Set ADTs implementation via:

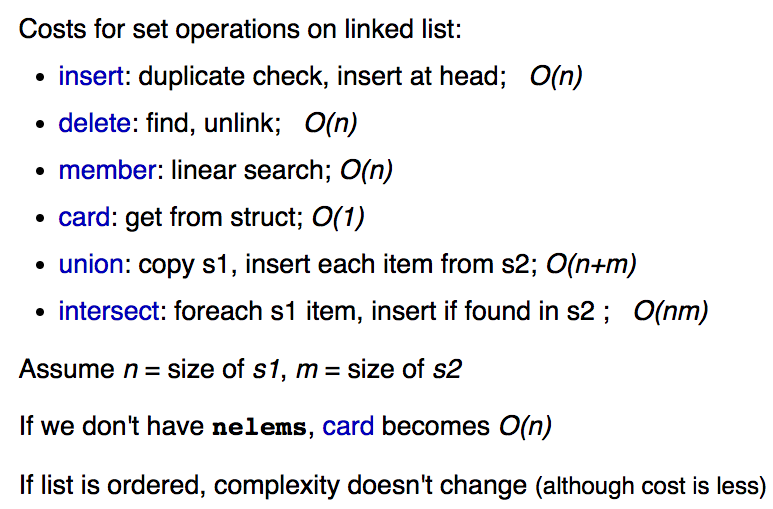
* // fixed size array, max # elements  
  **struct SetRep { int nelems; int elems[MAXELEMS]; }**
* // flexible array via. realloc()  
  **struct SetRep { int nelems; int \*elems; }**

**Sets as Linked Lists**

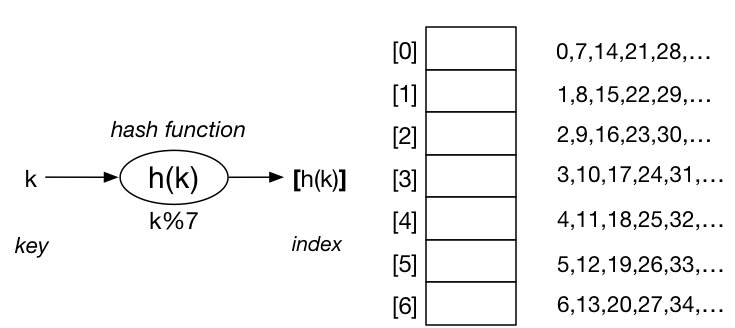
**Cardinality operations** are common to store # of elements, so you don’t need to do an O(n) traversal to count elts.

* However, this means higher space complexity / higher storage costs.
* This also means there will be extra work in incrementing / decrementing during other operations to update the counter

Efficient to have a single ptr to the **start AND end of the list**



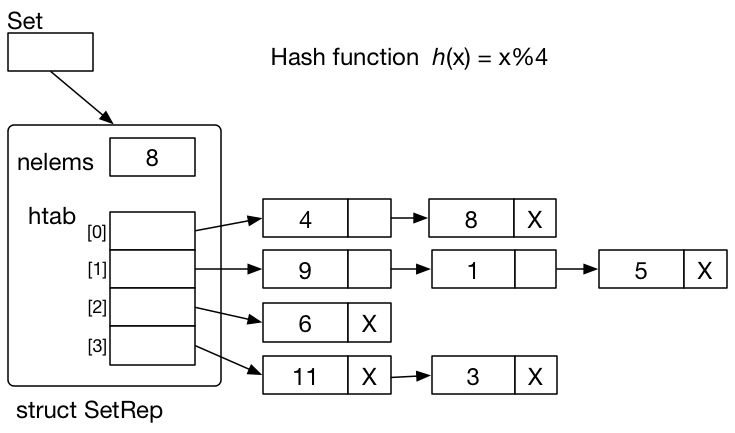
**Sets as Hash Tables**



**Hash Tables** = a way of indexing based off key values, which can be arbitrary values (numbers, strings etc.).

* Taking a key value and running it through a function that gives us an index into the Hash Table, telling us where the object belongs in the Hash Table.

If there are a large number of objects and a relatively small hash table, several objects with different key values could end up coming up with the same Hash Value, therefore going into the same spot in the Hash Table (which is bad).



E.g. insert val 4 into hash table.

* Compute h(4) = 4%4 = 0
* 4 belongs into htab[0] 🡪 added to list

E.g. insert val 12 into hash table.

* H(12) = 12%4 = 0
* 12 belongs into htab[0] 🡪 added to list

**STEP 1: Take key value  
STEP 2: Compute index into hash table  
STEP 3: Insert object into hash table elt**

**Real world example**: We have 1 million elements, but our Hash Function is extremely complex, so Hash Table elements may only contain 1 object each. This would have similar property to accessing an array.