

# Working with Advanced Data Transformations

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# Overview

**Understand the group by keyword and use aggregations on field values**

**Use joins to combine matching records from multiple relations**

**Use the union command to combine records together into one relation**

**Extract entities in bags into discrete records using the flatten command**

**Use real world data from the City of New York to perform analysis**

Demo

**Access and download the data for  
accident information for the City of New  
York**

# Grouping Records on the Same Key

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# Group By

ID	Product_ID	Quantity	Amount
o1	phone	1	199
o1	shoes	1	69
o2	book	2	22
o3	phone	1	149
o3	belt	2	19

**Tuple of fields**

# Group By

ID	Product_ID	Quantity	Amount
o1	phone	1	199
o1	shoes	1	69
o2	book	2	22
o3	phone	1	149
o3	belt	2	19

group orders by ID

# Group By

o1



o1	phone	1	199
----	-------	---	-----

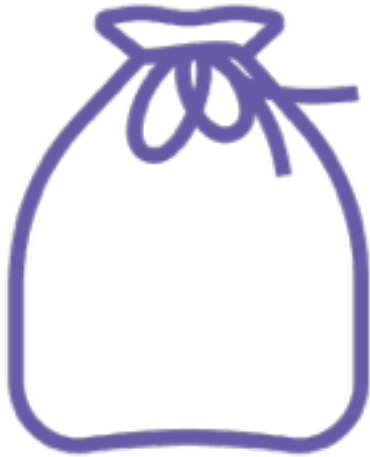
o1	shoes	1	69
----	-------	---	----

o2



o2	book	2	22
----	------	---	----

o3



o3	phone	1	149
----	-------	---	-----

o3	belt	2	19
----	------	---	----

## Group By

o1



o1	phone	1	199
o1	shoes	1	69

**All records with the same key are  
grouped into a bag**



Group By

o1



o1	phone	1	199
o1	shoes	1	69

group orders by ID creates a relation with 2 fields

key = field name “group”

value = bag with field name “orders”

# Demo



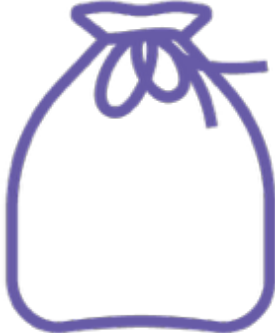
**Use the group by command on the collisions data in preparation to performing aggregation operations**

- group by reason for collisions across all boroughs
- group by collisions on a per borough basis

# Performing Aggregations on Grouped Records



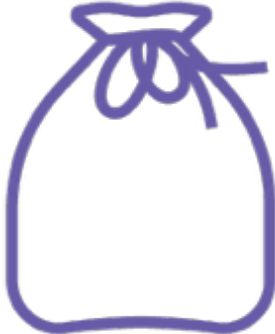
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# Aggregations on Groups

o1		o1	phone	1	199
		o1	shoes	1	69
o2		o2	book	2	22
o3		o3	phone	1	149
		o3	belt	2	19

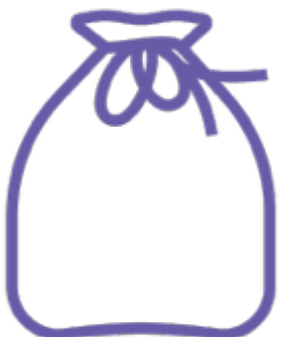
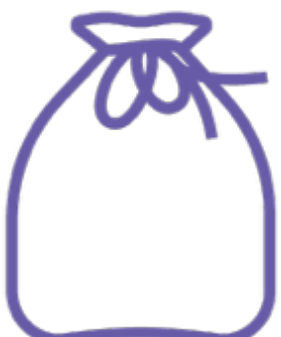
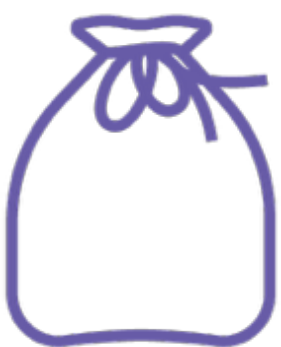
**Aggregations are UDFs which can be applied to field values from multiple records**

