## УНИВЕРСИТЕТ ИТМО

факультет программной инженерии и компьютерной техники

Направление подготовки 09.04.04 Системное и прикладное программное обеспечение

# Системное програмное обеспечение Семестр 1

Лабораторная работа No3

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#### 1. Goals

Implement linear code generation using certain instructions set by CFG analysis for a number of subprograms. Print mnemonic representation of the generated linear code to the output file as an assembly listing.

#### 2. Subtasks

- 1 Describe a virtual machine having instruction set and memory model according to your variant number
  - a. Learn the notation for the target architecture definitions
  - b. Describe VM according to your variant number
    - i. Describe a number of registers and memory banks
    - ii. Describe an instructions set: come up with an opcode encoding for each instruction, its

operands and a VM state changing operation description

- 1 Describe data movement and constant loading instructions
- 2 Describe arithmetic and logical instructions
- 3 Describe conditional and unconditional branching instructions
- 4 Describe IO-instructions using hidden registry store and an internal IO-port
- iii. Describe a number of mnemonics for the instructions set
- c. Prepare a script for assembled listing execution using target VM definition
  - i. Create a test listing using prepared mnemonics
  - ii. Use listing translator to generate binary module containing native code according to the target VM definition
  - iii. Execute binary module and retrieve its resulting output
  - iv. Ensure that all the prepared VM instructions work correctly
- 2 Choose and learn existing hardware VM ISA a. For a given VM
  - i. There should be an existing emulator (for example: qemu)
  - ii. There should be an existing toolchain (developer tools set): C compiler, assembler, disassembler, linker and debugger
- b. Approve your VM choice with the teacher
- c. Learn memory model and instruction set of the VM
- d. Learn how to use machine toolchain (building and launching programs in form of the assembly listing)
- e. Prepare a script for assembled listing execution using emulator
  - i. Create a test listing using VM instructions' mnemonics
  - ii. Use assembler and linker from the existing toolchain
  - iii. Execute binary module and get its resulting output

### 3. Solution description

**modules** this stage of the solution is included in 2 files, "asm.c"."asm.h", "localVariabl.c" and "localVariable.h"

In the "localVariable.h" header file, we define different structures to represent a linked list of variables. The variables have the structure varDeclaration which was defined in lab2.

More detail will be mentioned on that structure, and for full reference please refer to CПO2 report. The basic information about variables declaration was saved previously in the control flow graph as Instruction structures. Those variable declaration structures keep information about the name and the type of the variable, this type was divided into 2 categories: simple and array. The job of this module is to depth first traverse the CFG and create a linked list of all the local variables.

After this list of variables is compiled, the module asm.c starts its job of translating the variables and also in traversing the control flow graph to create the linear code.

#### 3.1 Modules Overview

The codebase is organized into several interrelated modules, each responsible for specific functionality. Below is an explanation of the primary modules and their purposes:

#### 3.1.1. Local variable list

This module defines and manages local variables used in the program. It consists of:

- varDeclaration: holds information about the variable such as its type and id. The type can be a simple type or array. Simple types are ones such as int, bool and long. And array types have 2 important fields which are the type of the elements (which could also be an array) and the number of those elements.
- Node\*: a pointer to the next variable.
- filename: name of the file in which the variable was declared.

In the module asm.c, the main function is translate(). This function invokes a number of different methods and is responsible for translating the cfg into a linear code.

#### 3.1.2 dynamic typing and variable translation

- local variables declarations are being translated first. a structure was chosen to represent a variable in the memory in order to retrieve and change meta data about its type at runtime.
- This structure for meta data follows this example

```
i_input2:
.type: dd 0x0 ; Offset for `type`
.value: dd 0x10 ; Offset for `value
```

- For each variable we keep two values which are its type and payload, local variables are stored on the stack frame while literals are stored on the stack (of the stack machine).
- The types and values are pushed to the stack and inside the arithmetic functions in assembly a check is done at runtime to make sure that the operation is done in a type appropriate way. The following snippet shows an example.

```
// Check if both operands are of type "number" (type == 2)
//both the type and value have been pushed to the stack
// Increment stack pointer to "pop" the first operand's pointer
  sp = sp + 2;
  let type1 = m stack:2[sp];
                               // First operand type (address of operand)
  sp = sp + 2;
                              // Type of first operand (e.g., 2 for number)
  let value1 = m_stack:2[sp];
 // Increment stack pointer to "pop" the second operand's pointer
  sp = sp + 2;
  let type2 = m_stack:2[sp]; // Second operand pointer
  sp = sp + 2;
  let value2 = m_stack:2[sp];
                                // Type of second operand
  // Check if both operands are of type "number" (type == 2)
  if (type1 == 2) && (type2 == 2) then {
         let r = value1 + value2;
```

#### 1.1.2 Translation implementation details

- **translate() Function**: this function distributes the traversal of the control flow graph into smaller functions that are responsible for translating:
  - Operation trees
  - While loops
  - o If conditions
  - o Literals and their types
- translateOT() Function: This is an important function. Its mission is to check the kinds of
  operations and operands and distribute them to the correct assembly operations. Since we are
  using a stack, the operands are pushed to the stack first and then they are translated into
  assembly functions such as wide\_add and wide\_store. These functions in assembly are
  responsible to type check the variables and literals in real time and also to store the values and
  types after the execution of the operations.

