Optimizing Hive Joins

Overview

Understand how joins are implemented in Hive

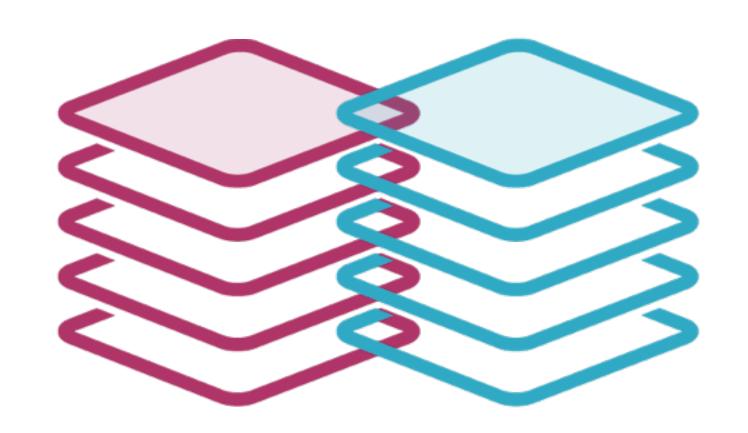
Optimize joins which involve large tables

Use semi-joins in place of IN/EXISTS subqueries

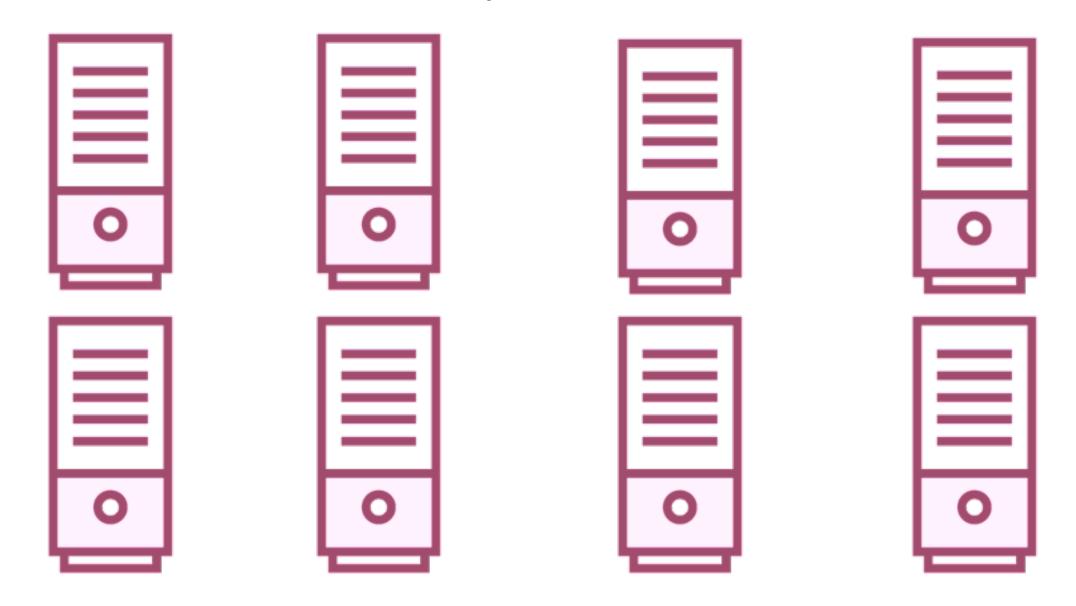
Optimize map-only joins

Join Operations as MapReduce Jobs

Join Operations



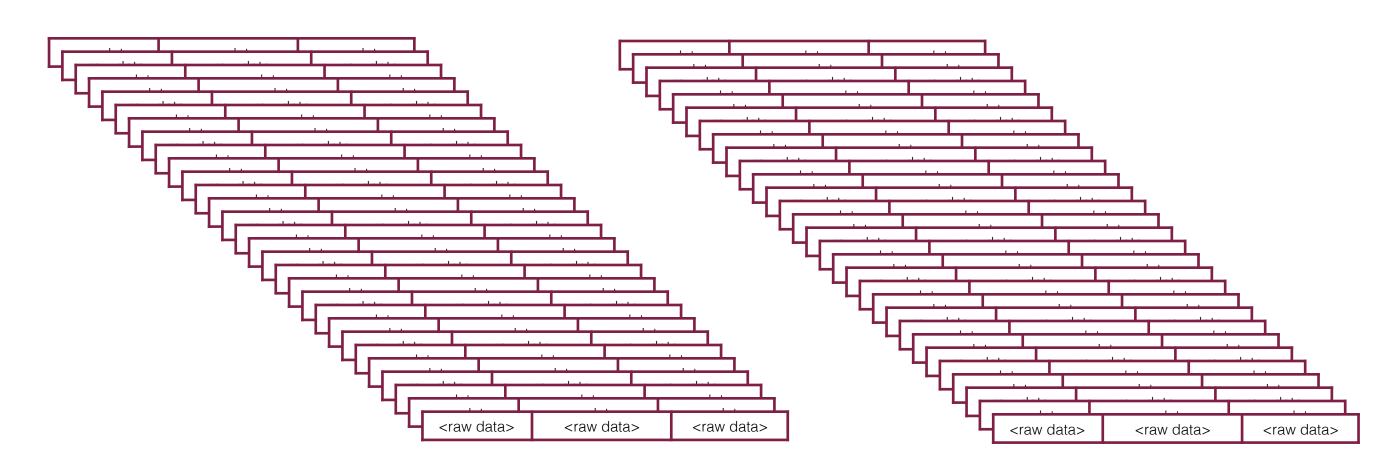
Join operations are MapReduce jobs under the hood



