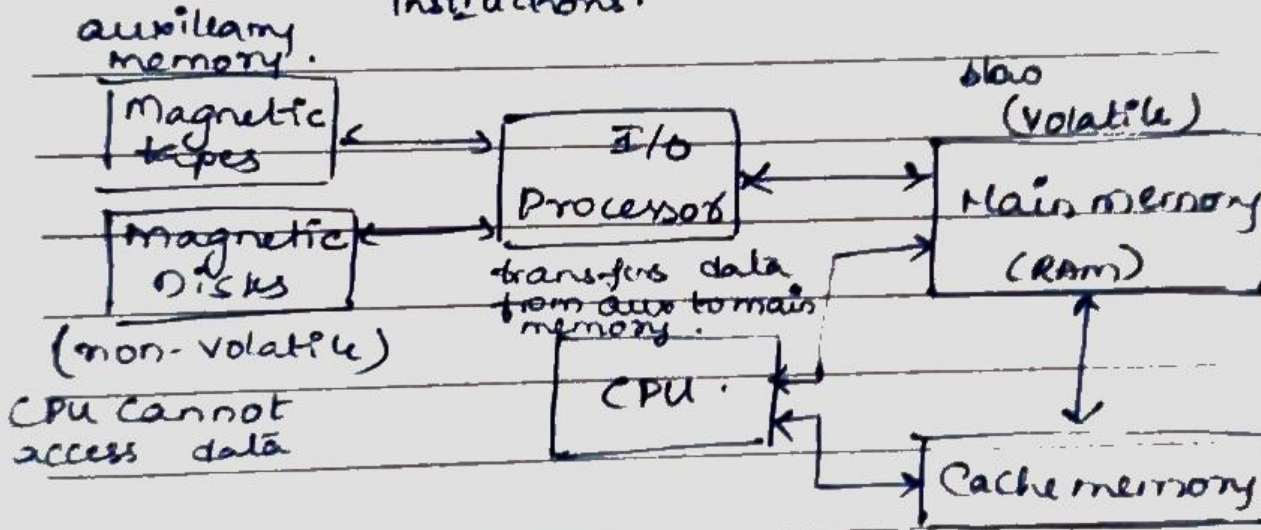


# Memory Hierarchy - Implement multi-processor multiprocessing

- \* Memory is essential part in any system.
- \* Useful to store information is computer

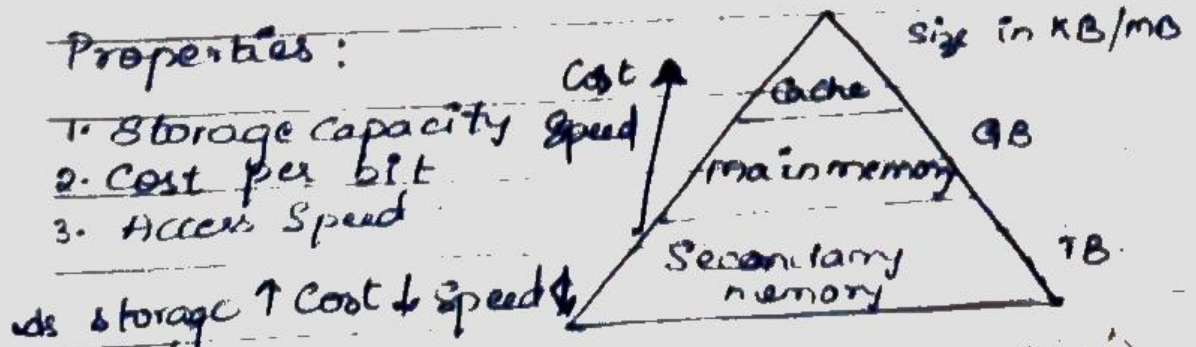
## Types of Memory

- \* Main memory - communicate with CPU.
  - \* Auxiliary " - (Secondary memory)
  - \* Cache memory
    - ↓
    - Store freq used instructions.
- Stores data permanent in computer.  
\* Backup purpose.



