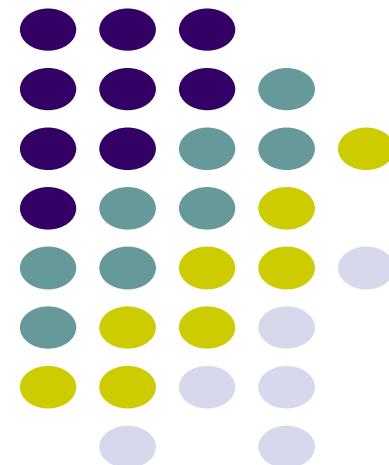


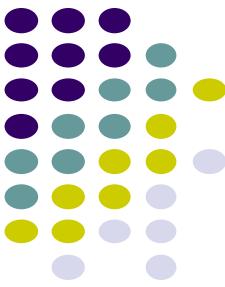
L29: TRIE和AC自动机

吉林大学计算机学院

谷方明

fmgu2002@sina.com



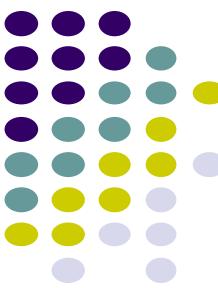
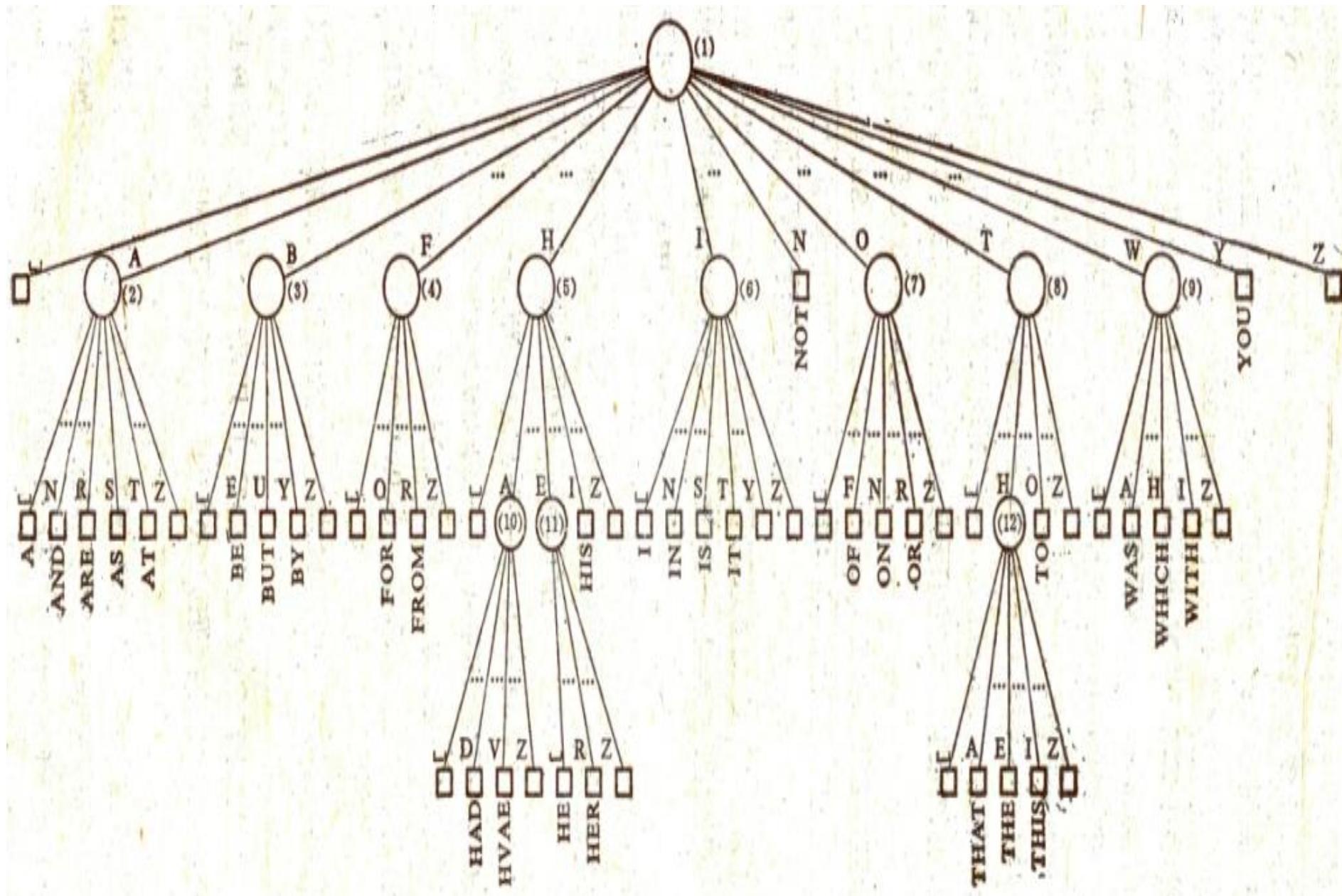