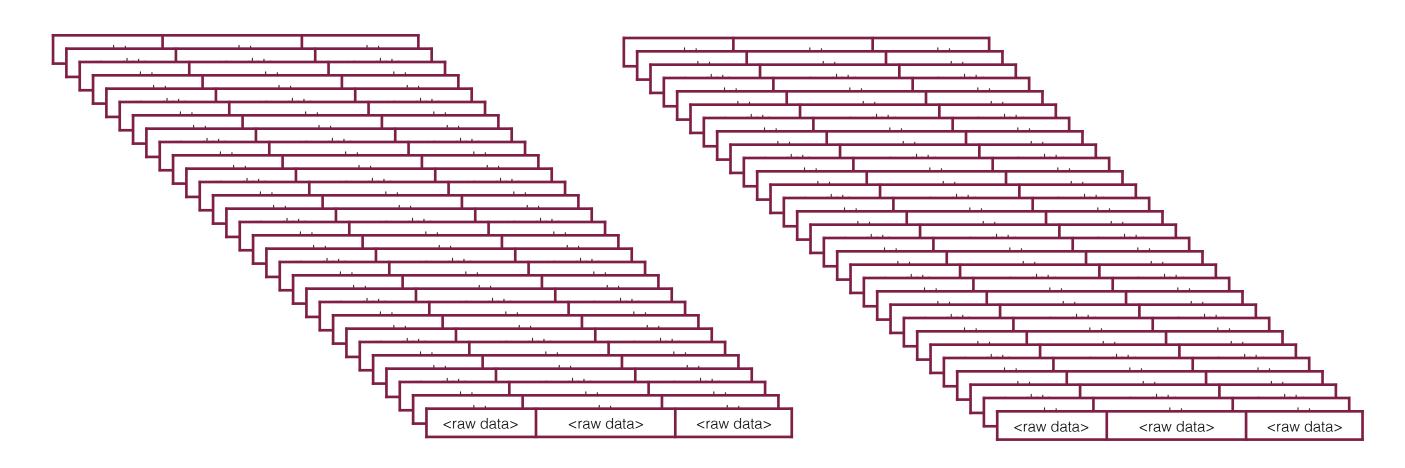
A programming paradigm which runs on a distributed system



Takes advantage of the inherent parallelism in data processing



Modern systems generate millions of records of raw data

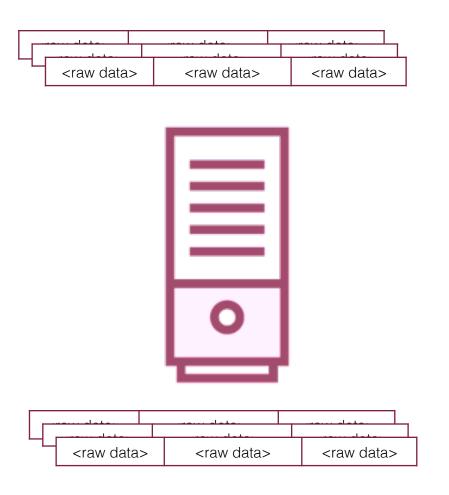


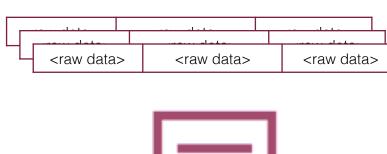
A task of this scale is processed in two stages

map

reduce

map







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reduce







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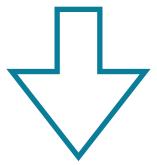


