

Binary Tree - Complete Notes

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Introduction

A **Binary Tree** is a hierarchical data structure where each node has at most two children, referred to as the left child and right child.

Basic Structure

```
class TreeNode {
    int data;
    TreeNode left;
    TreeNode right;

    TreeNode(int data) {
        this.data = data;
        this.left = null;
        this.right = null;
    }
}
```

Types of Binary Trees

1. Perfect Binary Tree

A binary tree is **perfect** if:

- All internal nodes have exactly two children
- All leaf nodes are at the same level
- The tree is completely filled up to the nth level

Properties:

- Total nodes = $2^{(h+1)} - 1$, where h is height
- Leaf nodes = 2^h

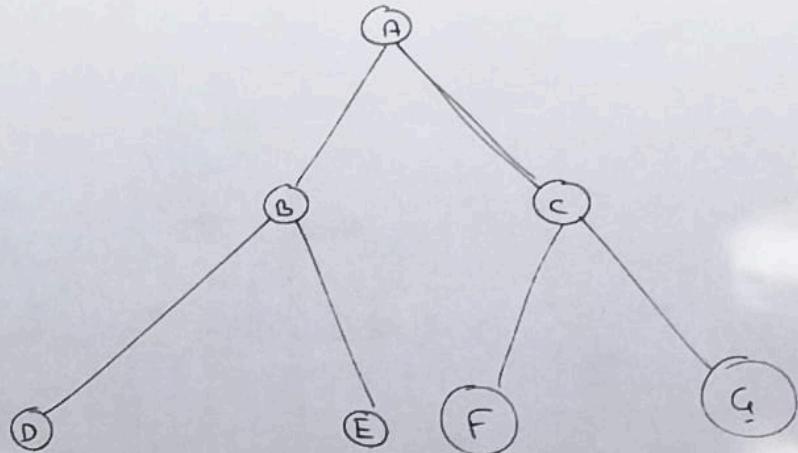
- Height = $\log_2(n+1) - 1$

Example:



(1) Perfect Binary Tree

(11)



