

Lab assignment solution

Problem 1:

- ① - Swap the pivot with the last element
- The pivot is 1 (given)

$$[1, 6, 2, 4, 3, 5] \rightarrow [5, 6, 2, 4, 3, 1]$$

- ② Place two index. i and j . on the first and second to last element.

$$\begin{array}{cccccc} [5, 6, 2, 4, 3, 1] \\ \uparrow \quad \quad \quad \uparrow \\ i \quad \quad \quad j \end{array}$$

- ③ move i to the right as long as the current element is less than the pivot element. move j to the right as long as the current element is more than the pivot.

$$\begin{array}{cccccc} [5, 6, 2, 4, 3, 1] \\ \uparrow \quad \uparrow \quad \uparrow \quad \uparrow \quad \uparrow \\ i_{j+} \quad j_3 \quad j_2 \quad j_1 \quad j_0 \end{array}$$

← moving j to the left 5 times.

$$\begin{array}{cccccc} [5, 6, 2, 4, 3, 1] \\ i_j \quad \quad \quad \# \end{array}$$

when j is about to pass i , stop the process and swap i with the pivot element.

$$[1, 6, 2, 4, 3, 5]$$

Problem 2:

$$3n/4 = 3 \cdot 9/4 = 6.75$$

For $x = 5$, $L = \{1, 1, 2, 3, 3, 4\}$, $R = \{6, 7\}$ and in each case $|L| < 6.75$, $|R| < 6.75$, therefore 5 is a good pivot

For $x = 6$ and 7 , $|R| > 6.75$, and therefore is not a good pivot

For $x = 4$, $R = \{5, 6, 7\}$, $L = \{1, 1, 2, 3, 3\}$, $|R| < 6.75$ and $|L| < 6.75$ therefore is a good pivot

For $x = 3$, $R = \{4, 5, 6, 7\}$, $L = \{1, 1, 2\}$ again is a good pivot

For $x = 2$, $R = \{3, 3, 4, 5, 6, 7\}$ and $L = \{1\}$ and is a good pivot point

For $x = 1$, $R = \{2, 2, 3, 3, 4, 5, 6, 7\}$ and $|R| > 6.75$ and is not a good pivot

- a. So to recap: x in $Bad = [1, 6, 7]$ are not a good pivot but x in $Good = \{2, 2, 3, 4, 5\}$ are good pivot, there are 3 bad pivot, and 5 good pivot.
- b. 5 out of 9 are good pivots, 5 is more than half of 9 so, yes, the number of pivot is more than half.

Problem 3:

1. **Valid recursion:** Base case is provided, i.e., when $n \leq 1$. And each recursion leads to base case because on each stage we are subtracting 1.
2. **Correct base case:** when $n \leq 0$, it is correctly returned because $0! = 1! = 1$
3. **Recursion steps correct:** Assuming the recursion gives the correct result for $1 \leq n \leq k$, for some k , it is then true that it gives the correct result for $n = k + 1$ because,
$$\text{recursiveFactorial}(k + 1) = (k + 1) * \text{recursiveFactorial}(k)$$

Problem 4:

```
Algorithm SubsetSum(S, k)
  Input: a set S of n elements, and a number k which is the sum
  Output: a subset of S whose elements sum up to k.
  elem <- S.getLast()
  if(k = 0)
    return emptySet()
  if(S.isEmpty() && k != 0)
    return null
  S.removeLast()
  if(elem > k)
    return SubsetSum(S, k)
  include <- SubsetSum(S, k)
  exclude <- SubsetSum(S, k - elem)
  if(include != null)
    include.addLast(elem)
  return include
return exclude
```