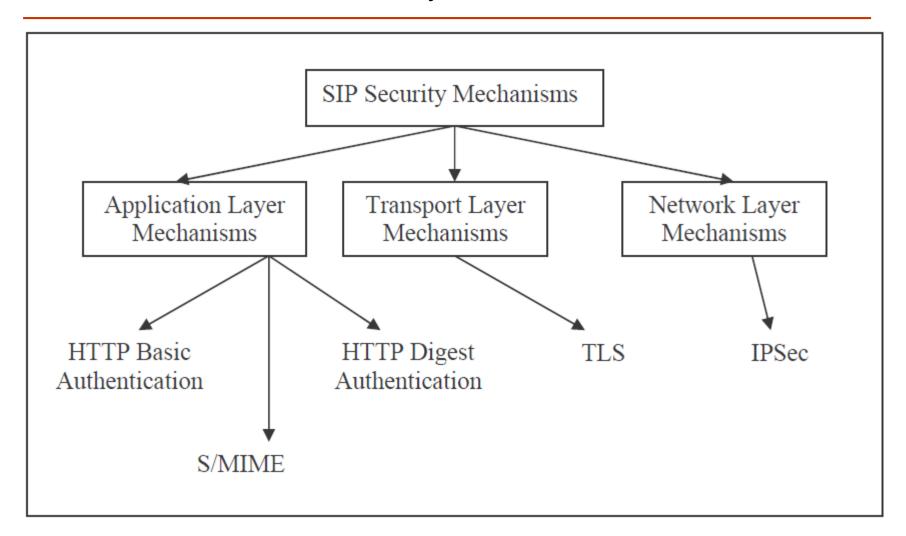
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- Authentication Goals
- Attack scenarios
- HTTP digest authentication
- Authentication scheme for a trusted domain
- Authentication Challenges

### SIP Security Architecture



#### Authentication Goals and Methods

- Identifying the source (sender)
  - Object ID in reality should be the same as the object claims to be
- Identifying the Receiver
  - Before sending sensitive data (such as personal information)
- Authentication methods
  - HTTP digest mechanism, transport layer mechanism
  - Secure Multipart Internet Mail Extensions (S/MIME)

#### **Attacks Scenarios**

- Hijacking registration of the user client
  - The client sends an REGISTER message to registrar server (includes IP address and contact information of the user)
  - IP Add. & Contact Info. can be used to hijack registration and forward incoming call to SIP phone of the hijacker
  - Hijacker can cheat registrar by first blocking original user's (DoS attacks to the user)

### Attacks Scenarios (cont.)

- Hijacking registration of the user client (cont.)
  - The attacker sends it's own registration messages to registrar
    - including the SIP contact information of the original user, but with the changed IP address
  - From the Registrar perspective
    - Only the IP address of the registered user is changed
    - SIP contact is associated to new IP address which is the hijacker's IP address
  - Then Registrar will forward incoming call request to the hijacker instead of the original user

# INVITE Spoofing

An example of spoofing the From, Via and Subject header fields in a INVITE request

INVITE sip:bob@biloxi.com SIP/2.0

Via: SIP/2.0/UDP somewhere.com

To: Bob <sip:bob@biloxi.com>

From: Alice <alice@atlanta.com>

Call-ID: a72bcf72f22d

Cseq: 100 INVITE

Contact: <sip:John@somewhere.com>

**Subject: Your Friend Alice** 

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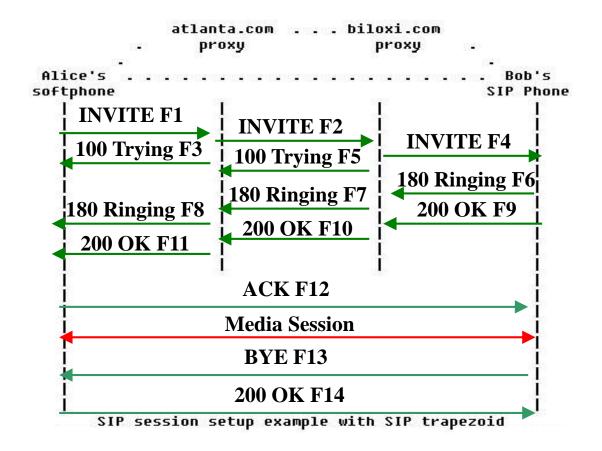
## Attacks Scenarios (Cont.)

