

## Lecture 15

\* Proposition 1 : Under assumption 1 the following hold for the average cost per stage problem:

a. The optimal average cost  $\lambda^*$  is the same for all initial states and together with some vector  $h^* = [h^*(1), \dots, h^*(n)]$  satisfies Bellman's equation

$$\lambda^* + h^*(i) = \min_{u \in U(i)} \left\{ g(i, u) + \sum_{j=1}^n p(i, u, j) h^*(j) \right\}, \quad i \in [n] \quad - ①$$

Furthermore, if  $\mu(i)$  attains the minimum in the above equation for all  $i$ , the stationary policy  $\mu$  is optimal. In addition, out of all vectors  $h^*$  satisfying ①, there is a unique vector with  $h^*(n) = 0$ .

b. If a scalar  $\lambda$  and a vector  $h = [h(1) \dots h(n)]^T$  satisfy ①, then  $\lambda$  is the average optimal cost per stage for each initial state.

c. Given a stationary policy  $\mu$  with corresponding average cost per stage  $\lambda_\mu$ , there is a unique vector  $h_\mu = [h_\mu(1), \dots, h_\mu(n)]^T$  such that  $h_\mu(n) = 0$  and

$$\lambda_\mu + h_\mu(i) = g(i, \mu(i)) + \sum_{j=1}^n p(i, \mu(i), j) h_\mu(j), \quad i \in [n] \quad - ②$$

Proof:

a. Let us denote

$$\bar{\lambda} := \min_{\mu} \frac{C_{nn}(\mu)}{N_{nn}(\mu)} \quad - (i)$$

Note that  $C_{nn}(\mu)$  and  $N_{nn}(\mu)$  are finite  $\forall \mu$  by assumption 1.

Also we have that

$$C_{nn}(\mu) - N_{nn}(\mu) \bar{\lambda} \geq 0 \quad - (ii)$$

with equality if  $\mu$  attains minimum in (i). By a previous proposition, the costs  $h^*(1), \dots, h^*(n)$  solve uniquely corresponding

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$$h^*(i) = \min_{u \in U(i)} \left\{ g(i, u) - \bar{\lambda} + \sum_{j=1}^{n-1} p(i, u, j) h^*(j) \right\}, \quad -(iii)$$

Since the transition probability  $p(i, u, n) = 0 \forall i, u$  in the assoc. SSPP. An optimal stationary policy should minimize

$$C_{nn}(\mu) - N_{nn}(\mu) \bar{\lambda}$$

to zero, so we have that  $h^*(n) = 0$ .

Thus,

$$h^*(i) + \bar{\lambda} = \min_{u \in U(i)} \left\{ g(i, u) + \sum_{j=1}^n p(i, u, j) \cdot h^*(j) \right\}, \quad i \in [n] \quad -(iv)$$

We will show that this relation implies  $\bar{\lambda} = \lambda^*$ .

Let  $\pi = \{\mu_0, \mu_1, \dots\}$  be any admissible policy, let  $N$  be a positive integer, and for all  $k = 0, \dots, N-1$ , define  $J_k(i)$  using the following recursion

$$J_0(i) = h^*(i), \quad i \in [n]$$

$$J_{k+1}(i) = g(i, \mu_{N-k-1}(i)) + \sum_{j=1}^n p(i, \mu_{N-k-1}(i), j) \cdot J_k(j), \quad i \in [n]$$

Note that  $J_N(i)$  is the  $N$ -stage cost of  $\pi$  when the starting state is  $i$  and the terminal cost function is  $h^*$ . By (iv)

$$\bar{\lambda} + J_0(i) \leq J_1(i), \quad i \in [n]$$

Using this relation, we have

$$\begin{aligned} & g(i, \mu_{N-2}(i)) + \bar{\lambda} + \sum_{j=1}^n p(i, \mu_{N-2}(i), j) J_0(i) \\ & \leq g(i, \mu_{N-2}(i)) + \sum_{j=1}^n p(i, \mu_{N-2}(i), j) J_1(j) \end{aligned}$$

and as  $J_0(i) = h^*(i)$

$$2\bar{\lambda} + h^*(i) \leq J_2(i)$$

Repeating this argument we obtain

$$k\bar{\lambda} + h^*(i) \leq J_k(i) \quad k = 0, \dots, N, \quad i \in [n]$$

Repeating this argument we obtain

$$k\bar{\lambda} + h^*(i) \leq J_k(i), \quad k=0, \dots, N, \quad i \in [n]$$

and in particular for  $k=N$ ,

$$\bar{\lambda} + \frac{h^*(i)}{N} \leq \frac{J_N(i)}{N}$$

Taking limit as  $N \rightarrow \infty$ ,

$$\bar{\lambda} \leq J_\pi(i)$$

for all admissible policy  $\pi$ , with equality if  $\pi$  is a stationary policy  $\mu$  such that  $\mu(i)$  attains minimum in (iv)  $\forall i, k$ .

