```
result := {R};

done := false;

compute F +;

while (not done) do

if (there is a schema R_i in result that is not in BCNF)

then begin

let \alpha \to \beta be a nontrivial functional dependency that

holds on R_i such that \alpha \to R_i is not in F +,

and \alpha \cap \beta = \emptyset;

result := (result -R_i) \cup (R_i - \beta) \cup (\alpha, \beta);

end

else done := true;
```

To check if a relation  $R_i$  in a decomposition of R is in BCNF,

- Either test R<sub>i</sub> for BCNF with respect to the restriction of F to R<sub>i</sub> (that is, all FDs in F<sup>+</sup> that contain only attributes from R<sub>i</sub>)
- or use the original set of dependencies F that hold on R, but with the following test:
  - for every set of attributes  $\alpha \subseteq R_i$ , check that  $\alpha^+$  (the attribute closure of  $\alpha$ ) either includes no attribute of  $R_{\Gamma}$  α, or includes all attributes of  $R_i$ .
  - If the condition is violated by some  $\alpha \to \beta$  in F, the dependency

 $\alpha \rightarrow (\alpha^+ - \alpha) \cap R_i$ 

can be shown to hold on  $R_i$ , and  $R_i$  violates BCNF.

We use above dependency to decompose R<sub>i</sub>

```
Let F_c be a canonical cover for F; i := 0;
for each functional dependency \alpha \to \beta in F_c do
 if none of the schemas R_i, 1 \le j \le i contains \alpha \beta
         then begin
                 R_i := \alpha \beta
            end
if none of the schemas R_i, 1 \le j \le i contains a candidate key for R
 then begin
            i := i + 1;

R_i := \text{any candidate key for } R_i;
         end
/* Optionally, remove redundant relations */
repeat
if any schema R_i is contained in another schema R_k
      then /* delete R_i */
        R<sub>j</sub> = R;;
i=i-1;
return (R<sub>1</sub>, R<sub>2</sub>, ..., R<sub>i</sub>)
```

Consider a set F of functional dependencies and the functional dependency  $\alpha \to \beta$  in F.

To test if attribute A  $\in \alpha$  is extraneous in  $\alpha$ 

- 1. compute  $({\alpha} A)^+$  using the dependencies in F
- 2. check that  $({\alpha} A)^+$  contains β; if it does, A is extraneous in α

To test if attribute  $A \in \beta$  is extraneous in  $\beta$ 

- 1. compute  $\alpha^+$  using only the dependencies in  $\mathsf{F}' = (F \{\alpha \to \beta\}) \cup \{\alpha \to (\beta A)\},$
- 2. check that  $\alpha^+$  contains A; if it does, A is extraneous in  $\beta$