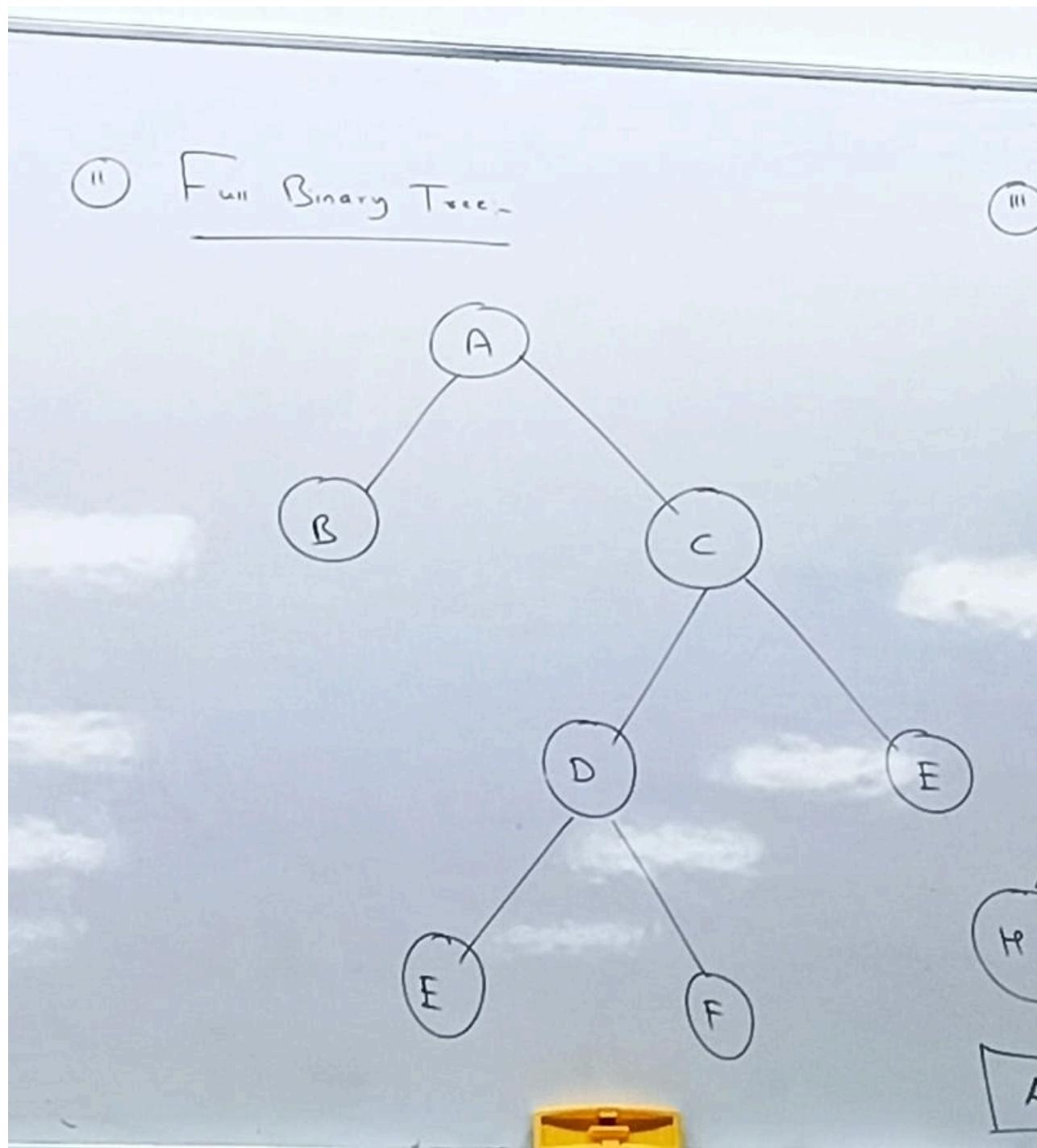
2. Full Binary Tree

A binary tree is **full** if every node has either:

- Zero children (leaf node), OR
- Two children (internal node)

Properties:

- If there are n nodes, then there are $(n+1)/2$ leaf nodes
- Number of internal nodes = $(n-1)/2$

Example:

3. Complete Binary Tree

A binary tree is **complete** if:

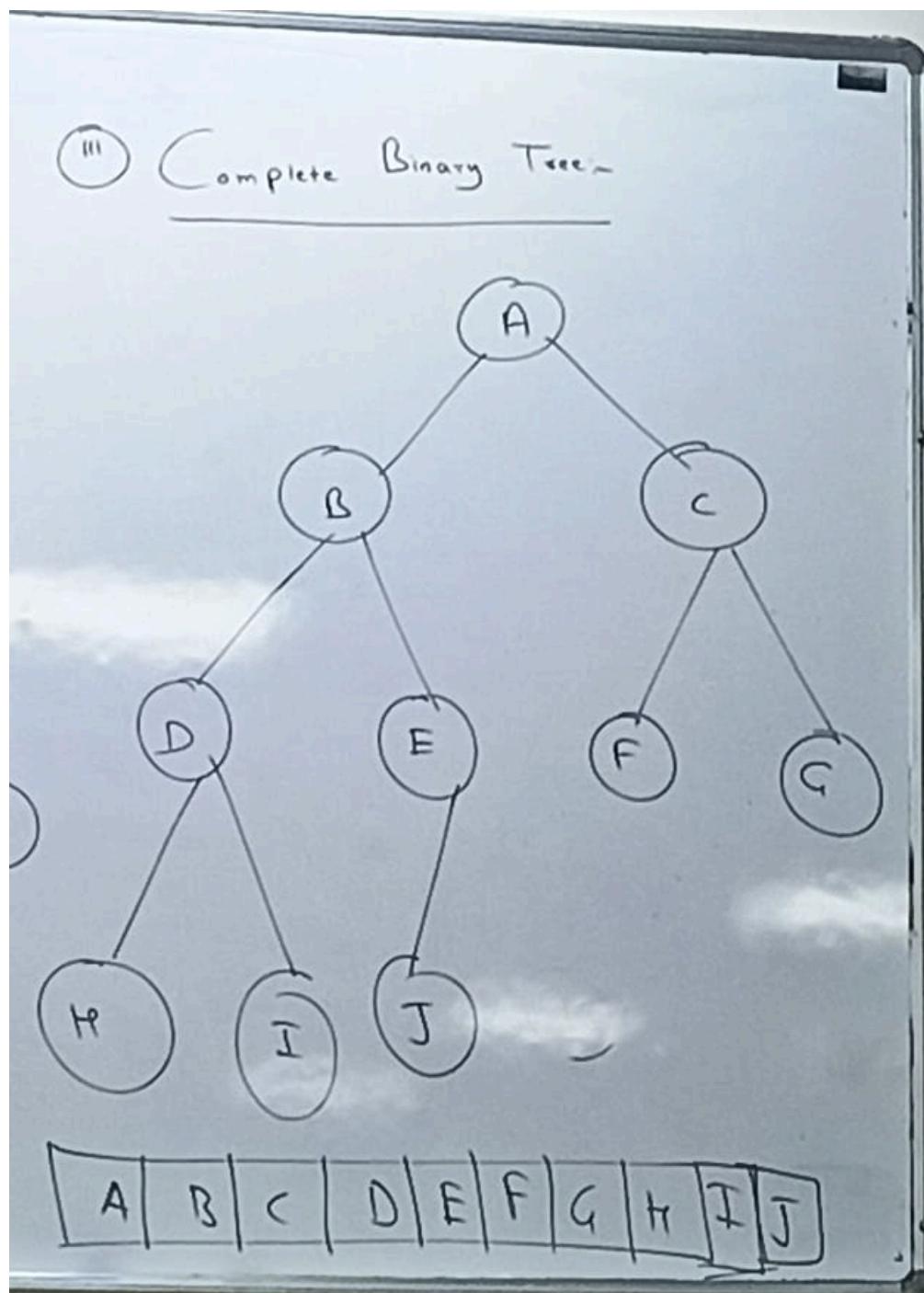
- All levels are completely filled except possibly the last level
- The last level is filled from left to right

- Used in heap data structures

Properties:

- Height = $\lfloor \log_2(n) \rfloor$
- Efficiently represented using arrays
- Parent of node at index i is at $\lfloor (i-1)/2 \rfloor$
- Left child at $2i + 1$, Right child at $2i + 2$