# Aggregations on Groups

o1		o1	phone	1	199
		o1	shoes	1	69
o2		o2	book	2	22
o3		o3	phone	1	149
		o3	belt	2	19



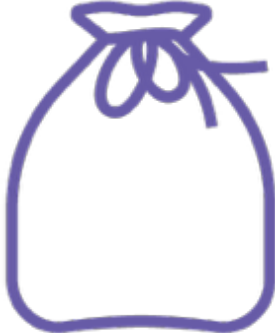
**COUNT()** the number of different products in each order

# Aggregations on Groups

o1		2
o2		1
o3		2

**COUNT()** the number of different  
products in each order

# Aggregations on Groups

o1		o1	phone	1	199
		o1	shoes	1	69
o2		o2	book	2	22
o3		o3	phone	1	149
		o3	belt	2	19

**SUM()** the total amount spent per order

# Aggregations on Groups

o1		268
o2		22
o3		168

**SUM()** the total amount spent per order



# Demo

**What kind of collision causes the most injuries in New York?**

- use the SUM() aggregation

**What boroughs have the most collisions?**

- use the COUNT() aggregation

# Join Operations in Pig

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# Joins

Name	Salary
Tom	1
John	1
Judy	150m



Name	Department
Judy	Google
Tom	GoogleX
John	Alphabet

# Joins

Name	Salary	Department
Tom	1	GoogleX
John	1	Alphabet
Judy	150m	Google

**Records from each relation matched on  
the join column**

# Joins

Name	Salary	Department
Tom	1	GoogleX
John	1	Alphabet
Judy	150m	Google

Pig provides support only for **equi-joins**

# Demo

**Perform join operations with 2 relations**

**Access individual fields from the joined relation using the :: operator**

# Types of Joins in Pig

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# Types of Joins

**Left Outer Join**

**Right Outer Join**

**Full Outer Join**

**Self Join**

**Cross Join**



# Types of Joins

**Left Outer Join**

Right Outer Join

Full Outer Join

Self Join

Cross Join

# Left Outer Join

Name	Salary
Tom	1
John	1
Judy	150m



Name	Department
Emily	Google
John	GoogleX
Tom	Alphabet

# Left Outer Join

Name	Salary
Tom	1
John	1
Judy	150m

**Every record on the left table will be present in the result**

- with a matching record
- padded with nulls

# Left Outer Join

Name	Salary	Department
Tom	1	Alphabet
John	1	GoogleX
Judy	150m	NULL

# Types of Joins

Left Outer Join

**Right Outer Join**

Full Outer Join

Self Join

Cross Join

# Right Outer Join

Name	Salary
Tom	1
John	1
Judy	150m



Name	Department
Emily	Google
John	GoogleX
Tom	Alphabet

# Right Outer Join

**Every record on the right table will be present in the result**

- with a matching record
- padded with nulls

Name	Department
Emily	Google
John	GoogleX
Tom	Alphabet

# Right Outer Join

Name	Salary	Department
Emily	NULL	Google
John	1	GoogleX
Tom	1	Alphabet



# Types of Joins

Left Outer Join

Right Outer Join

**Full Outer Join**

Self Join

Cross Join

# Full Outer Join

Name	Salary
Tom	1
John	1
Judy	150m



Name	Department
Emily	Google
John	GoogleX
Tom	Alphabet

# Full Outer Join

Name	Salary
Tom	1
John	1
Judy	150m



Name	Department
Emily	Google
John	GoogleX
Tom	Alphabet

**Records from both tables will be present in the result**

- with a matching record
- padded with nulls

# Full Outer Join

Name	Salary	Department
Emily	NULL	Google
John	1	GoogleX
Tom	1	Alphabet
Judy	150m	NULL