#### **Code Generation Process**

- Assembly Instructions: the assembly listing is written after the execution of this project
- RemoteTasks: is then used to assemble the listing, create a binary file and execute and debug this file.

#### 5. Results

As a result of the work done in this lab I was able to implement linear code translation unit with dynamic type checking. The memory banks used is a common ram and a stack. The result of this

lab allowed us to create a fully functioning program in our own language and execute it. First of all we can see the architecture created in the course of this lab work.

```
architecture spo {
/*
                      18
   case
   data word length
                     2
   spaces
                     common ram
*/
registers:
 storage ebp [16];
 storage esp [16];
 storage inp [16];
memory:
 range ram [0x0000..0xffff] { // Common RAM space for code/data
   cell = 16; // 8-bit cells
   endianess = little-endian; // Little endian
   granularity = 2; // 2-byte chunks
 range m stack [0x0000..0xffff] { // Stack space
   cell = 16;
   endianess = little-endian;
   granularity = 2;
instructions:
   encode imm16 field = immediate [16];
   encode reg field = register {
       esp = 0,
       \mathbf{ebp} = 1\};
 // Push immediate value onto the stack
 instruction push = { 0000 0000, imm16 as value } {
   m stack:2[sp] = value;
   sp = sp - 2; // Decrement stack pointer
   ip = ip + 3; // Increment instruction pointer
  // Pop value from the stack into outp
 instruction pop = { 0000 0001 } {
   sp = sp + 2;  // Increment stack pointer
ip = ip + 1;  // Increment instruction pointer
                    // Increment instruction pointer
 };
 instruction popOut = { 0000 0010 } {
   sp = sp + 2;
                     // Increment stack pointer
   outp = m_stack:2[sp]; // Retrieve value from stack to standard
   ip = ip + 1;  // Increment instruction pointer
```

```
};
// Add the top two integers on the stack
  instruction iadd = {0000 0011 } {
    sp = sp + 2;
                               // Increment stack pointer
    let x = m \text{ stack:2[sp]};
                                  // Pop first value into inp (top
value of the stack)
    sp = sp + 2;
                              // Increment stack pointer
   let y = m stack:2[sp];
                                 // Pop second value into outp (next
top value)
    let r = x + y; // Perform addition
   m_stack:2[sp] = r; // Push result back onto the stack
                             // Decrement stack pointer
    sp = sp - 2;
   ip = ip + 1;
                              // Increment instruction pointer
// Add the top two longs on the stack
instruction ladd = { 0000 0100 } {
                             // Increment stack pointer
    sp = sp + 2;
    let x = m stack:2[sp];
                              // Pop first value into inp (top
value of the stack)
    sp = sp + 2;
                              // Increment stack pointer
    let y = m_stack:2[sp];
                              // Pop second value into outp (next
top value)
                    // Perform addition
   let r = x + y;
   m_stack:2[sp] = r; // Push result back onto the stack
    sp = sp - 2;
                              // Decrement stack pointer
    ip = ip + 1;
                              // Increment instruction pointer
} ;
// store the type and the value , maybe I should push the type and
value on the top of the stack
//must test store
instruction st = \{ 0000 0101, imm16 as ptr \}
    sp=sp + 2;
    ram:2[ptr] = m stack:2[sp];
    //ram:1[ptr+1] = from>>8;
   ip = ip + 3;
};
//instruction r st = { 0000 1100, reg as target}
//{
// sp=sp + 2;
// target = m_stack:2[sp];
// ip = ip + 3;
//};
// I will use this to store the size of the array for array type
instruction a st = { 0000 0110, imm16 as ptr }
    sp=sp + 2;
```

```
ram:2[ptr+4] = m stack:2[sp];// we save 8 bytes for each variable ,
store the type in the first 4 bytes (32 bits from the stack) and the
value in the next
   ip = ip + 3;
};
// to create load I have to have different loads (iload, lload..) I
also need a general funciton to dispatch to them, but I also need a to
recieve the type
// and a way to extract it, and a way to handle the labels, in short I
need to know how I figuered out how to store it in metadata section
// were I keep all my variables and their types, or simple just how I
am storing so that I can load
 instruction jmp = { 0000 0111, imm16 as target } {
   ip = target;
 };
 instruction jz = \{0000\ 1000, imm16 \text{ as target}\}
                   // Increment stack pointer
   sp = sp + 2;
   let x = m \text{ stack:2[sp]};
                               // Pop first value into inp (top
value of the stack)
   if x == 0x0 then
     ip = target;
   else
     ip = ip + 3;
 };
 // for now, everywhere I pushed the value and then pushed the type to
the stack
instruction wide add ={0000 1001} {
   // Increment stack pointer to "pop" the first operand's pointer
   sp = sp + 2;
   of operand)
   sp = sp + 2;
                                   // Type of first operand (e.g.,
   let value1 = m stack:2[sp];
2 for number)
   // Increment stack pointer to "pop" the second operand's pointer
   sp = sp + 2;
   sp = sp + 2;
   // Check if both operands are of type "number" (type == 2)
