} &= \text{curr}.children[i] \\ \text{get}(\text{charAt}(i)) & \\ \text{else curr}.children.forEach(\text{curr}.children[i], \text{w}.charAt(i), n NC) & \\ \text{curr} &= \text{curr}.children.get(...) \end{aligned}$$

curr.isWord = T
Union (N l, N r)
if (l == 0) r copyTri(r)
else if (r == 0) r copyTri(l)
else {
 N curr = r NC
 bool hasBelow = F
 // get set of all children
 HS<ch> allC = r.getChildren(l.chil.keySet())
 allC.addAll(r.chil.keySet())
 for (char ch : allC)
 N uChil = trieUnion(l.chil.get(ch),
 r.chil.get(ch))
 if (uChil != null) // something found
 hasBelow = T
 curr.chil.put(ch, uChil)
 curr.isWord = l.isWord || r.isWord // in union?
 if (hasBelow || curr.isWord) // something to r
 r curr
 else r
intersection addAll/retainAll
curr.isWord = l.isWord && r.isWord

DraftTree

```
firstIndexof(TermC) a, Term key, Comp)
{low=1
 i mid=0
 i high=a.l-1
 c = comp.compare(a[mid].key)
 w(h-l+1)
 m=(l+h)/2
 if (c>0) h=m
 if (c<0) l=m
 if (comp.compare(a[(l+h)/2], key)==0)
    r (h+l)/2
 if (comp.compare(a[(l+h)/2+1], key)==0)
    r (h+l)/2+1
 else r -1
 &C].topLevelMatches(&pref, i, k)
 if (pref==0) error
 p&cT) pg=n P&cT...
 (k, nT.weightOrder())
 i lst = fIOF(-terms, ...
 nT.prefixOrder(pref, l))
 i last=Same
 if (lst==i) r n &C]
 T[] &= n T(last-l+1)
 F(i=lst, ind=0 → ind < ?)
 if (pg.size() < k)
    pg.add(t)
 else if (pg.parent().getW ...
 &t, optW)
    pg.remove()
    pg.add(t)
 i numbers = Math.min(k, pg.size())
 &C] ind=n &C(numbers)
 for (i=0 → numbers)
    ref[numbers-1-i]=pg.remove()...
    .getWord()
 r ref
```

```
pr i root
L<N> // 
pr cl N
pr &s one
pr L<N> child
pr N(&s)
m=&s
ch=&n AL<N>
pr &s[] draw((&s) par, &s[] names)
N=&n AL<N>()
pr (i=0 → par, l)
N.add(n NC.names[i])
pr (i=0 → par, l)
if (par[i]==-1) root=i
else N.get(par[i]).child[i]...
add(&n, get(par[i]).child[i])
L<&s> out = draw(N, get(root, F)
r out, toArea(&s[] par, l))
pr bool isLast(&n, i)
if (n.chil.size() == i) T
else ret F
pr L<&s> draw(&n, bool drawBar)
L<&s> l = n AL<&s>()
n.mre = "+" + n.mre
&s str
if (drawBar) str = "| "
else str = ""
l.add(n.mre)
for (i=0 → n.chil.size())
    L<&s> cList = draw...
    (n.chil.get(i), !isLast(n, i))
for (&s : cList)
    s = str + s
    l.add(s)
```

r l

