

Written Problems

Question 1

According to Dijkstra's algorithm, when the vertex is marked closed and assumes that the shortest path to it is found. If we have a open node such that the cost to the next node is minimal, the minimality will not change. This is because adding positive numbers to a vertex will not change this value. If we have negative numbers, the basic concept of Dijkstra's algorithm is not valid, and thus won't work.

Question 2

```
define delete_edge(graph, node1, node2) // using a dict representation of graphs
begin
    connections1 = graph.get[node1] // list of connections to that node
    connections2 = graph.get[node2] // list of connections to that node
    if node1 is in connections2 && node2 is in connections1 then
        delete node2 from connection1
        delete node1 from connection2
    else
        connection doesnt exist
end

define insert_edge(graph, node1, node2)
begin
    if graph.has_key(node1) == False && graph.has_key(node2) == False
        return // exit function
    else
        connections1 = graph.get[node1] // list of connections to that node
        connections2 = graph.get[node2] // list of connections to that node
        connections1.append(node2)
        connections2.append(node1)
end
```

Question 3

To make adjustments to the adjacency list, every connection list should have a corresponding node to which it is connected just once. In other words, if I have node a, and the connection list is [b,d,e] and when I want to represent node b, I do not include node a in the connection list as it would be redundant. Node with the higher order will be given priority can contain the connections to the subsequent nodes. Thus allowing us to delete any node at constant time, and also any edge at constant time.

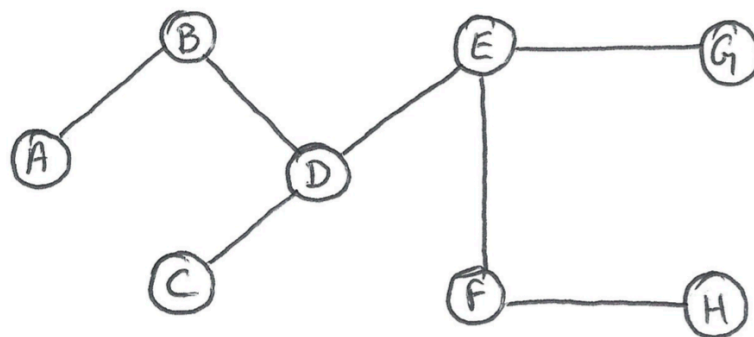
```

define delete_edge_const(graph, node1) // using a dict representation of graphs
begin
    connections = graph.get[node1] // list of connections to that node
    delete element at index 0 of connections
end

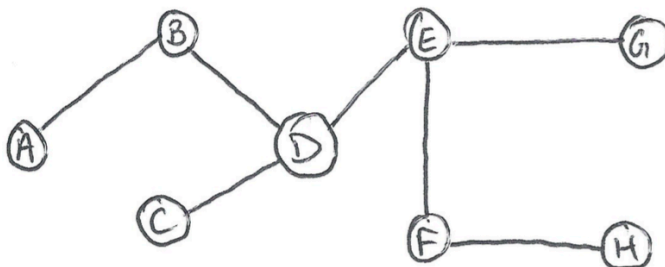
```

Question 4

a.) Using Prim's Algorithm.



b.) Using Kruskal's Algorithm.



Queue	
(C,d)	= 1
(e,f)	= 1
(f,h)	= 1
(a,b)	= 2
(b,d)	= 2
(d,e)	= 2
(e,g)	= 2
* (f,g)	= 2
(a,c)	= 3
(g,h)	= 3
(d,f)	= 4.

Starting at A (both DFS and BFS)

```
Aviks-MacBook-Pro:week_7 avikbag$ python graph_traversal8.py
Adjacency list for the graph used
A -> ['B', 'C']
B -> ['A', 'D']
C -> ['A', 'D']
D -> ['B', 'C', 'E', 'F']
E -> ['D', 'F', 'G']
F -> ['D', 'E', 'G', 'H']
G -> ['E', 'F', 'H']
H -> ['F', 'G']
DFS on the given graph (root = Node A)
-> A -> B -> D -> C -> E -> F -> G -> H
BFS on the given graph (root = Node A)
-> A -> B -> C -> D -> E -> F -> G -> H
```

Starting at D (both DFS and BFS)

```
Aviks-MacBook-Pro:week_7 avikbag$ python graph_traversal8.py
Adjacency list for the graph used
A -> ['B', 'C']
B -> ['A', 'D']
C -> ['A', 'D']
D -> ['B', 'C', 'E', 'F']
E -> ['D', 'F', 'G']
F -> ['D', 'E', 'G', 'H']
G -> ['E', 'F', 'H']
H -> ['F', 'G']
DFS on the given graph (root = Node D)
-> D -> B -> A -> C -> E -> F -> G -> H
BFS on the given graph (root = Node D)
-> D -> B -> C -> E -> F -> A -> G -> H
```