Applications of Spatial Coupling & Sparse Graph Codes for Sparse Recovery

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Outline

An (ℓ, r) LDPC Code

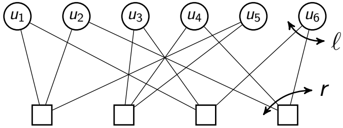
Parity-Check Matrix

$$H = \begin{pmatrix} 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 \end{pmatrix}$$

$$\ell = 2 \qquad r = 3$$

LDPC Code $\mathcal{C} = \{x : H \odot x = 0\}$

Tanner Graph



Compressed Representation

Belief Propagation Decoder & Threshold

Belief Propagation (BP)

- Popular choice
- Low-complexity
- Threshold: $\mathtt{h}^{\mathrm{BP}}$

Maximum a Posteriori (MAP)

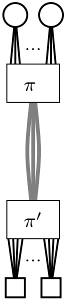
- Optimal Decoder
- Not Realizable
- Threshold: $\mathtt{h}^{\mathrm{MAP}}$

$$\mathtt{h}^{\mathrm{BP}} < \mathtt{h}^{\mathrm{MAP}}$$

Threshold Comparison

LDPC (ℓ, r)	Shannon $\mathtt{h}^{\mathrm{Sh}}$	AWGN $\mathtt{h}^{\mathrm{BP}}$	AWGN $\mathtt{h}^{\mathrm{MAP}}$	BSC $\mathtt{h}^{\mathrm{BP}}$	BSC $\mathtt{h}^{\mathrm{MAP}}$
$(3, 6)$	0.5000	0.4293	0.4794	0.4160	0.4681
$(4, 6)$	0.6667	0.5211	0.6645	0.5203	0.6633
$(5, 6)$	0.8333	0.5731	0.8333	0.5773	0.8333

(ℓ, r, N, w) Spatially-Coupled Ensemble



- An LDPC code of left-degree $\ell = 3$ and right-degree $r = 4$

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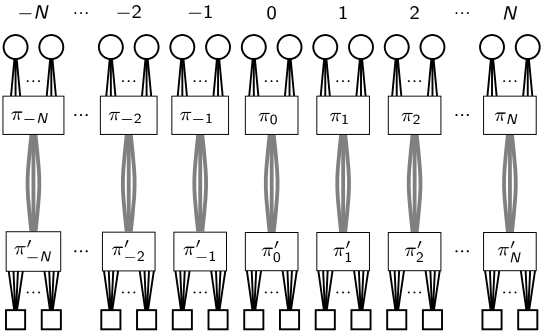
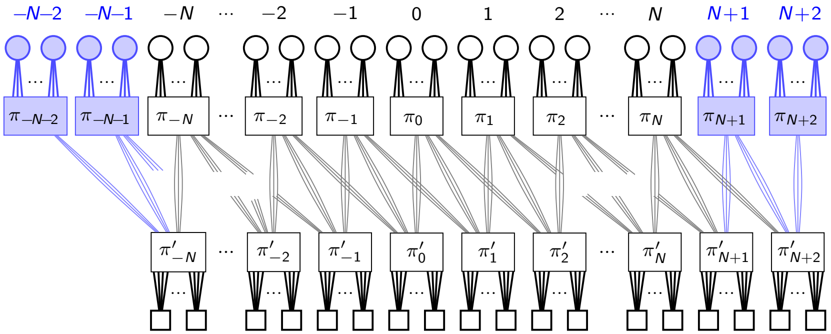
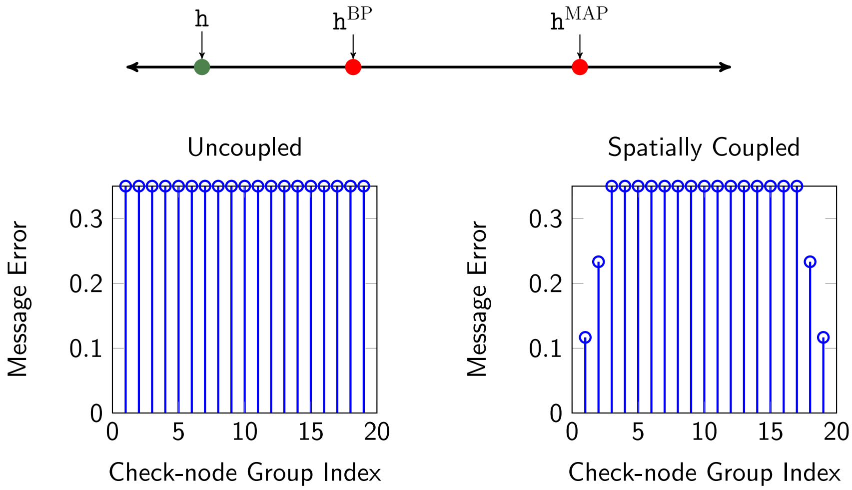


(ℓ, r, N, w) Spatially-Coupled Ensemble



- Shown for $\ell = 3$, $r = 4$, and $w = 3$
- Check-nodes at Section $\{i\}$ are connected to variable-nodes in Sections $\{i-(w-1), \ldots, i\}$
- Shown to have near optimal BP thresholds

Threshold Saturation via Spatial Coupling (1)



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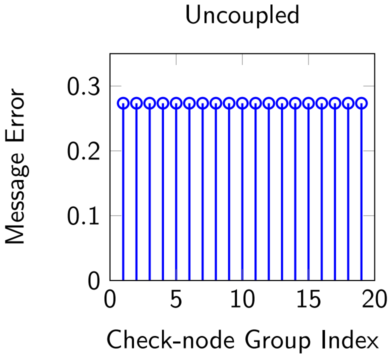

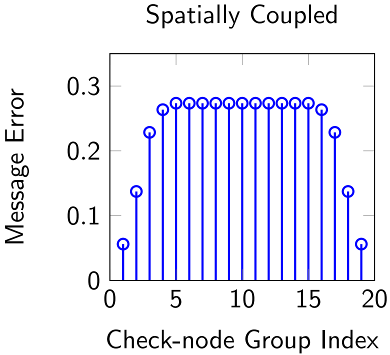

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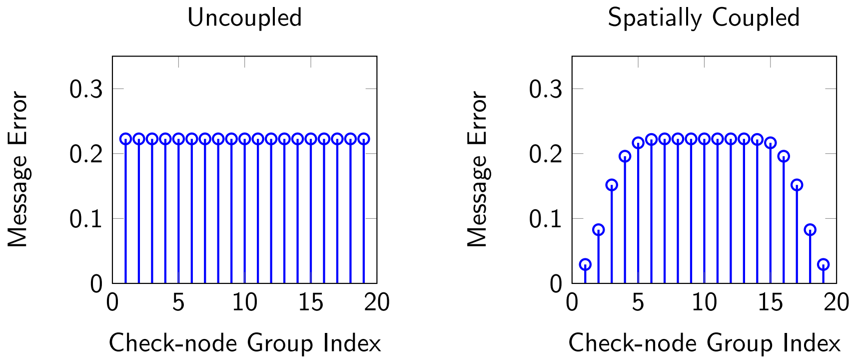

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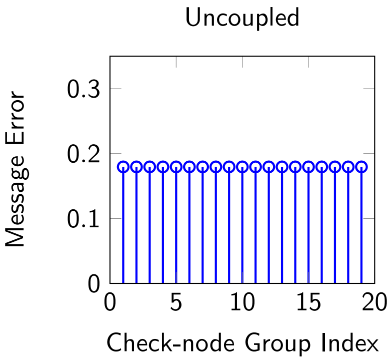

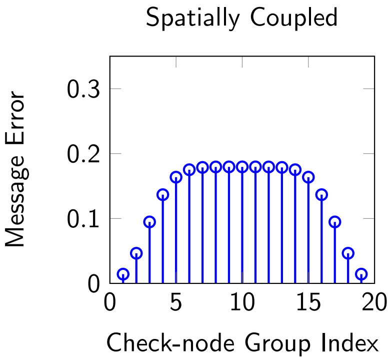

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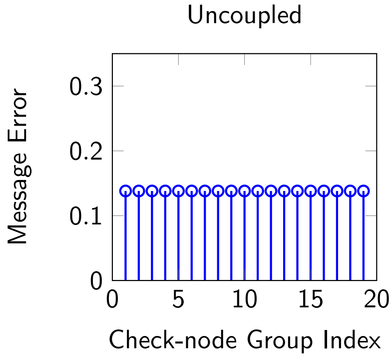

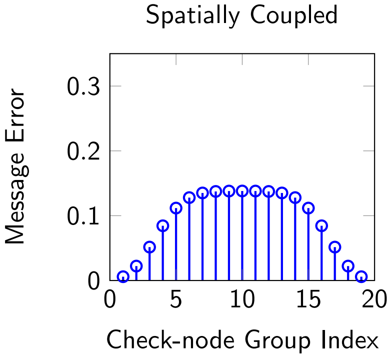

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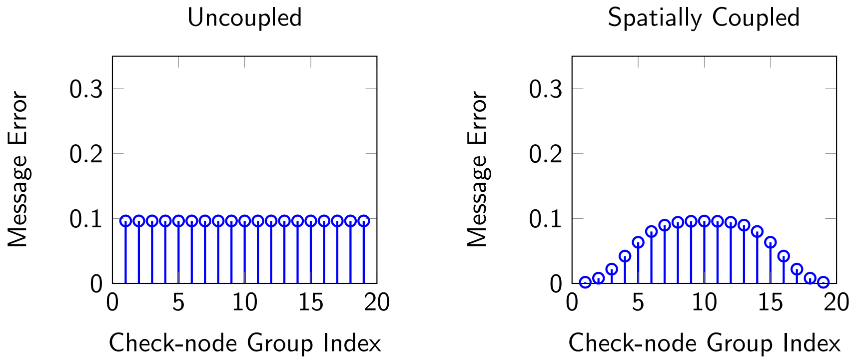

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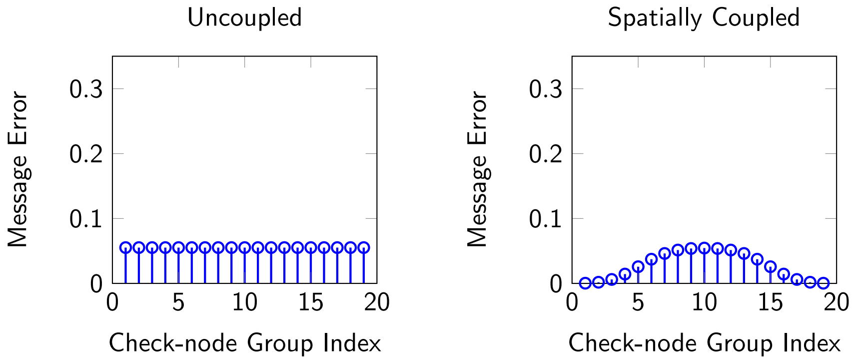

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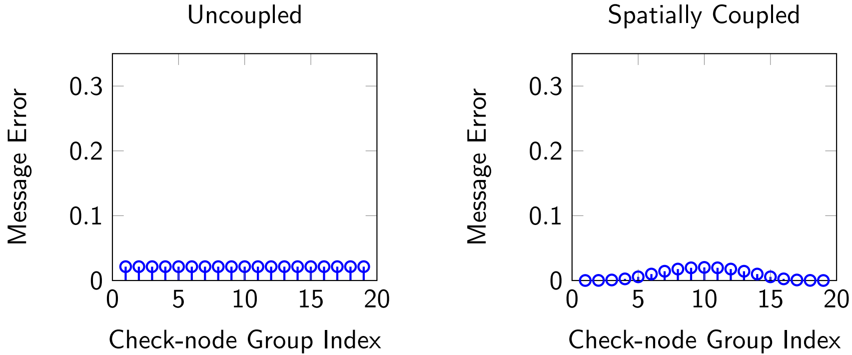

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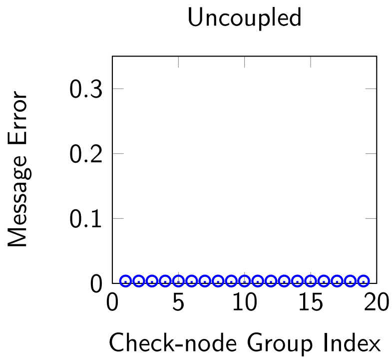
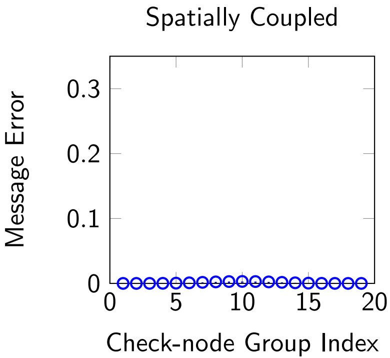
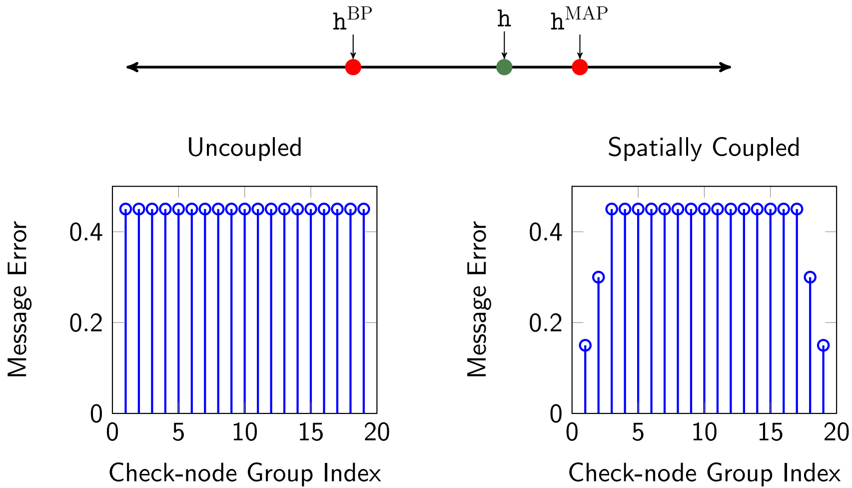




