

# Introduction to Parallel and Distributed Processing

## CUDA Reduce and Prefix

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# Classic Problem

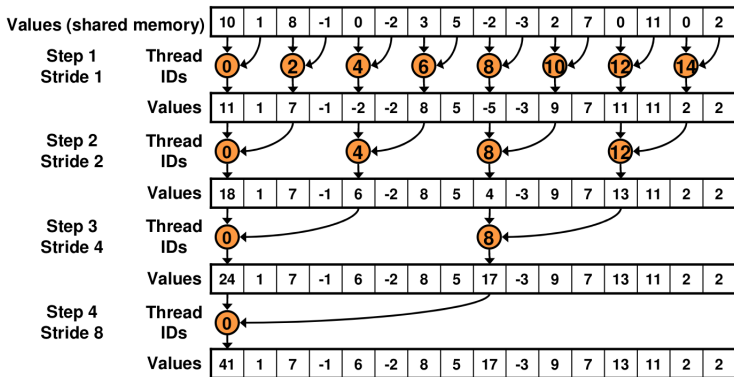
- Given vector  $X = (x_1, \dots, x_n)$  compute  $\sum_{i=1}^n x_i$
- Smells like a simple 1D problem...

# Direct Approach

```
1  __global__ void reduce0(int* gin, int* gout) {
2      extern __shared__ int sdata[];
3
4      int tid = threadIdx.x;
5      int i = blockIdx.x * blockDim.x + tid;
6
7      sdata[tid] = gin[i];
8      __syncthreads();
9
10     for (int s = 1; s < blockDim.x; s *= 2) {
11         if (tid % (2 * s) == 0) sdata[tid] += sdata[tid + s];
12         __syncthreads();
13     }
14
15     if (tid == 0) gout[blockIdx.x] = sdata[0];
16 } // reduce0
```

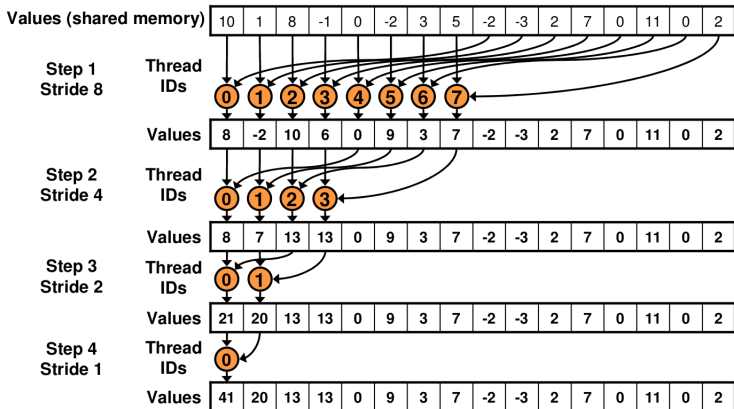
# Direct Approach – Problem

Divergent branching:



# Better Approach

Sequential addressing:



# Sequential Addressing

```
1  __global__ void reduce1(int* gin, int* gout) {
2      extern __shared__ int sdata[];
3
4      int tid = threadIdx.x;
5      int i = blockIdx.x * blockDim.x + tid;
6
7      sdata[tid] = gin[i];
8      __syncthreads();
9
10     for (int s = blockDim.x >> 1; s > 0; s >>= 1) {
11         if (tid < s) sdata[tid] += sdata[tid + s];
12         __syncthreads();
13     }
14
15     if (tid == 0) gout[blockIdx.x] = sdata[0];
16 } // reduce1
```

# Even Better Sequential Addressing

```
1  // we halve the number of blocks
2  __global__ void reduce2(int* gin, int* gout) {
3      __extern __shared__ int sdata[];
4
5      int tid = threadIdx.x;
6      int i = blockIdx.x * (blockDim.x * 2) + threadIdx.x;
7
8      sdata[tid] = gin[i] + gin[i + blockDim.x];
9      __syncthreads();
10
11     for (int s = blockDim.x >> 1; s > 0; s >>= 1) {
12         if (tid < s) sdata[tid] += sdata[tid + s];
13         __syncthreads();
14     }
15
16     if (tid == 0) gout[blockIdx.x] = sdata[0];
17 } // reduce2
```