# Types of Joins

Left Outer Join

Right Outer Join

Full Outer Join

**Self Join**

Cross Join

# Self Join

Name	Salary
Tom	1
John	1
Judy	150m



Name	Salary
Tom	1
John	1
Judy	150m

# Self Join

Name	Salary	Salary
Tom	1	1
John	1	1
Judy	150m	150m

# Types of Joins

Left Outer Join

Right Outer Join

Full Outer Join

Self Join

Cross Join



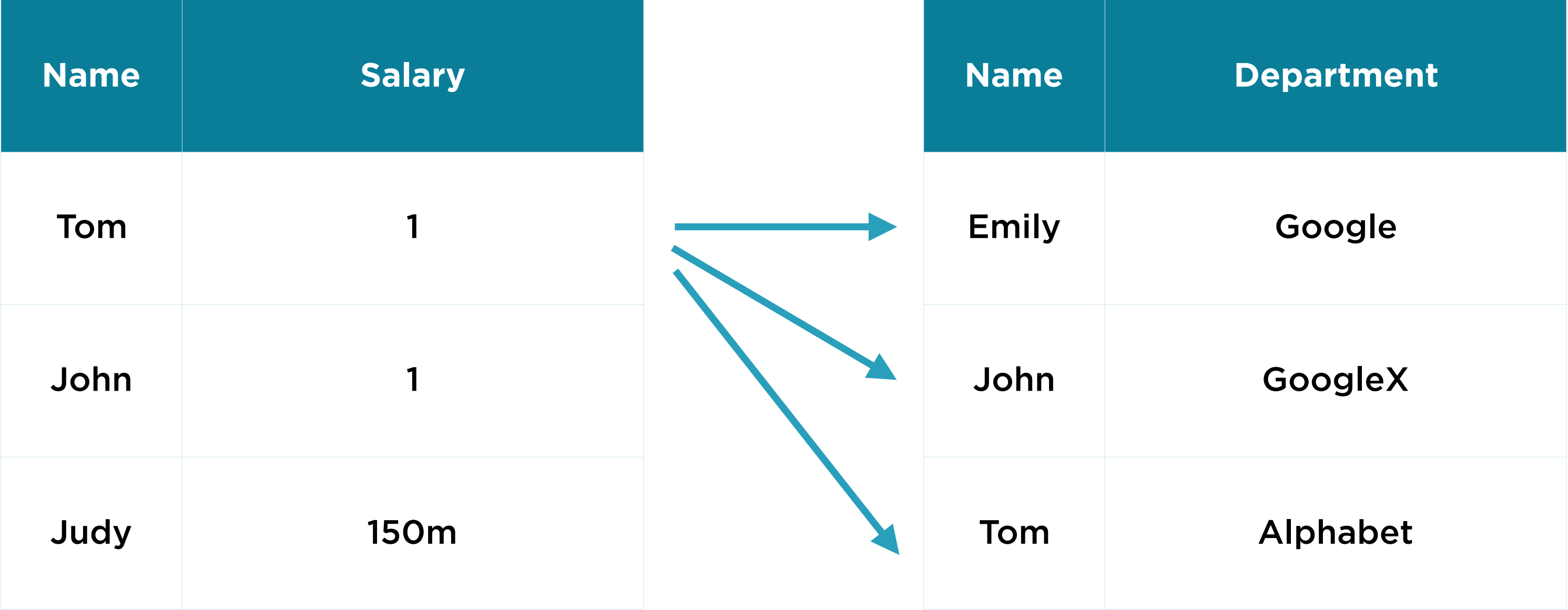
# Cross Join

Name	Salary
Tom	1
John	1
Judy	150m

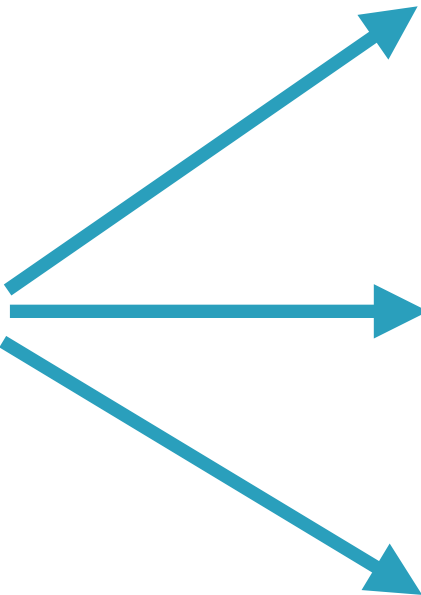


Name	Department
Emily	Google
John	GoogleX
Tom	Alphabet

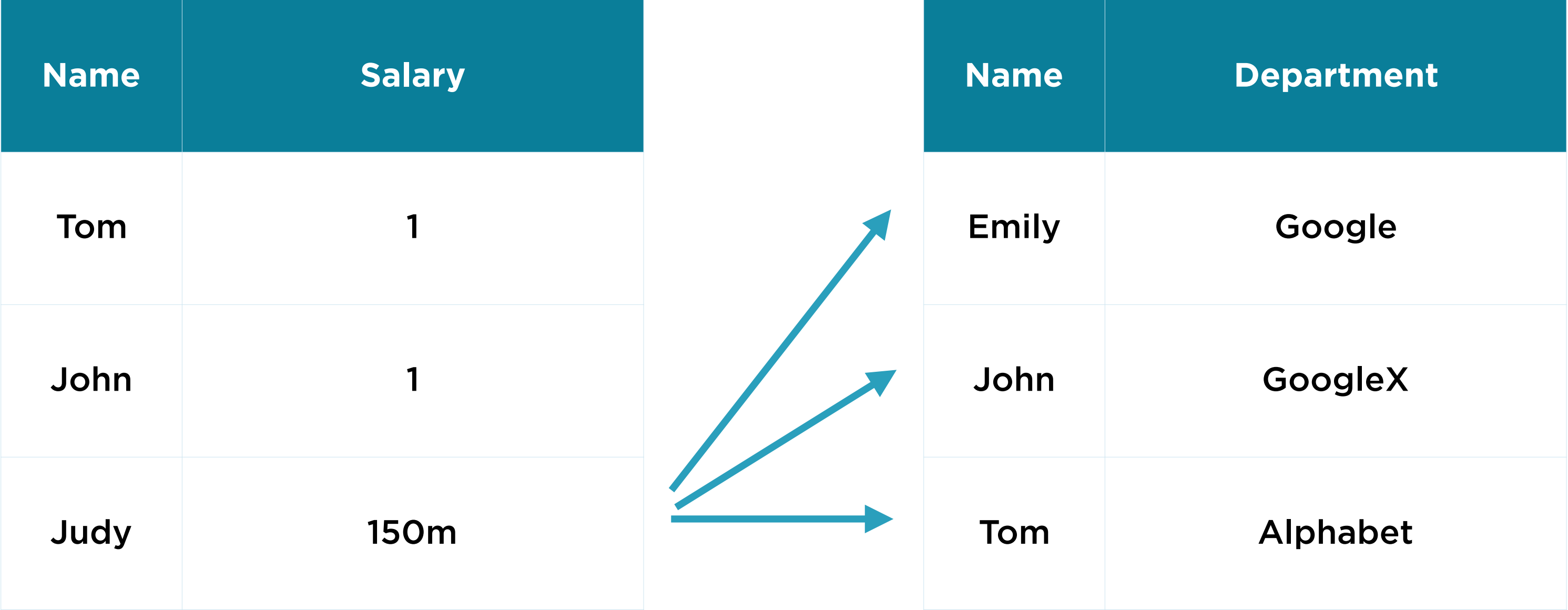
# Cross Join



# Cross Join

Name	Salary		Name	Department
Tom	1		Emily	Google
John	1		John	GoogleX
Judy	150m		Tom	Alphabet