map

An operation performed in parallel, on small portions of the dataset

map

One Record

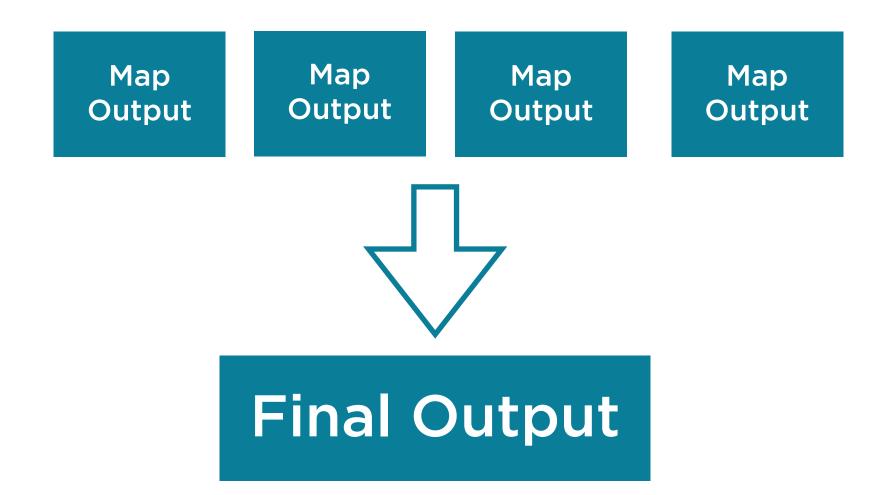


Key-Value Output

reduce

An operation to combine the results of the map step

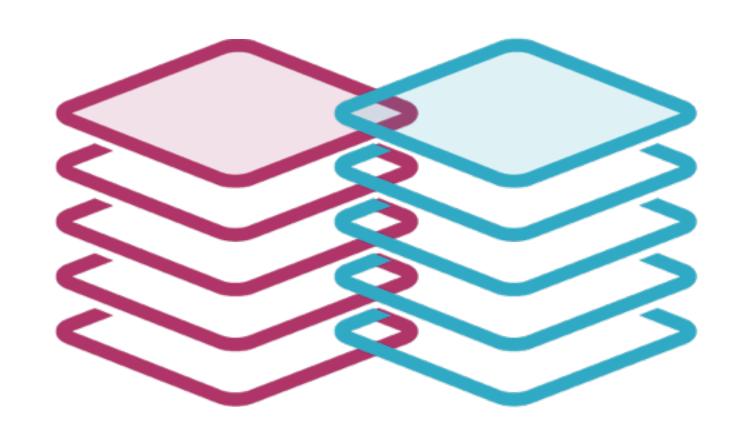
reduce



Map A step that can be performed in parallel

reduce A step to combine the intermediate results

Join Columns and MapReduce Jobs



Join combines records from two or more tables on the same column value

Trades

Symbol	Open	High	Low	Close	Day
GOOG	820	840	818	829	1-1-2017

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

select * from Names join Trades
on Names.Symbol = Trades.Symbol

Trades

Symbol	Open	High	Low	Close	Day
GOOG	820	840	818	829	1-1-2017

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

One join column = one MapReduce job

Trades

Symbol	Open	High	Low	Close	Day
GOOG	820	840	818	829	1-1-2017

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

```
select * from Names join Trades
on (Names.Symbol = Trades.Symbol)
join Revenues on (Names.Symbol = Revenues.Symbol)
```

Trades

Symbol	Open	High	Low	Close	Day
GOOG	820	840	818	829	1-1-2017

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues

Symbol	Name	Revenue
GOOG	Google	90B
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One join column = one MapReduce job

Trades

Symbol	Open	High	Low	Close	Day
GOOG	820	840	818	829	1-1-2017

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

```
select * from Names join Trades
on (Names.Symbol = Trades.Symbol)
join Revenues on (Names.Name = Revenues.Name)
```

Trades

Symbol	Open	High	Low	Close	Day
GOOG	820	840	818	829	1-1-2017

Names

Symbol	Name	
GOOG	Google	
AAPL	Apple	
MSFT	Microsoft	

Revenues

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

select * from Names join Trades

on (Names.Symbol = Trades.Symbol)

join Revenues on (Names.Name = Revenues.Name)

Trades

Symbol	Open	High	Low	Close	Day
GOOG	820	840	818	829	1-1-2017

Names

Symbol	Name	
GOOG	Google	
AAPL	Apple	
MSFT	Microsoft	

Revenues

Name	Revenue
Google	90B
Apple	215B
Microsoft	85B
	Google Apple

Two join columns = two MapReduce jobs

For faster queries...

Minimize the number of MapReduce jobs run

Demo

Join operations on 3 tables, with different columns in the join clause

Join Operations and Table Sizes

Trades 500GB

Symbol	Open	High	Low	Close
GOOG	820	840	818	829

Names 10MB

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues 100MB

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

```
select * from Names join Trades
on (Names.Symbol = Trades.Symbol)
join Revenues on (Names.Symbol = Revenues.Symbol)
```

Trades 500GB

Symbol	Open	High	Low	Close
GOOG	820	840	818	829

Names 10MB

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues 100MB

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

```
select * from Names join Trades
```

```
on (Names.Symbol = Trades.Symbol)
```

join Revenues on (Names.Symbol = Revenues.Symbol)

Trades 500GB

Symbol	Open	High	Low	Close
GOOG	820	840	818	829

Names 10MB

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues 100MB

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

Names 10MB

Trades
500GB

Revenues 100MB

Names

Trades

Revenues

10MB

500GB

100MB



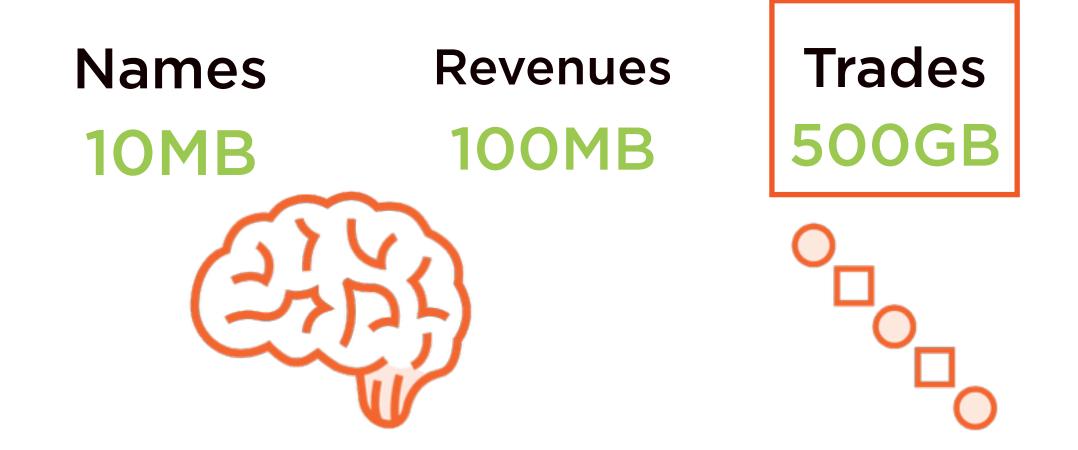
All tables, except for the last are held in memory



