

✅ Theory: LRU Page Replacement Algorithm

🔥 What is Page Replacement?

When a process tries to access a page that is not present in main memory (RAM), a page fault occurs. The OS must bring that page from disk into memory.

But if memory (frames) is full, OS must replace one of the existing pages using a page replacement algorithm.

🔥 LRU (Least Recently Used) Algorithm

- LRU replaces the page which has not been used for the longest time.
- It assumes past behavior predicts future behavior (i.e., page not used recently will not be used soon).
- It maintains a timestamp or counter to track when each page was last used.

✅ Advantages

- More efficient than FIFO, because it considers page usage history.
- Reduces page faults compared to FIFO.

❌ Disadvantages

- Requires extra storage and computation to maintain timestamps.
- Hardware support (counters/stack) may be required for efficiency.

✅ Code Explanation (Line-by-Line)

```
int frames, pages[100], frame[10], used[10], n, i, j, k;
```

```
int faults = 0, hits = 0, time = 0;
```

- frames → number of physical memory frames.
- pages[] → page reference string entered by user.
- frame[] → current pages stored in memory.
- used[] → stores timestamp of when each frame was last used.
- faults and hits → counters.
- time → logical counter to simulate timestamps.

```
printf("Enter number of frames: ");
```

```
scanf("%d", &frames);
```

Takes number of frames (e.g., 3 frames).

```
printf("Enter number of pages: ");
```

```
scanf("%d", &n);
```

Accepts length of reference string (e.g., 12 pages).

```
for(i = 0; i < n; i++)
```

```
    scanf("%d", &pages[i]);
```

Reads each page number into array.

✓ Initialization

```
for(i = 0; i < frames; i++)
```

```
    frame[i] = -1, used[i] = -1;
```

- -1 means empty frame.
 - used[] = -1 means no timestamp yet.
-

✓ Main Logic: Processing Each Page

```
for(i = 0; i < n; i++) {
```

```
    int page = pages[i], hit = 0;
```

- page → current page request
 - hit → becomes 1 if page already exists in memory
-

✓ Page Hit Checking

```
for(j = 0; j < frames; j++) {
```

```
    if(frame[j] == page) {
```

```
        hit = 1;
```

```
        used[j] = time++;
```

```
        hits++;
```

```
    }
```

```
}
```

- If page found → hit
 - Update timestamp of that frame → used[j] = time++
-

✓ Page Fault Handling

```
if(!hit) {
```

```

int pos = 0;
for(j = 0; j < frames; j++)
    if(frame[j] == -1 || used[j] < used[pos])
        pos = j;

```

- If page not found → Page Fault
- Select frame with least used timestamp → $used[j] < used[pos]$
- If frame is empty (-1) → directly use it

```

frame[pos] = page;
used[pos] = time++;
faults++;
}

```

- Insert new page into selected frame
- Update timestamp
- Increase fault count

✅ Displaying Output

```

printf("%d\t", page);
for(k = 0; k < frames; k++)
    printf("%s ", frame[k] == -1 ? "-" : (char [3]){frame[k]+'0',0});

```

Prints current memory state after every page request.

```

if(hit)
    printf("\t\tHit (%d)\n", page);
else
    printf("\t\tPage Fault\n");

```

Displays whether page was hit or fault.

✅ Final Output

```

printf("\nTotal Page Faults: %d", faults);
printf("\nTotal Page Hits: %d\n", hits);

```