Threshold Saturation via Spatial Coupling (2)



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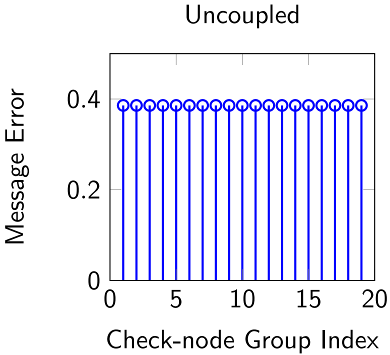

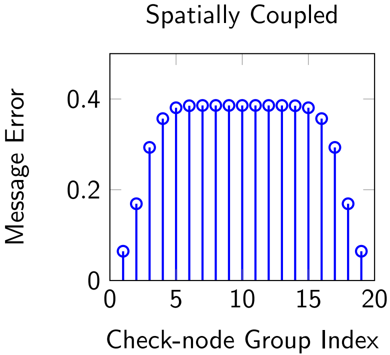

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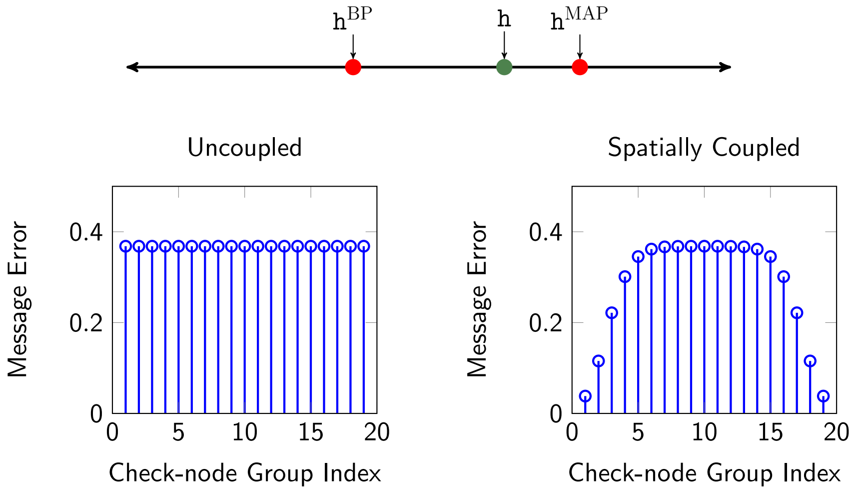

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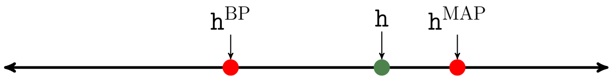

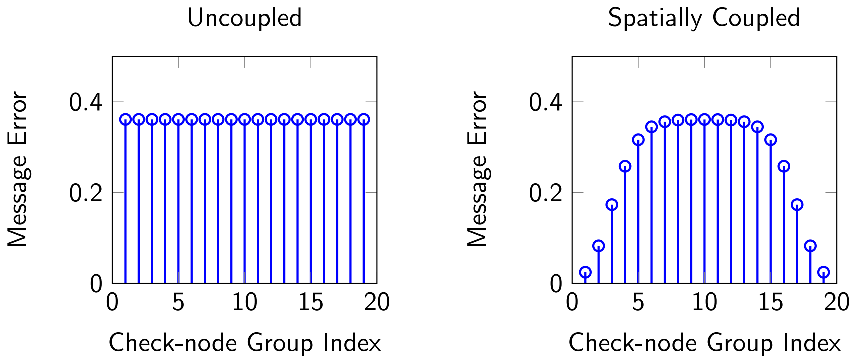

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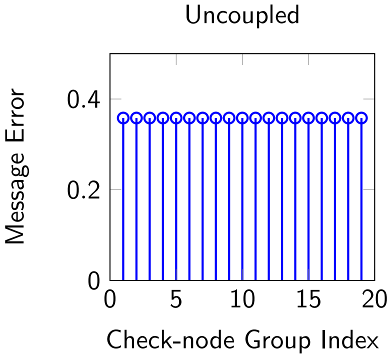

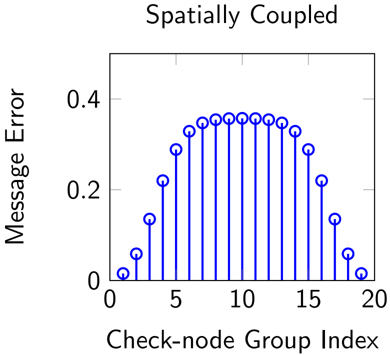

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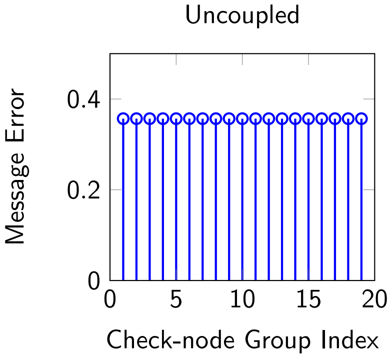

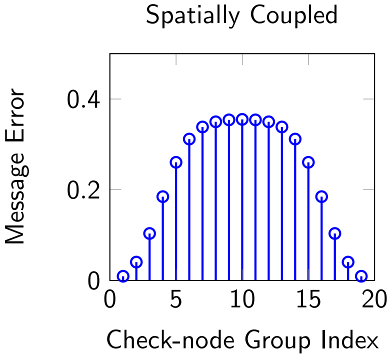

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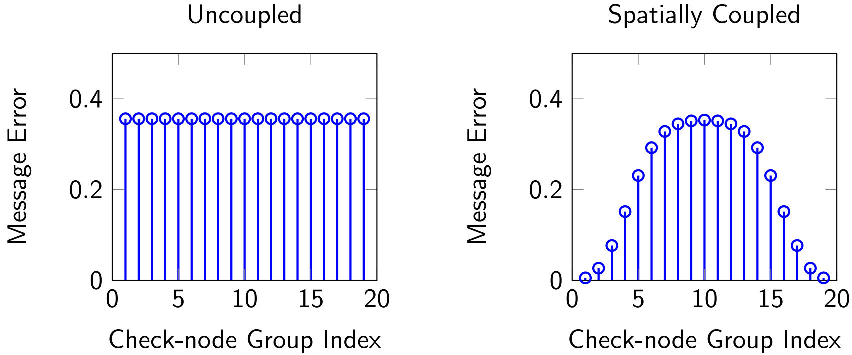

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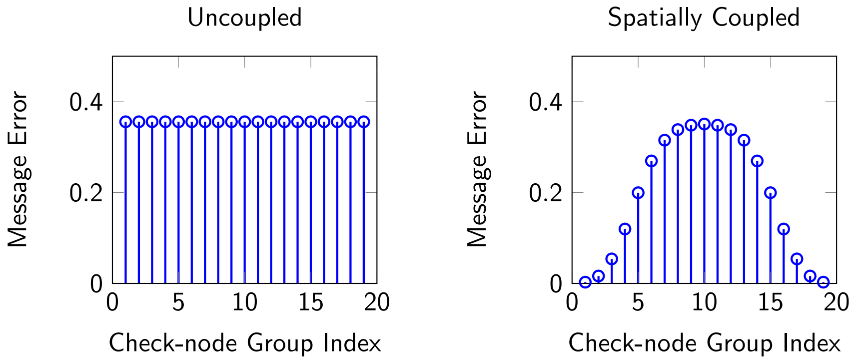
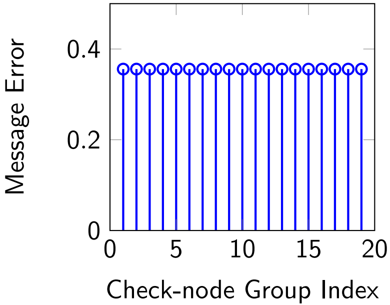
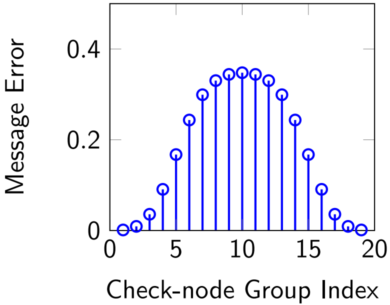


Threshold Saturation via Spatial Coupling (2)

Uncoupled



Spatially Coupled



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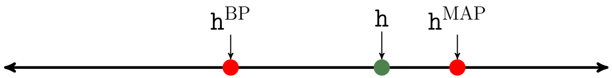

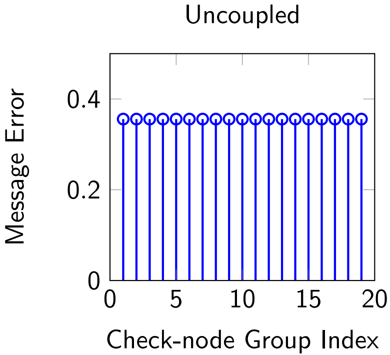

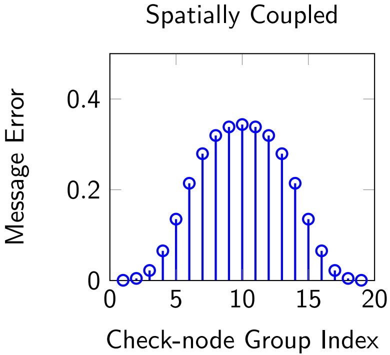

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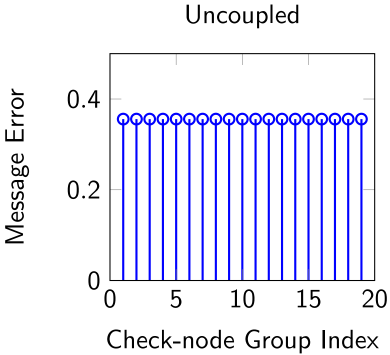

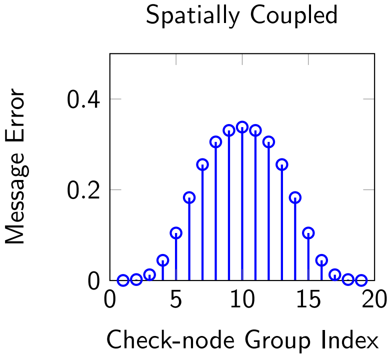

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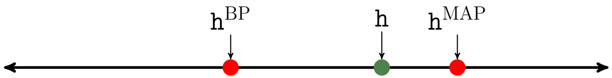

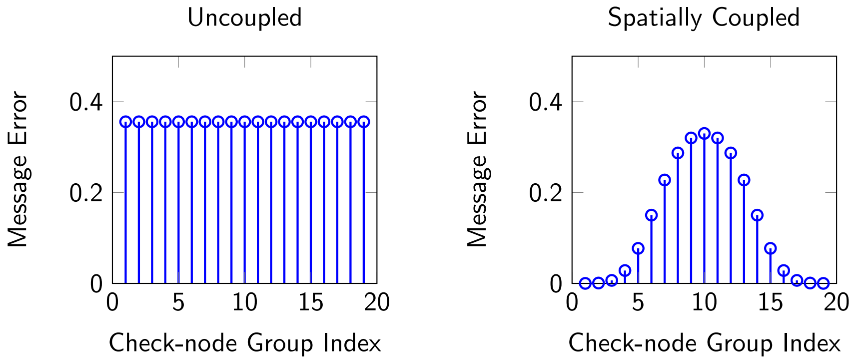

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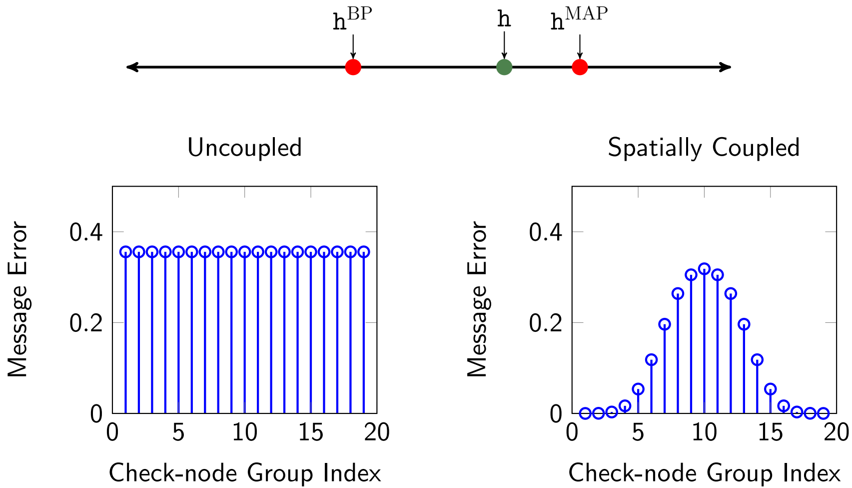