## Properties:

1. Storage capacity
2. Cost per bit
3. Access Speed



## Main memory

- \* Central storage unit in computer system
- \* relatively large and fast memory
- \* Store prog and data during comp oper
- \* Major technology involved is based on Semiconductor Integrated Circuits

### Integrated Circuit RAM chips

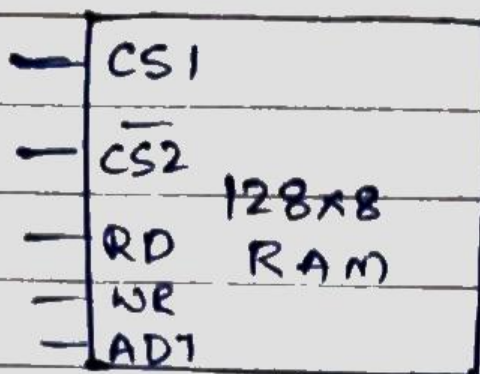
- static (expensive)
  - 1 bit to store binary info
  - extremely fast
- Dynamic
  - Stores info in form of electric charges applied to capacitors → provide with help of transistor
  - ⇒ Refresher
  - reduces power consumption
  - Stores large data

### Primary memory

- RAM (volatile)
- ROM (non-volatile)
  - ↓
  - Boot strap loader Pgm.



# RAM chip



RAM chip stores 128 words

and each word stores 8 bits.

$$128 = 2^7$$

8bit data bus

7 bits are required to store address

## 2 chip select

4 combination  $\rightarrow$  we get

depending on no. of RAM CS will be there or vice versa

eg; if we take 4 ram chips

$$4 = 2^2 \Rightarrow 2 \Rightarrow \text{chipselects}$$

$$8 = 2^3 = 3 \Rightarrow \text{chipselects}$$

CS1	CS2	RD	WR	memory-function	data bus
0	0	x	x	Inhibit	High impedance
0	1	x	x	"	"
1	0	x	x	"	"
1	0	0	1	write	input data to RAM
1	0	1	0	Read	output data from RAM
1	1	x	x	Inhibit	high impedance

# Associative Memory

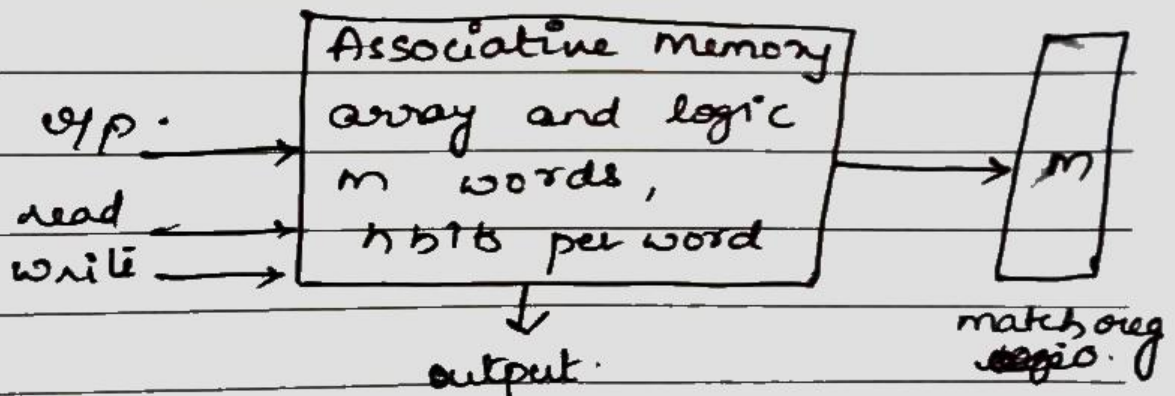
also called as CAM (Content addressable memory)

CPU executes instr based on address.  
but " " " " " Content

Y Am? → because time required to search an item or execute an item can be substantially reduced