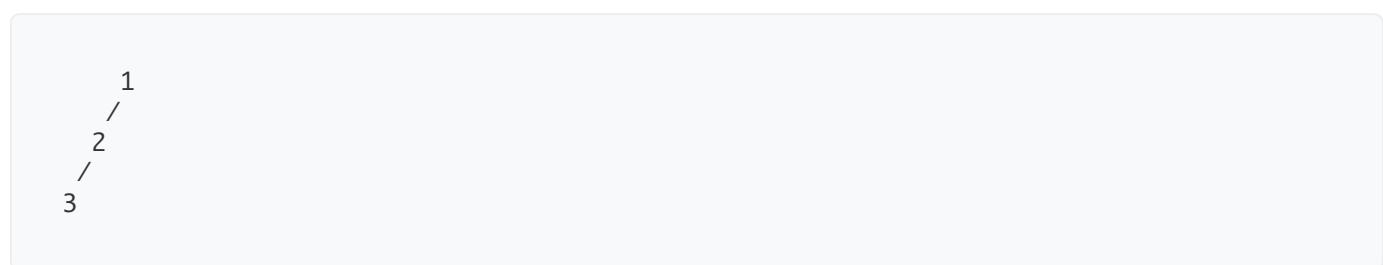
Example:



4. Degenerate (Skewed) Binary Tree

Every parent node has only one child. Performance degrades to $O(n)$ like a linked list.

Example (Left Skewed):



5. Balanced Binary Tree

A binary tree where the height difference between left and right subtrees is at most 1 for every node.
Examples: AVL Tree, Red-Black Tree.

Tree Traversal Methods

1. Inorder Traversal (Left → Root → Right)

Visits nodes in ascending order for BST.

```
public void inorderTraversal(TreeNode root) {
    if (root == null) return;

    inorderTraversal(root.left);
    System.out.print(root.data + " ");
    inorderTraversal(root.right);
}
```

Time Complexity: O(n) **Space Complexity:** O(h) - recursion stack, where h is height

2. Preorder Traversal (Root → Left → Right)

Used to create a copy of the tree or get prefix expression.

```
public void preorderTraversal(TreeNode root) {
    if (root == null) return;

    System.out.print(root.data + " ");
    preorderTraversal(root.left);
    preorderTraversal(root.right);
}
```

Time Complexity: O(n) **Space Complexity:** O(h)

3. Postorder Traversal (Left → Right → Root)

Used to delete the tree or get postfix expression.

```
public void postorderTraversal(TreeNode root) {
    if (root == null) return;

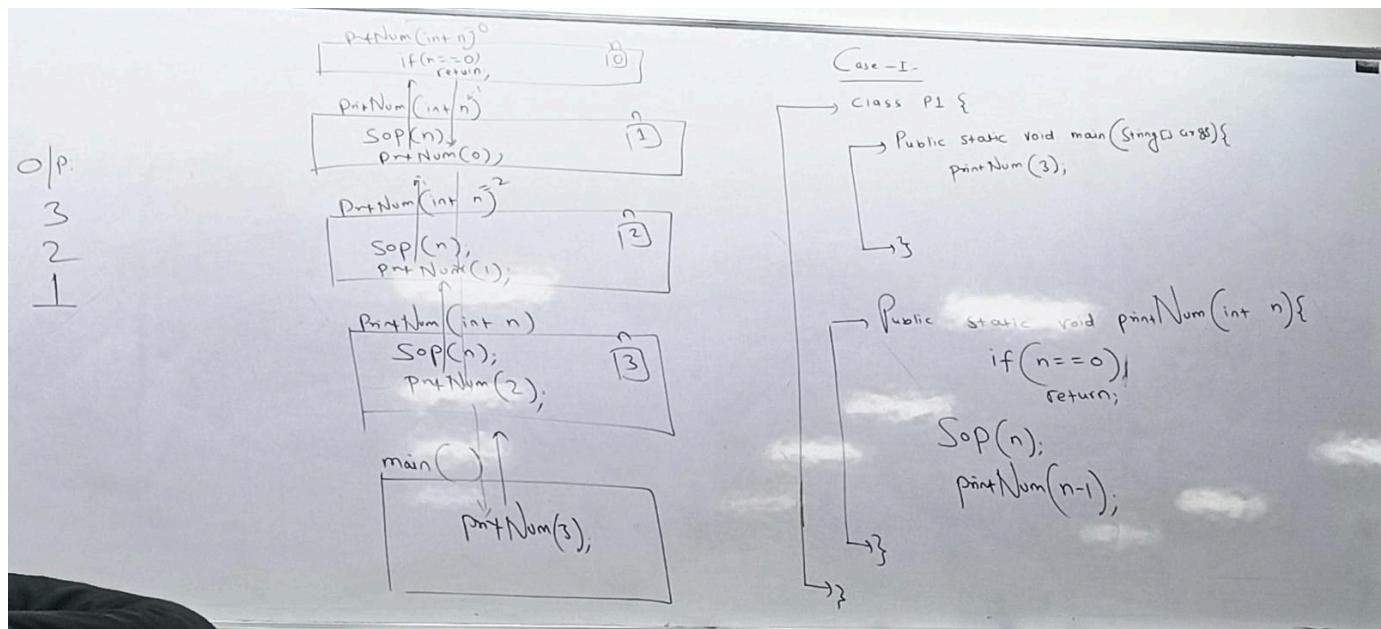
    postorderTraversal(root.left);
    postorderTraversal(root.right);
    System.out.print(root.data + " ");
}
```