   if (type1 == 2) && (type2 == 2) then {
       let r = value1 + value2;
                                           // Perform addition
       // Push the result back onto the stack
       m stack:2[sp] = r;
       sp = sp - 2;
                                 // Decrement stack pointer
       m stack:2[sp]= 2; // push the type
       sp=sp-2;
   } else {
       // Push error code (-1) onto the stack if types don't match
       m \text{ stack:2[sp]} = -1;
```

```
sp = sp - 2;
                                    // Decrement stack pointer
       m \text{ stack:2[sp]} = -1;
       sp=sp -2; // push type error
   }
   // Increment instruction pointer to the next instruction
   ip = ip + 1;
} ;
instruction wide store = { 0000 1010 }
   sp=sp + 2;
   let ptr = m stack:2[sp];// TO
   sp=sp + 2;
   let type=m stack:2[sp];//pop type
   sp=sp + 2;
   let value=m stack:2[sp];//pop value
   ram:2[ptr] = type;// store type
   ram:2[ptr+4] = value;// store value
   ip = ip + 1;
};
instruction wide sub ={0000 1011} {
   // Increment stack pointer to "pop" the first operand's pointer
   sp = sp + 2;
   let type1 = m stack:2[sp];
                                    // First operand type (address
of operand)
   sp = sp + 2;
   2 for number)
   // Increment stack pointer to "pop" the second operand's pointer
   sp = sp + 2;
                                 // Second operand pointer
   let type2 = m stack:2[sp];
   sp = sp + 2;
                                      // Type of second operand
   let value2 = m stack:2[sp];
   // Check if both operands are of type "number" (type == 2)
   if (type1 == 2) && (type2 == 2) then {
                                              // Perform addition
       let r = value1 - value2;
       // Push the result back onto the stack
       m \text{ stack:2[sp]} = r;
       sp = sp - 2;
                                     // Decrement stack pointer
       m stack:2[sp]= 2; // push the type
       sp= sp -2;
    } else {
       // Push error code (-1) onto the stack if types don't match
       m \text{ stack:2[sp]} = -1;
       sp = sp - 2;
                                    // Decrement stack pointer
       m stack:2[sp] = -1;
       sp=sp -2; // push type error
   // Increment instruction pointer to the next instruction
   ip = ip + 1;
```

```
instruction wide mult ={0000 1100} {
   // Increment stack pointer to "pop" the first operand's pointer
   sp = sp + 2;
  of operand)
   sp = sp + 2;
   2 for number)
   // Increment stack pointer to "pop" the second operand's pointer
   sp = sp + 2;
   let type2 = m stack:2[sp];  // Second operand pointer
   sp = sp + 2;
   // Check if both operands are of type "number" (type == 2)
   if (type1 == 2) && (type2 == 2) then {
      let r = value1 * value2;
                                    // Perform addition
      // Push the result back onto the stack
      m \text{ stack:2[sp]} = r;
                            // Decrement stack pointer
      sp = sp - 2;
      m stack:2[sp]= 2; // push the type
      sp= sp -2;
   } else {
      // Push error code (-1) onto the stack if types don't match
      m_stack:2[sp] = -1;
      sp = sp - 2;
                            // Decrement stack pointer
     m stack:2[sp] = -1;
     sp=sp -2; // push type error
   // Increment instruction pointer to the next instruction
   ip = ip + 1;
};
instruction wide div ={0000 1101} {
   // Increment stack pointer to "pop" the first operand's pointer
   sp = sp + 2;
  of operand)
   sp = sp + 2;
   2 for number)
   // Increment stack pointer to "pop" the second operand's pointer
   sp = sp + 2;
   sp = sp + 2;
   let value2 = m stack:2[sp];
                              // Type of second operand
   // Check if both operands are of type "number" (type == 2)
   if (type1 == 2) && (type2 == 2) then {
      let r = value1 / value2;
                                    // Perform addition
```

```
// Push the result back onto the stack
        m \text{ stack:2[sp]} = r;
        sp = sp - 2;
                                        // Decrement stack pointer
        m stack:2[sp]= 2; // push the type
        sp= sp -2;
    } else {
        // Push error code (-1) onto the stack if types don't match
        m \text{ stack: 2[sp]} = -1;
        sp = sp - 2;
                                       // Decrement stack pointer
        m \text{ stack:2[sp]} = -1;
        sp=sp -2; // push type error
    // Increment instruction pointer to the next instruction
    ip = ip + 1;
} ;
  instruction jgt = { 0000 1111, imm16 as target } {
                                // Increment stack pointer
    sp = sp + 2;
    let x = m \text{ stack:2[sp]};
                                     // Pop first value into inp (top
value of the stack)
    if (x >> 15 == 0x0) && (x != 0x0) then
     ip = target;
    else
      ip = ip + 3;
    instruction jge = { 0001 0000, imm16 as target } {
                                 // Increment stack pointer
    sp = sp + 2;
    let x = m \text{ stack:2[sp]};
                                    // Pop first value into inp (top
value of the stack)
    if (x >> 15 == 0x0) then
      ip = target;
   else
     ip = ip + 3;
  } ;
      instruction jlt = { 0001 0001, imm16 as target } {
                                // Increment stack pointer
    sp = sp + 2;
    let x = m \text{ stack:2[sp]};