字符串查找

口 例：有 31 个最常用的英文单词的集合如下

A , AND , ARE , AS , AT , BE , BUT , BY , FOR , FROM , HAD ,
HAVE , HE , HER , HIS , I , IN , IS , IT , NOT , OF , ON , OR ,
THAT , THE , THIS , TO , WAS , WHICH , WITH , YOU

给定一个单词，查询是否在最常用单词集合中

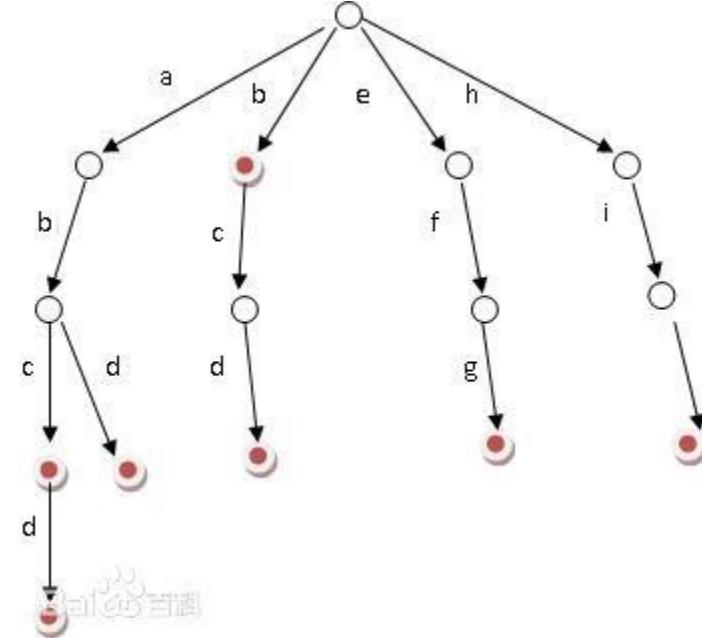


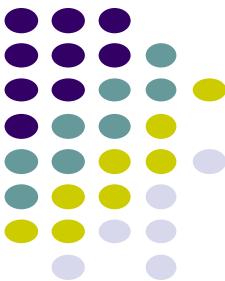
M叉检索树



字典树

- 单词查找树、 TRIE树等
- 字典树建立在字符集 Σ 上，
 $|\Sigma| = M$
- 字典树是一种M叉树
 - ✓ 根节点不包含字符，根节点外的每一个节点都包含一个字符
 - ✓ 每个节点的所有子节点包含的字符都不相同。
 - ✓ 根节点到某一节点的路径上的字符连起来，为该节点对应的字符串；