# Cross Join



# Cross Join

Name	Salary	Name	Department
Tom	1	Emily	Google
John	1	John	GoogleX
Judy	150m	Tom	Alphabet
Tom	1	Emily	Google
John	1	John	GoogleX
Judy	150m	Tom	Alphabet
Tom	1	Emily	Google
John	1	John	GoogleX
Judy	150m	Tom	Alphabet

# Demo

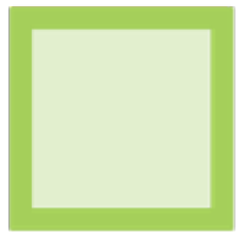
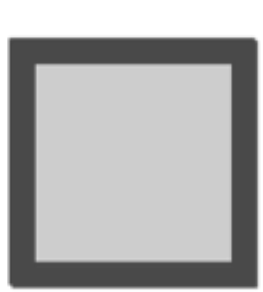
## **Implement join operations in Pig**

- left outer join
- self join
- cross join

# Unions in Pig

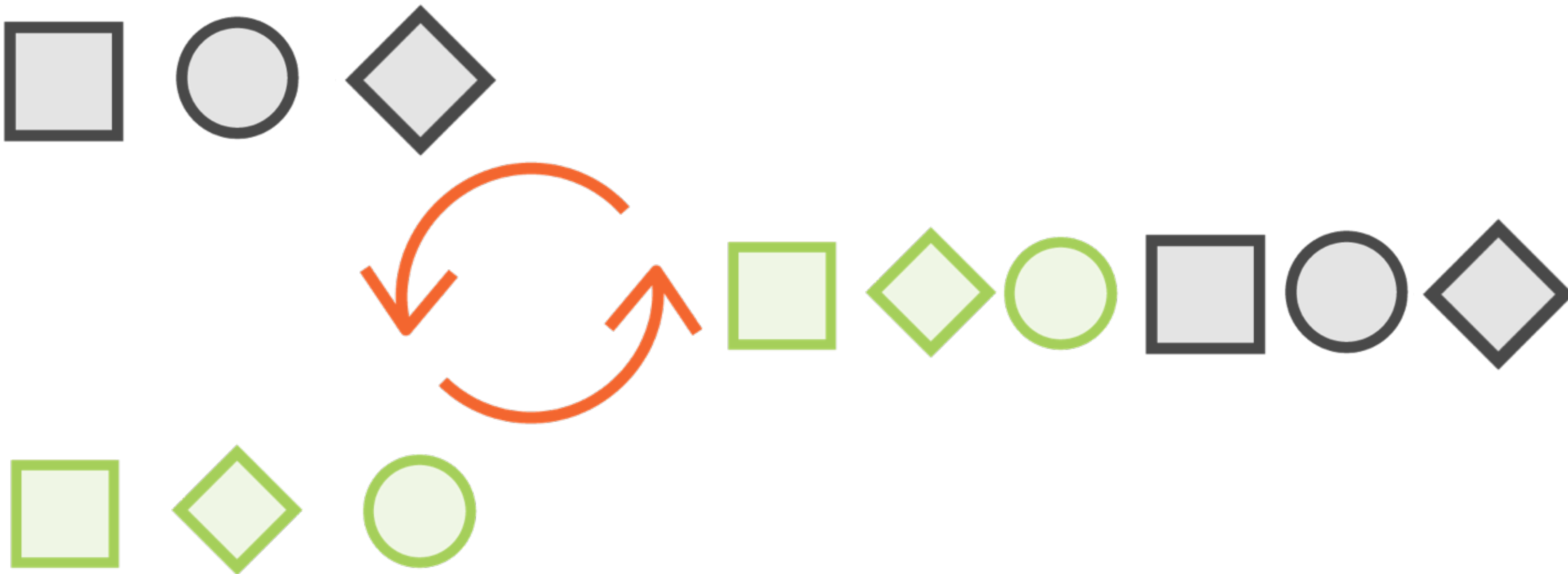
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Union