- SIP calls forwarding
  - SIP calls forwarding to possibly rogue server
    - Exploiting the lack of identifying response messages
- Attack process: SIP request is send to proxy server of the receiver
  - The attacker hijacks the request message and uses the header fields from that request
  - Then, attacker assigns the response with status code 302 ("moved temporarily")
  - Message with headers stolen from request message (for example INVITE, incoming call is forwarded to attacker)

## Attacks Scenarios (Cont.)

- The two attack scenarios have implications on security and privacy
  - Attacker could get personal and confidential data of the users
  - Rogue party could mislead end users
  - Forward calls to malicious entity
  - Existing call parties could be conferencing with rogue party (interception)
- Two main security mechanisms
  - Authentication
    - To prevent attackers from modifying and/or replaying SIP requests and responses
  - Encryption
    - To ensure confidentiality

#### Basic Call Flow



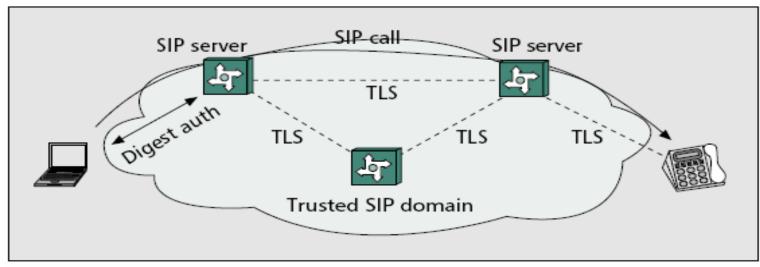
### Security Mechanisms

- End-to-end versus hop-by-hop
  - End-to-end security use SIP mechanisms to ensure security
  - Hop-by-hop relies on the security provided by the network
- Examples of hop-by-hop mechanism are Transport Layer Security (TLS) and Internet Protocol Security (IPsec)
- SIP introduces HTTP digest authentication and usage of S/MIME extensions

### SIP Methods and Authentication

| Method   | Purpose                        | Supports Authentication |
|----------|--------------------------------|-------------------------|
| INVITE   | Initiate a session             | Yes                     |
| ACK      | Acknowledge session initiation | No                      |
| OPTIONS  | Query server capabilities      | No                      |
| BYE      | Terminate a session            | Yes                     |
| CANCEL   | Cancel a pending request       | Yes                     |
| REGISTER | Register a user's location     | Yes                     |

# Security Mechanisms (cont.)

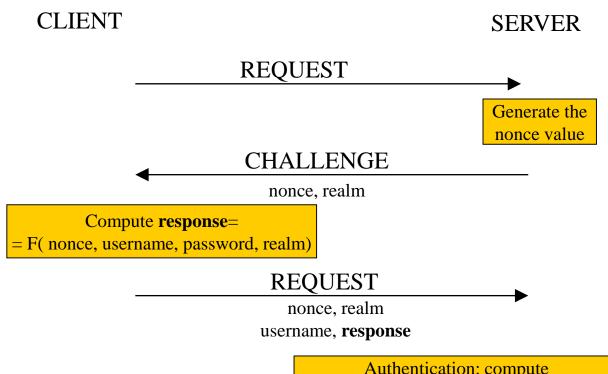


An example of a trusted network scenario

# HTTP digest authentication

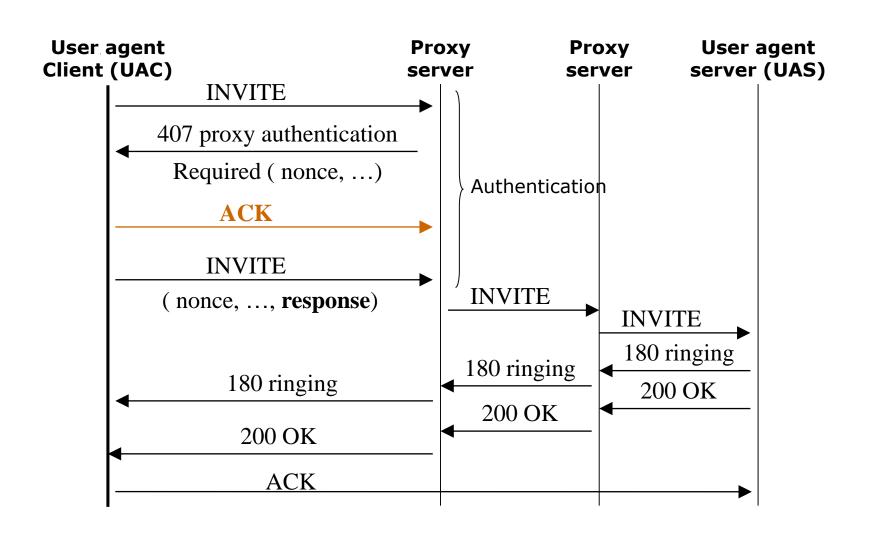
- HTTP digest is a challenge-response protocol
  - Nonce value is used in challenging the target
  - The response: a checksum of the username, password, nonce value
- The usage: users to users or users to proxies
  - Not proxies to proxies
- The security between proxies relies on other mechanisms
  - TLS or IPsec

#### The Authentication Procedure



Authentication: compute
F( nonce, username, password, realm)
And compare with **response** 

### The Authentication Procedure (cont.)



## HTTP digest authentication

- When a server receives a request message from the client (such as INVITE)
  - it may challenge the sender of the message
    - The server sends an response message containing a nonce value and a realm
    - The response is actually an error message requesting authentication
    - The realm in the message is the digest algorithm used in this challenge