存储结构

□ 每个结点存储M个儿子即可

□ 顺序存储

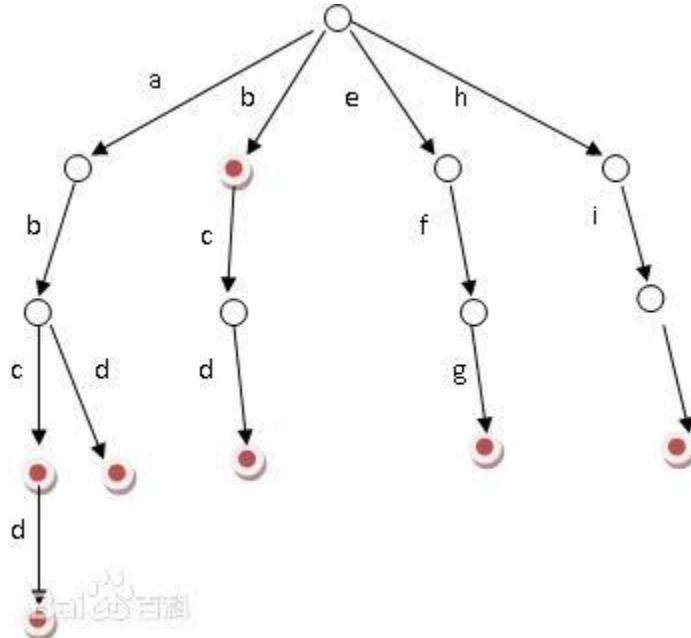
✓ $O(M^{MAXL})$, 浪费空间

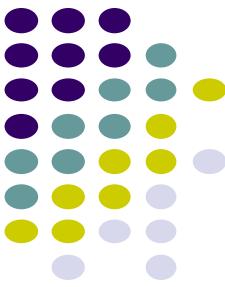
□ 链接存储

✓ $O(N * MAXL * M)$

✓ 共用前缀，一般达不到。

✓ 可用静态链表实现。

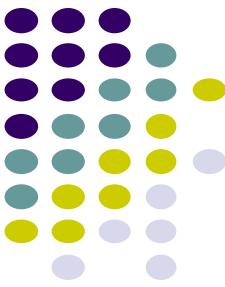




静态链表（结构体数组）

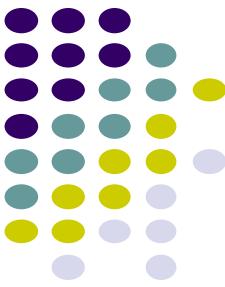
```
struct Node{  
    int son [ 26 ];  
    int flag; //单词结束标志  
};
```

```
Node trie[ MAXN ];  
int sp;
```



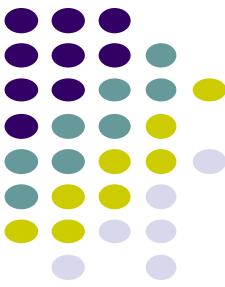
插入操作

```
//设根节点为0
void insertWord(char s[]) {
    int i,p=0,len=strlen(s),c;
    for(i=0; i<len; i++) {
        c=s[i]-'a';
        if(!trie[p].son[c]) trie[p].son[c] = ++sp;
        p = trie[p].son[c];
    }
    trie[p].flag = 1;
}
```



查询操作

```
//设根节点为0
int query(char s[]){
    int i,p=0,len=strlen(s),c;
    for(i=0;i<len;i++){
        c=s[i]-'a';
        if(!trie[p].son[c]) return 0;
        p = trie[p].son[c];
    }
    return trie[p].flag == 1;
}
```