The last table is streamed from disk to the job



Inefficient to keep large tables in memory



Re-order the tables in the join so the largest table is at the end

Trades 500GB

Symbol	Open	High	Low	Close
GOOG	820	840	818	829

Names 10MB

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues 100MB

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

```
select * from Names join Revenues
on (Names.Symbol = Revenues.Symbol)
join Trades on (Names.Symbol = Trades.Symbol)
```

Trades 500GB

Symbol	Open	High	Low	Close
GOOG	820	840	818	829

Names 10MB

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Revenues 100MB

Symbol	Name	Revenue
GOOG	Google	90B
AAPL	Apple	215B
MSFT	Microsoft	85B

```
select * from Names join Revenues
```

```
on (Names.Symbol = Revenues.Symbol)
join Trades on (Names.Symbol = Trades.Symbol)
```

For faster queries...

Specify the largest table at the very end

```
select /*+ streamtable(Trades) */
Names.Symbol, Trades.High, Revenues.Revenue
from Names join Trades
on (Names.Symbol = Trades.Symbol)
join Revenues on (Names.Symbol = Revenues.Symbol)
```

The Streamtable Keyword

Specify which table to stream in a join operation Stream the largest table, do not hold it in memory

Join Optimizations with Bucketing and Partitioning

Orders

ID	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35
03	8	1	33
04	5	2	69
05	1	1	123
06	6	1	99
07	2	2	24
80	3	2	20

Products

ID	Name	Cost
1	iPhone	599
2	Doll	35
3	Shoes	33
4	Jeans	69
5	Skates	123
6	Make Up	99
7	Book	24
8	Belt	20

Join Orders and Products to get the names of the products users have bought

Ord	ers
-----	-----

ID	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35
03	8	1	33
04	5	2	69
05	1	1	123
06	6	1	99
07	2	2	24
80	3	2	20

Products

ID	Name	Cost
7	iPhone	599
8	Doll	35
3	Shoes	33
1	Jeans	69
6	Skates	123
5	Make Up	99
4	Book	24
2	Belt	20

Product ID is the join column

Orders Products

ΪD	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35
03	8	1	33
04	5	2	69
05	1	1	123
06	6	1	99
07	2	2	24
08	3	2	20

ID	Name	Cost
7	iPhone	599
8	Doll	35
3	Shoes	33
1	Jeans	69
6	Skates	123
5	Make Up	99
4	Book	24
2	Belt	20

Need to scan the entire dataset to find the corresponding row

Orders

ID	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35
03	8	1	33
04	5	2	69
05	1	1	123
06	6	1	99
07	2	2	24
80	3	2	20

Bucket the Products table on the ID column

Products

ID	Name	Cost
6	Skates	123
3	Shoes	33
ID	Name	Cost
7	iPhone	599
1	Jeans	69
4	Book	24
ID	Name	Cost
2	Belt	20
8	Doll	35
5	Make Up	99

Orders

İD	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35
03	8	1	33
04	5	2	69
05	1	1	123
06	6	1	99
07	2	2	24
08	3	2	20

Products

ID	Name	Cost
6	Skates	123
3	Shoes	33
ID	Name	Cost
7	iPhone	599
1	leans	69
•	 	" "
4	Book	24
4 ID	3 3 3.1 1 3	-
	Book	24
ID	Book	24 Cost

Scan a much smaller dataset to access each row

Orders

ID	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35
03	8	1	33
04	5	2	69
05	1	1	123
06	6	1	99
07	2	2	24
08	3	2	20

Products

ID	Name	Cost
6	Skates	123
3	Shoes	33
ID	Name	Cost
7	iPhone	599
1	Jeans	69
4	Book	24
ID	Name	Cost
2	Belt	20
8	Doll	35
5	Make Up	99

Faster joins

Partitioning

Join optimizations would work the same way

Reduce the dataset to scan to find the corresponding row

For faster queries...

Use bucketing or partitioning on the join columns

Left Semi-joins Instead of Subqueries

Any join query that requests rows from the left row source based on the existence of rows in the right row source without including data from the right row source in the final result and without duplicating rows from the left row source is a logical left semi join.

http://sqlity.net/

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Any join query that requests rows from the left row source based on the existence of rows in the right row source without including data from the right row source in the final result and without duplicating rows from the left row source is a logical left semi join. http://sqlity.net/

Names Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					
MSFT					

```
select names.symbol
from names left semi join trades
on
(names.symbol = trades.symbol);
```

Names Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					
MSFT					

Symbol	Name
GOOG	Google
MSFT	Microsoft

```
select symbol from names
where symbol in
(select symbol from trades);
```

The IN/NOT IN Keywords

Check for existence of a column value in the subquery

```
select symbol from names
where symbol in
(select symbol from trades);
```

The IN/NOT IN Keywords

Only one column can be selected in the subquery

```
select names.symbol
from names
where exists
(select trades.symbol from trades
where names.symbol = trades.symbol);
```

The EXISTS/NOT EXISTS Keywords

Check if even one result is returned from the subquery

```
select names.symbol
from names
where exists
(select trades.symbol from trades
where names.symbol = trades.symbol);
```