                                    // Pop first value into inp (top
value of the stack)
    if (x >> 15 == 0x1) then
      ip = target;
    else
     ip = ip + 3;
  };
        instruction jle = { 0001 0010, imm16 as target } {
    sp = sp + 2;
                                 // Increment stack pointer
    let x = m \text{ stack:2[sp]};
                                     // Pop first value into inp (top
value of the stack)
    if (x >> 15 == 0x1) || (x == 0x0) then
      ip = target;
    else
      ip = ip + 3;
```

```
};
 // Halt execution
 instruction hlt = { 1111 1111 } {
   // Execution halts
 };
   instruction write = { 0001 0011 } {
   sp = sp + 2;
   value of the stack)
   outp = x; // Retrieve value from stack to standard output
   ip = ip + 1;  // Increment instruction pointer
 };
   instruction read = { 0001 0100 } {
   sp = sp + 2;
   let ptr= m stack:2[sp];
   ram:2[ptr+4] = inp;
   ip = ip + 1;  // Increment instruction pointer
 };
   instruction jmp emp = { 0001 0101, imm16 as ptr } {
   sp = sp + 2;
                             // Increment stack pointer
   let target = ram:2[ptr];
   ip = target;
 } ;
 instruction wide mod ={0001 0110} {
    sp = sp + 2;
                             // First operand type (address
   let type1 = m stack:2[sp];
of operand)
   sp = sp + 2;
   let value1 = m stack:2[sp];  // Type of first operand (e.g.,
2 for number)
   // Increment stack pointer to "pop" the second operand's pointer
   sp = sp + 2;
                               // Second operand pointer
   let type2 = m stack:2[sp];
   sp = sp + 2;
                                     // Type of second operand
   let value2 = m stack:2[sp];
   // Check if both operands are of type "number" (type == 2)
   if (type1 == 2) && (type2 == 2) then {
                                             // Perform addition
       let r = value1 % value2;
       // Push the result back onto the stack
       m \text{ stack:2[sp]} = r;
                                    // Decrement stack pointer
       sp = sp - 2;
       m stack:2[sp]= 2; // push the type
       sp= sp -2;
   } else {
       // Push error code (-1) onto the stack if types don't match
       m \text{ stack: 2[sp]} = -1;
       sp = sp - 2;
                                   // Decrement stack pointer
       m \text{ stack:2[sp]} = -1;
       sp=sp -2; // push type error
```

```
// Increment instruction pointer to the next instruction
    ip = ip + 1;
};
  instruction push sf =\{0001\ 0111\ ,\ imm16\ as\ value\}\ \{
    ram[esp]= value;
    esp = esp - 4;
                     // Decrement stack pointer
    ip = ip + 3;
                   // Increment instruction pointer
  };
  instruction pop sf ={0001 1000} {
    esp = esp + 4;
                          // Increment stack pointer
    ip = ip + 1;
                       // Increment instruction pointer
};
  instruction call = { 0001 1010, imm16 as ptr, 0000 0000 } {
    esp = esp + 4;
    ram:2[esp] = ip + 4; // save return point
    esp=esp + 4;
    ram:2[esp]= ebp; // save the old ebp
    ebp = esp;
    ip = ptr;
  instruction ret = { 0001 1011 } {
        //if ebp != 0 then {
            esp = ebp;
            ebp = ram:2[esp] ;
            esp = esp - 4;
            ip = ram:2[esp];
            esp = esp - 4;
            // }
        //ip = ip + 4;
    // ip=ip + 1;
};
  instruction load = { 0001 1100 } {
        sp= sp + 2; // pop the pointer to the value in ram
        let ptr= m stack:2[sp];
        let type = ram:2[ptr];
        let value = ram:2[ptr + 4];
        m stack[sp]= value;
        sp = sp - 2; //push the value
        m stack[sp] = type;
        sp = sp - 2; //push the type
        ip=ip+1;
};
instruction add reg imm = { 0001 1101, imm16 as value } {
        if value ==1 then {
        m stack:2[sp] = ebp + 4;// the size of the cell is 2 bytes (16
bits) and I need 4 for value and 4 for type
        else {
        m \text{ stack:2[sp]} = ebp + 8 * (value - 1) + 4 ;
        sp = sp - 2; // Decrement stack pointer
```

```
ip=ip+3;
};
instruction sub reg imm = { 0001 1110, imm16 as value } {
        if value ==1 then {
        m stack:2[sp] = ebp - 4 * value + 8;
        else {
        m_stack:2[sp] = ebp - (8 * (value - 1)) + 4;
        sp = sp - 2; // Decrement stack pointer
        ip=ip+3;
};
instruction init ={ 0001 1111,imm16 as target }{
    ebp=target;
    esp=ebp;
    ip=ip + 3;
};
mnemonics:
  mnemonic init for init(target) "{1}";
  mnemonic push (value) "{1}";
  mnemonic pop() ;
  mnemonic popOut();
  mnemonic iadd();
  mnemonic ladd() ;
  mnemonic jump for jmp(target) "{1}";
  mnemonic jz(target) "{1}";
  mnemonic jgt(target) "{1}";
  mnemonic jge(target) "{1}";
  mnemonic jlt(target) "{1}";
  mnemonic jle(target) "{1}";
  mnemonic store for st(ptr) "{1}";
  mnemonic astore for a st(ptr) "{1}";
  mnemonic wide add();
  mnemonic wide store();
  mnemonic wide sub();
  mnemonic wide mult();
  mnemonic wide div();
  mnemonic wide mod();
  mnemonic read();
  mnemonic write();
  mnemonic jmp emp(ptr) "{1}";
  mnemonic push sf (value) "{1}";
  mnemonic pop sf ();
  mnemonic call ( ptr) "{1}" ;
  mnemonic ret() ;
  mnemonic load() ;
  mnemonic add s for add reg imm(value) "{1}" ;
  mnemonic sub_s for sub_reg_imm(value) "{1}";
```

```
mnemonic hlt();
}
Architecture.pdsl
```