# HTTP digest authentication (cont.)

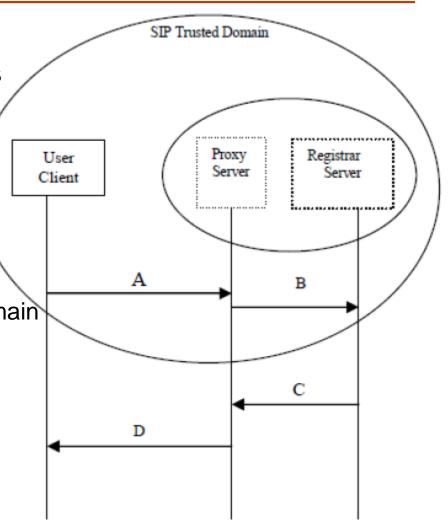
- Nonce value is used in challenging the target
  - The initiator, client, of the request receives the response
    - Computes the response value with nonce value received in challenge and with a username and a secret password
    - The secret password is known by both the client and the server
  - The client sends back the original request message with the computed response value, username, nonce value and realm

#### Authentication scheme- trusted SIP domain

 The outbound proxy for the domain is at the edge of the domain

 Authentication of the user client has to be handled by that proxy

 Proxy and Registrar servers needs to have public key certificates, which are authorized by authorities of that domain



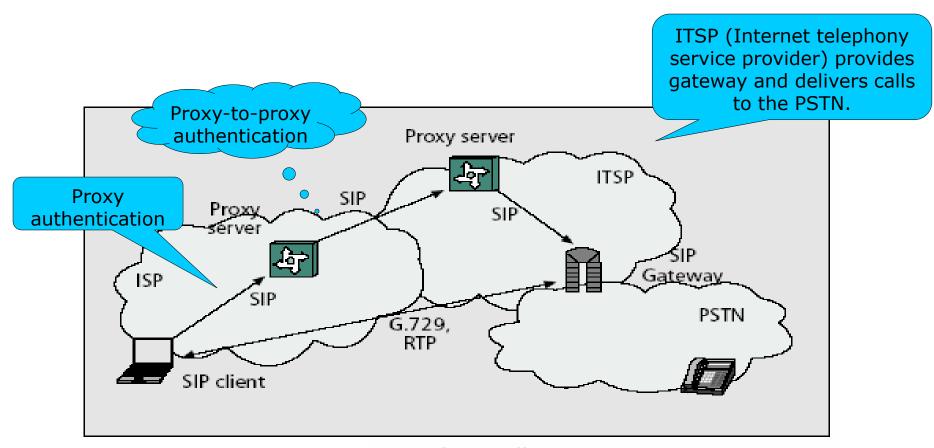
- The user client registers itself to the registrar server
  - Sends inside the registration message the identity associated to the user client
- The registrar server creates a large Secret number (N) when it receives the registration message
  - Replies to client with a value, which is computed with the N value and the identity of the user client
  - This value is the password for the user client
- Then the registrar server computes a number r, which is generated with identities of the user client and registrar server
- The user client initiating the call is required to authenticate itself with the outbound proxy of the trusted SIP domain

- The user client sends request with the parameters as follows
  - User client creates a secret random number R With the password, received in registration
  - The client generates a number n by computing it with the r and the password
  - User client computes a timestamp and a temporary key K
    - K is created with the timestamp and the password
    - The secret random number R is then encrypted with key K
    - Then user client sends the parameter A
    - A consists of n, R encrypted using K, identity of registrar and timestamp

- The Proxy server receives the request A with the parameters
  - Compares the timestamp in A and current time in order to verify that the message is in acceptable timeframe
  - Proxy server verifies the user client with registrar server using the proxy certificate and the parameters received in message A
  - Message B is sent by the proxy to registrar
  - Registrar replies by sending a message C to proxy server, if the user client is identified and authorized in this domain
  - The proxy server then sends to the user client a temporary certificate, which
    is valid until the timestamp associated with the certificate expires
  - The proxy encrypts the certificate with session key, which is then used for all signalling traffic

- User client receives the message D containing the encrypted certificate
  - Got a temporary certificate and the session key to continue the call establishment
- Call is established between calling user client and called user's server
  - The identity information is shared
  - The server of the called user verifies the received certificate and if it is valid
  - it saves the session key and allows call to be established to called user client

# An Example Scenario of a SIP-Based IP Telephony Service



PC-to-phone call.

# Challenges in SIP authentication

- Two major weaknesses in HTTP digest authentication
  - Lack of securing all headers and parameters in SIP
  - The requirement of pre-existing user configuration on servers, which does not scale well
- The authentication scheme for a trusted SIP domain
  - Usage of several hash computations and server certificates
    - Additional load (overhead) and decreases the overall performance of the server
    - If the load increases, the server comes more vulnerable to denial of service attacks

### Summary

- Authentication methods
  - HTTP digest
  - Usage of S/MIME in SIP (to be presented)
  - Scheme for a trusted domain
- Security trade-off: Overhead and Performance