



Introduction to Parallel and Distributed Processing

Shared Memory

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Suggested Reading

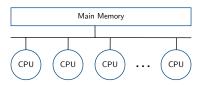
- Cilk+ reference guide https://software.intel.com/en-us/node/522579
- General resources https://www.cilkplus.org/





Shared Memory Architecture

- A type of parallel architecture where multiple processors share main memory (e.g. modern multi-core processors)
- Software perspective: threads share address space







Why Shared Memory

- Currently the most accessible type of parallel architecture
- Always at hand, why not use it
- Perceived easy parallelism





Why Shared Memory

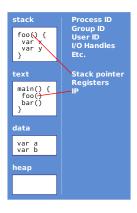


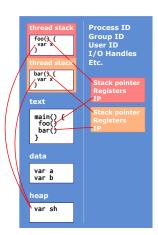




Programming Shared Memory

 Create multiple threads that communicate by reading/writing shared variables









Programming Shared Memory

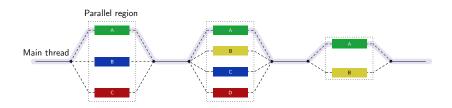
- Working with threads is too low-level for practical purposes scheduling, synchronization
- At the high level two types of parallelism:
 - Task parallelism (MIMD): pool of threads executes tasks, each task is a set of instructions on some data
 - Data parallelism: pool of threads acts on different parts of data with the same set of instructions
- Ideally, we want to describe our tasks and data transformations, without worrying about low-level realization





Fork-Join Pattern

- General strategy to express and execute task parallelism
- Main model for OpenMP, Cilk+, Intel TBB, etc.

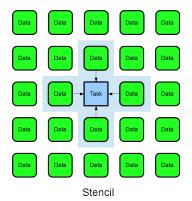






Stencil Pattern

 Basic data parallel pattern: data structured (e.g. in 1D, 2D), update to data element depends on its neighborhood





Cilk/Cilk+ API

- Provide a minimal set of C/C++ language extensions to enable shared memory parallelism (task and data)
- Focus on a parallel algorithms not implementation
- Efficient realization left to a compiler and run-time
- Comes with a powerful work-stealing scheduler and lock-free hyperobjects
- Deterministic Cilk+ code has serial semantics



Cilk+ Support

- Open source version in GCC 4.8 to 6.9, Clang/LLVM some support
- Commercial release by Intel (Parallel Studio XE, etc.)
- Using in GCC:

```
ı g++ −std=c++11 −fcilkplus −03 ...
```



Spawn and Sync

- The Cilk+ standard defines three new keywords: cilk_spawn, cilk_sync and cilk_for
- These keywords are all what is needed to fork and join tasks

```
int fib(int n) {
   if (n < 2) return n;

// create child task
   int a = cilk_spawn fib(n - 1);
   int b = fib(n - 2);

// wait for all child tasks in function cilk_sync;

return a + b;
// implicit cilk_sync</pre>
```



Spawn and Sync

- Arguments passed to "spawned" function are evaluated before invocation
- Pass-by-references and pass-by-pointer arguments must be valid at least till cilk_sync
- Implicit cilk_sync follows after destructors
- It is a programmer job to avoid race conditions more to folow





Simple Parallel Postorder

```
struct node {
        node* left;
        node* right;
   };
4
5
   node* build tree() { ... }
   int release tree(node* p) { ... }
   void process(const node* p) { ... }
9
   void visit(const node* p) {
10
        if (p->left != 0) visit(p->left);
11
        if (p->right != 0) visit(p->right);
12
        process(p);
13
14
15
   int main(int argc, char* argv[]) {
16
        node* tree = build tree();
17
        visit(tree);
18
        return release tree(tree);
19
20
```



Simple Parallel Postorder

```
struct node {
        node* left;
2
        node* right:
   };
4
5
6
   node* build tree() { ... }
7
   int release tree(node* p) { ... }
   void process(const node* p) { ... }
9
   void visit(const node* p) {
10
        if (p->left != 0) cilk spawn visit(p->left);
11
        if (p->right != 0) visit(p->right):
12
        cilk sync;
13
        process(p);
14
15
16
   int main(int argc, char* argv[]) {
17
        node* tree = build tree();
18
        visit(tree);
19
20
        return release tree(tree);
21
```





[f1,l1) and [f2,l2) are sorted ranges





• [f1,l1) and [f2,l2) are sorted ranges



```
template <typename InIter, typename OutIter>
   void merge dc(InIter f1, InIter l1, InIter f2, InIter l2,
2
                  OutIter res) {
3
     if ((l1 - f1) + (l2 - f2) < 5) {
4
        return merge(f1, l1, f2, l2, res);
5
6
7
     InIter mid1, mid2;
8
     if ((l2 - f2) < (l1 - f1)) {
10
       mid1 = f1 + ((l1 - f1) >> 1);
11
       mid2 = std::lower bound(f2, l2, *mid1);
12
13
     else {
14
       mid2 = f2 + ((l2 - f2) >> 1);
15
       mid1 = std::lower bound(f1, l1, *mid2);
16
17
18
     merge dc(f1, mid1, f2, mid2, res);
19
20
     merge dc(mid1, l1, mid2, l2, res + (mid1 - f1) + (mid2 - f2));
   } // merge dc
```



```
template <typename InIter, typename OutIter>
   void merge dc(InIter f1, InIter l1, InIter f2, InIter l2,
2
                  OutIter res) {
3
     if ((l1 - f1) + (l2 - f2) < MERGE LIMIT) {
4
        return merge(f1, l1, f2, l2, res);
5
6
7
     InIter mid1, mid2;
8
     if ((l2 - f2) < (l1 - f1)) {
10
       mid1 = f1 + ((l1 - f1) >> 1);
11
       mid2 = std::lower bound(f2, l2, *mid1);
12
13
     else {
14
       mid2 = f2 + ((l2 - f2) >> 1);
15
       mid1 = std::lower bound(f1, l1, *mid2);
16
17
18
     cilk spawn merge dc(f1, mid1, f2, mid2, res);
19
     merge dc(mid1, l1, mid2, l2, res + (mid1 - f1) + (mid2 - f2));
20
   } // merge dc
```



```
template <typename InIter, typename OutIter>
   void merge sr(InIter f1, InIter l1, InIter f2, InIter l2,
2
                  OutIter res) {
3
     while ((l1 - f1) + (l2 - f2) >= MERGE LIMIT) {
4
       InIter mid1, mid2;
       if ((l2 - f2) < (l1 - f1)) {
6
         mid1 = f1 + ((l1 - f1) >> 1):
7
         mid2 = std::lower bound(f2, l2, *mid1);
8
9
       else {
10
         mid2 = f2 + ((l2 - f2) >> 1);
11
         mid1 = std::lower bound(f1, l1, *mid2);
12
13
       cilk spawn merge_sr(f1, mid1, f2, mid2, res);
14
        res += (mid1 - f1) + (mid2 - f2);
15
       f1 = mid1;
16
       f2 = mid2:
17
18
     merge(f1, l1, f2, l2, res);
19
20
   } // merge sr
```



Efficiency Analysis

• If n is the total number of elements to merge:

$$T_1 = 2T_1(\frac{n}{2}) + O(\log(n)) = \Theta(n)$$

$$T_{\infty} = T_{\infty}(\frac{n}{2}) + O(\log(n)) = \Theta(\log^2(n))$$

$$S_{\infty} = O\left(\frac{n}{\log^2 n}\right)$$



For Fun

Build parallel sorting routine starting with parallel merge