#### 5.1 Simple calculator

And we used our language to create a simple calculator. This calculator expects 2 numbers as its input and an operation + - / \*. The calculator uses loops to enter the input from the console one byte by the other. An if statement is used to direct the output to a different result depending on what operation the user want to apply to the numbers. A separate function is created to output the result from ASCII encoding, that is continuously dividing and taking the modulo of the number with 10. The output is then printed in the console. Here we can see this code example according to variant 3.

```
function calculator()
  dim i, sum, tmp1, tmp2 as int
  tmp1 = 10;
  tmp2=0;
  scan(i);
  sum = i - 48;
  scan(i);
  while i != 32
    tmp1 = i - 48;
    tmp2 = tmp1 * 10;
    sum = sum + tmp2;
    scan(i);
  wend
  dim i1, sum1, tmp11, tmp21, res, x as int
  tmp11 = 10;
  tmp21= 0;
  scan(i1);
  sum1 = i1 - 48;
  scan(i1);
  while i1 != 32
    tmp11= i1 - 48;
    tmp21 = tmp11 * 10;
    sum1 = sum1 + tmp21;
    scan(i1);
  wend
  scan(x);
  if x == 43 then
    res= sum + sum1;
  else
```

```
if x ==45 then
      res= sum - sum1;
     else
      if x ==47 then
        res = sum / sum1;
      else
        if x ==42 then
          res = sum * sum1;
        end if
       end if
    end if
  end if
  print_value(res);
end function
function print_value(res)
  dim nextLine, revertedNum, tmp as int
  nextLine = 10;
  revertedNum = 0;
  while res != 0
  revertedNum = revertedNum * 10 ;
  tmp = res % 10;
  tmp=tmp;
  revertedNum= revertedNum + tmp;
  res = res / 10;
  wend
  while revertedNum != 0
  tmp = revertedNum % 10;
  tmp=tmp + 48;
  print(tmp);
  revertedNum = revertedNum / 10;
  wend
end function
                                  Input.txt
```