(A) Augment Register → Contains the content  
↓

(K) Key Register → Comparison purpose  
↓



Note:

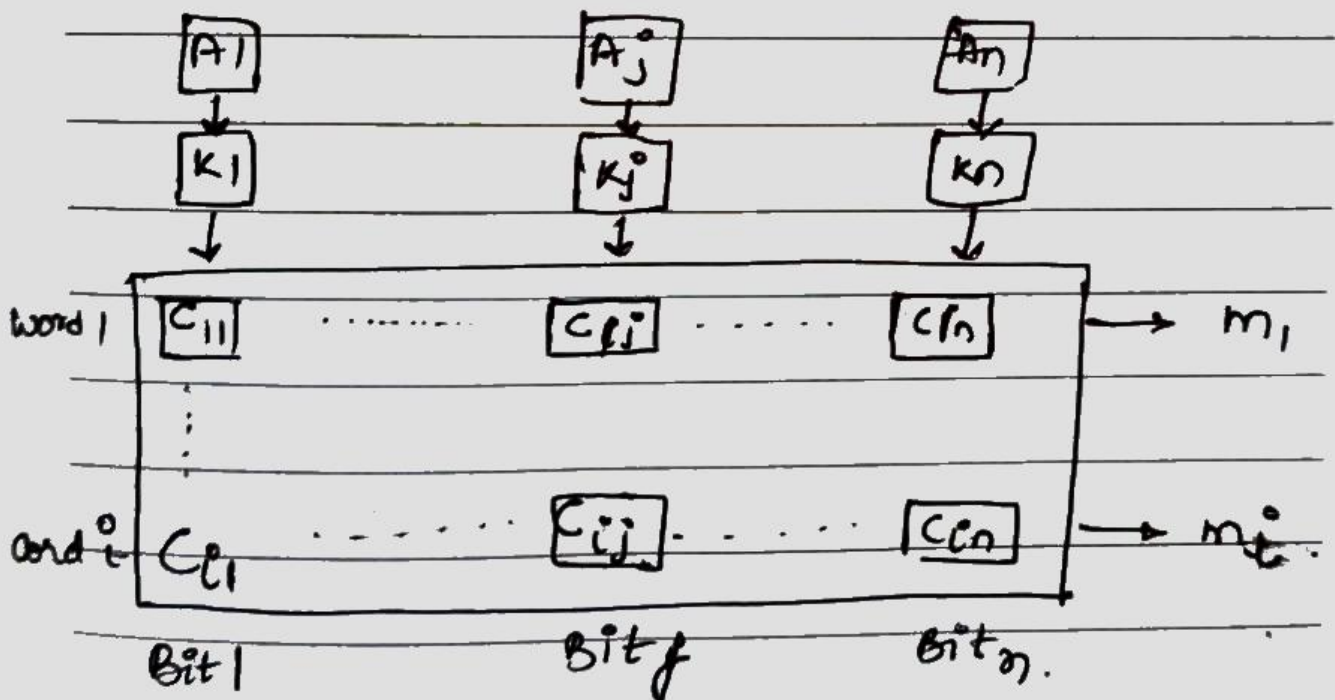
1. The content in augment reg is compared with data in associative memory only if the key register contains 1's.

→ If key registers contains 0's comparison is not done.

2. Size of match logic depends on the no. of words in associative memory.  
for each word - match i/o has to be maintained