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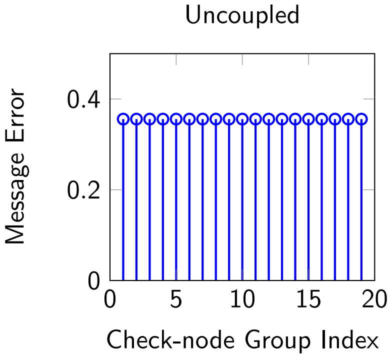

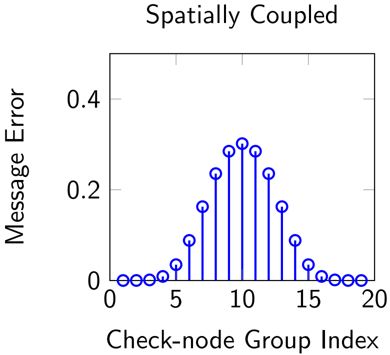

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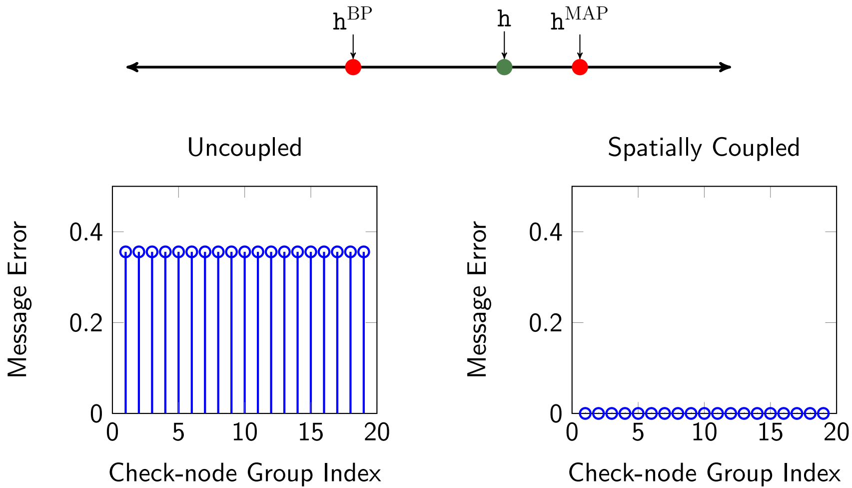
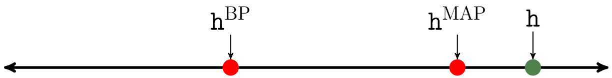
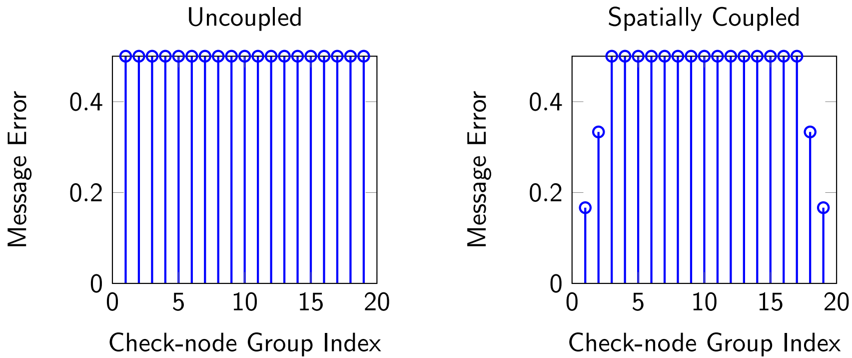


Threshold Saturation via Spatial Coupling (3)





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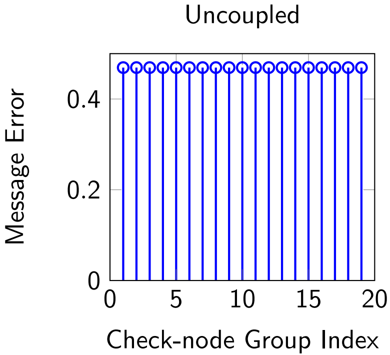

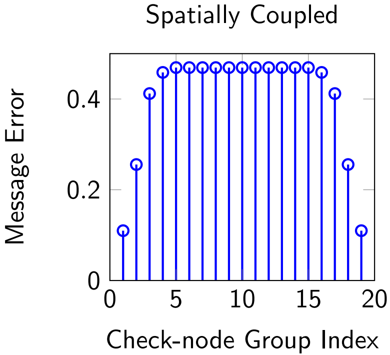

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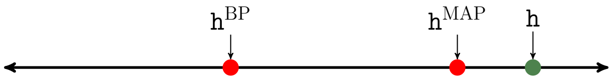

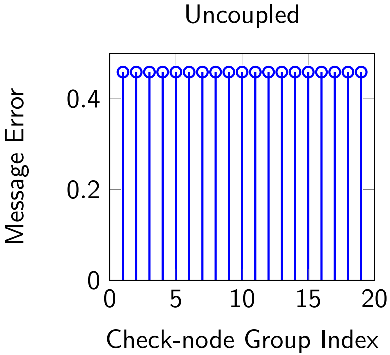

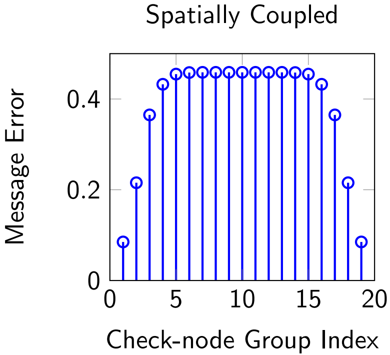

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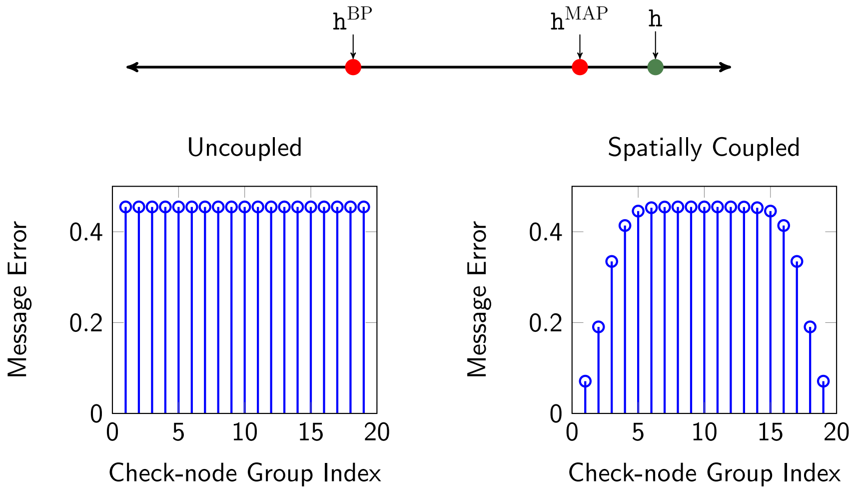

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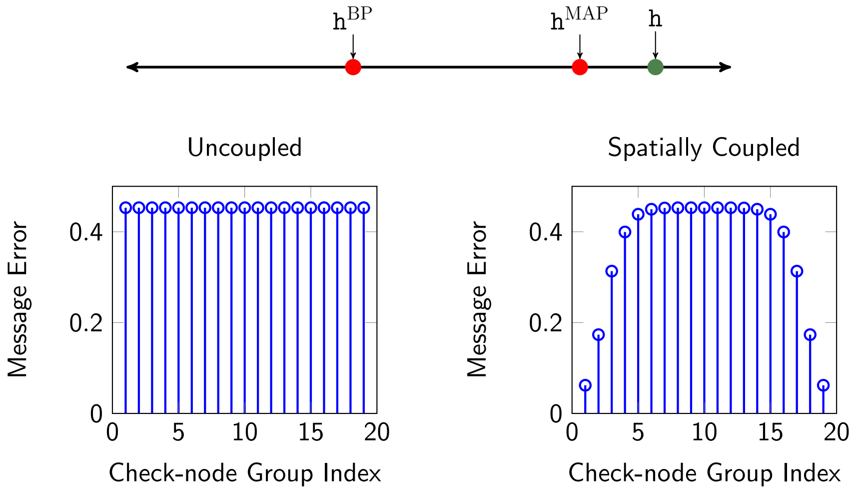

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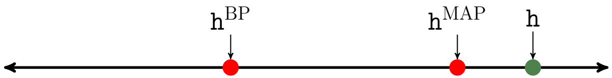

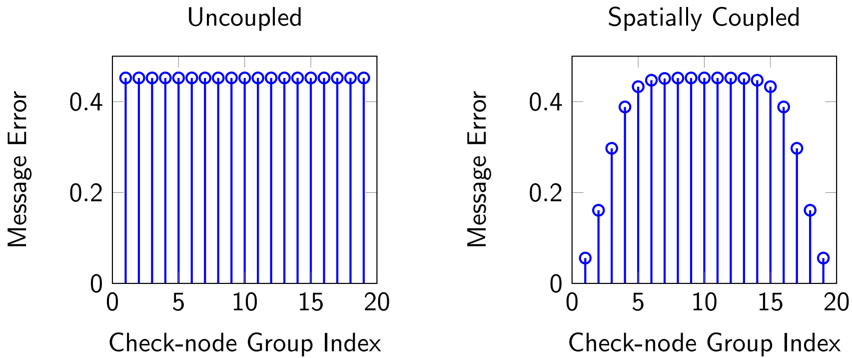

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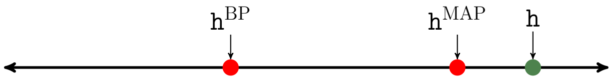

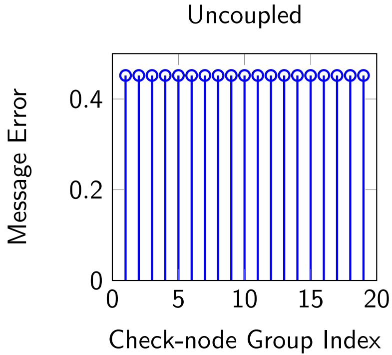

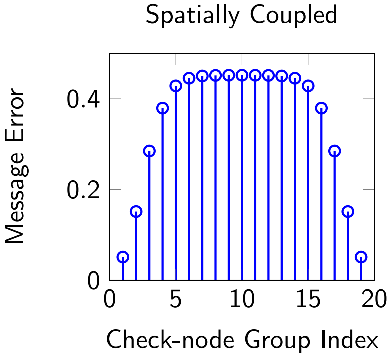

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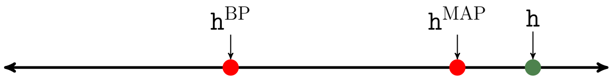
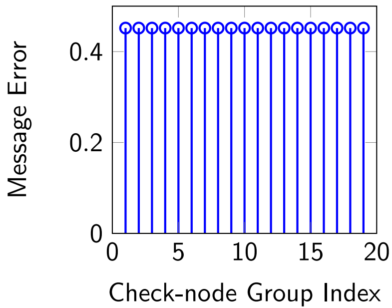
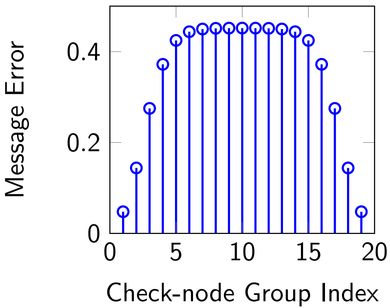


Uncoupled



Spatially Coupled



Threshold Saturation Result

MAP Performance with a BP Decoder!