Time Complexity: O(n) **Space Complexity:** O(h)

Traversal Cases Summary

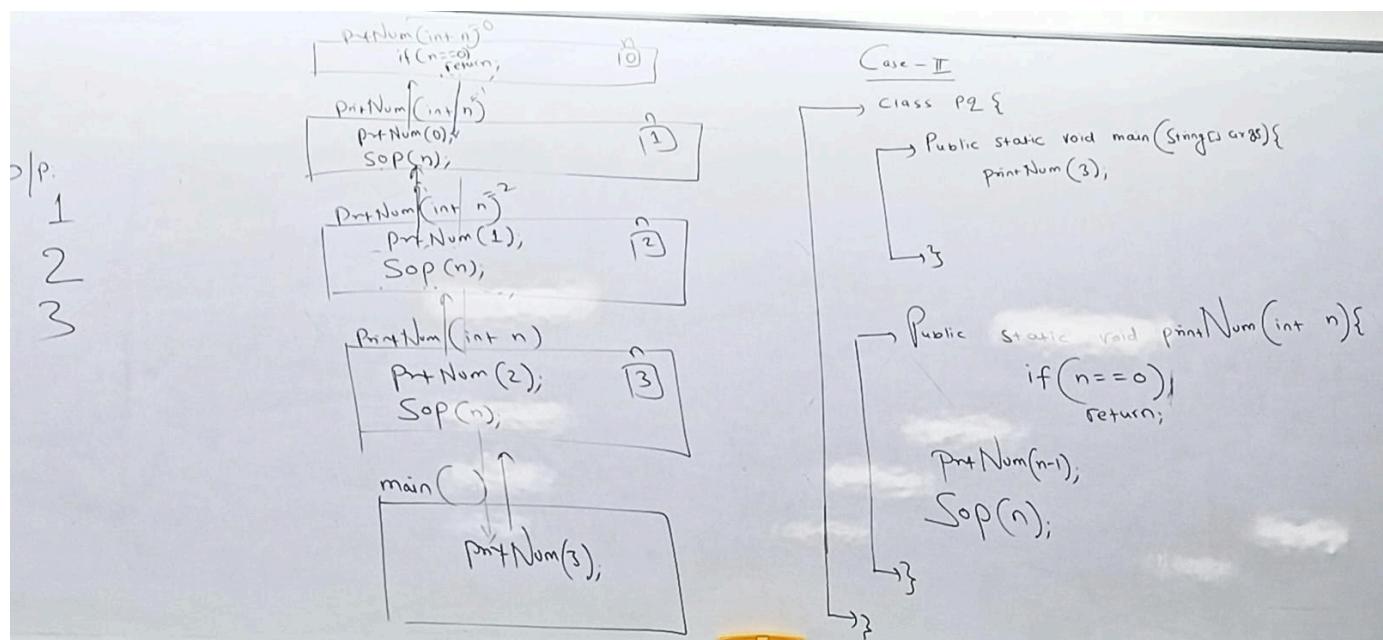
Case 1: Print First, Then Recurse (Preorder)

```
void case1(TreeNode root) {
    if (root == null) return;
    System.out.print(root.data + " ");
    case1(root.left);           // Then recurse left
    case1(root.right);         // Then recurse right
}
```



Case 2: Recurse First, Then Print (Postorder)

```
void case2(TreeNode root) {
    if (root == null) return;
    case2(root.left);           // Recurse left first
    case2(root.right);         // Then recurse right
    System.out.print(root.data + " ");
}
```

