

Lab 5

1. Show all steps of In-Place QuickSort in sorting the array [1, 6, 2, 4, 3, 5] when doing first partition. Use leftmost values as pivots.
2. In our average case analysis of QuickSort, we defined a *good self-call* to be one in which the pivot x is chosen so that number of elements $< x$ is less than $3n/4$, and also the number of elements $> x$ is less than $3n/4$. We call an x with these properties a *good pivot*. When n is a power of 2, it is not hard to see that at least half of the elements in an n -element array could be used as a good pivot (exactly half if there are no duplicates). For this exercise, you will verify this property for the array $A = [5, 1, 4, 3, 6, 2, 7, 1, 3]$ (here, $n = 9$). Note: For this analysis, use the version of QuickSort in which partitioning produces 3 subsequences L, E, R of the input sequence S .
 - a. Which x in A are good pivots? In other words, which values x in A satisfy:
 - i. the number of elements $< x$ is less than $3n/4$, and also
 - ii. the number of elements $> x$ is less than $3n/4$
 - b. Is it true that at least half the elements of A are good pivots?
3. Prove the following recursive factorial algorithm is correct.

Algorithm recursiveFactorial(n)

Input: A non-negative integer n

Output: $n!$

if ($n = 0 \parallel n = 1$) **then**

return 1

return $n * \text{recursiveFactorial}(n-1)$

4. Review of SubsetSum Problem: Given a set $S = \{s_0, s_1, s_2, \dots, s_{n-1}\}$ of positive integers and a non-negative integer k , find a subset T of S so that the sum of the integers in T equals k or indicate no such subset can be found.

We have already seen a brute force solution to this problem in an earlier lab. In this exercise, you are going to come up with a recursive solution for SubsetSum. Write the pseudo code for your algorithm.

Hint:

We are seeking a $T \subseteq S = \{s_0, s_1, \dots, s_{n-2}, s_{n-1}\}$ whose sum is k . Such a T can be found if and only if one of the following is true:

- (1) A subset T_1 of $\{s_0, s_1, \dots, s_{n-2}\}$ can be found whose sum is k , OR
- (2) A subset T_2 of $\{s_0, s_1, \dots, s_{n-2}\}$ can be found whose sum is $k - s_{n-1}$

If (1) holds, then the desired set T is T_1 . If (2) holds, the desired set T is $T_2 \cup \{s_{n-1}\}$.