分析

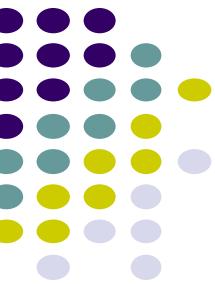
□ 时间复杂度

- ✓ 插入: $O(N * MAXL)$
- ✓ 查找: $O(MAXL)$

□ 空间复杂度

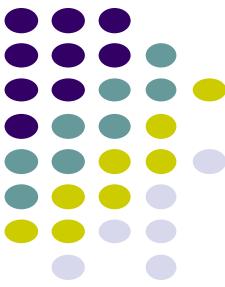
- ✓ $O(N * MAXL) * O(M)$

□ 时间效率比Hash快，但浪费空间



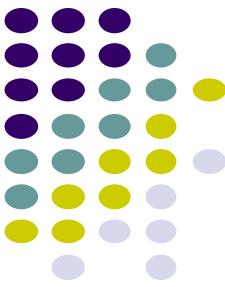
Trie树应用广泛

- 串的快速检索
- 最长公共前缀
- “串”排序
- 辅助结构



多模匹配

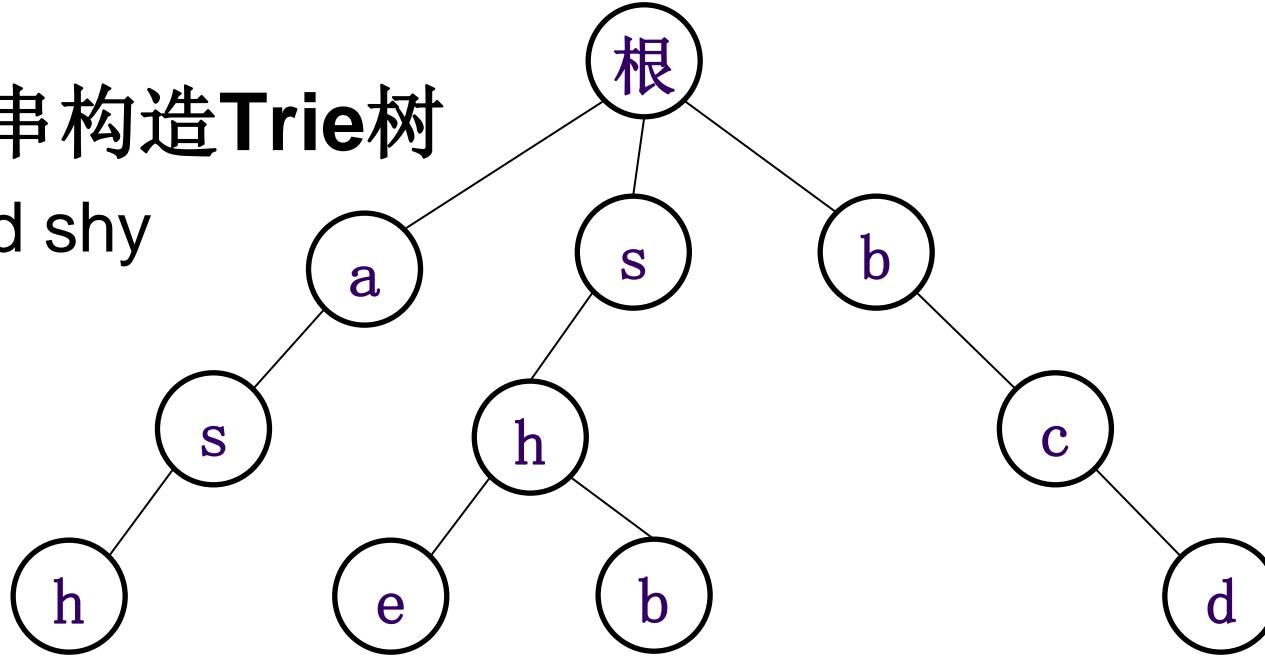
- 文本串: **ashecd**
- 模式串: **ash she bcd shb**
- 询问: 模式串在文本串中出现的次数
- K遍KMP算法?
 - ✓ 效率低