The EXISTS/NOT EXISTS Keywords

The subquery has to reference the parent query i.e. should be correlated

It is much more efficient to replace the IN/EXISTS subqueries with a left semi-join

```
select names.symbol
from names left semi join trades
on
(names.symbol = trades.symbol);
```

Choose records from the left table

```
select names.symbol
from names left semi join trades
on
(names.symbol = trades.symbol);
```

Choose records from the left table

```
select names.symbol
from names left semi join trades
on
(names.symbol = trades.symbol);
```

Choose records from the left table

Based on the existence of rows in the right table

```
select names.symbol
from names left semi join revenues
on
(names.symbol = revenues.symbol
and
names.name = revenues.name);
```

Allows specifying multiple columns in the join

Not allowed when we use the IN keyword

Left Semi-join vs. In/Exists

Left Semi-join

In/Exists

Only scans the right table till a match is found

Can specify matches on multiple columns

Bucketing or partitioning the right table can improve performance

Needs to the scan the entire subquery table

IN allows only one column to be selected in the subquery

No advantage of a bucketed or partitioned right table

For faster queries...

Use the **left semi-join** rather than in/exists keywords with subqueries

Demo

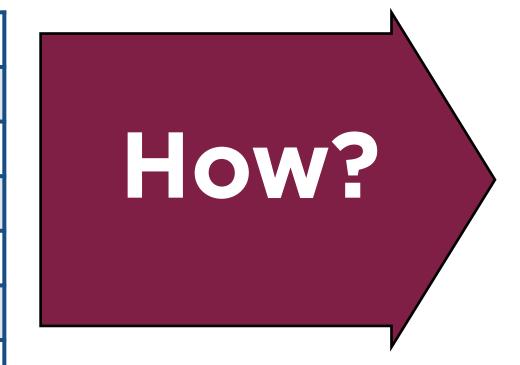
The left semi-join as opposed to the in/exists subqueries

The Anatomy of a MapReduce Program

Counting Word Frequencies

Consider a large text file

Twinkle twinkle little star		
How I wonder what you are		
Up above the world so high		
Like a diamond in the sky		
Twinkle twinkle little star		
How I wonder what you are		