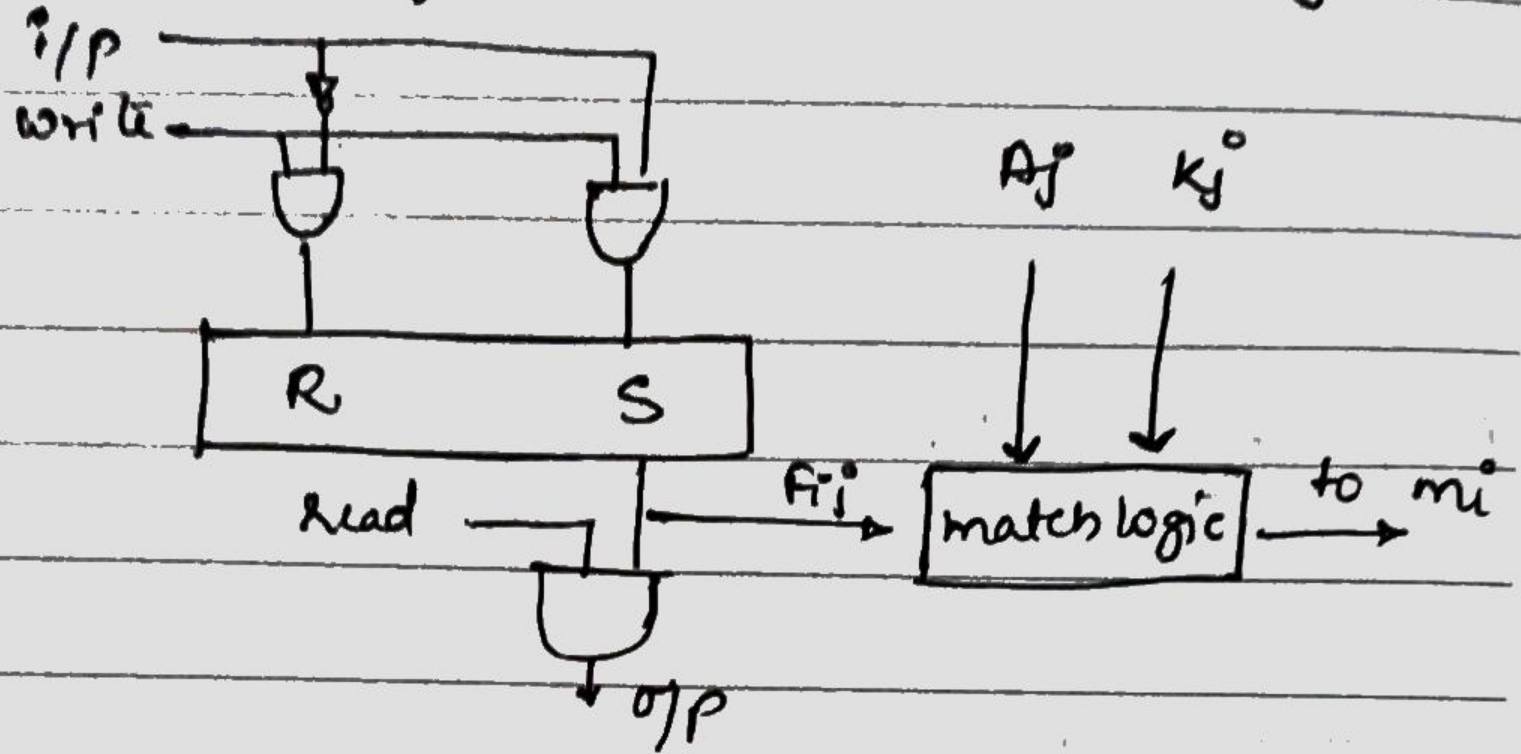
Eg; A      101      111100  
K      111      000000

word 1      100      111100      no match  
word 2      101      000001      match.





one cell of associative memory.



SR flt is used

S - Set → accepts through and gate.  
R - Reset -

# Cache Memory

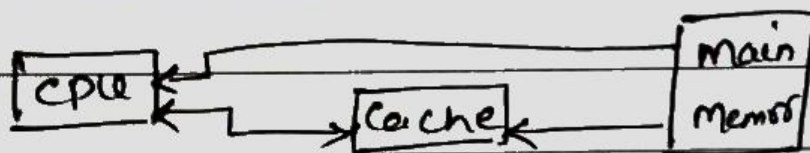
- extremely fast
- high accessibility
- expensive
- small in size

Property; - locality of reference

(freq executed instructions in the user prog)