Union



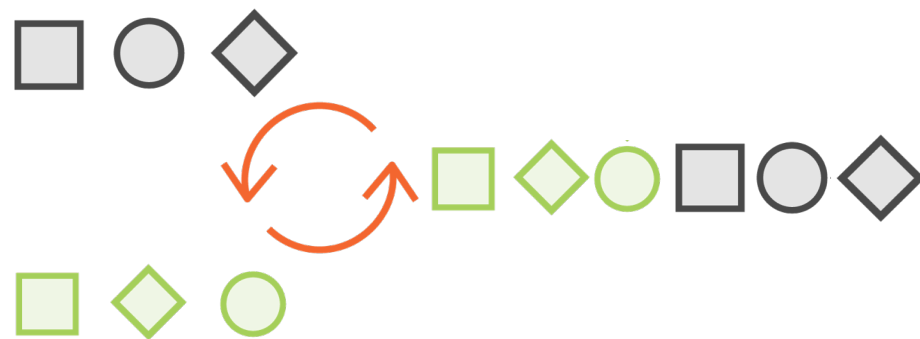
# Union

The relations involved in a union should have:

- the same number of fields
- compatible schema

**Does not preserve** the order of tuples

**Preserves** duplicates



# Demo

**Implement a union between 2 relations  
which have the same schema**

# Unions with Different Schemas

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R1 : (a1 : long, a2 : long)

R2 : (b1 : long, b2 : long, b3 : long)

R1 union R2 : null

---

## Union When Schema Is Mismatched

**R1: (a1: long, a2: long)**

R2: (b1: long, b2: long, b3: long)

R1 union R2: null

---

## Union When Schema Is Mismatched

R1 : (a1 : long, a2 : long)

**R2 : (b1 : long, b2 : long, b3 : long)**

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## Union When Schema Is Mismatched

R1 : (a1 : long, a2 : long)

R2 : (b1 : long, b2 : long, b3 : long)

**R1 union R2: null**

---

# Union When Schema Is Mismatched



R1 : (a1 : long, a2 : long)

R2 : (b1 : (x : int, y : int), b2 : long)

R1 union R2 : (a1 : bytearray, a2 : long)

---

Union When Schema Types Are Not the Same

R1 : (a1: long, a2: long)

R2 : (b1: (x: int, y: int), b2: long)

R1 union R2 : (a1: bytearray, a2: long)

---

# Union When Schema Types Are Not the Same

R1: (a1: long, **a2: long**)

R2: (b1: (x: int, y: int), **b2: long**)

R1 union R2: (a1: bytearray, **a2: long**)

---

## Union When Schema Types Are Not the Same

R1: (a1: long, a2: bytearray, a3: int)

R2: (b1: float, b2: chararray, b3: bytearray)

R1 union R2: (a1: float, a2: chararray, a3: int)

---

Compatible Types Will Be Cast to Higher Type

R1: (**a1: long**, a2: bytearray, a3: int)

R2: (**b1: float**, b2: chararray, b3: bytearray)

R1 union R2: (**a1: float**, a2: chararray, a3: int)

---

## Compatible Types Will Be Cast to Higher Type

**double > float > long > int > bytearray**

R1: (a1: long, **a2: bytearray**, a3: int)

R2: (b1: float, **b2: chararray**, b3: bytearray)

R1 union R2: (a1: float, **a2: chararray**, a3: int)

---

## Compatible Types Will Be Cast to Higher Type

**double > float > long > int > bytearray**

**tuple | bag | map | chararray > bytearray**

R1: (a1: long, a2: bytearray, a3: **int**)

R2: (b1: float, b2: chararray, b3: **bytearray**)

R1 union R2: (a1: float, a2: chararray, a3: **int**)

---

## Compatible Types Will Be Cast to Higher Type

**double > float > long > int > bytearray**

**tuple | bag | map | chararray > bytearray**

R1: (a1: long, a2: bytearray, a3: int)

R2: (b1: float, b2: chararray, b3: bytearray)

R1 union R2: (a1: float, a2: chararray, a3: int)