Word	Frequency
above	14
are	20
how	21
star	22
twinkle	32

MapReduce Flow

Twinkle twinkle little star

How I wonder what you are



Up above the world so high

Like a diamond in the sky



Each partition is given to a different process i.e. to mappers

Twinkle twinkle little star

How I wonder what you are



MapReduce Flow

Twinkle twinkle little star

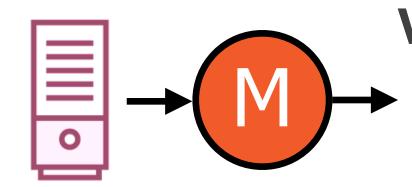
How I wonder what you are

■

Each mapper works in parallel

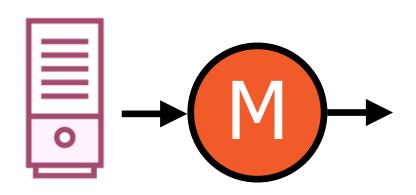
Up above the world so high

Like a diamond in the sky



Twinkle twinkle little star

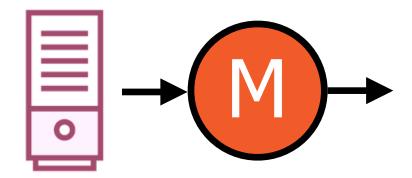
How I wonder what you are



Map Flow

Twinkle twinkle little star

How I wonder what you are

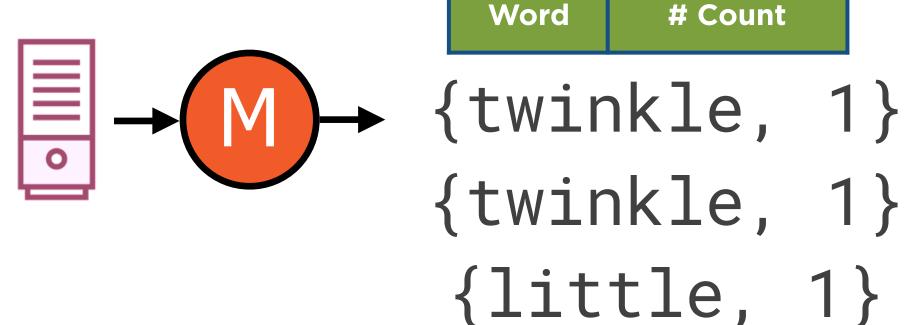


Within each mapper, the rows are processed serially

Map Flow

Twinkle twinkle little star

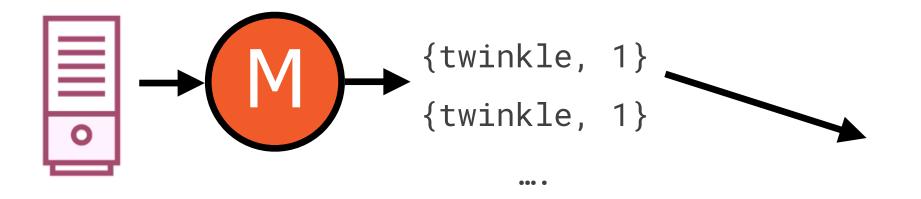
How I wonder what you are

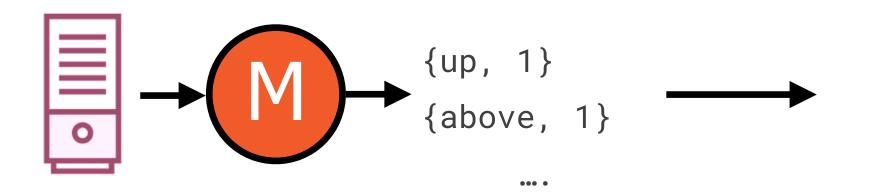


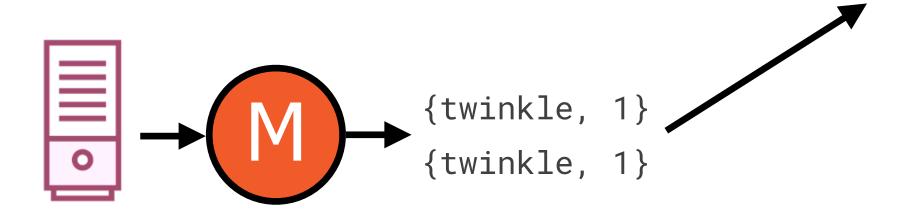
{star, 1}

Each row emits {key, value} pairs

Reduce Flow

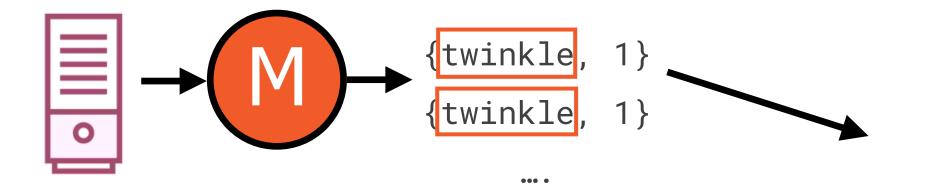


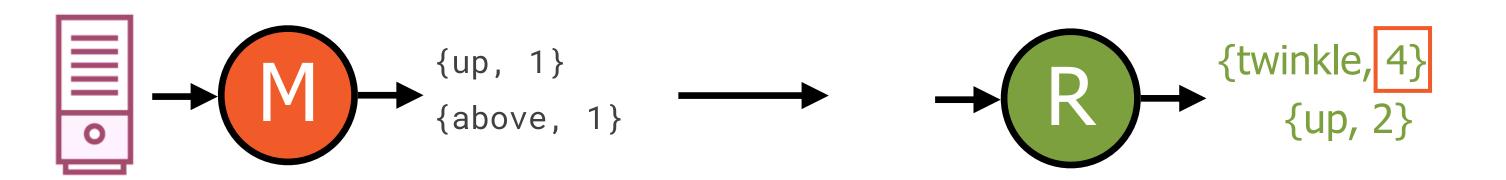


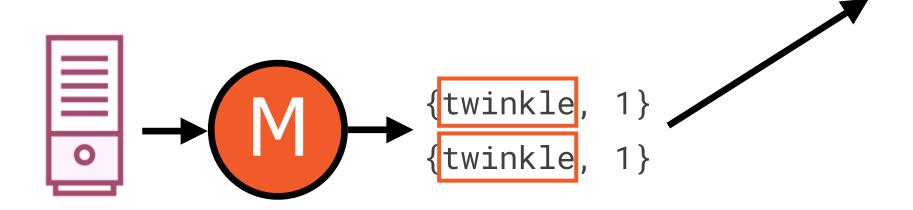


The results are passed on to another process i.e. a reducer

Reduce Flow







The reducer combines the values with the same key

Key Insight Behind MapReduce

Many data processing tasks can be expressed in this form

Orders

ID	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35
03	8	1	33
04	5	2	69
05	1	1	123
06	6	1	99
07	2	2	24
80	3	2	20

Products

ID	Name	Cost
1	iPhone	599
2	Doll	35
3	Shoes	33
4	Jeans	69
5	Skates	123
6	Make Up	99
7	Book	24
8	Belt	20

select * from Orders join Products
on Orders.ProductID = Products.ID

Orders

ID	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35



Products

ID	Name	Cost
1	iPhone	599
2	Doll	35



The mapper operates on each row of the tables

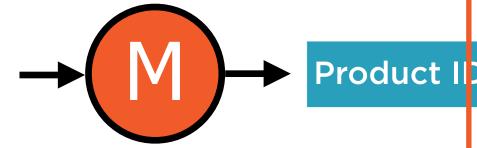
Orders

ID	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35
	Р	roducts	
ID	Name	9	Cost
1	iPhon	е	599
2	Doll		35

The join column is the key

Orders

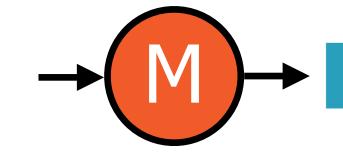
ID	Product ID	Quantity	Amount
01	4	1	599
02	7	1	35