It follows that

$$\bar{\lambda} = \min_{\pi} J_\pi(i) = \lambda^*, \quad i \in [n]$$

Equation (iv) with  $h^*(n)=0$  is equivalent to equation (iii).

Since the solution of (iii) is unique, the same is true for (iv) with  $h^*(n)=0$ .

\* b and c are home work!

"  $\square$  "

### \* Value Iteration :

The natural one: Initialize iterate  $J_0$  arbitrarily. For  $k=0, 1, 2, \dots$

Compute  $J_{k+1}$  iteratively as follows:

$$J_{k+1}(i) = \min_{u \in U(i)} \left\{ g(i, u) + \sum_{j=1}^n P(i, u, j) J_k(j) \right\}, \quad i \in [n]$$

If it is natural to expect that  $\frac{J_k(i)}{k}$  should converge to  $\lambda^*$  as  $k \rightarrow \infty$ , i.e.,

$$\lim_{k \rightarrow \infty} \frac{J_k(i)}{k} = \lambda^* \quad \forall i.$$

To show this, let us define the recursion

$$J_{k+1}^*(i) = \min_{u \in U(i)} \left\{ g(i, u) + \sum_{j=1}^n P(i, u, j) J_k^*(j) \right\}, \quad i \in [n]$$

with  $J_0^*(i) = h^*(i)$ ,  $\forall i \in [n]$ .

By induction it is easy to see that

$$J_k^*(i) = k\lambda^* + h^*(i), \quad \forall i \in [n]$$

It can also be shown that

$$|J_k(i) - J_k^*(i)| \leq \max_{j \in [n]} |J_0(j) - h^*(i)|, \quad i \in [n].$$

$$\text{So, } |J_k(i) - k\lambda^*| \leq \max_{j \in [n]} |J_0(j) + h^*(j)| + \max_{j \in [n]} |h^*(j)|, \quad i \in [n]$$

So,  $\frac{J_k(i)}{k}$  converges to  $\lambda^*$

\* But see that  $\lambda^*$  is finite hence as  $k \rightarrow \infty, J_k(i) \rightarrow \infty \forall i \in [n]$ .  
So there is a computational problem!

\* Also we have no information about the differential cost vector  $h^*$ .

\* Relative Value Iteration:

- To overcome the problems in previous section we subtract a constant (cleverly) to get  $h^*$  as well as  $\lambda^*$ .

Consider the algorithm:

$$h_k(i) = J_k(i) - J_k(s), \quad i = 1, \dots, n, \quad s \in [n] \text{ be fixed.}$$

Then

$$\begin{aligned} h_{k+1}(i) &= J_{k+1}(i) - J_{k+1}(s) \\ &= \min_{u \in U(i)} \left\{ g(i, u) + \sum_{j=1}^n p(i, u, j) J_k(j) \right\} \\ &\quad - \min_{u \in U(s)} \left\{ g(s, u) + \sum_{j=1}^n p(s, u, j) J_k(j) \right\} \\ &= \min_{u \in U(i)} \left\{ g(i, u) + \sum_{j=1}^n p(i, u, j) (J_k(j) - J_k(s)) \right\} \\ &\quad - \min_{u \in U(s)} \left\{ g(s, u) + \sum_{j=1}^n p(s, u, j) (J_k(j) - J_k(s)) \right\} \\ &= \min_{u \in U(i)} \left\{ g(i, u) + \sum_{j=1}^n p(i, u, j) h_k(j) \right\} \\ &\quad - \min_{u \in U(s)} \left\{ g(s, u) + \sum_{j=1}^n p(s, u, j) h_k(j) \right\} \end{aligned}$$

If can be seen that if relative value iteration converges to some vector  $h$ , then we have

to some vector  $h$ , then we have  $\downarrow$

$$\lambda + h(i) = \min_{u \in U(i)} \left\{ g(i, u) + \sum_{j=1}^n p(i, u, j) \cdot h(j) \right\}$$

with  $h(s) = 0$  and

$$\lambda = \min_{u \in U(s)} \left\{ g(s, u) + \sum_{j=1}^n p(s, u, j) \cdot h(j) \right\}$$