

2.71

Problem

We want to create a data structure where 4 signed bytes are packed into a 32-bit unsigned int (`typedef unsigned packed_t`). We need to be able to extract a desired byte from the packed “word”.

Visualization

```
1100 0011 1010 0101 0011 1100 0101 1010  <- unsigned int
+---+---+ +---+---+ +---+---+ +---+---+
|         |         |         |
Byte 3    Byte 2    Byte 1    Byte 0
```

Answers

A. Faulty implementation

The code is trying to bit shift the word in such a way that the desired byte is located at the last 8 bits of the shifted word.

```
int xbyte(packed_t word, int bytenum) {
    int shifted_word = (word >> (bytenum << 3)); // shift word by (bytenum * 8)
    return (shifted_word & 0xFF); // extract only last 8 bits of shifted_word
}
```

The issue lies in how the last 8 bits are extracted. Since the shifted word is always bitwise and’ed with 0xFF, the padding is always filled with zeros. This is undesired since our byte should represent a *signed* int, which if negative, should be padded by ones.

B. Correct implementation

```
int xbyte(packed_t word, int bytenum) {
    // shift word so that byte occupies the 1st 8 bits (leftmost side)
    int left_shifted_word = (word << (3 - (bytenum << 3)));

    // shift byte to rightmost side (padded by ones if neg. else zeros)
    return (left_shifted_word >> 24);
}
```

2.82

Problem

The following is the initial code for the problem:

```

// ints are 32 bits for this problem
int x = random();
int y = random();

// unsigned ints are 32 bits for this problem
unsigned ux = (unsigned) x;
unsigned uy = (unsigned) y;

```

Given the expressions for each problem, determine if they are *always* true. If yes, show a mathematical proof, else give a counterexample.

Answers

A. $(x < y) == (-x > -y)$ False, if x is $TMIN$, then the first expression is always true for any value of y that is greater than $TMIN$. However, in that case, $-x$ is $TMIN$, which any value of $-y$ cannot be less than.

For example:

| x | y | $(x < y)$ | $(-x > -y)$ |
|--------|---|-----------|-------------|
| $TMIN$ | 0 | 1 | 0 |

B. $((x + y) \ll 4) + (y - x) == 17*y + 15*x$ True. If one side overflows the other overflows as well.

$$\begin{aligned}
 ((x + y) \ll 4) + (y - x) &= 17y + 15x \\
 (x + y) \ll 4 &= 16y + 16x \\
 &= 16(y + x) \\
 &= 2^4(y + x)
 \end{aligned}$$

C. $\sim x + \sim y + 1 == \sim(x + y)$ True, by using definition of negation: $-x == \sim x + 1$.

$$\begin{aligned}
 \sim x + \sim y + 1 &= (-x - 1) + (-y - 1) + 1 \\
 &= -(x + y) - 1 \\
 &= \sim(x + y)
 \end{aligned}$$

D. $(ux - uy) == \text{-(unsigned)}(y - x)$ True. The explicit cast serves no purpose because it only alters the interpretation of the bits. Using the definition of negation, the result of negating $\text{-(unsigned)}(y - x)$ and $(y - x)$ would be the same. Therefore, we can distribute in the negative.

$$\begin{aligned} -(\text{unsigned})(y - x) &= -(y - x) && \text{(same bit pattern)} \\ &= x - y \end{aligned}$$

Since we are doing a comparison, both sides are interpreted as unsigned, therefore

$$\begin{aligned} ux - uy &= x - y \\ &= -(\text{unsigned})(y - x) \end{aligned}$$

E. $((x \gg 2) \ll 2) \leq x$ True

$$\begin{aligned} ((x \gg 2) \ll 2) &= \text{floor}(x/4) \cdot 4 \\ &\leq (x/4) \cdot 4 \\ &= x \end{aligned}$$

2.83

Problem

Consider binary numbers with a representation of the form $0.y \ y \ y \ y \dots$ where y is a k -bit sequence. Let Y represent the decimal value for y .

For example:

1. $1/3 \rightarrow 0.01010101\dots$ ($y = 01$, $k = 2$, $Y = 1$)
2. $1/5 \rightarrow 0.00110011\dots$ ($y = 0011$, $k = 4$, $Y = 3$)

Answers

A. Find the expression for the value of these binary numbers in terms of Y and k

The expression to find the base-10 value of the given form of a binary number is the following:

$$\sum_{n=1}^{\infty} Y \left(\frac{1}{2^k} \right)^n$$

This can be written as a infinite geometric series, which has the following property:

$$\sum_{n=0}^{\infty} ar^n = \frac{a}{1-r}$$

Re-arranging our equation so that we can exploit this, we get:

$$\left[\sum_{n=0}^{\infty} Y \left(\frac{1}{2^k} \right)^n \right] - Y$$

Finally, we get:

$$\frac{Y}{1 - \frac{1}{2^k}} - Y = \left(\frac{Y}{2^k - 1} \right)$$

B. Numeric Values of Examples

(a) $y = 101$

- $k = 3$
- $Y = 5$
- **numeric value** = $\frac{5}{6}$

(b) $y = 0110$

- $k = 4$
- $Y = 6$
- **numeric value** = $\frac{6}{15}$

(c) $y = 010011$

- $k = 6$
- $Y = 19$
- **numeric value** = $\frac{19}{65}$