# What About Parallel Prefix?

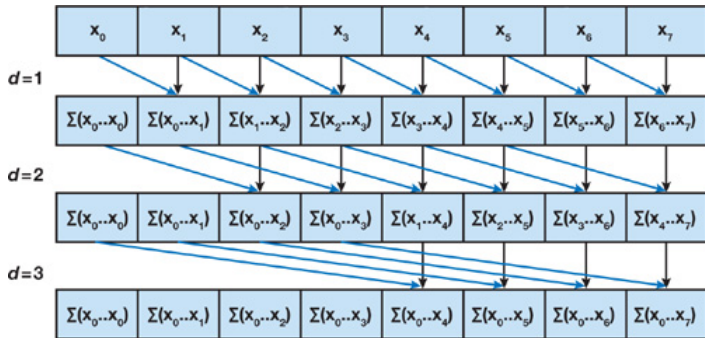
- Again, 1D problem PREFIX\_SUM

**Input:**  $[x_0, x_1, \dots, x_{n-1}]$

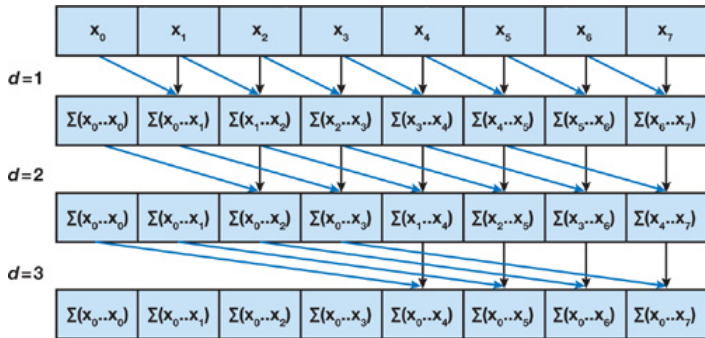
```
1: for  $d = 1 \dots \log(n)$  do  
2:   for  $j = 1 \dots n$  pardo  
3:     if  $j \geq 2^d$  then  
4:        $x[j] = x[j - 2^{d-1}] + x[j]$ 
```



# What About Parallel Prefix?



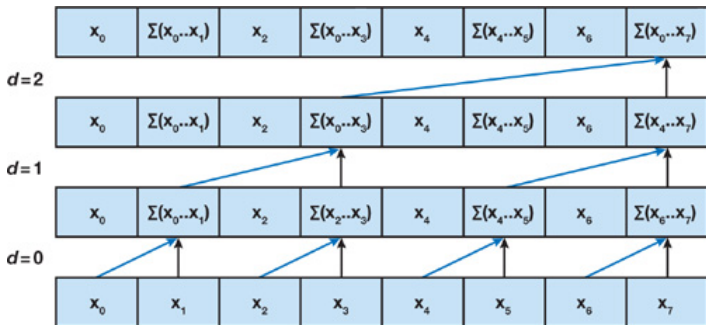
# What About Parallel Prefix?



Problem: total work  $O(n \log(n))$

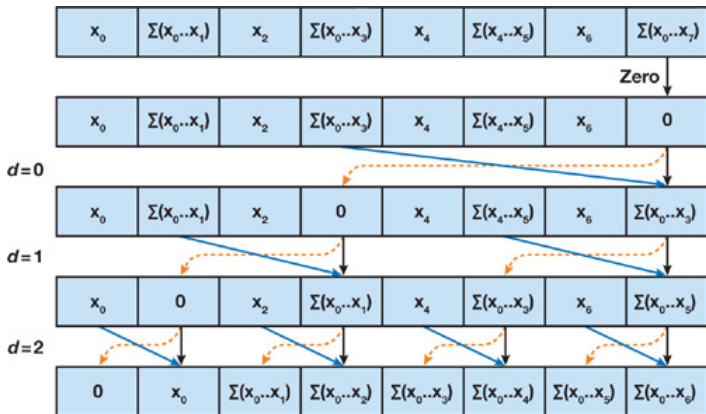
# Going Work Efficient

- We resolve to our old friend :-)
- Decompose problem into two phases: up-sweep (reduction) and down-sweep (update)

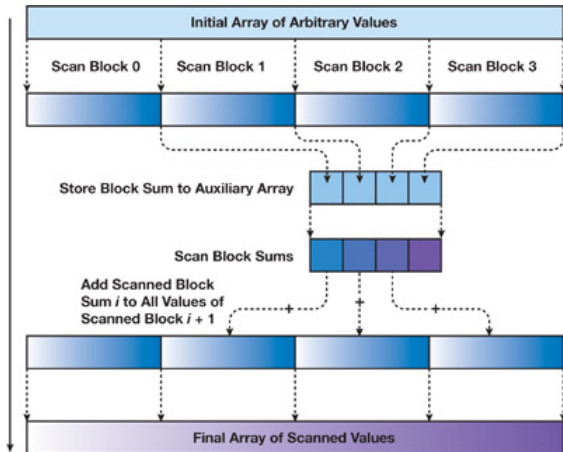


# Going Work Efficient

- We resolve to our old friend :-)
- Decompose problem into two phases: up-sweep (reduction) and down-sweep (update)



# Working on Large Arrays



# For Fun

- Modify the prefix code to work with large arrays.