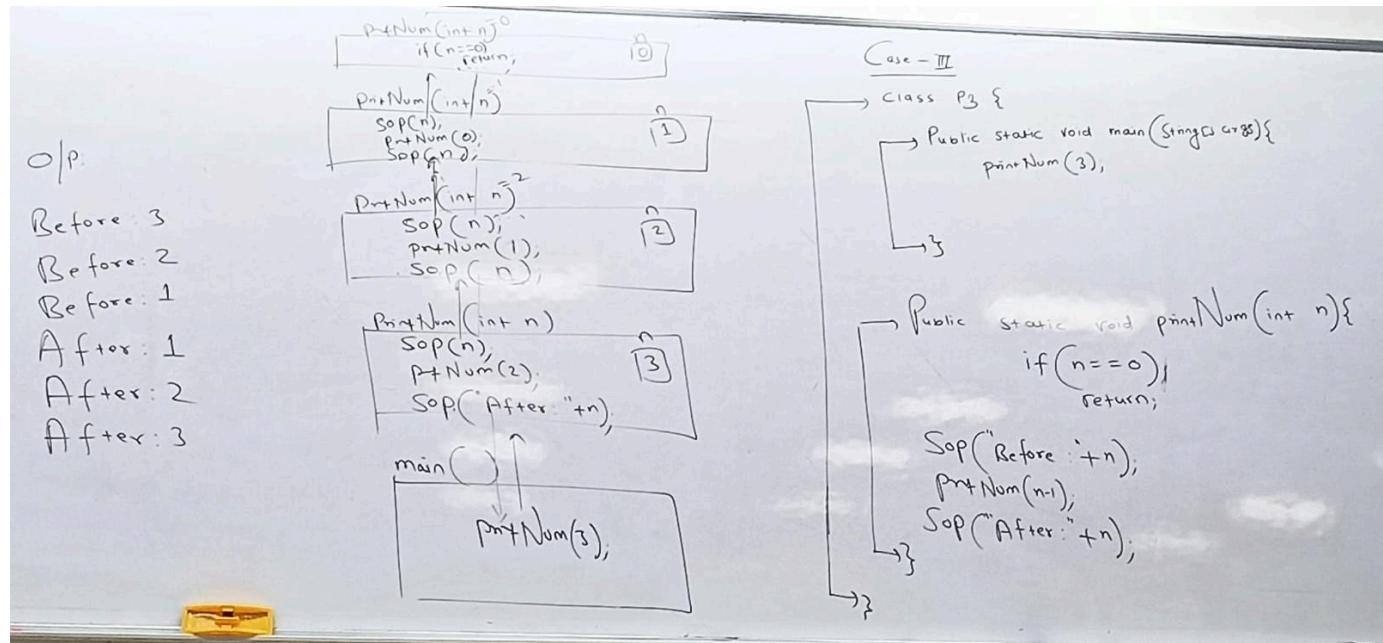


Case 3: Print Before and After Recursion (Inorder with extras)

```

void case3(TreeNode root) {
    if (root == null) return;
    System.out.print(root.data + " "); // Print before
    case3(root.left); // Recurse left
    case3(root.right); // Recurse right
    System.out.print(root.data + " "); // Print after
}

```



Binary Search Tree (BST)

A BST is a binary tree where:

- Left subtree contains only nodes with keys less than the node's key
- Right subtree contains only nodes with keys greater than the node's key

- Both subtrees are also BSTs

Search in BST (Recursive)

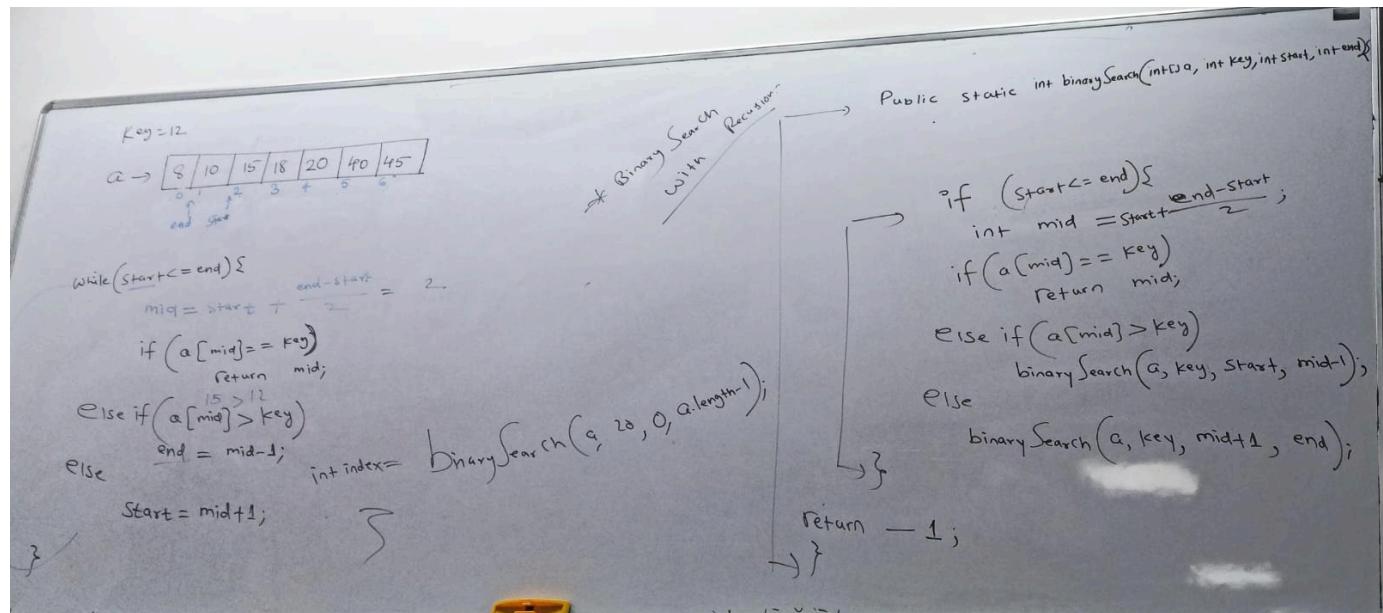
```

public TreeNode search(TreeNode root, int key) {
    // Base case: root is null or key is present at root
    if (root == null || root.data == key) {
        return root;
    }

    // Key is greater than root's key
    if (key > root.data) {
        return search(root.right, key);
    }

    // Key is smaller than root's key
    return search(root.left, key);
}