Product ID ID, Quantity, Amount

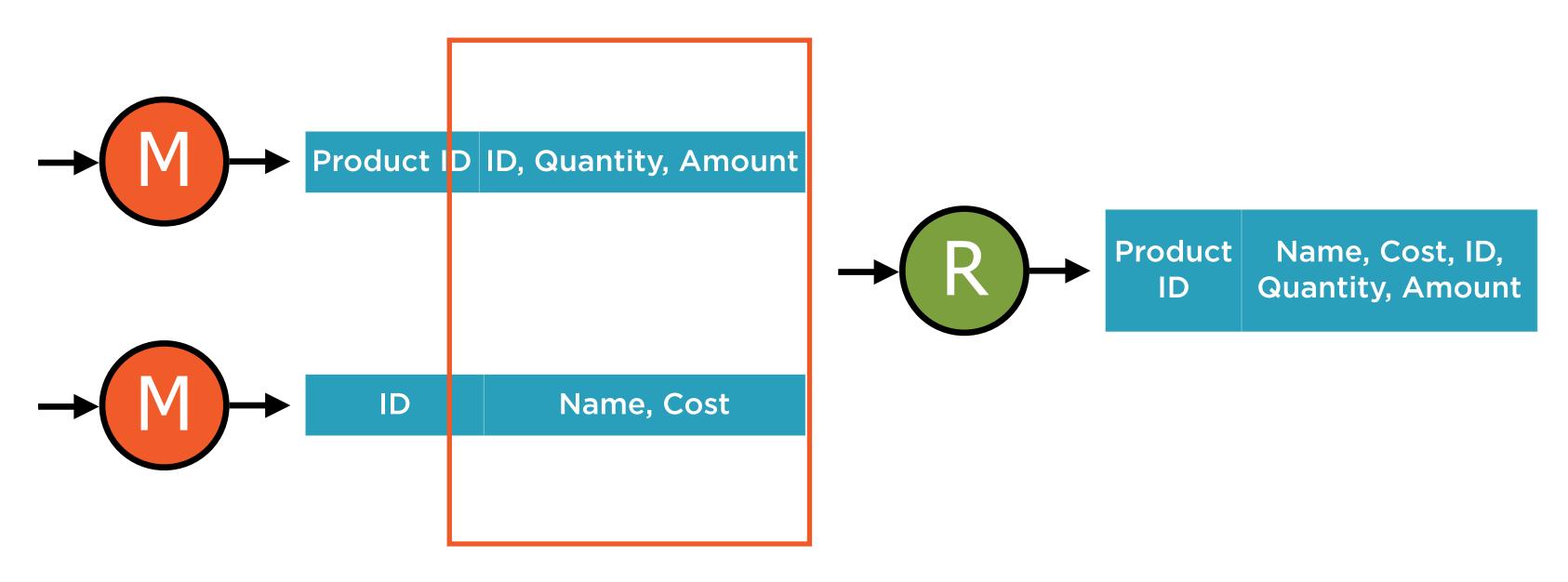
Products

ID	Name	Cost
1	iPhone	599
2	Doll	35



ID Name, Cost

The remaining columns are values



The reducer combines all columns which have the same key



MapReduce operations have 2 phases of processing





Certain queries can be structured to have no reduce phase



Such joins are called map-side joins

More performant



Improves processing time

Reduces data transfer between machines in the cluster

Reduces operations such as shuffle and sort between map and reduce phases



We'd like joins to be map-side joins if possible

Conditions for Map-side Joins

Conditions for Map-side Joins



All tables, except one, are small enough to be held in memory

Tables are bucketed on the join columns and Table1 Buckets = N * Table2 Buckets

Conditions for Map-side Joins



All tables, except one, are small enough to be held in memory

Tables are bucketed on the join columns and Table1 Buckets = N * Table2 Buckets

All Tables but One Are Small

Names				
Symbol	Name			
GOOG	Google			
AAPL	Apple			
MSFT	Microsoft			

Trades

Symbol	Open	High	Low	Close	Day
GOOG					

Left table is the smaller table

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Trades

Symbol	Open	High	Low	Close	Day
GOOG					

Only rows which have a match in both the left and right table

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Trades

```
SymbolOpenHighLowCloseDayGOOG
```

select * from Names join Trades
on (Names.Symbol = Trades.Symbol)

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

A copy of the smaller table is stored in a hash table like structure

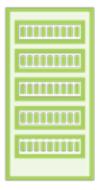
Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



This hash table is copied to each mappers' local disk

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Trades

Symbol	Open	High	Low	Close	Day
GOOG					

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Trades

Symbol	Open	High	Low	Close	Day
AAPL					

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Trades

Symbol	Open	High	Low	Close	Day
MSFT					

Parts of the larger table are distributed to each mapper

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Mappers run on the entire Names and parts of the Trades table

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



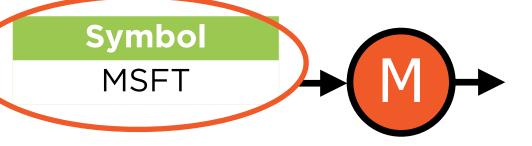
Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Combine with those rows which are available on the mapper

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Symbol	Other Columns
GOOG	

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Symbol	Other Columns
AAPL	

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Symbol	Other Columns
MSFT	

The output of all the mappers forms the final output

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	
GOOG	→(M)→

Symbol	Other Columns
GOOG	

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Symbol	Other Columns
AAPL	

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Symbol	Other Columns
MSFT	

No reducer needed

The left table is the smaller table

Inner Joins

Are possible as map-only joins

Names Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					

All rows from the left table are in the result

- with a matching row
- padded with nulls

Names Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					

select * from Names left outer join Trades
on (Names.Symbol = Trades.Symbol)

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



The smaller table is copied to each mappers's disk

Names

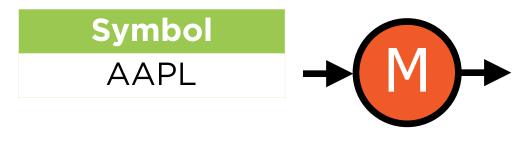
Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol GOOG M

Is the row not present in this chunk or not present in the entire table?