∴ Avg memory access time is reduced.

placed in between CPU and main memory



performance can be analyzed by hit ratio

$$\text{hit ratio} = \frac{\text{no. of hits}}{\text{hits + miss}}$$

mapping → process through which the words in main memory are transferred to cache memory

3 types of mapping

- 1) Associative
- 2) Direct
- 3) Set Associative

I. associative mapping:  $\uparrow$  no of words

assume CPU size is  $32K \times 12$

$$2^5 \times 2^{10} \\ = 2^{15}$$

$\downarrow$   
each word contains  
no of bits.

15  $\rightarrow$  address generated by CPU.

Thus, Each address generated is stored in  
augment register.

Cache memory size  $512 \times 12$ .

CPU addr



Augment register

addr	data
01000	3450
00777	3710
23451	7734

address of augment reg is compared with  
associative mapping of cache and instructions  
are executed.

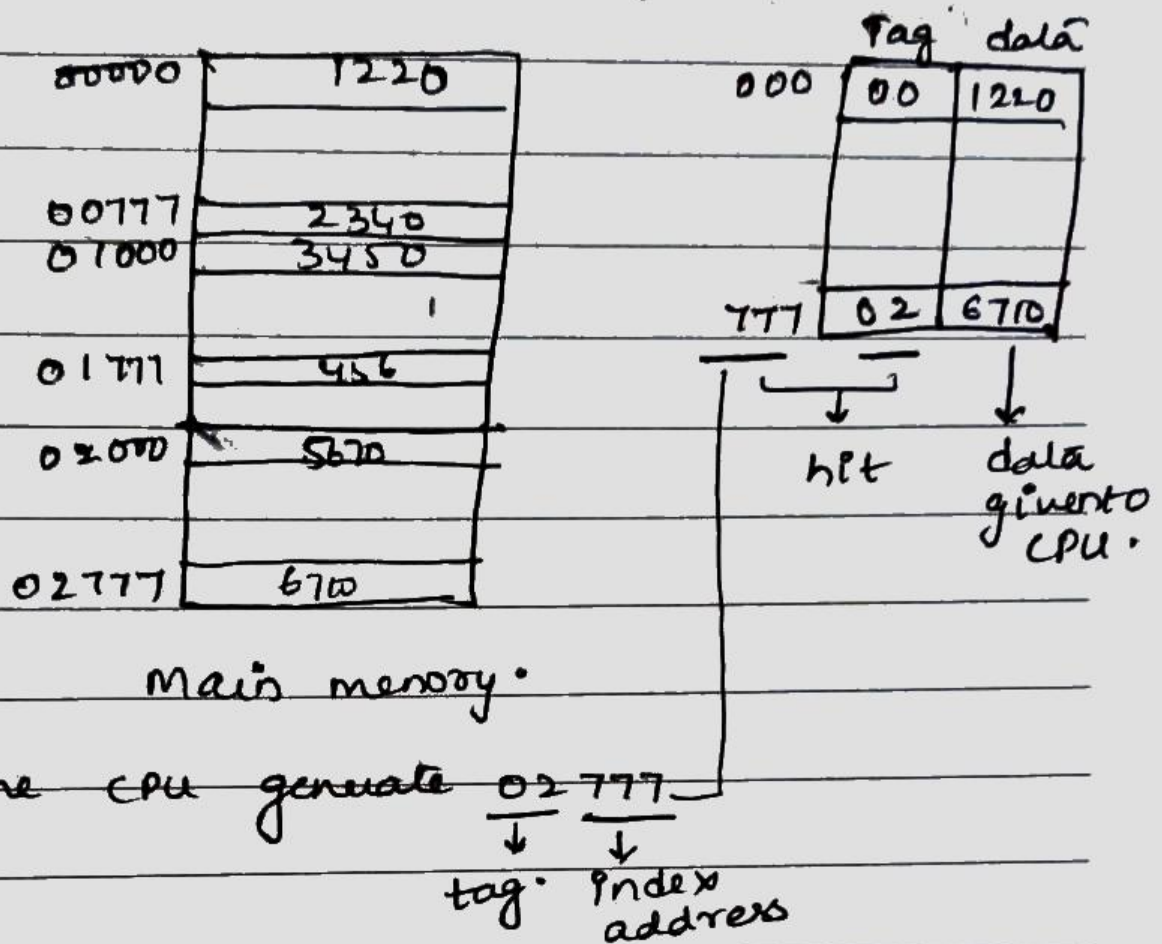


# I. Direct mapping

CPU generated address

6 bits 9 bits.

Tag	Index
-----	-------



Problem Several address can be generated by CPU with same index address but different tags.

Hence, hit ratio is drastically ↓

To overcome this set associative mapping

## III set associative mapping

In single index, multiple words are ~~taken~~ stored.

	tag	data	tag	data
000	00	1220	01	1234
777	02	6710	04	4523

# Auxillary memory (Secondary memory)

↓  
low cost

High storage capacity

low access speed

Eg; magnetic tapes, magnetic Disks, optical disks.  
(HDD) (CD)

\* depends on magnetic, electronics, optical & electro mechanical system

Access time = Seek time + Transfer time

↓  
position read/write  
head to location  
where data is residing