For large N, w $\quad \mathtt{h}_c^{\mathrm{BP}} = \mathtt{h}^{\mathrm{MAP}}$

SC-LDPC (ℓ, r)	Shannon $\mathtt{h}^{\mathrm{Sh}}$	AWGN $\mathtt{h}_c^{\mathrm{BP}}$	BSC $\mathtt{h}_c^{\mathrm{BP}}$
$(3, 6)$	0.5000	0.4794	0.4681
$(4, 6)$	0.6667	0.6645	0.6633
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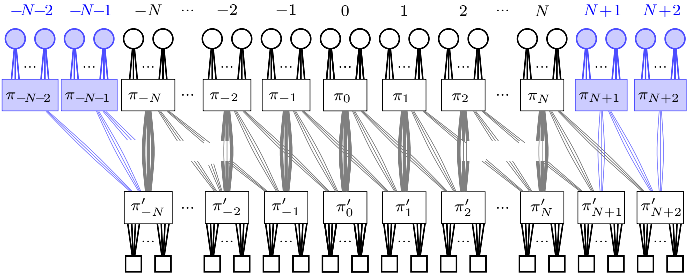

$-N-2$
$-N-1$
$-N$
$\cdots$
-2
-1
0
1
2
$\cdots$
N
$N+1$
$N+2$
π_{-N-2}
π_{-N-1}
π_{-N}
π_{-2}
π_{-1}
π_0
π_1
π_2
π_N
π_{N+1}
π_{N+2}
π'_{-N}
π'_{-2}
π'_{-1}
π'_0
π'_1
π'_2
π'_N
π'_{N+1}
π'_{N+2}

Rate loss for finite N and w

SC-LDPC (ℓ, r, N, w)	Shannon $\mathtt{h}^{\mathrm{Sh}}$	AWGN $\mathtt{h}_{\mathrm{c}}^{\mathrm{BP}}$	BSC $\mathtt{h}_{\mathrm{c}}^{\mathrm{BP}}$
$(3, 6, 10, 3)$	0.5434	0.4794	0.4681
$(3, 6, 20, 3)$	0.5222	0.4794	0.4681
$(3, 6, 30, 3)$	0.5149	0.4794	0.4681
$(4, 6, 10, 3)$	0.7245	0.6645	0.6633
$(4, 6, 20, 3)$	0.6963	0.6645	0.6633
$(4, 6, 30, 3)$	0.6866	0.6645	0.6633
$(5, 6, 10, 3)$	0.9056	0.8333	0.8333
$(5, 6, 20, 3)$	0.8704	0.8333	0.8333
$(5, 6, 30, 3)$	0.8582	0.8333	0.8333

Pros & Cons

Pros

- Significant improvment in thresholds
- Achieves capacity under simple BP decoding
- Universality - works for all channels models!

Cons

- Need large blocklengths to leverage the gains

Outline

Lattices and Lattice Codes

Lattice

A lattice of dimension n is a discrete subgroup of $\mathbb{R}^n$ isomorphic to $\mathbb{Z}^n$

$$\Lambda = \{\mathbf{Gz}, \mathbf{z} \in \mathbb{Z}^n\}$$

Lattices and Lattice Codes

Lattice

A lattice of dimension n is a discrete subgroup of $\mathbb{R}^n$ isomorphic to $\mathbb{Z}^n$

$$\Lambda = \{\mathbf{Gz}, \mathbf{z} \in \mathbb{Z}^n\}$$

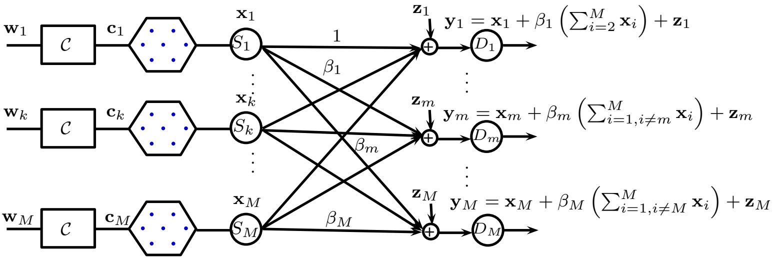
- Efficient structures for
  - Mathematics: sphere packing and sphere covering problems
  - Information Theory: channel coding & quantization
- Single user Gaussian channel - Erez and Zamir
- Coding with side information - Wyner-Ziv and Costa, Zamir, Erez and Shamai
- Secrecy - He and Yener
- Dirty multiple access channel - Philosof, Khisti, Erez and Zamir

"Lattices are everywhere" by Ram Zamir

Prior Work

New perspectives for dealing with interference:

- Interference alignment - Sridharan, Jafarian, Vishwanath and Jafar



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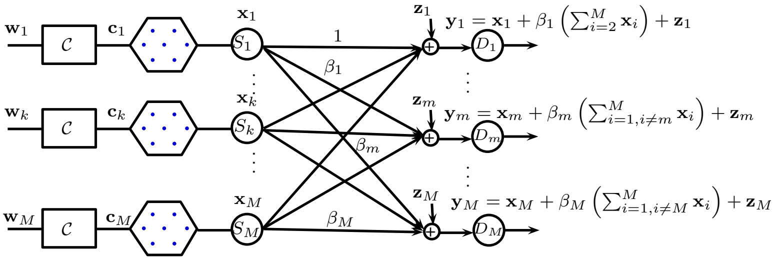

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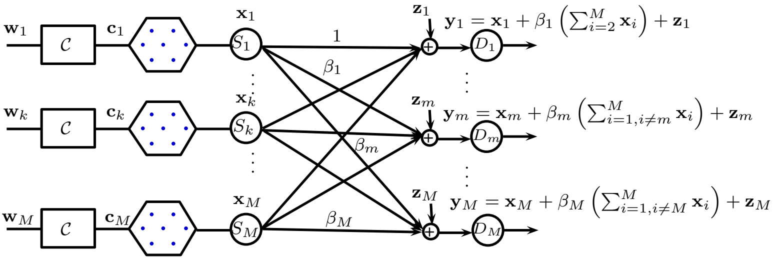

- Above schemes are all based on lattices good for channel coding

Goodness of Lattices for Channel Coding

- Voronoi region $\mathcal{V}$ of a lattice Λ, $\mathcal{V} := \{\mathbf{x} : \|\mathbf{x}\| \leq \|\mathbf{x} - \mathbf{c}\| \quad \forall \mathbf{c} \in \Lambda\}$
- Fundamental volume of Λ, $V(\Lambda)$: $\mathrm{Vol}(\mathcal{V})$
- Let a lattice point $\lambda \in \Lambda$ is trasnmitted via AWGN channel of variance σ^2
- Volume-to-noise ratio(VNR) of Λ:

$$\mathrm{VNR} = \frac{V(\Lambda)^{2/n}}{2\pi e \sigma^2}$$

- $P(\Lambda, \sigma^2) := \Pr(d(\lambda, \lambda + \mathbf{z}) \geq d(\lambda', \lambda' + \mathbf{z}))$ for some $\lambda' \in \Lambda$

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Poltyrev Goodness for Channel Coding

For any VNR>1 $\exists\{\Lambda_n\}$ such that $P(\Lambda_n, \sigma^2) \to 0$ as $n \to \infty$.

- *Poltyrev*-good lattices are at the core of such lattice coding schemes

Objective

Motivating questions

- All the existing results were based on Construction-A.
  - Linear codes over increasing field sizes and their ML decoding

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Motivating questions

- All the existing results were based on Construction-A.
  - Linear codes over increasing field sizes and their ML decoding
- Is this construction fundamental to good lattices?
- Can we work with just binary codes under practical decoding schemes?

Main Result

Codes over $\mathbb{F}_2$ and BP decoding suffice

- Recall Forney et al's result - based on nested random binary linear codes
- Propose capacity-achieving nested SC LDPC ensemble
- Construct lattices using Construction-D, based on the above ensemble
- Show existence of sequence of lattices that are *Poltyrev*-good under BP

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Applications

- Apply proposed lattices to Symmetric Interference Channel
- Can be applied to other problems which adopt Construction A lattices

Construction D with L levels

- Barnes and Sloane '83, Forney, Chung and Trott '00, Yan, Ling, Wu ' 13
- Choose $G_1 \subseteq \ldots \subseteq G_L$ where G_l is a gen matrix of code $\mathcal{C}_l$ over $\mathbb{F}_2$.

$$\left[\begin{array}{ccc} \cdots & g_L & \cdots \\ \cdots & \vdots & \cdots \\ \hline \cdots & g_{k_2} & \cdots \\ \cdots & \vdots & \cdots \\ \hline \cdots & g_{k_1} & \cdots \\ \cdots & \vdots & \cdots \\ \cdots & g_2 & \cdots \\ \cdots & g_1 & \cdots \end{array}\right]$$

G_1 $\quad$ G_2 $\quad \ldots \quad$ G_L

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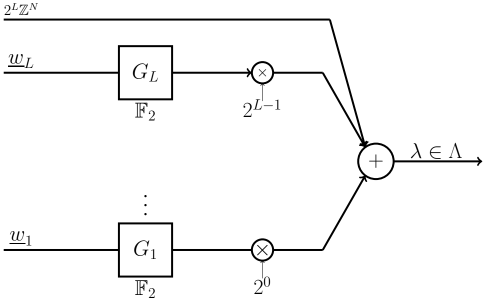
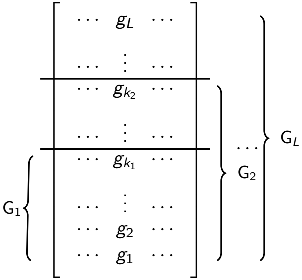
- $\underline{\lambda} = \underline{w}_1\mathbf{G}_1 + 2\underline{w}_2\mathbf{G}_2 \ldots + 2^{L-1}\underline{w}_{L-1}\mathbf{G}_{L-1} + 2^L\mathbb{Z}^N \in \Lambda$





Multi-Level Decoding(Successive Cancellation)

- $\underline{y} = \boxed{\underline{w}_1\mathbf{G}_1 + 2\underline{w}_2\mathbf{G}_2 \ldots + 2^{L-1}\underline{w}_{L-1}\mathbf{G}_{L-1} + 2^L\mathbb{Z}^N} + \underline{n}$

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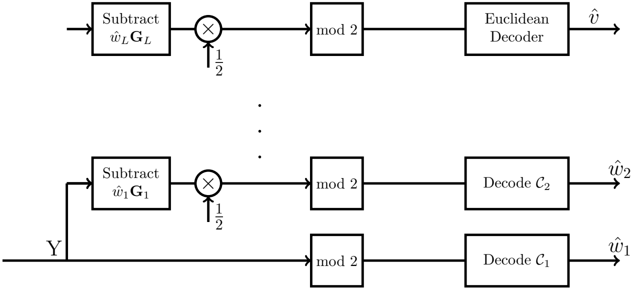
- $\underline{y} \bmod 2 = [\underline{w}_1\mathbf{G}_1 + \underline{n}] \bmod 2 = \underline{w}_1 \odot \mathbf{G}_1 + \boxed{\underline{n} \bmod 2}$
- Decode $\underline{w}_1$, reconstruct $\underline{w}_1\mathbf{G}_1$ and subtract from $\underline{y}$



Theorem (Forney, Trott & Chung)

There exists a sequence of Construction D lattices based on $\mathcal{C}_1 \subseteq \mathcal{C}_2 \ldots \subseteq \mathcal{C}_L$ such that the VNR $\rightarrow 1$ and the $Pr(\lambda, \sigma^2) \rightarrow 0$.

- Take L large enough.
- It's sufficient that $\mathcal{C}_i$ at each level is capacity achieving for the mod-2 AWGN channel.

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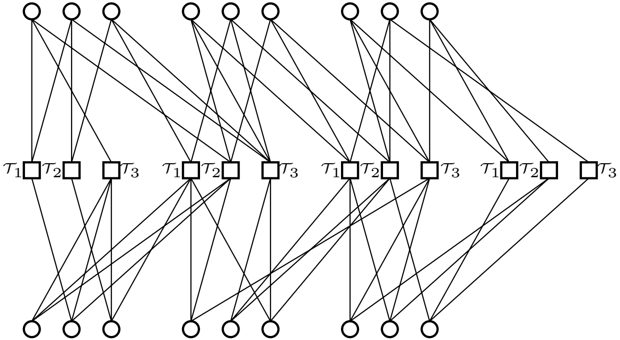
Objective:

- Capacity achieving nested code constructions, preferably under BP decoding.

Proposed Nested Spatially-Coupled LDPC Ensemble

1. Begin with a (d_v^1, d_c) SC LDPC code. For ex, $(d_v^1 = 3, d_c = 6, L = 3, w = 2)$.
2. Group check nodes into type $\mathcal{T}_k$, $k \in \{1, \ldots, d_v^1\}$



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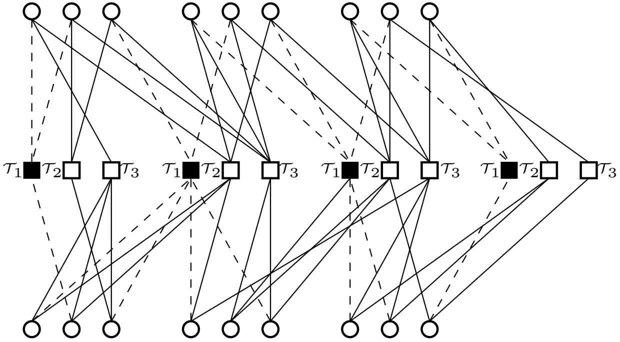

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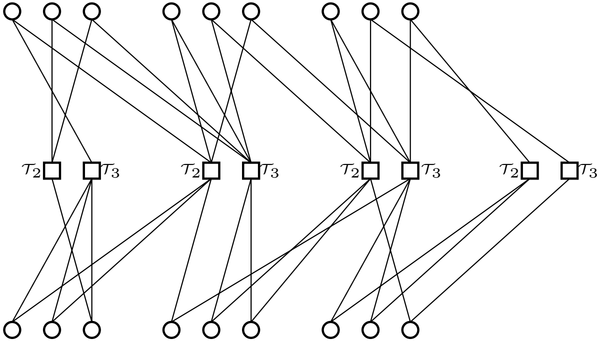
4. Results in a super-code that is a (d_v^2, d_c) SC LDPC code.



