## 1. Brute Force Approach:

## • Approach:

- Iterate over each character in the string.
- o For each character, remove it and check if the resulting string is a palindrome.
- o If any of the resulting strings is a palindrome, return true; otherwise, return false.

```
class Solution {
public:
    bool validPalindrome(string s) {
        for (int i = 0; i < s.length(); i++) {</pre>
            string modified = s.substr(0, i) + s.substr(i + 1);
            if (isPalindrome(modified)) {
                return true;
        }
        return isPalindrome(s); // Check if the original string is already
    }
private:
    bool isPalindrome(const string& s) {
        int i = 0, j = s.length() - 1;
        while (i < j) {
            if (s[i] != s[j]) {
                return false;
            i++;
            j--;
        return true;
};
```

## • Complexity:

- o **Time Complexity:** O(n^2) due to checking for palindrome after removing each character.
- Space Complexity: O(n) for creating substrings.

#### 2. Better Approach:

#### Approach:

o Use two pointers starting from the beginning and end of the string.

 If a mismatch is found, check if either the substring obtained by removing the left character or the one by removing the right character is a palindrome.

```
class Solution {
public:
    bool validPalindrome(string s) {
        int i = 0, j = s.length() - 1;
        while (i < j) {
            if (s[i] != s[j]) {
                return isPalindrome(s, i + 1, j) || isPalindrome(s, i, j -
1);
            i++;
        return true;
    }
private:
    bool isPalindrome(const string& s, int i, int j) {
        while (i < j) {
            if (s[i] != s[j]) {
                return false;
            i++;
        return true;
};
```

## • Complexity:

- **Time Complexity:** O(n), since we traverse the string at most twice.
- Space Complexity: O(1), as no extra space is required.

## 3. Optimal Approach:

# Approach:

- o The optimal approach is essentially the same as the better approach.
- The better approach already achieves the most efficient solution, ensuring minimal time complexity by checking both possibilities when a mismatch occurs.

```
• • •
class Solution {
public:
    bool validPalindrome(string s) {
        int i = 0, j = s.length() - 1;
        while (i < j) {
            if (s[i] != s[j]) {
                return isPalindrome(s, i + 1, j) || isPalindrome(s, i, j -
1);
            }
            i++;
        return true;
    }
private:
    bool isPalindrome(const string& s, int i, int j) {
        while (i < j) {
            if (s[i] != s[j]) {
                return false;
            i++;
        return true;
};
```

# Complexity:

- Time Complexity: O(n).
- Space Complexity: O(1).