# Report

# **TLB Replacement Policy**

- LRU
  - sys\_config.txt

```
TLB Replacement Policy: LRU
Page Replacement Policy: FIFO
Frame Allocation Policy: GLOBAL
Number of Processes: 2
Number of Virtual Page: 128
Number of Physical Frame: 64
```

o analysis.txt

```
Process A, Effective Access Time = 164.758

Process A, Page Fault rate = 0.723

Process B, Effective Access Time = 163.709

Process B, Page Fault rate = 0.665
```

- RANDOM
  - sys\_config.txt

```
TLB Replacement Policy: RANDOM
Page Replacement Policy: FIFO
Frame Allocation Policy: GLOBAL
Number of Processes: 2
Number of Virtual Page: 128
Number of Physical Frame: 64
```

analysis.txt

```
Process A, Effective Access Time = 165.200
Process A, Page Fault rate = 0.723
Process B, Effective Access Time = 163.333
Process B, Page Fault rate = 0.665
```

 若範例input來看,LRU和RANDOM並沒有造成太大的差距,稍微可以看出使用RANDOM的時候Process A的TLB miss比較多,Process B的TLB miss比較少,因為TLB miss 但Page hit的時候會多一次查詢時

### 間,但不會增加Page Fault

• 造成這個結果可能是因為input file並沒有很規律的重複reference相同的page,如果process一直 reference差不多的page的時候,LRU的效果會比reference好

### Page Replacement Policy

- FIFO
  - sys\_config.txt

```
TLB Replacement Policy: LRU
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Frame Allocation Policy: GLOBAL
Number of Processes: 2
Number of Virtual Page: 128
Number of Physical Frame: 64
```

o analysis.txt

```
Process A, Effective Access Time = 164.758

Process A, Page Fault rate = 0.723

Process B, Effective Access Time = 163.709

Process B, Page Fault rate = 0.665
```

- CLOCK
  - sys\_config.txt

```
TLB Replacement Policy: LRU
Page Replacement Policy: CLOCK
Frame Allocation Policy: GLOBAL
Number of Processes: 2
Number of Virtual Page: 128
Number of Physical Frame: 64
```

o analysis.txt

```
Process A, Effective Access Time = 164.758
Process A, Page Fault rate = 0.723
Process B, Effective Access Time = 163.709
Process B, Page Fault rate = 0.665
```

• 兩者看起來並沒有差異,因此我準備了LOCAL的來做比對

#### • FIFO

sys\_config.txt

```
TLB Replacement Policy: LRU
Page Replacement Policy: FIFO
Frame Allocation Policy: LOCAL
Number of Processes: 2
Number of Virtual Page: 128
Number of Physical Frame: 64
```

o analysis.txt

```
Process A, Effective Access Time = 164.980
Process A, Page Fault rate = 0.774
Process B, Effective Access Time = 163.144
Process B, Page Fault rate = 0.700
```

### CLOCK

sys\_config.txt

```
TLB Replacement Policy: LRU
Page Replacement Policy: CLOCK
Frame Allocation Policy: LOCAL
Number of Processes: 2
Number of Virtual Page: 128
Number of Physical Frame: 64
```

o analysis.txt

```
Process A, Effective Access Time = 164.980
Process A, Page Fault rate = 0.774
Process B, Effective Access Time = 163.522
Process B, Page Fault rate = 0.694
```

• 使用clock 的話有機會使Page fault rate在local的時候變小,如果頻繁接觸某幾個page的話,clock policy可以保留頻繁接觸的frame但fifo不行

# Frame Allocation Policy

- GLOBAL
  - sys\_config.txt

```
TLB Replacement Policy: LRU
Page Replacement Policy: FIFO
Frame Allocation Policy: GLOBAL
```

Number of Processes: 2 Number of Virtual Page: 128 Number of Physical Frame: 64

### o analysis.txt

```
Process A, Effective Access Time = 164.758

Process A, Page Fault rate = 0.723

Process B, Effective Access Time = 163.709

Process B, Page Fault rate = 0.665
```

#### LOCAL

### sys\_config.txt

```
TLB Replacement Policy: LRU
Page Replacement Policy: FIFO
Frame Allocation Policy: LOCAL
Number of Processes: 2
Number of Virtual Page: 128
Number of Physical Frame: 64
```

### analysis.txt

```
Process A, Effective Access Time = 164.980
Process A, Page Fault rate = 0.774
Process B, Effective Access Time = 163.144
Process B, Page Fault rate = 0.700
```

- page fault的話global會有比local還要好的效果,因為global能夠根據當下process執行的狀況來分配適當的page,找到process比較夠用的page數量
- 但因為effective access rate受到TLB的影響比較多,所以影響沒有那麼大
- 另外,綜合frame policy和page policy,在範例測資中相同TLBpolicy下,frame policy對Page fault rate 的影響會比較大