Lattice Design based on the proposed Nested SC LDPC ensemble

1. For a given σ, compute the capacity of the mod-2 AWGN channel at each level:

$$\underline{y}_i = \underline{w}_i \mathbf{G}_i + \frac{1}{2^{i-1}}\underline{n} \mod 2 = \underline{w}_i \odot \mathbf{G}_i + \boxed{\frac{1}{2^{i-1}}\underline{n} \mod 2}$$

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Lemma

Given nested binary linear codes $\mathcal{C}_1 \subseteq \mathcal{C}_2 \subseteq \ldots \subseteq \mathcal{C}_r$ there exists nested generator matrices for these codes.

Proposed Ensemble is Capacity achieving

Theorem

Each code ensemble in the proposed nested Spatially-Coupled LDPC ensemble is capacity achieving.

Proof.

- Show that the mod 2 AWGN channel is BMS.
- Each derived protograph has the same spatially coupled structure.
- The proof follows from Kudekar & Urbanke, Kumar & Pfister's results.

□

Proposed Lattices are Poltyrev-Good

Theorem

There exists a sequence of SC LDPC lattices with $VNR(\Lambda, \sigma^2) \to 1$ *for which, under multistage BP decoding,* $\mathbb{E}\left[P(\lambda, \sigma^2)\right] \to 0$ *as* $w, L, M \to \infty$.

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Proof.

- The proposed nested ensemble achieve capacity.
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□

- Binary codes and more importantly practical BP decoding suffices.
- Practically we observe that two levels of coding gets you lattices very close to Poltyrev limit.

Design Example of Poltyrev-Good Lattice

Target error probability $P(2^L\mathbb{Z}^n, \sigma_L^2) = 10^{-4}$ in the uncoded level $\implies \sigma_L = 0.08$

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Cap	0.99	0.57	0.02
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   - Note $P(4\mathbb{Z}^n, \sigma^2) \approx nP(4\mathbb{Z}, \sigma^2)$
   - We fix $n = 2 \times 10^5$

(d_c, d_v^1, d_v^2)	(L,w)	$P(4\mathbb{Z}, \sigma^2)$	$\sigma_{\max}$	VNR	$\text{VNR}_{\text{rate-loss}}$
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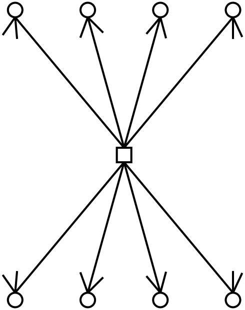
(60, 26, 3)	(72, 12)	5×10^{-10}	0.3200	0.482dB	0.927dB
(60, 27, 3)	(64, 9)	5×10^{-10}	0.3203	0.57dB	0.951dB

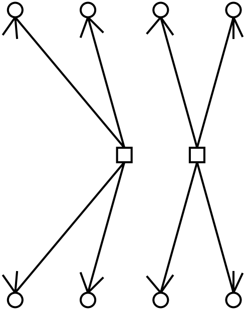
Alternate Nested SC LDPC ensemble

- Derive a lower rate code by "splitting the checks"
- Consider a $(3, 8)$ code

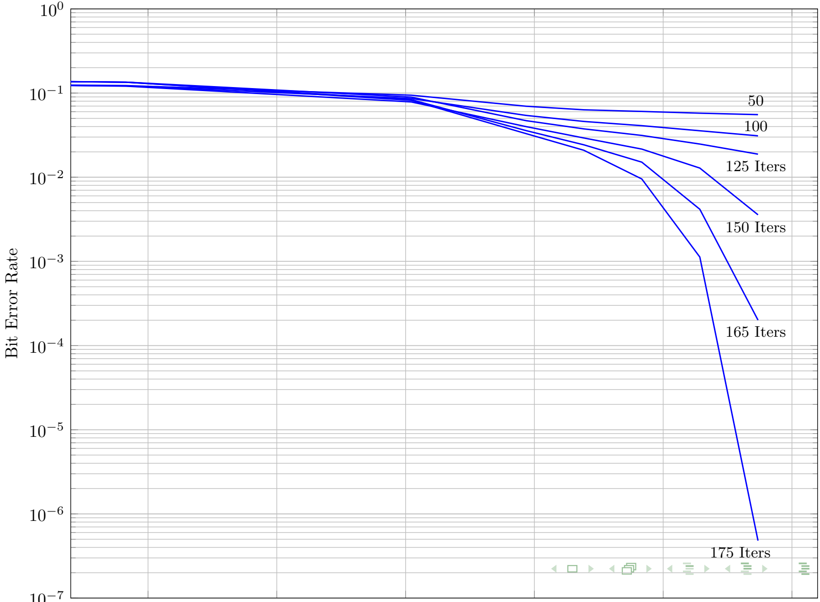


Alternate Nested SC LDPC ensemble

- Derive a lower rate code by “splitting the checks”
- Consider a $(3, 8)$ code
- Split each check into “two” checks to derive a $(3, 4)$ sub-code
- Easy to prove that resulting code is from the $(3, 4)$ SC LDPC ensemble



Simulation Results



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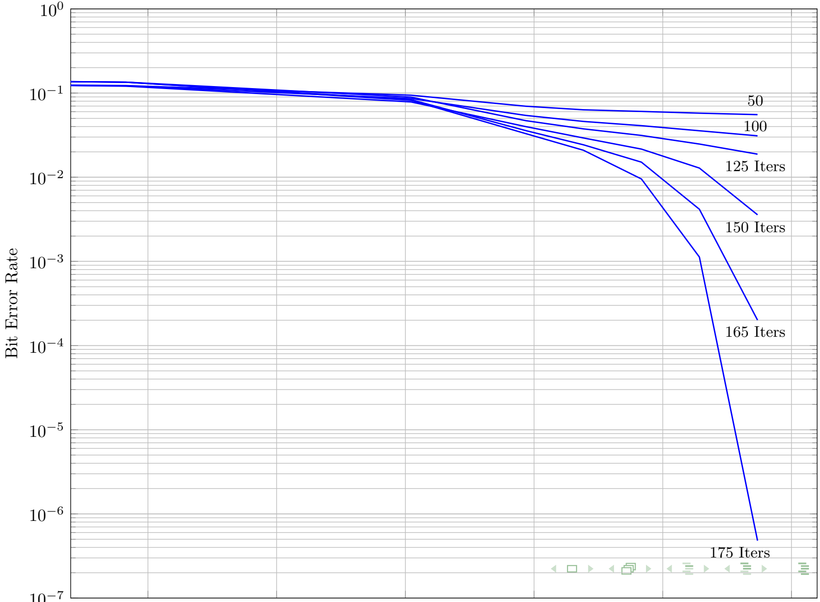
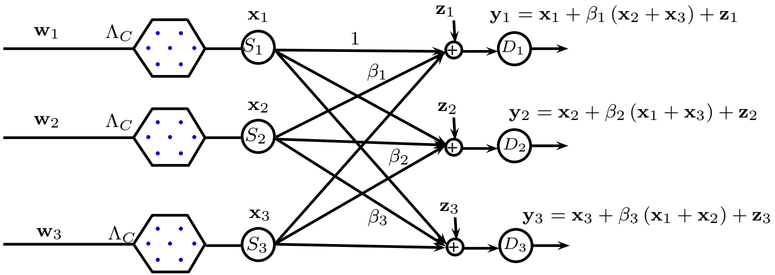


3-User Symmetric Interference Channel



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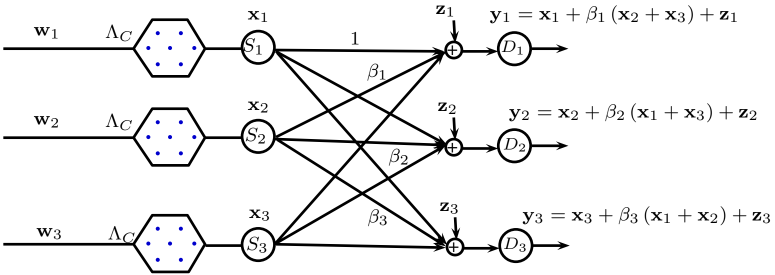


- $\mathbf{x}_i \in \Lambda_C \triangleq \Lambda \cap \mathbb{Z}_4^N$ is transmitted.

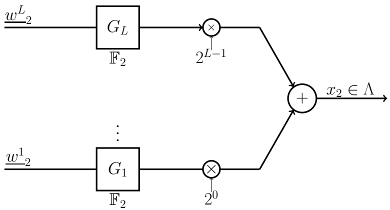
Symmetric Interference Channel - Decoding Sums

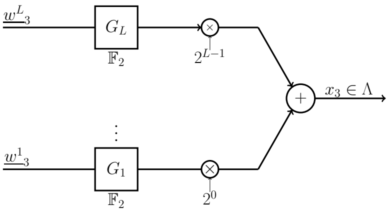
Interference at Destination 1:

$$\begin{aligned}\mathbf{x}_2 + \mathbf{x}_3 &= (\underline{w}_2^1 + \underline{w}_3^1)\mathbf{G}_1 + 2(\underline{w}_2^2 + \underline{w}_3^2)\mathbf{G}_2 + 4\mathbf{k}_{23} \\ &= (\underline{w}_2^1 \oplus \underline{w}_3^1)\mathbf{G}_1 + 2(\underline{c}_{23}^1 \oplus \underline{w}_2^2 \oplus \underline{w}_3^2)\mathbf{G}_2 + 4(\underline{c}_{23}^2 + \mathbf{k}_{23})\mathbf{Z}\end{aligned}$$

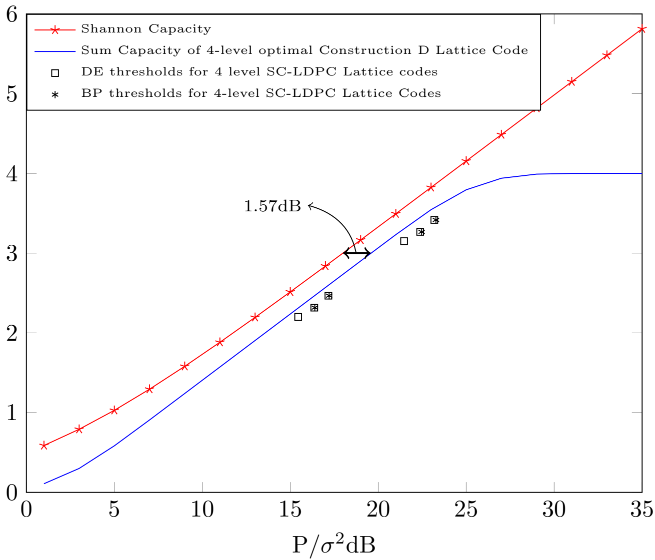
where the carry overs are

$$\begin{aligned}\underline{c}_{23}^1 &= 0.5\left(\underline{w}_1^1 + \underline{w}_1^2 - \underline{w}_1^1 \oplus \underline{w}_1^2\right), \\ \underline{c}_{23}^2 &= 0.5\left(\underline{c}_{23}^1 + \underline{w}_2^1 + \underline{w}_2^2 - \underline{c}_{23}^1 \oplus \underline{w}_2^1 \oplus \underline{w}_2^2\right)\end{aligned}$$





Achievable Information Rates



Concluding Remarks

- Multilevel constructions - efficient ways to decode integer combinations
- Need capacity achieving nested codes
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- Need capacity achieving nested codes
- Multilevel construction is provably good under message passing decoding
- Coding schemes based on binary LDPC codes and iterative decoding suffice

Outline

Lossy Source Coding Problem

$$X^n = (X_1, \cdots, X_n),\ X_i \sim \text{Bernoulli}(\tfrac{1}{2})$$

Binary code $\mathcal{C} = (n, k)$, rate $R = k/n$

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Lossy Source Coding

- Compress X^n to $\hat{X}^n \in \mathcal{C}$
- Min. Hamming distortion

$$D = \frac{1}{n} \sum_{i=1}^{n} \mathbb{E}|X_i - \hat{X}_i|$$

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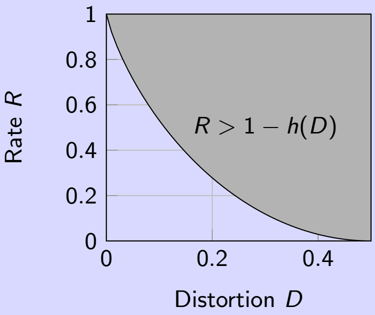
- Rate-Distortion theory:

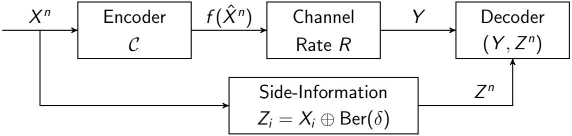
$$R > 1 - h(D)$$

- $h(\cdot)$ is binary entropy function

$$h(D) = -D\log_2 D - (1-D)\log_2(1-D)$$



Side-Information Problems: Wyner-Ziv



Wyner-Ziv Formulation

- Side-information Z^n about X^n
- Decoder additionally has Z^n
- Say $Z_i = X_i \oplus \mathrm{Ber}(\delta)$

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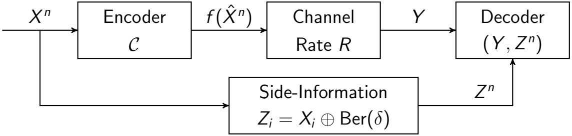

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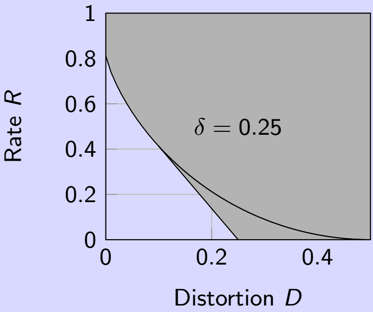
- Wyner-Ziv theory:

$$R > l.c.e\{h(D * \delta) - h(D), (\delta, 0)\}$$

- $D * \delta = D(1 - \delta) + \delta(1 - D)$



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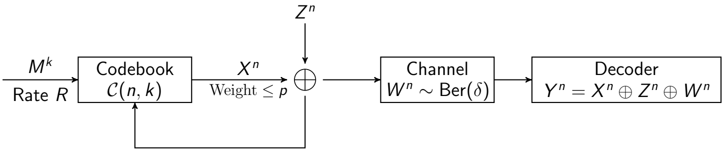

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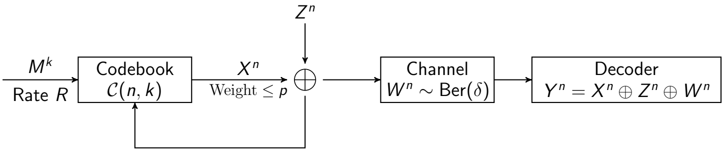
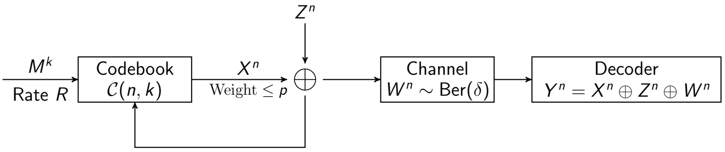


Gelfand-Pinsker Formulation

- Message M^k encoded to $X^n \in \mathcal{C}$ with $\frac{1}{n}\sum_{i=1}^{n} \mathbb{E}[X_i] \le p \le \frac{1}{2}$
- Side-information Z^n is available only at the encoder

Side-Information Problems: Gelfand-Pinsker

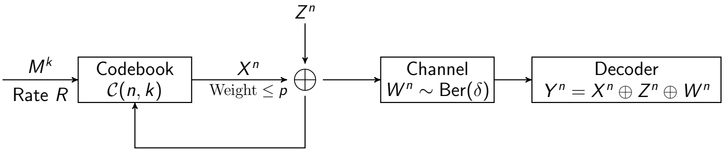


Gelfand-Pinsker Formulation

- Message M^k encoded to $X^n \in \mathcal{C}$ with $\frac{1}{n}\sum_{i=1}^{n} \mathbb{E}[X_i] \leq p \leq \frac{1}{2}$
- Side-information Z^n is available only at the encoder
- The output at the decoder is

$$Y^n = X^n \oplus Z^n \oplus W^n, \quad \{W_i\} \sim \mathrm{Ber}(\delta)$$

Side-Information Problems: Gelfand-Pinsker



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- Capacity region by Gelfand-Pinsker:

$$R < h(p) - h(\delta)$$

Main Result

Objective

- Construct low-complexity coding schemes that achieve the complete rate regions of Wyner-Ziv and Gelfand-Pinsker
  - Low-complexity encoding and decoding

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- Message-passing algorithms have non-negligible gap

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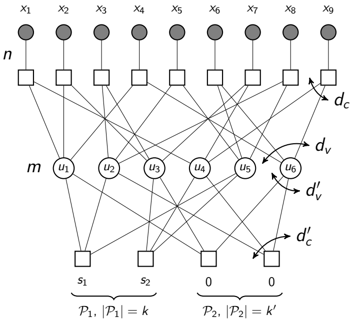
Idea

- Wainwright et al. used compound LDGM/LDPC codes with optimal encoding/decoding
- Message-passing algorithms have non-negligible gap
- Remedy via Spatial-Coupling
  - Channel coding in coupled compound codes (Kasai et al.)
  - Lossy source coding with spatially-coupled LDGM (Aref et al.)
  - Encoding with compound codes has additional challenges

Compound LDGM/LDPC Codes



- Codebook $\mathcal{C}(n, m-k-k')$
- Message constraints

$$u_1 \oplus u_2 \oplus u_5 = s_1, \quad u_1 \oplus u_3 \oplus u_6 = 0$$

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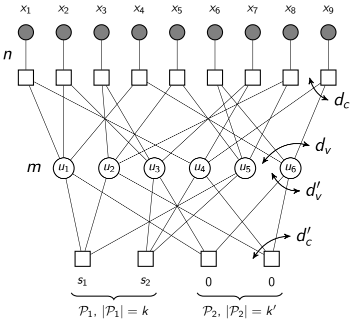

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Key Properties

- Compound code is
  - a good source code under optimal encoding
  - a good channel code under optimal decoding
- LDGM code is
  - a good source code under optimal encoding
  - (side note) LDGM code is not a good channel code

Good Code

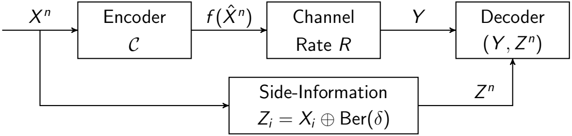
"Good" source code

- Rate of the code is $R = 1 - h(D) + \varepsilon$
- When this code is used to optimally encode $\mathrm{Ber}(\frac{1}{2})$
- The average Hamming distortion is at most D

"Good" channel code

- Rate of the code is $R = 1 - h(\delta) - \varepsilon$
- When this code is used for channel coding on $\mathrm{BSC}(\delta)$
- Message est. under optimal decoding with error at most ε

Coding Scheme: Wyner-Ziv



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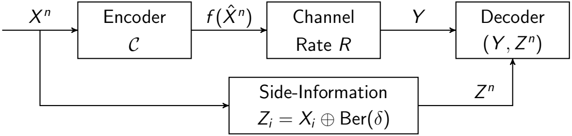

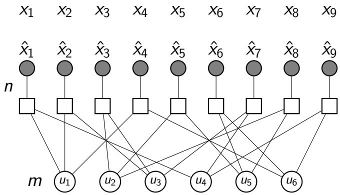

- Encode X^n to $\hat{X}^n$ using LDGM w/ Distortion $\approx D$

$\frac{m}{n} \approx 1 - h(D) + \varepsilon$

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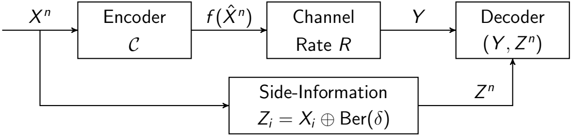
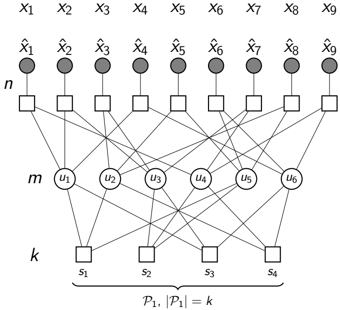




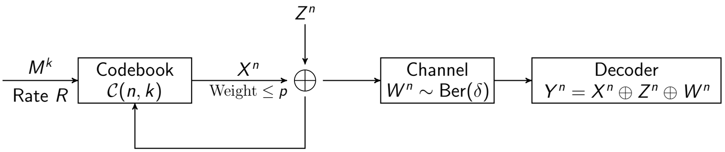
$\frac{m}{n} \approx 1 - h(D) + \varepsilon \qquad \frac{m-k}{n} \approx 1 - h(D * \delta) + \varepsilon$

- Encode X^n to $\hat{X}^n$ using LDGM w/ Distortion $\approx D$
- Compute & transmit s_i's
  $R = \frac{k}{n} \approx h(D * \delta) - h(D)$
- Decoder has Z^n:

$$\begin{aligned} Z_i &= X_i \oplus \mathrm{Ber}(\delta) \\ &\approx \hat{X}_i \oplus \mathrm{Ber}(D) \oplus \mathrm{Ber}(\delta) \\ &= \hat{X}_i \oplus \mathrm{Ber}(D * \delta) \end{aligned}$$

- Decode $\hat{X}^n$ from Z^n & s_i

Coding Scheme: Gelfand-Pinsker



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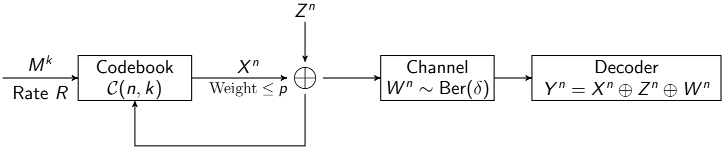

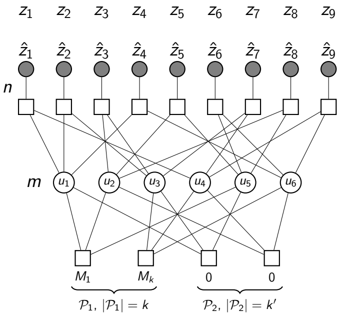

$\frac{m-k-k'}{n} \approx 1 - h(p) + \varepsilon \qquad \frac{m-k'}{n} \approx 1 - h(\delta) + \varepsilon$

- With message M^k, encode Z^n to $\hat{Z}^n$ (Distortion $\approx p$)
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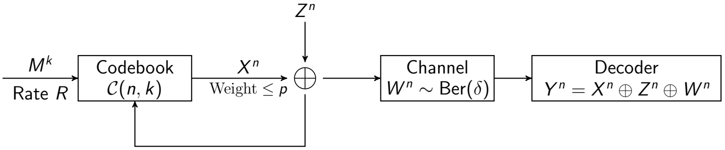

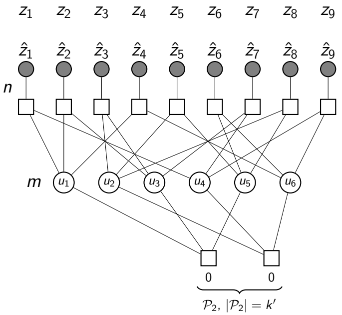

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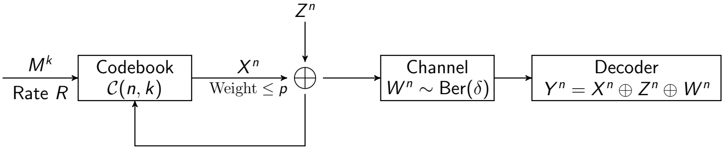

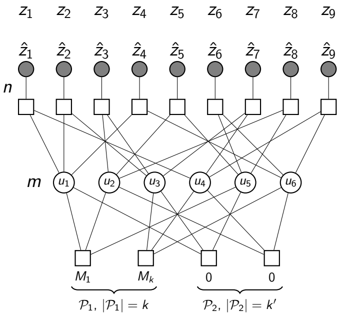

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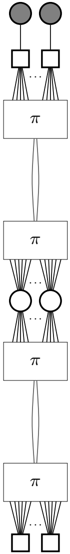
- Decode $\hat{Z}^n$ and compute M^k
- $R = \frac{k}{n} \approx h(p) - h(\delta)$

Remarks

- Need codes that are simultaneously good for channel and source coding
- Use message-passing algorithms instead of optimal
- Use spatial-coupling for goodness of codes under message-passing