```



Time Complexity: $O(h)$ - $O(\log n)$ for balanced, $O(n)$ for skewed **Space Complexity:** $O(h)$ - recursion stack

Search in BST (Iterative)

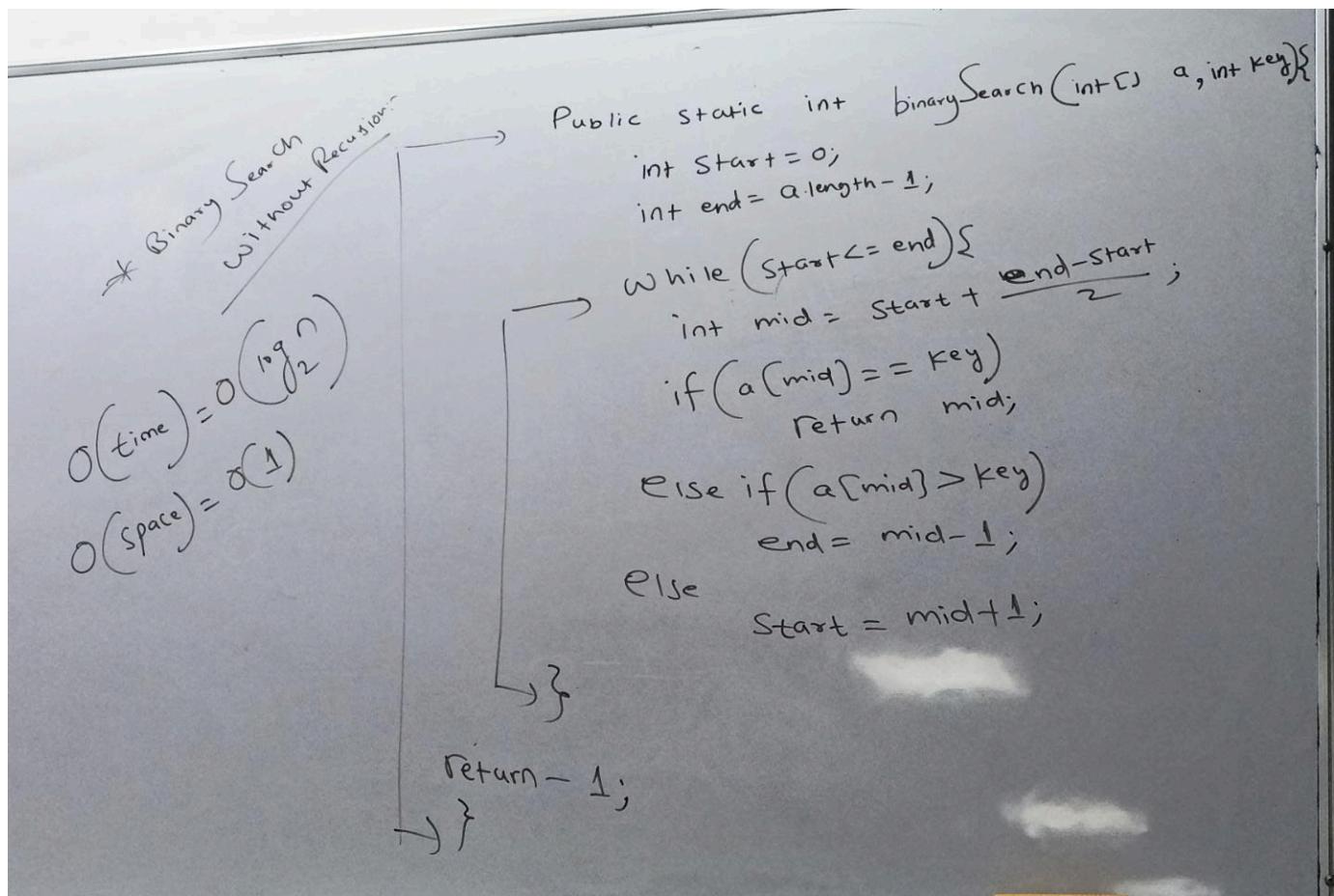
```

public TreeNode searchIterative(TreeNode root, int key) {
    TreeNode current = root;

    while (current != null && current.data != key) {
        if (key > current.data) {
            current = current.right;
        } else {
            current = current.left;
        }
    }

    return current;
}

```



Time Complexity: $O(h)$ **Space Complexity:** $O(1)$

Complexity Analysis

Time Complexities Summary

Operation	BST (Balanced)	BST (Skewed)	General Binary Tree
Search	$O(\log n)$	$O(n)$	$O(n)$
Insert	$O(\log n)$	$O(n)$	$O(n)$
Delete	$O(\log n)$	$O(n)$	$O(n)$

Operation	BST (Balanced)	BST (Skewed)	General Binary Tree
Traversal	$O(n)$	$O(n)$	$O(n)$
Height	$O(n)$	$O(n)$	$O(n)$

Space Complexities

Operation	Space Complexity	Notes
Recursive Traversal	$O(h)$	Recursion stack
Iterative Traversal	$O(w)$	Queue for level order
Search (Recursive)	$O(h)$	Stack space
Search (Iterative)	$O(1)$	No extra space

Where:

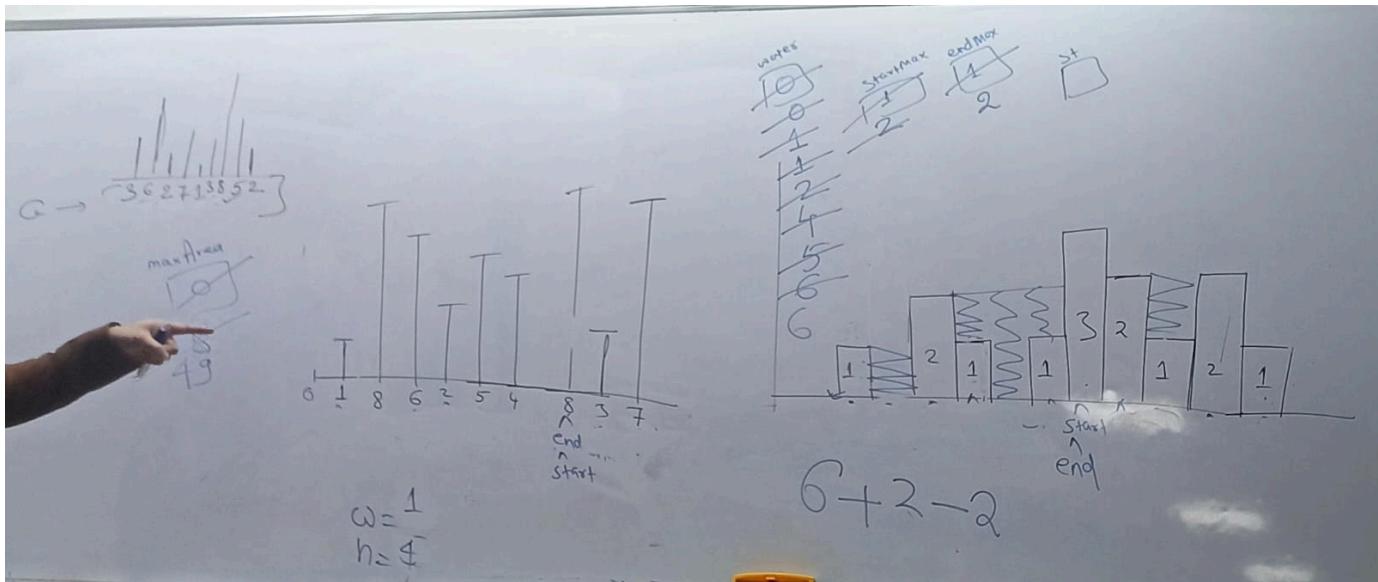
- n = number of nodes
 - h = height of tree
 - w = maximum width of tree
-

Best Practices

1. **Always check for null** before accessing node properties
 2. **Use iterative approaches** when possible to save stack space
 3. **For BST operations**, maintain the BST property
 4. **Choose balanced trees** (AVL, Red-Black) for guaranteed $O(\log n)$ operations
 5. **Use level order traversal** for problems requiring level-wise processing
 6. **Consider space-time tradeoffs** when choosing recursive vs iterative
-

Common Patterns

1. **Two Pointer**: Left and right child pointers
 2. **Recursion**: Most tree problems solved recursively
 3. **DFS**: Preorder, Inorder, Postorder
 4. **BFS**: Level order traversal using queue
 5. **Backtracking**: Path finding problems
 6. **Divide and Conquer**: Split problem into left and right subtrees
-



35. Search Insert Position

Solved

Example 1:

Input: nums = [1,3,5,6], target = 5
Output: 2

Example 2:

Input: nums = [1,3,5,6], target = 2
Output: 1

Example 3:

Input: nums = [1,3,5,6], target = 7
Output: 4

Constraints:

- 18.3K
- 397
- 97 Online

Code

```

1 class Solution {
2     public int searchInsert(int[] nums, int target) {
3         int start = 0;
4         int end = nums.length - 1;
5
6         while (start <= end) {
7             int mid = start + (end - start) / 2;
8
9             if (nums[mid] == target)
10                return mid;
11             else if (nums[mid] > target)
12                 end = mid - 1;
13             else
14                 start = mid + 1;
15         }
16         return start;
17     }
18 }
```

Java Auto

Saved Ln 18, Col 2

Testcase | Test Result

Accepted Runtime: 0 ms

Case 1 Case 2 Case 3

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leetcode.com/problems/binary-search/ Binary Search - LeetCode Search Insert Position - LeetCode

Problem List Submit Premium

Description Editorial Solutions Submissions

704. Binary Search

Solved

Easy Topics Companies

Given an array of integers `nums` which is sorted in ascending order, and an integer `target`, write a function to search `target` in `nums`. If `target` exists, then return its index. Otherwise, return `-1`.

You must write an algorithm with $O(\log n)$ runtime complexity.