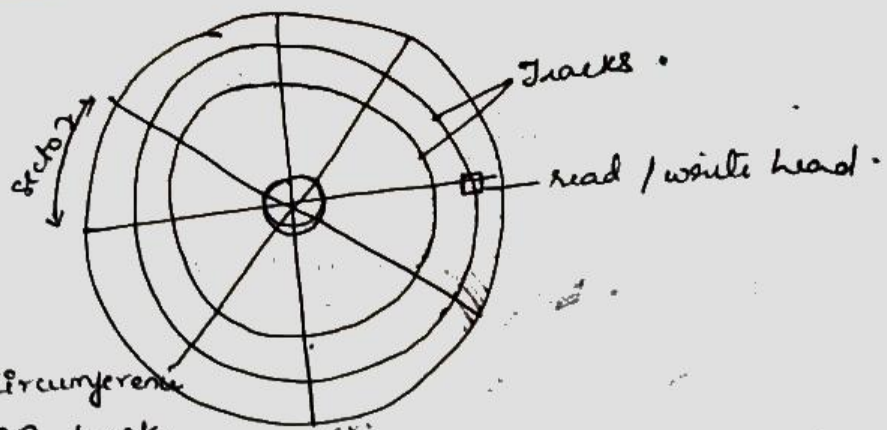
↓  
time to transfer data  
to the destination

transfer rate = no. of characters/words that device can transfer per second

## Magnetic disks

- made of circular plates (made with metal/plastic) and coated with magnetized material.
- many disks are stacked on one spindle.
- read/write heads available on each surface.
- all disks rotate together at high speed.
- Hence, bits are stored on magnetized surface.
- Concentric circles are called as tracks and these are divided into sectors.





- \* track near circumference is longer than tracks near centre of disk.
- \* If bits are recorded with equal density, some tracks will have more bits than others.
- \* Hence, variable recording density  $\Rightarrow$  preferred

### Read/Write Process:

- 1) min quantity of info which can be transferred is a sector
  - 2) single read/write head used for each surface.
- $\therefore$  when track address is given, read/write head is moved to that particular location (specified track).

OR

Separate read/write heads can be also maintained for each track on each surface.  $\therefore$  address bits can select particular track head.

### Magnetic Tape

- $\rightarrow$  strip of plastic coated with magnetic recording medium.
- $\rightarrow$  Bits are recorded as magnetic spots on tape along several tracks.

1 character - 7/9 bits are used.

multiple read/write heads are used to read data as a seq of characters.

- $\rightarrow$  Information recorded in blocks are called as records

• Each record has id (bit pattern).

fixed/variable

Hence addressed through the bit pattern.