AC自动机

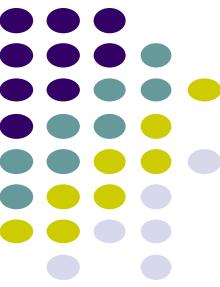
- 用多个模式串构造Trie树

- ✓ ash she bcd shy



- 文本串在Trie树上进行匹配

- ✓ ashecd



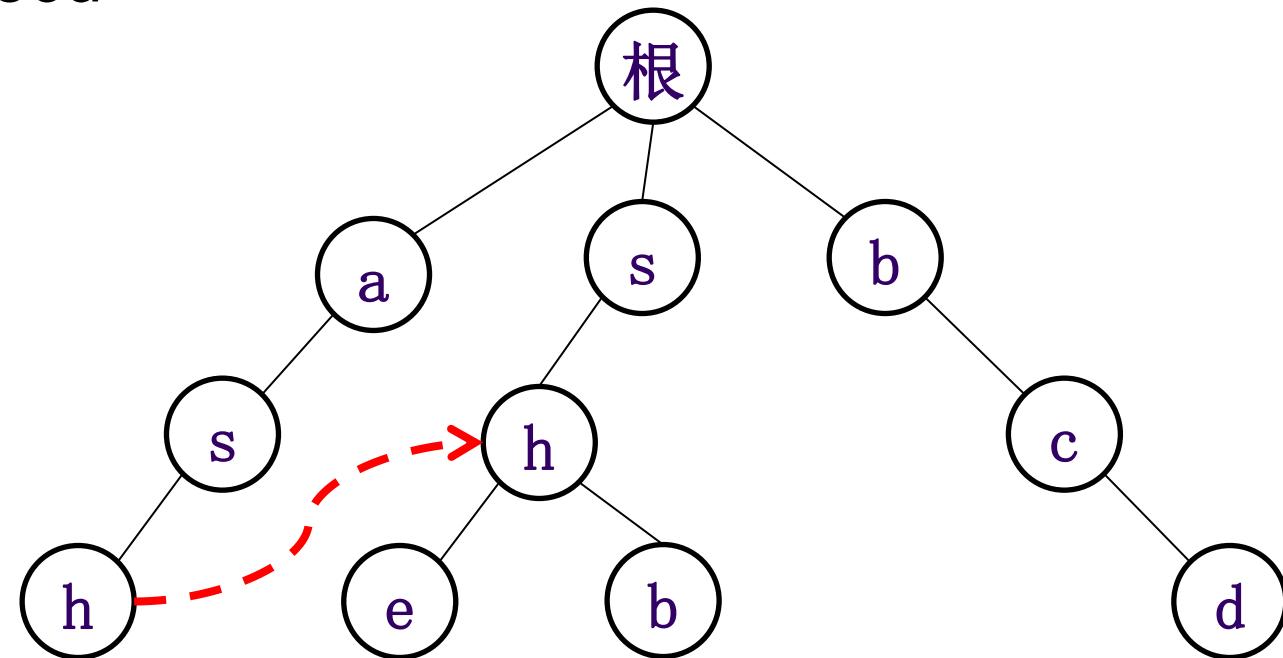
fail指针

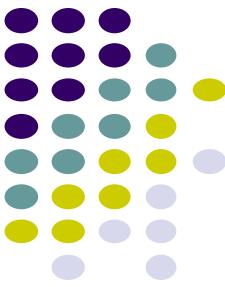
□ 类似KMP，引入fail指针

- ✓ 例：文本串ashecd

□ fail指针的意义

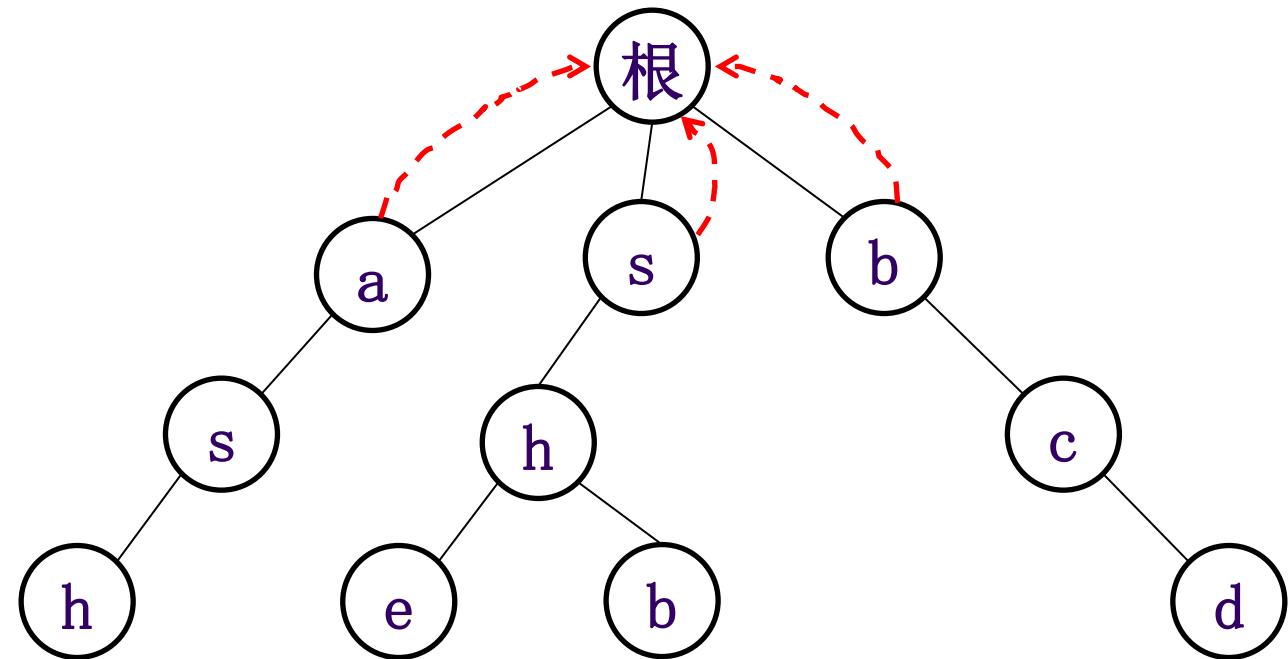
- ✓ 最长公共后缀

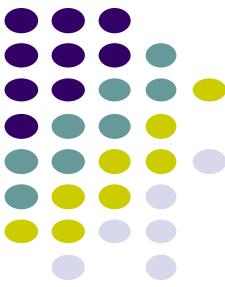




fail指针的建立

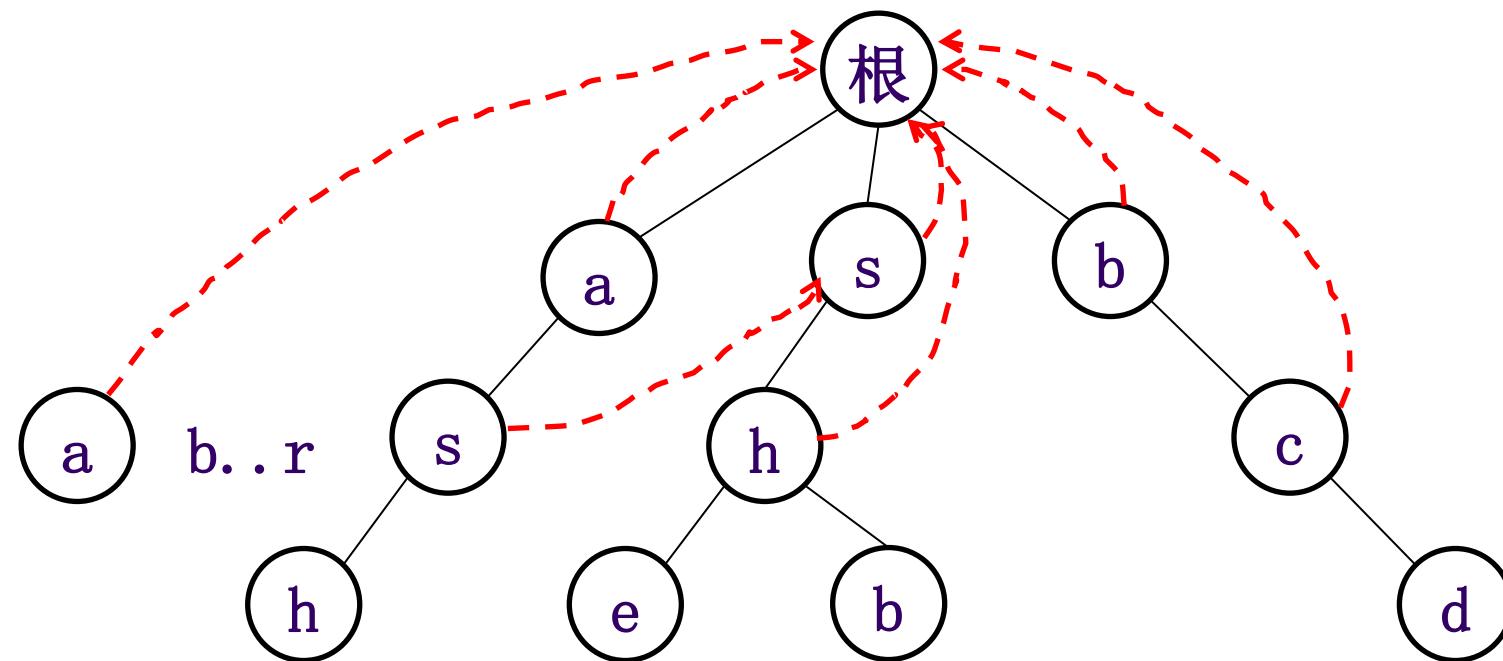
