

Java Multithreading for Senior Engineering Interviews / ... / LongAdder

## LongAdder

Learn why LongAdder can be a better choice for computing statistics in a multithreaded environment than AtomicLong.

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## Overview

Consider a scenario where you are hosting an API service and need to maintain a running count of all the invocations of a particular API endpoint. You could use Atomictong for the occasion as a counter but if you expect several concurrent requests hitting your endpoint, then multiple threads will attempt to increment the Atomictong instance. In such a scenario where the intention is to compute statistics a better approach is to use the LongAdder class rather than the Atomictong.

AtomicLong internally uses the compare and set instruction to perform the increment operation. In case of high contention, that is multiple threads attempting to increment an instance, a single thread wins the race to perform the increment operation while the rest have to retry until all of them succeed one by one. AtomicLong doesn't use locks and threads don't suspend accessing it. We discuss non-blocking synchronization to implement thread-safe algorithms and cons of locking in the Atomic Classes section.

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LongAdder overcomes high contention by keeping an array of counts, which can be thought of as variants of AtomicLong, and each one of them can be incremented atomically and independently of the other. The contention is spread out across the array resulting in increased throughput. Though the increased throughput comes at the cost of using more space. When the final value is requested, the sum of the array is calculated and returned by invoking the sum() method. Note that sum() isn't an atomic snapshot and any writes by threads while sum() is executing may not be reflected in the result. An accurate result is returned when concurrent updates don't occur.

The counts are of a non-public type cell which is a variant of AtomicLong. The array is referred to as the *table of Cells* and it grows in size as contention increases but its maximum size is capped by the number of CPUs.

## **Example**

The widget below runs a crude performance test between LongAdder and AtomicLong. We create ten threads that increment an instance of either class a million times. The sum at the end should come to be ten million. From the test set-up we can see that the instance of each class is highly contended for. From the test output we can see that LongAdder performs much better than AtomicLong.