Example 1:

Input: `nums = [-1,0,3,5,9,12]`, `target = 9`
Output: 4
Explanation: 9 exists in `nums` and its index is 4

Example 2:

Input: `nums = [-1,0,3,5,9,12]`, `target = 2`
Output: -1
Explanation: 2 does not exist in `nums` so return -1

Constraints:

- $1 \leq \text{nums.length} \leq 10^4$

13.1K 193 99 Online

Code Java Auto

```
1 class Solution {
2     public int search(int[] nums, int target) {
3         int start = 0;
4         int end = nums.length - 1;
5         int mid = 0;
6
7         while (start <= end) {
8             mid = start + (end - start) / 2;
9
10            if (nums[mid] == target)
11                return mid;
12            else if (nums[mid] > target)
13                end = mid - 1;
14            else
15                start = mid + 1;
16        }
17        return -1;
18    }
19}
20}
```

Saved Ln 20, Col 2

Testcase Test Result

Accepted Runtime: 0 ms

Case 1 Case 2

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leetcode.com/problems/best-time-to-buy-and-sell-stock/ Best Time to Buy and Sell - LeetCode Binary Search - LeetCode Search Insert Position - LeetCode

Problem List Submit Premium

Description Editorial Solutions Submissions

121. Best Time to Buy and Sell Stock

Solved

Easy Topics Companies

You are given an array `prices` where `prices[i]` is the price of a given stock on the `ith` day.

You want to maximize your profit by choosing a **single day** to buy one stock and choosing a **different day in the future** to sell that stock.

Return the **maximum profit** you can achieve from this transaction. If you cannot achieve any profit, return `0`.

Example 1:

Input: `prices = [7,1,5,3,6,4]`
Output: 5
Explanation: Buy on day 2 (price = 1) and sell on day 5 (price = 6), profit = 6-1 = 5.
Note that buying on day 2 and selling on day 1 is not allowed because you must buy before you sell.

Example 2:

Input: `prices = [7,6,4,3,1]`
Output: 0

34.7K 727 270 Online

Code Java Auto

```
1 class Solution {
2     public int maxProfit(int[] prices) {
3         int buy = prices[0];
4         int profit = 0;
5
6         for (int price : prices) {
7
8             if (price < buy)
9                 buy = price;
10            else if (price - buy > profit)
11                profit = price - buy;
12        }
13        return profit;
14    }
15}
```

Saved Ln 15, Col 2

Testcase Test Result

Accepted Runtime: 0 ms

Case 1 Case 2

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Maximum Depth of Binary Tree | Search Insert Position - LeetCode

leetcode.com/problems/maximum-depth-of-binary-tree/ Problem List Submit

Description Editorial Solutions Submissions

104. Maximum Depth of Binary Tree Solved

Easy Topics Companies

Given the `root` of a binary tree, return its maximum depth.

A binary tree's **maximum depth** is the number of nodes along the longest path from the root node down to the farthest leaf node.

Example 1:

```

10 *      this.val = val;
11 *      this.left = left;
12 *      this.right = right;
13 *     }
14 */
15 class Solution {
16     public int maxDepth(TreeNode root) {
17         if (root == null)
18             return 0;
19         int left = maxDepth(root.left);
20         int right = maxDepth(root.right);
21
22         return 1 + Math.max(left, right);
23     }
24 }
25 
```

Saved Ln 1, Col 1

Testcase Test Result

Accepted Runtime: 0 ms

Case 1 Case 2

14K 198 84 Online

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Minimum Depth of Binary Tree | Search Insert Position - LeetCode

leetcode.com/problems/minimum-depth-of-binary-tree/description/ Problem List Submit

Description Editorial Solutions Submissions

111. Minimum Depth of Binary Tree Solved

Easy Topics Companies

Given a binary tree, find its minimum depth.

The minimum depth is the number of nodes along the shortest path from the root node down to the nearest leaf node.

Note: A leaf is a node with no children.

Example 1:

```

14 */
15 */
16 class Solution {
17     public int minDepth(TreeNode root) {
18         if (root == null)
19             return 0;
20
21         int left = minDepth(root.left);
22         int right = minDepth(root.right);
23
24         if (left == 0 || right == 0)
25             return 1 + right + left;
26
27         return 1 + Math.min(left, right);
28     }
29 } 
```

Saved Ln 1, Col 1

Testcase Test Result

You must run your code first

7.7K 134 27 Online

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leetcode.com/problems/container-with-most-water/ Container With Most Water - LeetCode Search Insert Position - LeetCode

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Description Editorial Solutions Submissions

11. Container With Most Water

Solved

Medium Topics Companies Hint

You are given an integer array `height` of length `n`. There are `n` vertical lines drawn such that the two endpoints of the `ith` line are $(i, 0)$ and $(i, height[i])$.

Find two lines that together with the x-axis form a container, such that the container contains the most water.

Return the maximum amount of water a container can store.

Notice that you may not slant the container.

Example 1:

Accepted Runtime: 0 ms

Case 1 Case 2

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The Great Mantra: Hare Krishna avinash-01 - LeetCode Profile Search Insert Position - LeetCode

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Description Editorial Solutions Submissions

42. Trapping Rain Water

Solved

Hard Topics Companies

Given `n` non-negative integers representing an elevation map where the width of each bar is `1`, compute how much water it can trap after raining.

Example 1:

Input: height = [0,1,0,2,1,0,1,3,2,1,2,1]
Output: 6
Explanation: The above elevation map (black section) is represented by array [0,1,0,2,1,0,1,3,2,1,2,1]. In this case, 6 units of rain water (blue section) are being trapped.

Example 2:

Accepted Runtime: 0 ms

Case 1 Case 2