As a result of executing the translation module we get the following listing.

```
[section ram, code]
store ret.value
push 2
store i_input2
push 2
```

```
store sum_input2
       push 2
       store tmp1_input2
       push 2
       store tmp2_input2
       push 2
       store i1_input2
       push 2
       store sum1_input2
       push 2
       store tmp11_input2
       push 2
       store tmp21_input2
       push 2
       store res_input2
       push 2
       store x_input2
       push 2
       store nextLine_input2
       push 2
       store revertedNum_input2
       push 2
       store tmp_input2
calculator:
       push label 0
       push tmp1_input2
       wide_store
       push label_1
       push tmp2_input2
       wide_store
       push i_input2
       read
       push label_2
       push i_input2
       wide_sub
       store label_3.type
       store label_3.value
       push label_3
       push sum_input2
       wide_store
       push i_input2
       read
               ;while
label_4:
       push label_5
       push i_input2
       wide_sub
       pop
```

```
jz label_6
       push 1
               ;false branch
       jump label_7
label 6:
       push 0
               ;true branch
label_7:
       jz label_8
               ;while body
       push label_9
       push i_input2
       wide_sub
       store label_10.type
       store label_10.value
       push label 10
       push tmp1_input2
       wide_store
       push label_11
       push tmp1_input2
       wide_mult
       store label_12.type
       store label_12.value
       push label_12
       push tmp2_input2
       wide_store
       push tmp2_input2
       push sum_input2
       wide_add
       store label_13.type
       store label_13.value
       push label 13
       push sum_input2
       wide_store
       push i_input2
       read
       jump label_4
label_8:
               ;end while
       push label_14
       push tmp11_input2
       wide_store
       push label_15
       push tmp21_input2
       wide_store
       push i1_input2
       read
       push label_16
```

```
push i1_input2
       wide_sub
       store label_17.type
       store label_17.value
       push label 17
       push sum1_input2
       wide_store
       push i1_input2
       read
               ;while
label_18:
       push label_19
       push i1_input2
       wide_sub
       pop
       jz label 20
       push 1
               ;false branch
       jump label_21
label_20:
       push 0
               ;true branch
label_21:
       jz label_22
               ;while body
       push label_23
       push i1_input2
       wide_sub
       store label_24.type
       store label_24.value
       push label_24
       push tmp11_input2
       wide_store
       push label_25
       push tmp11_input2
       wide_mult
       store label_26.type
       store label_26.value
       push label_26
       push tmp21_input2
       wide store
       push tmp21_input2
       push sum1_input2
       wide_add
       store label_27.type
       store label_27.value
       push label_27
       push sum1_input2
```

```
wide_store
       push i1_input2
       read
       jump label_18
label 22:
               ;end while
       push x_input2
       read
       push label_28
       push x_input2
       wide_sub
       pop
       jz label_29
       push 0
               ;false branch
       jump label_30
label_29:
       push 1
               ;true branch
label_30:
               ;if
       jz label_32
               ;then
       push sum1_input2
       push sum_input2
       wide_add
       store label_33.type
       store label_33.value
       push label_33
       push res_input2
       wide_store
       push res_input2
       push label_34
       store ret.value
       jump print_value
label_34:
       push halt
       store ret.value
       jump label_31
label_32:
               ;else
       push label_35
       push x_input2
       wide_sub
       pop
       jz label_36
       push 0
               ;false branch
```

```
jump label_37
label_36:
       push 1
               ;true branch
label_37:
               ;if
       jz label_39
               ;then
       push sum1_input2
       push sum_input2
       wide_sub
       store label_40.type
       store label_40.value
       push label_40
       push res_input2
       wide store
       push res_input2
       push label_41
       store ret.value
       jump print_value
label_41:
       push halt
       store ret.value
       jump label_38
label_39:
               ;else
       push label_42
       push x_input2
       wide_sub
       pop
       jz label_43
       push 0
               ;false branch
       jump label_44
label_43:
       push 1
               ;true branch
label_44:
               ;if
       jz label_46
               ;then
       push sum1_input2
       push sum_input2
       wide_div
       store label_47.type
       store label_47.value
       push label_47
       push res_input2
```

```
wide_store
       push res_input2
       push label_48
       store ret.value
       jump print_value
label_48:
       push halt
       store ret.value
       jump label_45
label_46:
               ;else
       push label_49
       push x_input2
       wide_sub
       pop
       jz label_50
       push 0
               ;false branch
       jump label_51
label_50:
       push 1
               ;true branch
label_51:
               ;if
       jz label_52
               ;then
       push sum1_input2
       push sum_input2
       wide_mult
       store label_53.type
       store label_53.value
       push label_53
       push res_input2
       wide_store
       push res_input2
       push label_54
       store ret.value
       jump print_value
label_54:
       push halt
       store ret.value
               ;endif
label_52:
       jump label_45
               ;endif
label_45:
       jump label_38
               ;endif
```

```
label_38:
       jump label_31
               ;endif
label_31:
       jmp_emp ret.value
print_value:
       push label_55
       push nextLine_input2
       wide_store
       push label_56
       push revertedNum_input2
       wide_store
               ;while
label_57:
       push label_58
       push res input2
       wide_sub
       pop
       jz label_59
       push 1
               ;false branch
       jump label_60
label_59:
       push 0
               ;true branch
label_60:
       jz label_61
               ;while body
       push label_62
       push revertedNum_input2
       wide_mult
       store label_63.type
       store label_63.value
       push label_63
       push revertedNum_input2
       wide_store
       push label_64
       push res_input2
       wide_mod
       store label_65.type
       store label_65.value
       push label_65
       push tmp_input2
       wide_store
       push tmp_input2
       push tmp_input2
       wide_store
       push tmp_input2
```

```
push revertedNum_input2
       wide_add
       store label_66.type
       store label_66.value
       push label 66
       push revertedNum_input2
       wide_store
       push label_67
       push res_input2
       wide div
       store label_68.type
       store label_68.value
       push label_68
       push res_input2
       wide_store
       jump label_57
label_61:
               ;end while
               ;while
label_69:
       push label 70
       push revertedNum_input2
       wide_sub
       pop
       jz label_71
       push 1
               ;false branch
       jump label_72
label_71:
       push 0
               ;true branch
label 72:
       jz label_73
               ;while body
       push label_74
       push revertedNum_input2
       wide_mod
       store label_75.type
       store label_75.value
       push label_75
       push tmp_input2
       wide_store
       push label_76
       push tmp_input2
       wide_add
       store label_77.type
       store label_77.value
       push label_77
```

```
push tmp_input2
        wide_store
        push tmp_input2
        write
        push label 78
        push revertedNum_input2
        wide_div
        store label_79.type
        store label_79.value
        push label 79
        push revertedNum_input2
        wide store
        jump label_69
label 73:
                ;end while
        jmp emp ret.value
        jump halt
ret:
  .type: dd 0x0; Offset for `type`
  .value: dd "; Offset for `value`
i_input2:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for 'value'
sum input2:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for 'value'
tmp1_input2:
  .type: dd 0x0; Offset for `type`
  .value: dd 0x10; Offset for `value`
tmp2 input2:
  .type: dd 0x0; Offset for `type`
  .value: dd 0x10; Offset for 'value'
i1_input2:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for 'value'
sum1_input2:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for `value`
tmp11_input2:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for 'value'
tmp21_input2:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for `value`
res_input2:
  .type: dd 0x0; Offset for `type`
  .value: dd 0x10; Offset for 'value'
x_input2:
```

