

# Parallel FFT

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## Algorithm:

### Recursive Implementation:

```
 $X_0, \dots, X_{N-1} \leftarrow \text{ditfft2}(x, N, s):$           DFT of  $(x_0, x_s, x_{2s}, \dots, x_{(N-1)s})$ :  
  if  $N = 1$  then                                trivial size-1 DFT base case  
     $X_0 \leftarrow x_0$   
  else  
     $X_0, \dots, X_{N/2-1} \leftarrow \text{ditfft2}(x, N/2, 2s)$       DFT of  $(x_0, x_{2s}, x_{4s}, \dots)$   
     $X_{N/2}, \dots, X_{N-1} \leftarrow \text{ditfft2}(x+s, N/2, 2s)$   DFT of  $(x_s, x_{s+2s}, x_{s+4s}, \dots)$   
    for  $k = 0$  to  $N/2-1$                                 combine DFTs of two halves into full DFT:  
       $t \leftarrow X_k$   
       $X_k \leftarrow t + \exp(-2\pi i k/N) X_{k+N/2}$   
       $X_{k+N/2} \leftarrow t - \exp(-2\pi i k/N) X_{k+N/2}$   
    endfor  
  endif
```

In the above method, it is difficult to parallelize the above code. Iterative approaches can be easily parallelized.

### Iterative Approach of Cooley-Tukey FFT:

```

algorithm iterative-fft is
  input: Array  $a$  of  $n$  complex values where  $n$  is a power of 2
  output: Array  $A$  the DFT of  $a$ 

  bit-reverse-copy( $a, A$ )
   $n \leftarrow a.length$ 
  for  $s = 1$  to  $\log(n)$ 
     $m \leftarrow 2^s$ 
     $\omega_m \leftarrow \exp(-2\pi i/m)$ 
    for  $k = 0$  to  $n-1$  by  $m$ 
       $\omega \leftarrow 1$ 
      for  $j = 0$  to  $m/2 - 1$ 
         $t \leftarrow \omega A[k + j + m/2]$ 
         $u \leftarrow A[k + j]$ 
         $A[k + j] \leftarrow u + t$ 
         $A[k + j + m/2] \leftarrow u - t$ 
         $\omega \leftarrow \omega \omega_m$ 

  return  $A$ 

```

In this method, we are doing bit-reverse copy, pseudocode for this is :

```

algorithm bit-reverse-copy( $a, A$ ) is
  input: Array  $a$  of  $n$  complex values where  $n$  is a power of 2,
  output: Array  $A$  of size  $n$ 

   $n \leftarrow a.length$ 
  for  $k = 0$  to  $n - 1$ 
     $A[\text{rev}(k)] = a[k]$ 

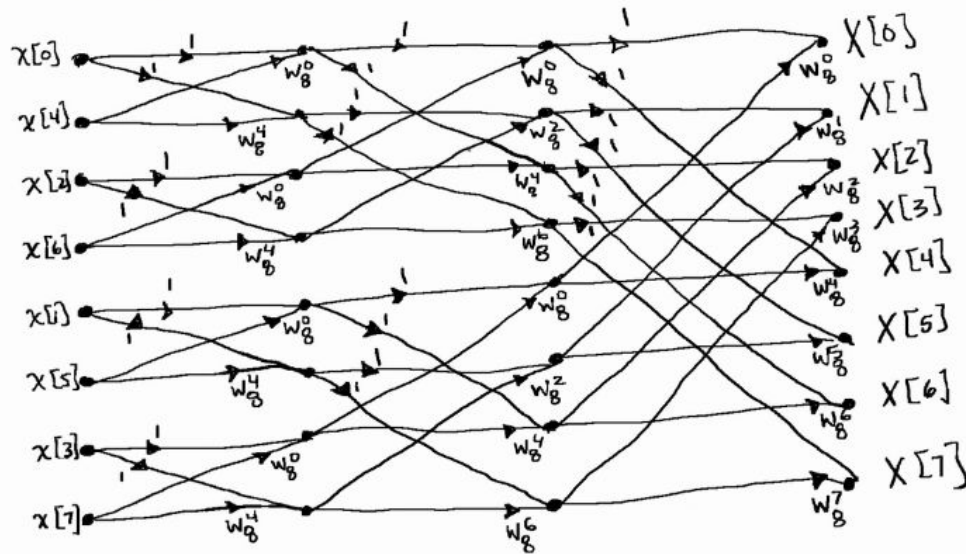
```

This method can be easily parallelized as it is an iterative approach.

Butterfly representation of above approach is:

Example:  $N=8=2^3$

6



## Comparison:

For  $N = 131072$

- For recursive code, time taken is 0.025 sec
- For iterative code, time taken is 0.021 sec
- For omp code,
  - Threads = 4, static schedule, time taken is 0.017 sec
  - Threads = 4, guided schedule, time taken is 0.014 sec
  - Threads = 6, static schedule, time taken is 0.016 sec
  - Threads = 6, guided schedule, time taken is 0.013 sec
  - Threads = 2, static schedule, time taken is 0.019 sec
  - Threads = 2, guided schedule, time taken is 0.017 sec

For  $N = 262144$

- For recursive code, time taken is 0.030
- For iterative code, time taken is 0.024
- For omp code,
  - Threads = 4, static schedule, time taken is 0.019
  - Threads = 4, guided schedule, time taken is 0.016
  - Threads = 6, static schedule, time taken is 0.017

- Threads = 6, guided schedule, time taken is 0.015
- Threads = 2, static schedule, time taken is 0.022
- Threads = 2, guided schedule, time taken is 0.020

## **Observation**

- From the above experiment, we can observe that the time taken by recursive approach is more when compared to iterative approach
- Also, we can observe that the Parallelised iterative approach is faster than other two methods.
- We can also observe that if we increase the threads, speed of execution also increases.
- Also we can observe that, guided works better than static.