Spatially-Coupled Compound LDGM/LDPC Codes



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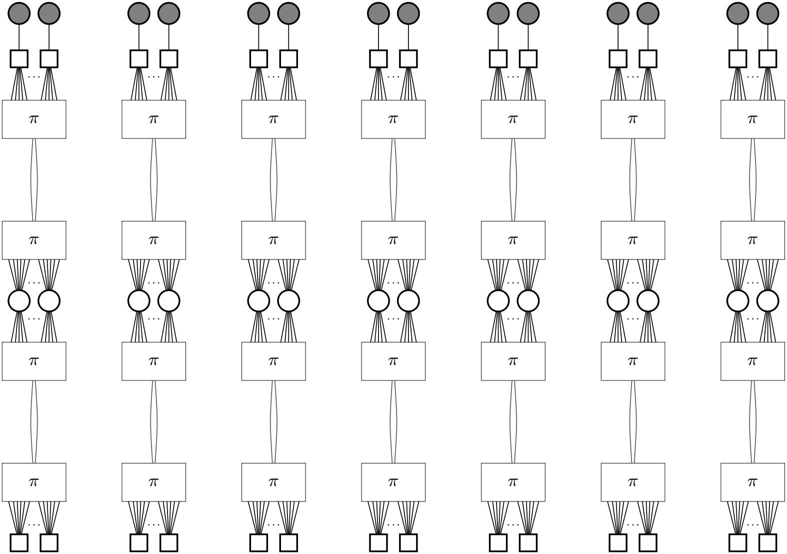

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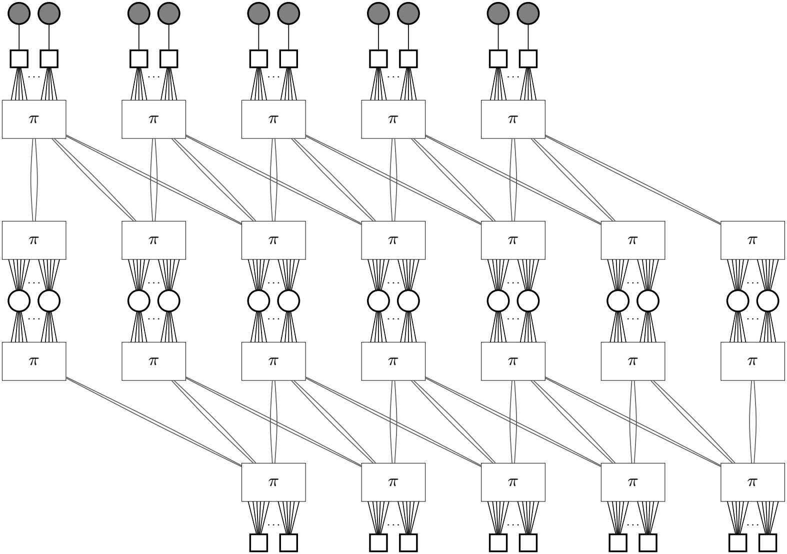

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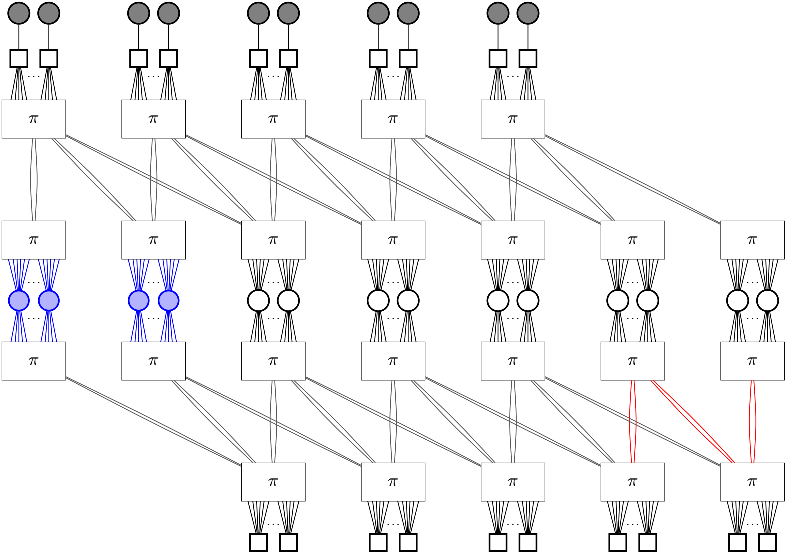
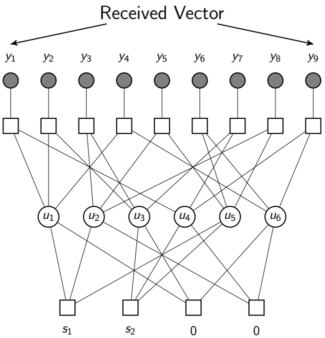
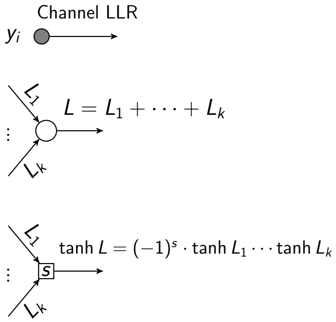


Decoding in Spatially-Coupled Compound Codes

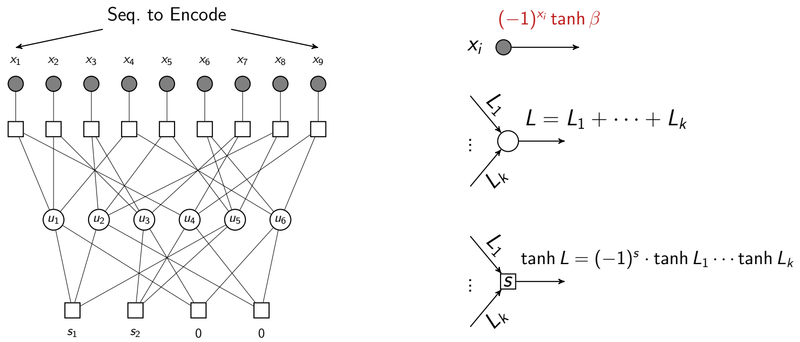




Remarks

- Standard message-passing algorithm
- Threshold saturation proven for SC compound codes on BEC
- Empirically observed for BMS channels

Encoding in Spatially-Coupled Compound Codes



Remarks

- Inverse temperature parameter β
- Message-passing rules are the same
- However, a crucial decimation step is needed

Encoding in SC Compound Codes: BPGD Algorithm

Encoding in SC Compound Codes: Remarks

- Randomization in setting u_{i^*} is crucial
- BPGD applied to uncoupled code always failed
- Spatially-coupled structure is crucial for successful encoding
  - In addition, distortion is close to optimal thresholds
  - Does not encode if decimated from both left and right
  - Does not encode if both left and right boundary is set to 0

Encoding in SC Compound Codes: Numerical Example

Block length (n)	4-cycles	Attempts $1/2/3/4/\geq 5$
9000	yes	5/3/5/2/35
9000	no	21/12/5/3/9
27000	no	35/15/0/0/0
45000	no	40/9/0/0/1
63000	no	44/6/0/0/0
81000	no	50/0/0/0/0

Remarks

- # Attempts to encode 50 seq. in $(6,3)$ LDGM / $(3,6)$ LDPC
- $L = 20,\ w = 4,\ \beta = 0.65,\ T = 10$
- Removing 4-cycles dramatically improves success
- How much do 6-cycles matter?

Numerical Results: Wyner-Ziv

LDGM (d_v, d_c)	LDPC (d'_v, d'_c)	(L, w)	(D_*, δ_*)	(D, δ)
$(6, 3)$	(3,6)	(20,4)	(0.111,0.134)	(0.1174,0.122)
$(8, 4)$	(3,6)	(20,4)	(0.111,0.134)	(0.1149,0.120)
$(10, 5)$	(3,6)	(20,4)	(0.111,0.134)	(0.1139,0.122)

Remarks

- D_* and δ_* are calculated based on the rate of the respective code:

$$D_* = h^{-1}(1 - R1) \qquad \delta_* = h^{-1}(1 - R2)$$

- $n \approx 140000$, $\beta = 1.04$, $T = 10$

Numerical Results: Gelfand-Pinsker

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Concluding Remarks

Conclusion

- Spatially-coupled codes achieve the rate regions of Wyner-Ziv and Gelfand-Pinsker problems
- Coupling structure is also crucial
  - to achieve optimum thresholds
  - for encoding to succeed with decimation

Open Questions

- Effect of degree profiles, short-cycles on encoding success
- Precise trade-offs with polar codes

Outline

Write-Once Memories

Flash Memory

- In typical flash memory, changing from 0 to 1 is easy
- Resetting 1 to 0 requires rewriting whole block
- Write-once memories model such storage systems

Write-Once Memories

0	0	0	0	0	0	0	0	0

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Binary Write-Once Memories

- $0 \longrightarrow 1$ is allowed

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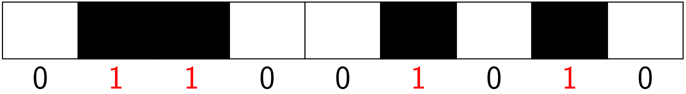

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Binary Write-Once Memories

- $0 \longrightarrow 1$ is allowed
- $1 \longrightarrow 0$ is forbidden

Capacity Region (I) - Noiseless

Message

0	0	1	0	0	1	0	1	0

Write-Once Memory without Noise

- In 1982, Rivest and Shamir gave first WOM codes
  - 2 bits in 2 writes with only 3 cells
- Only about $nt/\log(t)$ cells required to store n bits for t writes

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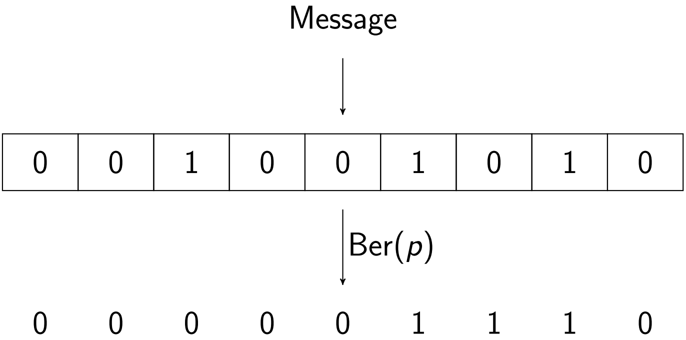
- Only about $nt/\log(t)$ cells required to store n bits for t writes
- In 1985, Heegard gave the capacity for t-write system
- For a 2-write system, it is

$$\left\{(R_1, R_2) \mid 0 \le R_1 < h(\delta),\ 0 \le R_2 < 1 - \delta\right\}$$

Capacity Region (II) - Read Errors



Write-Once Memory with Read Errors

- Different from write errors
- $Y = X \oplus \text{Ber}(p)$, where $\text{Ber}(p)$ denotes the Bernoulli noise
- Capacity region is unknown

Main Result

Objective

- Construct low-complexity coding schemes that achieve the capacity region of the WOM system
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$$R_1 < h(\delta) - h(p), \qquad R_2 < 1 - \delta - h(p).$$

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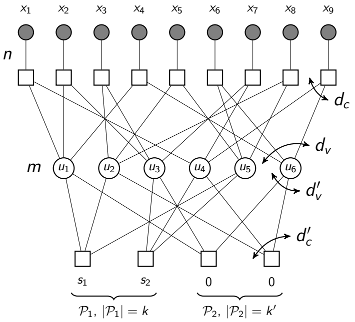
Idea

- Use compound LDGM/LDPC codes
- Encoding for second write is erasure quantization
- Use spatial coupling with message-passing

Compound LDGM/LDPC Codes



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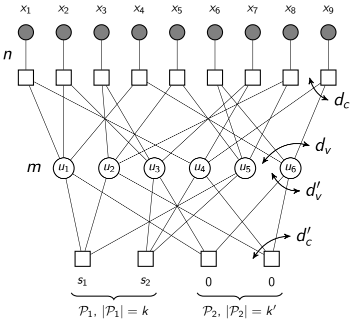

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Key Properties of Compound Codes

- a natural coset decomposition: $\mathcal{C} = \bigcup_{s^k \in \{0,1\}^k} \mathcal{C}(s^k)$
- achieves capacity over eras. chan. under MAP (when $m = n$)
- a good source code under optimal encoding
- a good channel code under optimal decoding

Good Code

"Good" source code

- Rate of the code is $R = 1 - h(\delta) + \varepsilon$
- When this code is used to optimally encode $\mathrm{Ber}(\frac{1}{2})$
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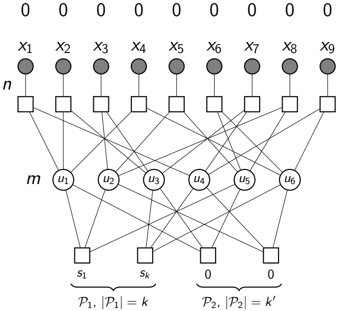
Coding Scheme for 2-write WOM: First Write

$$R_1 < h(\delta) - h(p)$$



$\frac{m-k-k'}{n} \approx 1 - h(\delta)$ $\frac{m-k'}{n} \approx 1 - h(p)$

- With message s^k, encode 0^n to x^n (Distortion $\approx \delta$)
- Store x^n

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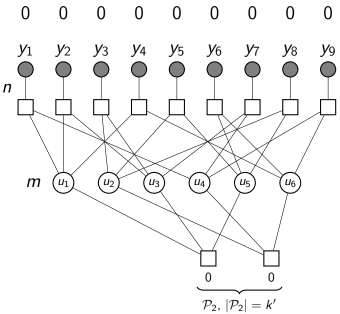