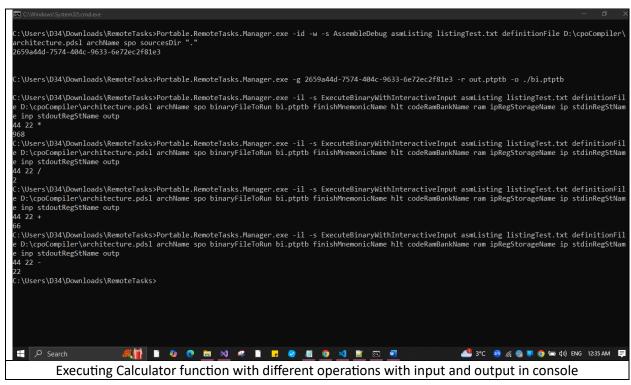
```
.type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for 'value'
nextLine input2:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for 'value'
revertedNum_input2:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for 'value'
tmp_input2:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x10; Offset for 'value'
halt:
        hlt
label 0:
  .type: dd 0x2; Offset for 'type'
  .value: dd 10; Offset for 'value'
label_1:
  .type: dd 0x2; Offset for 'type'
  .value: dd 0; Offset for `value`
label 2:
  .type: dd 0x2; Offset for 'type'
  .value: dd 48; Offset for 'value'
label 3:
  .type: dd 0x0; Offset for `type`
  .value: dd 0x0; Offset for `value`
label 5:
  .type: dd 0x2; Offset for 'type'
  .value: dd 32; Offset for 'value'
label 9:
  .type: dd 0x2; Offset for 'type'
  .value: dd 48; Offset for 'value'
label 10:
  .type: dd 0x0; Offset for `type`
  .value: dd 0x0; Offset for 'value'
label 11:
  .type: dd 0x2; Offset for `type`
  .value: dd 10; Offset for 'value'
label_12:
  .type: dd 0x0; Offset for `type`
  .value: dd 0x0; Offset for `value`
label 13:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label 14:
  .type: dd 0x2; Offset for 'type'
  .value: dd 10; Offset for 'value'
label_15:
  .type: dd 0x2; Offset for 'type'
```

```
.value: dd 0; Offset for `value`
label_16:
  .type: dd 0x2; Offset for 'type'
  .value: dd 48; Offset for 'value'
label 17:
  .type: dd 0x0; Offset for `type`
  .value: dd 0x0; Offset for 'value'
label 19:
  .type: dd 0x2; Offset for 'type'
  .value: dd 32; Offset for 'value'
label 23:
  .type: dd 0x2; Offset for 'type'
  .value: dd 48; Offset for 'value'
label 24:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label_25:
  .type: dd 0x2; Offset for 'type'
  .value: dd 10; Offset for 'value'
label 26:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label 27:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for `value`
label 28:
  .type: dd 0x2; Offset for 'type'
  .value: dd 43; Offset for 'value'
label 33:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label 35:
  .type: dd 0x2; Offset for `type`
  .value: dd 45; Offset for 'value'
label 40:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label_42:
  .type: dd 0x2; Offset for 'type'
  .value: dd 47; Offset for 'value'
label 47:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label 49:
  .type: dd 0x2; Offset for 'type'
  .value: dd 42; Offset for 'value'
label_53:
  .type: dd 0x0; Offset for 'type'
```

```
.value: dd 0x0; Offset for 'value'
label_55:
  .type: dd 0x2; Offset for 'type'
  .value: dd 10; Offset for 'value'
label 56:
  .type: dd 0x2; Offset for 'type'
  .value: dd 0; Offset for `value`
label 58:
  .type: dd 0x2; Offset for 'type'
  .value: dd 0; Offset for 'value'
label_62:
  .type: dd 0x2; Offset for 'type'
  .value: dd 10; Offset for 'value'
label 63:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label_64:
  .type: dd 0x2; Offset for 'type'
  .value: dd 10; Offset for 'value'
label_65:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label 66:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for `value`
label 67:
  .type: dd 0x2; Offset for 'type'
  .value: dd 10; Offset for 'value'
label 68:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label 70:
  .type: dd 0x2; Offset for `type`
  .value: dd 0; Offset for `value`
label 74:
  .type: dd 0x2; Offset for 'type'
  .value: dd 10; Offset for 'value'
label_75:
  .type: dd 0x0; Offset for `type`
  .value: dd 0x0; Offset for 'value'
label 76:
  .type: dd 0x2; Offset for 'type'
  .value: dd 48; Offset for 'value'
label 77:
  .type: dd 0x0; Offset for 'type'
  .value: dd 0x0; Offset for 'value'
label_78:
  .type: dd 0x2; Offset for 'type'
```

```
.value: dd 10; Offset for `value`
label_79:
.type: dd 0x0; Offset for `type`
.value: dd 0x0; Offset for `value`
listingTest.txt
```

And as a result of assembling and executing this listing according to our architecture we get the following results



#### 5.1 Fibonacci

The input we gave our module was supposed to output the Fibonacci series which it calculates recursively (until 11).

```
function fib(n as int)
 dim r, l, k,m as int
 if n < 2 then
     r = 1;
 end if
 if n \ge 2 then
     r = fib(n-2) + fib(n-1);
 end if
 r;
end function
function printf(res as int)
   dim nextLine, revertedNum, tmp as int
   nextLine = 10;
   revertedNum = 0;
   while res != 0
        revertedNum = revertedNum * 10 ;
       tmp = res % 10;
```

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```
revertedNum= revertedNum + tmp;
        res = res / 10;
    wend
    while revertedNum != 0
        tmp = revertedNum % 10;
        tmp=tmp + 48;
        print(tmp);
        revertedNum = revertedNum / 10;
    wend
    print(10);
end function
function main()
  dim i as int
  i = 0;
 while i < 11
     printf(fib(i));
     i = i + 1;
   wend
end function
Input.txt
```

```
[section ram, code]
       init CodeEnd
       jump main
fib:
       sub_s 1
       wide_store
       pop_sf
       pop_sf
       pop_sf
       pop_sf
       pop_sf
       pop_sf
       pop_sf
       pop_sf
       push 0
       push 2
       add_s 1
       wide_store
       push 0
       push 2
       add_s 2
       wide_store
       push 0
       push 2
       add_s 3
       wide_store
```

```
push 0
       push 2
       add_s 4
       wide_store
       push 2
       push 2
       sub_s 1
       load
       wide_sub
       pop
       jlt label_1
       push 0
               ;false branch
       jump label_2
label_1:
       push 1
               ;true branch
label_2:
               ;if
       jz label_3
               ;then
       push 1
       push 2
       add_s 1
       wide_store
       jump label_3
label_3:
       push 2
       push 2
       sub_s 1
       load
       wide_sub
       pop
       jge label_6
       push 0
               ;false branch
       jump label_7
label_6:
       push 1
               ;true branch
label_7:
               ;if
       jz label_8
               ;then
       push 1
       push 2
       sub_s 1
       load
```