---

## Compatible Types Will Be Cast to Higher Type

**double > float > long > int > bytearray**

**tuple | bag | map | chararray > bytearray**



R1: (a1:(x:long, y:int), a2:{(n:float, m:chararray)})

R2: (b1:(g:chararray, h:float), b3:{(n:int, m:long)})

R1 union R2: (a1: (), a2: {()})

---

## Different Inner Types

**The union may result in an empty complex type**

R1: (a1:(x:long, y:int), a2:{(n:float, m:chararray)})

R2: (b1:(g:chararray, h:float), b3:{(n:int, m:long)})

R1 union R2: (a1: (), a2: {()})

---

## Different Inner Types

**The union may result in an empty complex type**

R1: (a1:(x:long, y:int), a2:{(n:float, m:chararray)})

R2: (b1:(g:chararray, h:float), b3:{(n:int, m:long)})

R1 union R2: (a1: (), a2: {})

---

## Different Inner Types

**The union may result in an empty complex type**

# Union Onschema for Schema Mismatches

---

```
R1 : (a1 : long, a2 : chararray)
R2 : (b1 : long, b2 : float, b3 : bytearray)
```

```
union onschema R1, R2
```

```
U : (a1 : long, a2 : chararray, b2 : float, b3 : bytearray)
```

---

# Union Onschema Combines Schemas

```
R1 : (a1 : long, a2 : chararray)
R2 : (b1 : long, b2 : float, b3 : bytearray)
```

```
union onschema R1, R2
```

```
U : (a1 : long, a2 : chararray, b2 : float, b3 : bytearray)
```

---

# Union Onschema Combines Schemas

```
R1 : (a1 : long, a2 : chararray)
R2 : (b1 : long, b2 : float, b3 : bytearray)
```

```
union onschema R1, R2
```

```
U : (a1 : long, a2 : chararray, b2 : float, b3 : bytearray)
```

---

## Union Onschema Combines Schemas

```
R1 : (a1 : long, a2 : chararray)
R2 : (b1 : long, b2 : float, b3 : bytearray)
```

```
union onschema R1, R2
```

```
U : (a1 : long, a2 : chararray, b2 : float, b3 : bytearray)
```

---

# Union Onschema Combines Schemas



# Demo

**Implement union on schema between 2 relations which have only a few columns with matching schema**

# The Flatten Function

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# Flatten

User_ID	Username	Products_Bought
u123	John	{(phone), (book), (shoes), (shirt)}
u876	Jill	{(speakers)}
u654	Nina	{(handbag), (book)}

**The flatten function is applied to a bag of tuples**

# Flatten

User_ID	Username	Products_Bought
u123		{(handbag), (book), (shoes), (shirt)}
u876		{(handbag), (book), (shoes), (shirt)}
u654	Nina	{(handbag), (book)}

**The products each user has bought is specified as a bag**

# Flatten

User_ID	Username	Products_Bought
u123	John	{(phone), (book), (shoes), (shirt)}
u876	Jill	{(speakers)}
u654	Nina	{(handbag), (book)}

**Flattening a bag makes entity in the bag a separate record**

# Flatten

User_ID	Username	Products
u123	John	phone
u123	John	book
u123	John	shoes
u123	John	shirt
u876	Jill	speakers
u654	Nina	handbag
u654	Nina	book

# Flatten

User_ID	Username	Products
u123	John	phone
u123	John	book
u123	John	shoes
u123	John	shirt
u876	Jill	speakers
u654	Nina	handbag
u654	Nina	book

# Flatten

User_ID	Username	Products
u123	John	phone
u123	John	book
u123	John	shoes
u123	John	shirt
u876	Jill	speakers
u654	Nina	handbag
u654	Nina	book



Flattening an empty bag  
results in **null**

# Demo

**Use the flatten function with a bag of tuples**

# Summary

**Used advanced Pig transformations such as:**

- group by and aggregations
- join operations
- union operations
- flatten command

**Analyzed real world data from the City of New York**