□ 第一层

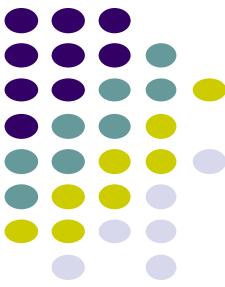




fail指针的建立

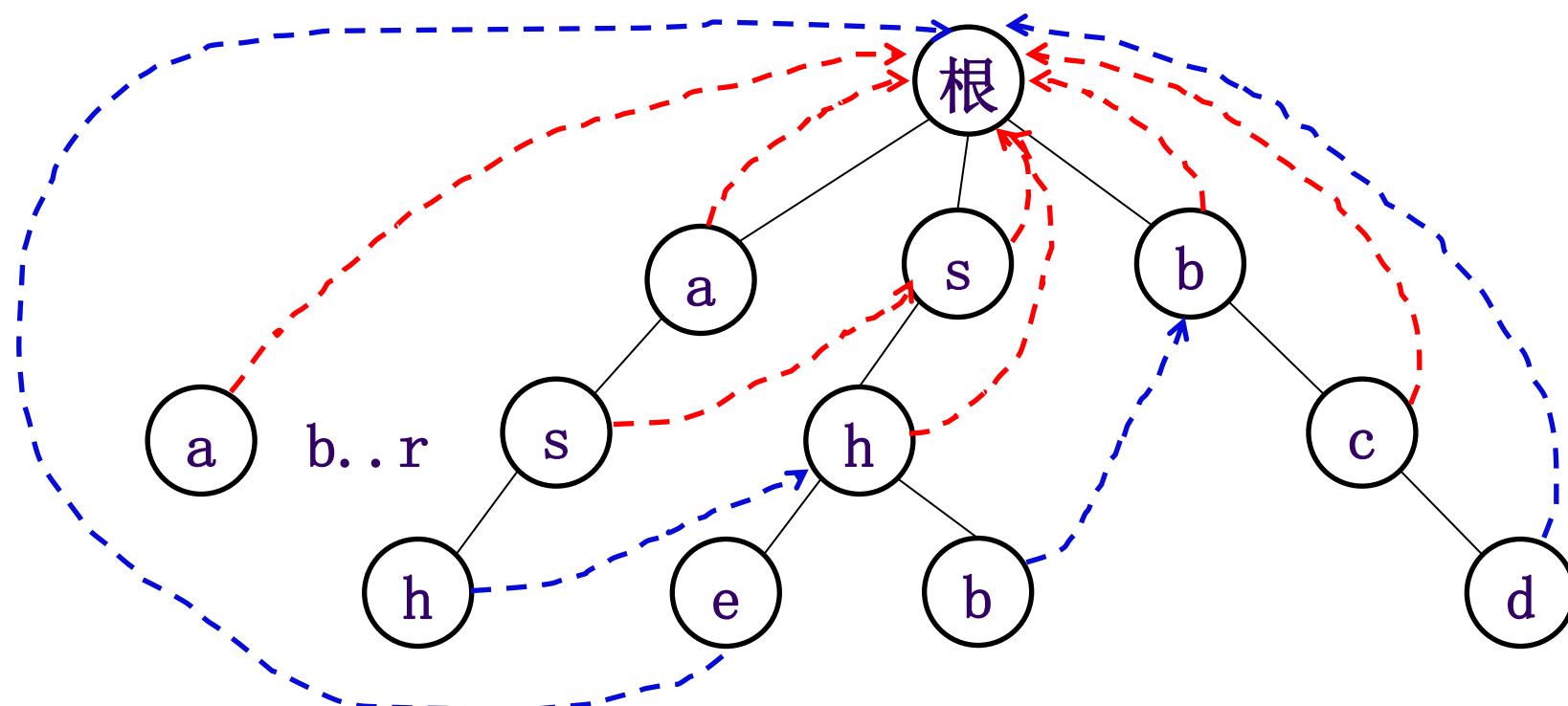
- 第二层
- BFS: a 和 fail(a) 具有相同的儿子s

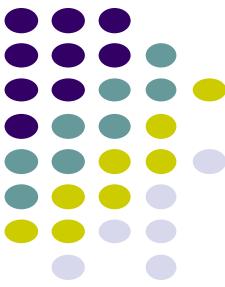




fail指针的建立

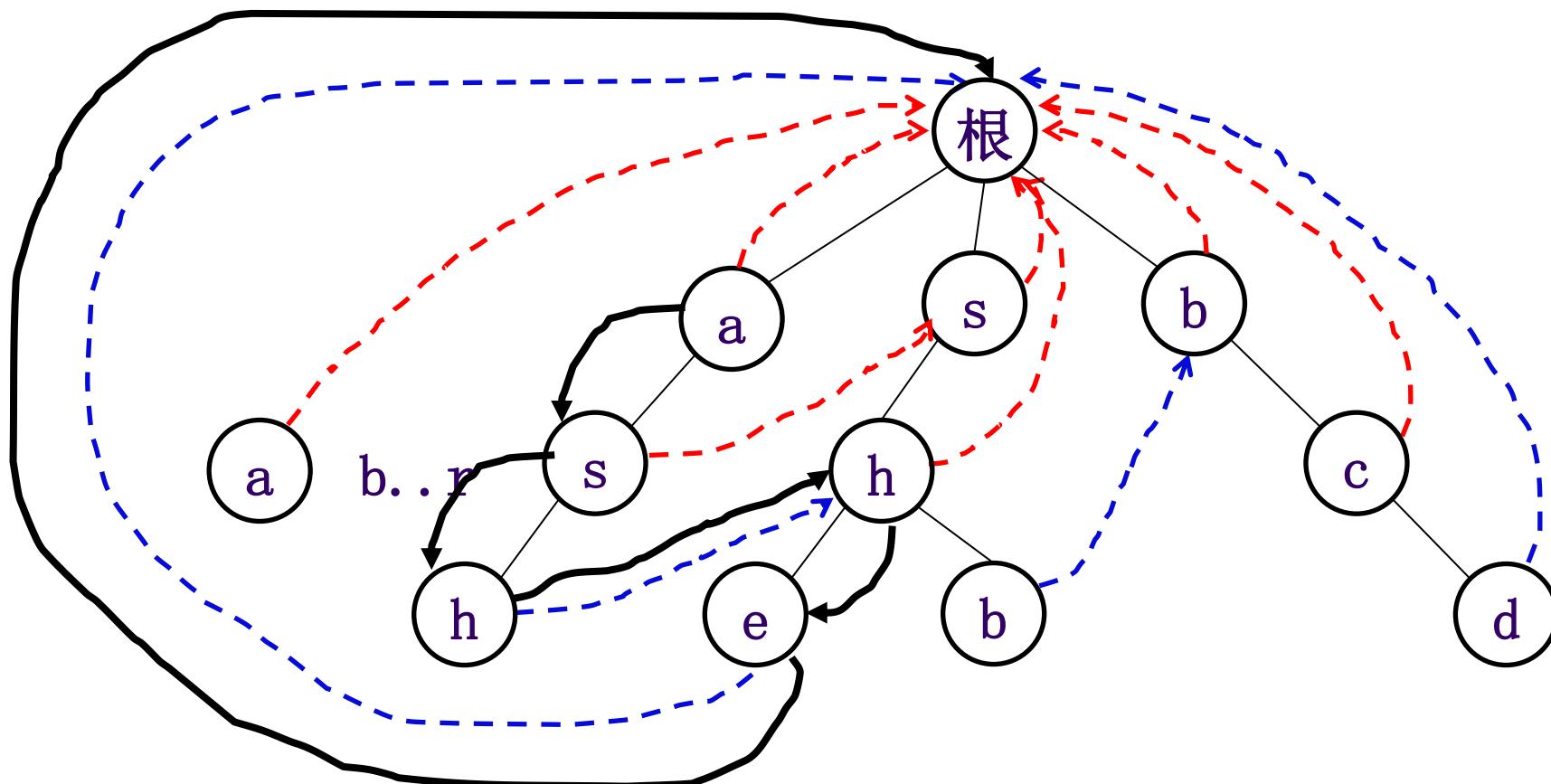
□ 第三层

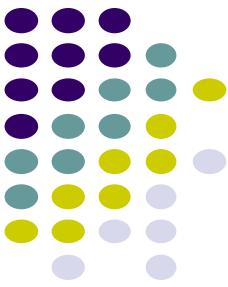




匹配过程

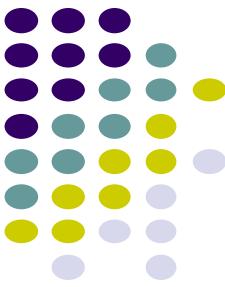
- 文本串 ashecd



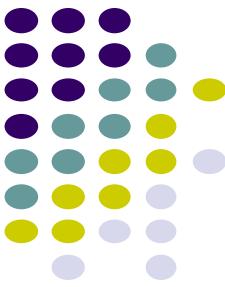


建立fail指针参考实现（朴素版）

```
void getFail()
{
    int i,now;
    queue<int> q;
    for(i=0;i<26;i++)
        if(trie[0][i]){
            fail[trie[0][i]]=0;
            q.push(trie[0][i]);
        }
}
```



```
while(!q.empty()){
    now=q.front();q.pop();
    for(i=0;i<26;i++)
        if(trie[now][i]){
            fail[trie[now][i]]=trie[fail[now]][i];
            q.push(trie[now][i]);
        }else{
            trie[now][i]=trie[fail[now]][i];
        }
}
}//时间复杂度O(|S|), |S| =Σ |sil|
```



多模匹配参考代码（朴素版）

```
int query(char s[]) {  
    int now=0,len=strlen(s),i,j,ans=0;  
    for(i=0;i<len;i++){  
        now = trie[now][s[i]-'a'];  
        for(j=now; j&&cntWord[j]!=-1;j=fail[j]){  
            ans += cntWord[j];  
            cntWord[j] = -1;  
        }  
    }  
    return ans;  
} // 时间复杂度O(|T|)
```