```
wide_sub
       pop_sf
       pop_sf
       call fib
       push 2
       push 2
       sub_s 1
       load
       wide_sub
       pop_sf
       pop_sf
       call fib
       wide_add
       add_s 1
       wide_store
       jump label_8
label_8:
       add_s 1
       load
               ;endif
               ;endif
       ret
printf:
       sub_s 1
       wide_store
       pop_sf
       pop_sf
       pop_sf
       pop_sf
       pop_sf
       pop_sf
       push 0
       push 2
       add_s 1
       wide_store
       push 0
       push 2
       add_s 2
       wide_store
       push 0
       push 2
       add_s 3
       wide_store
       push 10
       push 2
       add_s 1
       wide_store
       push 0
```

```
push 2
       add_s 2
       wide_store
               ;while
label_15:
       push 0
       push 2
       sub_s 1
       load
       wide_sub
       pop
       jz label_17
       push 1
               ;false branch
       jump label_18
label_17:
       push 0
               ;true branch
label_18:
       jz label_19
               ;while body
       push 10
       push 2
       add_s 2
       load
       wide_mult
       add_s 2
       wide_store
       push 10
       push 2
       sub_s 1
       load
       wide_mod
       add_s 3
       wide_store
       add_s 3
       load
       add_s 2
       load
       wide_add
       add_s 2
       wide_store
       push 10
       push 2
       sub_s 1
       load
       wide_div
       sub_s 1
```

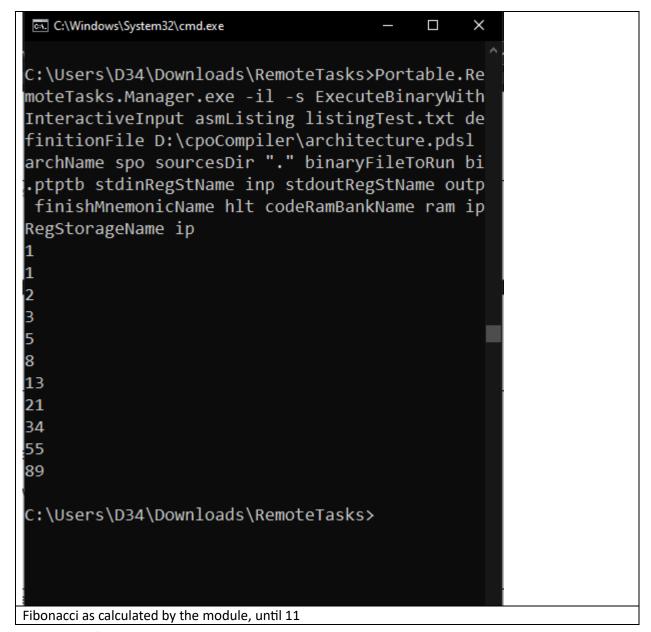
```
wide_store
       jump label_15
label_19:
              ;end while
              ;while
label_23:
       push 0
       push 2
       add_s 2
       load
       wide_sub
       pop
       jz label_25
       push 1
              ;false branch
       jump label_26
label_25:
       push 0
              ;true branch
label_26:
       jz label_27
              ;while body
       push 10
       push 2
       add_s 2
       load
       wide_mod
       add_s 3
       wide_store
       push 48
       push 2
       add_s 3
       load
       wide_add
       add_s 3
       wide_store
       add_s 3
       load
       pop
       write
       push 10
       push 2
       add_s 2
       load
       wide_div
       add_s 2
       wide_store
       jump label_23
```

```
label_27:
               ;end while
       push 10
       push 2
       pop
       write
       ret
main:
       pop_sf
       pop_sf
       push 0
       push 2
       add_s 1
       wide_store
       push 0
       push 2
       add_s 1
       wide_store
               ;while
label_33:
       push 11
       push 2
       add_s 1
       load
       wide_sub
       pop
       jlt label_35
       push 0
               ;false branch
       jump label_36
label_35:
       push 1
               ;true branch
label_36:
       jz label_37
               ;while body
       add_s 1
       load
       pop_sf
       pop_sf
       call fib
       pop_sf
       pop_sf
       call printf
       push 1
       push 2
       add_s 1
       load
```

```
wide_add
add_s 1
wide_store
jump label_33
label_37:
;end while
jump halt
halt:
hlt
CodeEnd:

Listing in assembly, generated by our module
```

And finally we got the following output



#### 6. Conclusions

I believe the results of this lab work align with the goals. The architecture description of the instruction set of our virtual machine was created successfully in alignment with our variant to include: arithmetic functions, branching, data movement and IO handeling. The module responsible for translating the control flow graphs' information to linear code was created successfully. And the result of this work was executed successfully to solve a meaningful task.

During the process of completing this lab assignment I learned what a virtual machine instruction set is, how the instructions are performed at a low level, what stages the compiler takes to reach an executable file, how to translate a code into a linear code, what dynamic and static typing is and learned a lot about real life architectures in the process of learning how to implement this task.

7. Link to the project on gitlab https://gitlab.se.ifmo.ru/Aveena/cpoCompiler