



Introduction to Parallel and Distributed Processing

One-sided communication

Jaroslaw 'Jaric' Zola

http://www.jzola.org/





Problem with Point-to-Point Communication

- Point-to-point communication in MPI (send/recv) requires that both ranks participate
- Hence, the delay on one end slows down the rank on the other end
- This becomes problematic, especially when dealing with irregular communication patterns





One-sided Communication Idea (RMA)

- Each rank exposes part of its memory to other ranks
- Other ranks can directly read/write to this memory
- No (explicit) synchronization required!





One-sided Communication Advantage

- Certain communication patterns much easier to express (e.g. recall shifting)
- Modern networks (e.g. IB) support RDMA, and hence provide great performance



Exposing Memory

- Ranks collectively declare window: memory that will be accessible to all ranks in the communicator
- Memory must be properly aligned (helper API provided)

```
#include <mpi.h>
 3
     int main(int argc, char* argv[]) {
      MPI Init(&argc, &argv):
6
      int* buf:
      MPI Alloc mem(1024 * sizeof(int), MPI INFO NULL, &buf);
      MPI Win win;
10
      MPI Win create(buf, 1024 * sizeof(int), sizeof(int),
                      MPI INFO NULL. MPI COMM WORLD. &win):
11
12
13
      // ...
14
15
      MPI Win free(&win):
      MPI Free mem(buf);
16
17
18
       return MPI Finalize():
19
     } // main
```





Moving Data

- Primitives for putting, getting and modifying data in a window
- We talk about origin and target for data (instead of sender/receiver)
- Some important caveats
 - No guaranteed ordering of basic operations
 - Basic concurrent writes to the same location undefined



Getting Data

```
#include <iostream>
    #include <mpi.h>
3
 4
    int main(int argc, char* argv[]) {
      MPI Init(&argc, &argv);
      int* buf:
      MPI Alloc mem(1024 * sizeof(int), MPI INFO NULL, &buf):
9
10
      MPI Win win;
11
      MPI Win create(buf, 1024 * sizeof(int), sizeof(int),
12
                      MPI INFO NULL, MPI COMM WORLD, &win);
13
14
      int x[2]:
15
      MPI Win lock(MPI LOCK EXCLUSIVE, 0, 0, win);
      // get from buf + 1021 at target 0
16
      MPI_Get(x, 2, MPI_INT, 0, 1021, 2, MPI INT, win);
17
      MPI Win unlock(0, win):
18
19
20
      MPI Win free(&win);
21
      MPI Free mem(buf):
22
23
      return MPI Finalize();
24
    } // main
```



Putting Data

```
#include <iostream>
    #include <mpi.h>
3
    int main(int argc, char* argv[]) {
 4
      MPI Init(&argc, &argv);
      int* buf:
      MPI Alloc mem(1024 * sizeof(int), MPI INFO NULL, &buf):
9
10
      MPI Win win;
11
      MPI Win create(buf, 1024 * sizeof(int), sizeof(int),
12
                      MPI INFO NULL, MPI COMM WORLD, &win);
13
14
      int x[2]:
15
      MPI Win lock(MPI LOCK EXCLUSIVE, 0, 0, win);
      // put from x into buf + 1021 at target 0
16
17
      MPI Put(x, 2, MPI INT, 0, 1021, 2, MPI INT, win);
18
      MPI Win unlock(0, win):
19
20
      MPI Win free(&win);
21
      MPI Free mem(buf):
22
23
      return MPI Finalize();
24
    } // main
```



Other Goodies

```
#include <iostream>
2
    #include <mpi.h>
4
    int main(int argc, char* argv[]) {
5
      MPI Init(&argc, &argv);
6
 7
      int* buf:
8
      MPI_Alloc_mem(1024 * sizeof(int), MPI_INFO NULL, &buf);
9
10
      MPI Win win:
      MPI Win create(buf, 1024 * sizeof(int), sizeof(int),
11
                      MPI INFO NULL, MPI_COMM_WORLD, &win);
12
13
14
      MPI Win fence(0, win):
      MPI Accumulate(buf, 1024, MPI INT, 0, 0, 1024, MPI INT, MPI SUM, win);
15
16
      MPI Win fence(0, win):
17
18
      MPI Win free(&win);
19
      MPI Free mem(buf);
20
21
      return MPI Finalize();
22
    } // main
```



For Fun

• Implement a simple distributed hash table using one-sided communication.