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



No way to tell

The left table is the smaller table

Inner Joins Left Outer Joins

Are possible as map-only joins

Right Outer Join

Names Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					

All rows from the right table are in the result

- with a matching row
- padded with nulls

Right Outer Join

Names Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					

select * from Names right outer join Trades
on (Names.Symbol = Trades.Symbol)

Right Outer Join

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol GOOG M

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



The entire left table is present to check for matches

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



The left table is the smaller table

Inner Joins

Left Cuter Joins

Right Outer Joins

Are possible as map-only joins

Full Outer Join

Names Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					

All rows from the both tables are in the result

- with a matching row
- padded with nulls

Full Outer Join

Names

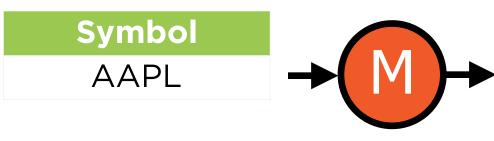
Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol GOOG M

Should a row be padded with nulls or does it have a match?

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft



No way to tell

The left table is the smaller table

Inner Joins

Left Outer Joins

Right Outer Joins

Full Outer Joins

Are possible as map-only joins

Right Table Is Small

Trades

Symbol	Open	High	Low	Close	Day
GOOG					

Names

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

What is possible as a map-only join now is different

The right table is the smaller table

Inner Joins
Left Outer Joins
Right Outer Joins
Full Outer Joins

Are possible as map-only joins

Conditions for Map-side Joins



All tables, except one, are small enough to be held in memory

Tables are bucketed on the join columns and Table1 Buckets = N * Table2 Buckets

Conditions for Map-side Joins



All tables, except one, are small enough to be held in memory

Tables are bucketed on the join columns and Table1 Buckets = N * Table2 Buckets

Names				
Symbol	Name			
GOOG	Google			
AAPL	Apple			
MSFT	Microsoft			



Join column = Symbol

Bucketed column = Symbol

Names

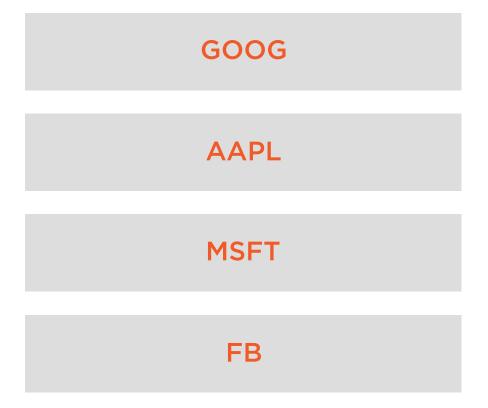
Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					

Trades

GOOG, AAPL

MSFT, FB



Names Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					

GOOG, AAPL

MSFT, FB



Names

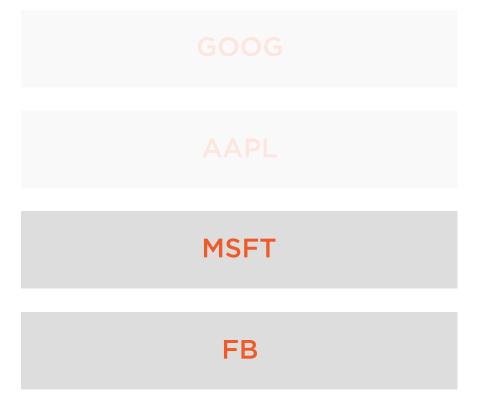
Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					

GOOG, AAPL

MSFT, FB



Bucket Map Join

Names

Trades

Symbol	Name
GOOG	Google
AAPL	Apple
MSFT	Microsoft

Symbol	Open	High	Low	Close	Day
GOOG					

GOOG, AAPL

MSFT, FB

GOOG

AAPL

MSFT

FB

set hive.optimize.bucketmapjoin = true;

Enable Bucket Map Join

This is not enabled by default in Hive

Bucket map joins are more efficient when tables are...

Bucketed on the join column

Sorted on the join column

Have the same number of buckets

Rows can be joined using merge-sort

Can only be used for equi-joins

```
set hive.input.format =
org.apache.hadoop.hive.ql.io.BucketizedHiveInputFormat;
set hive.optimize.bucketmapjoin = true;
set hive.optimize.bucketmapjoin.sortedmerge = true;
```

Enable Bucket Map Join with Sort-merge

Format of the file that is read in from disk into Hive

```
set hive.input.format =
org.apache.hadoop.hive.ql.io.BucketizedHiveInputFormat;
set hive.optimize.bucketmapjoin = true;
set hive.optimize.bucketmapjoin.sortedmerge = true;
```

Enable Bucket Map Join with Sort-merge

Format of the file that is read in from disk into Hive

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```

Enable Bucket Map Join with Sort-merge

Format of the file that is read in from disk into Hive

```
select /*+ mapjoin(trades) */ names.symbol, high, day
from names join trades on
(names.symbol = trades.symbol)
```

Explicitly Perform a Map-join

Read the table specified completely on the mapper node

Summary

A deep understanding of how joins work in Hive

Faster joins on large tables

Optimized semi-joins in place of IN/ EXISTS subqueries

Understood under what conditions joins are map-only and how to optimize them