## Introduction to Parallel Computer Architecture Gaussian Elimination

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This assignment is due February 8th (11:59PM), 2017. You may work on it in a team of up to two people.

Consider the problem of solving a system of linear equations of the form

In matrix notation, the above system is written as Ax = b where A is a dense  $n \times n$  matrix of coefficients such that  $A[i,j] = a_{i,j}$ , b is an  $n \times 1$  vector  $[b_0,b_1,\ldots,b_{n-1}]^T$ , and x is the desired solution vector  $[x_0,x_1,\ldots,x_{n-1}]^T$ . From here on, we will denote the matrix elements  $a_{i,j}$  and  $x_i$  by A[i,j] and x[i], respectively. A system of equations Ax = b is usually solved in two stages. First, through a set of algebraic manipulations, the original system of equations is reduced to an upper triangular system of the form

We write the above system as Ux = y, where U is a unit upper-triangular matrix, that is, one where the subdiagonal entries are zero and all principal diagonal entries are equal to one. More formally, U[i,j] = 0 if i > j, otherwise  $U[i,j] = u_{i,j}$ , and furthermore, U[i,i] = 1 for  $0 \le i < n$ . In the second stage of solving a system of linear equations, the upper-triangular system is solved for the variables in reverse order, from x[n-1] to x[0] using a procedure called back-substitution.

A serial implementation of a simple Gaussian elimination algorithm is shown in the next page. The algorithm converts the system of linear equations Ax = b into a unit upper-triangular system Ux = y. We assume that the matrix U shares storage with A and overwrites the upper-triangular

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1: procedure GAUSS_ELIMINATE(A, b, y)
2: int i, j, k;
3: for k := 0 to n - 1 do
      for i := k + 1 to n - 1 do
         A[k, j] := A[k, j]/A[k, k];
                                        /* Division step. */
5:
      end for
6:
      y[k] := b[k]/A[k, k];
7:
8:
      A[k, k] := 1;
      for i := k + 1 to n - 1 do
9:
         for j := k + 1 to n - 1 do
10:
            A[i, j] := A[i, j] - A[i, k] \times A[k, j];
                                                     /* Elimination step. */
11:
         end for
12:
         b[i] := b[i] - A[i, k] \times y[k];
13:
         A[i, k] := 0;
14:
15:
      end for
16: end for
```

portion of A. So, the element A[k,j] computed in line 5 of the code is actually U[k,j]. Similarly, the element A[k,k] that is equated to 1 in line 8 is U[k,k]. Also, our program assumes that  $A[k,k] \neq 0$  when it is used as a divisor in lines 5 and 7. So, our implementation is numerically unstable, though it should not be a concern for this assignment. For k ranging from 0 to n-1, the Gaussian elimination procedure systematically eliminates the variable x[k] from equations k+1 to n-1 so that the matrix of coefficients becomes upper-triangular. In the  $k^{\text{th}}$  iteration of the outer loop (line 3), an appropriate multiple of the  $k^{\text{th}}$  equation is subtracted from each of the equations k+1 to n-1

You are asked to develop a multi-threaded implementation of GAUSS\_ELIMINATE using pthreads. The program given to you accepts no arguments. The CPU computes the reference solution which is compared with the result provided by the multi-threaded implementation. If the solutions match, the application will print out "Test PASSED" to the screen before exiting. Edit the gauss\_eliminate\_using\_pthreads() function in gauss\_eliminate.c to complete the functionality of Gaussian elimination using pthreads. The source files for this question are available in a zip file called qauss\_code.zip. E-mail the files needed to run your code as a single zip file called gauss.zip. Provide a brief report describing how you designed your program (use code or pseudocode if that helps the discussion) and the amount of speedup obtained over the reference version for matrices of the following sizes:  $1024 \times 1024$ ,  $2048 \times 2048$ , and  $4096 \times 4096$  elements, when using 4, 8, 16, and 32 threads.