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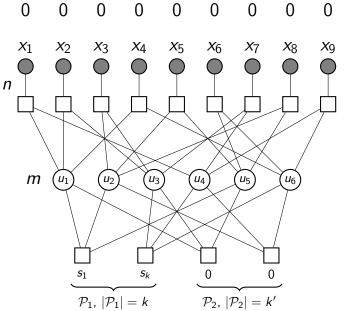

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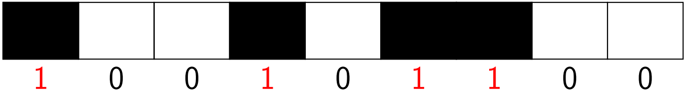
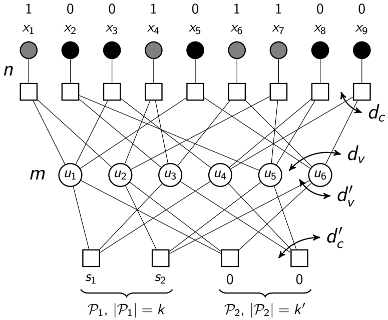
$$y_i = x_i \oplus \mathrm{Ber}(p)$$

- Dec. x^n and compute s^k
- $R_1 = \frac{k}{n} \approx h(\delta) - h(p)$

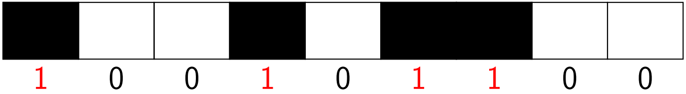
Coding Scheme for 2-write WOM: Second Write

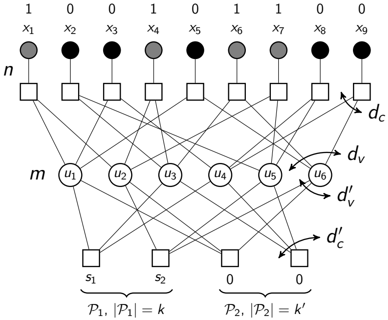




- Need to find a consistent codeword in $\mathcal{C}(s^k)$

Coding Scheme for 2-write WOM: Second Write





- Need to find a consistent codeword in $\mathcal{C}(s^k)$
- Closely related to Binary Erasure Quantization (BEQ)
- En Gad, Huang, Li and Bruck (ISIT 2015)

Binary Erasure Quantization

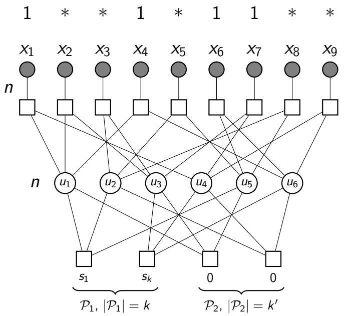
- Quantize a sequence in $\{0, 1, *\}^n$ to $x^n \in \mathcal{C} \subset \{0, 1\}^n$
  - 0's and 1's should match exactly
  - $*$'s can take either 0 or 1
- Can map the second write of 2-write WOM to BEQ
  - Map 0's to $*$'s and keep 1's
  - Quantize to codeword in $\mathcal{C}(s^k)$
- BEQ is the dual of decoding on binary erasure channel
  - Martinian and Yedidia (Allerton 2003)
  - Can quan. all seq. with erasure pattern $e^n \in \{0, 1\}^n$ to $\mathcal{C}$

    $$\Updownarrow$$

    Chan. dec. for $\mathcal{C}^{\perp}$ can correct all vectors with eras. $1^n \oplus e^n$
- Choose a good (dual) code $\mathcal{C}(s^k)$

Coding Scheme for 2-write WOM: Second Write

$$R_2 < 1 - \delta - h(p)$$

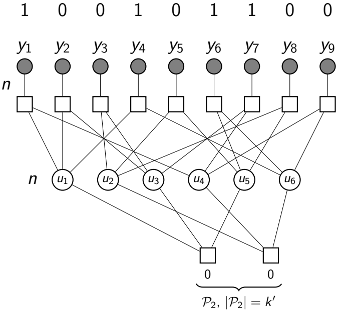


$\frac{n-k-k'}{n} \approx \delta$ $\quad$ $\frac{n-k'}{n} \approx 1 - h(p)$

- Change 0's to $*$'s
- With message s^k, encode seq. to $\mathcal{C}(s^k)$

Coding Scheme for 2-write WOM: Second Write

$$R_2 < 1 - \delta - h(p)$$



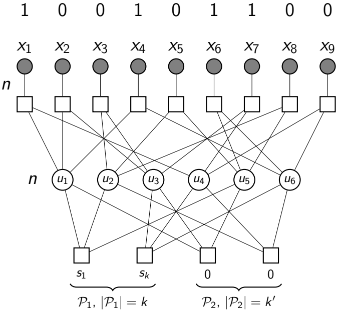
$$\frac{n-k'}{n} \approx 1 - h(p)$$

- Change 0's to $*$'s
- With message s^k, encode seq. to $\mathcal{C}(s^k)$
- Decoder has

$$y_i = x_i \oplus \mathrm{Ber}(p)$$

Coding Scheme for 2-write WOM: Second Write

$$R_2 < 1 - \delta - h(p)$$



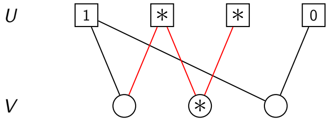
$\frac{n-k-k'}{n} \approx \delta$ $\qquad$ $\frac{n-k'}{n} \approx 1 - h(p)$

- Change 0's to $*$'s
- With message s^k, encode seq. to $\mathcal{C}(s^k)$
- Decoder has

$$y_i = x_i \oplus \mathrm{Ber}(p)$$

- Dec. x^n and compute s^k
- $R_2 = \frac{k}{n} \approx 1 - \delta - h(p)$

Iterative Erasure Quantization Algorithm

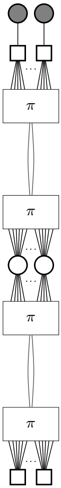


- Peeling type encoder

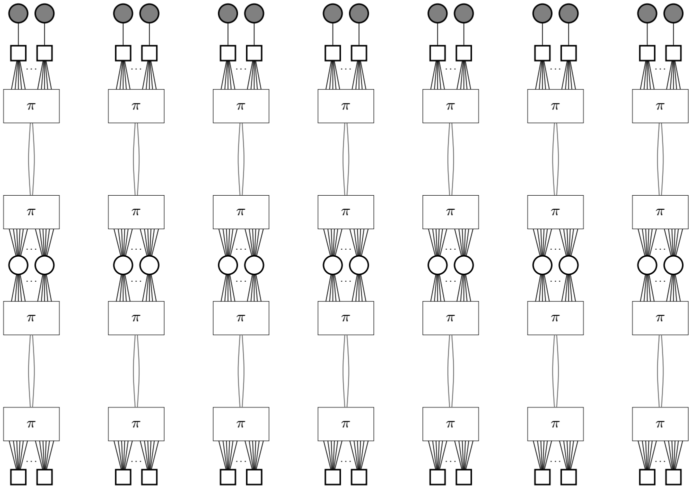
Remarks

- Need codes that are simultaneously good for channel/source coding and erasure quantization
- Use message-passing algorithms instead of optimal
- Use spatial-coupling for goodness of codes under message-passing

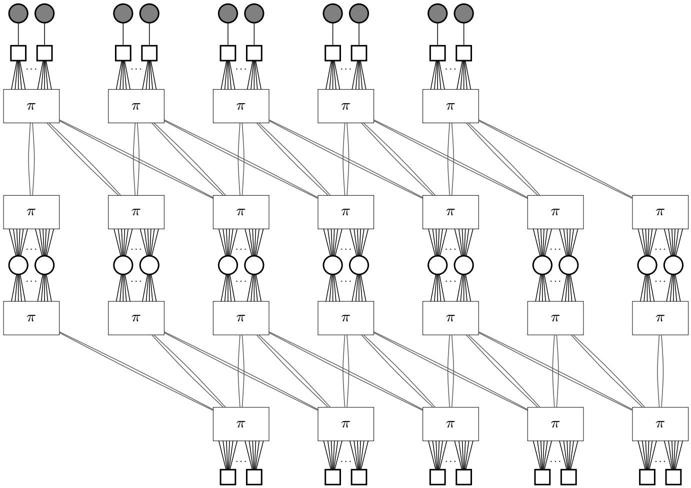
Spatially-Coupled Compound LDGM/LDPC Codes



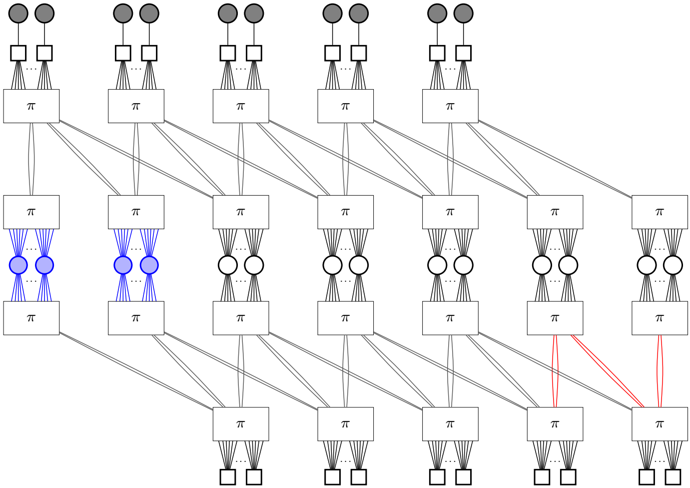
Spatially-Coupled Compound LDGM/LDPC Codes



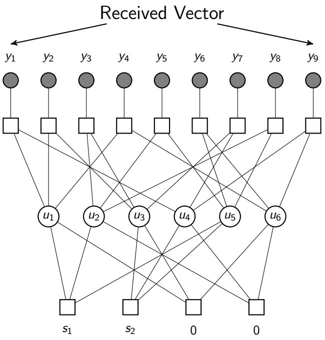
Spatially-Coupled Compound LDGM/LDPC Codes

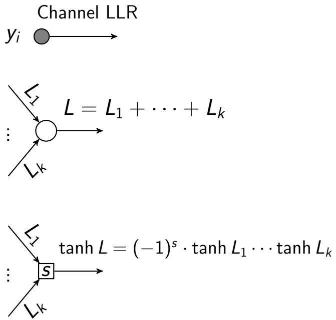


Spatially-Coupled Compound LDGM/LDPC Codes



Decoding in Spatially-Coupled Compound Codes





Remarks

- Standard message-passing algorithm
- Threshold saturation proven for SC compound codes on BEC
- Empirically observed for BMS channels

Numerical Results: Noiseless WOM

LDGM/LDPC (d_v, d_c, d'_v, d'_c)	δ^*	δ $w = 2$	δ $w = 3$	δ $w = 4$
$(3, 3, 3, 6)$	0.500	0.477	0.492	0.494
$(3, 3, 4, 6)$	0.333	0.294	0.324	0.326
$(3, 3, 5, 6)$	0.167	0.095	0.156	0.158
$(4, 4, 3, 6)$	0.500	0.461	0.491	0.492
$(4, 4, 4, 6)$	0.333	0.278	0.323	0.325
$(4, 4, 5, 6)$	0.167	0.086	0.155	0.159
$(5, 5, 3, 6)$	0.500	0.436	0.488	0.491
$(5, 5, 4, 6)$	0.333	0.260	0.320	0.324
$(5, 5, 5, 6)$	0.167	0.079	0.154	0.159

Remarks

- δ^* is the Shannon threshold
- $L = 30$, Single system length ≈ 24000

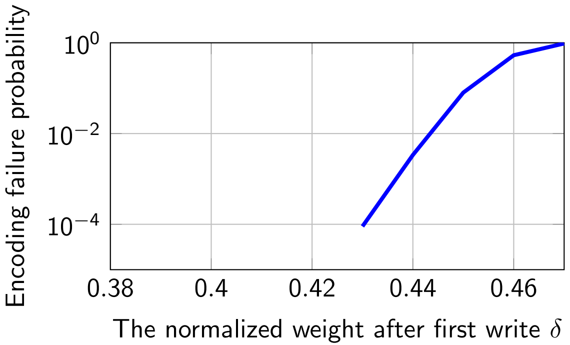
Numerical Results: WOM with Read Errors

LDGM/LDPC (d_v, d_c, d'_v, d'_c)	w	(δ^*, p^*)	(δ, p)
$(3, 3, 4, 6)$	3	$(0.333, 0.0615)$	$(0.321, 0.0585)$
$(3, 3, 4, 8)$	3	$(0.500, 0.0417)$	$(0.490, 0.0387)$
$(3, 3, 6, 8)$	4	$(0.250, 0.0724)$	$(0.239, 0.0684)$
$(4, 4, 4, 6)$	4	$(0.333, 0.0615)$	$(0.324, 0.0585)$
$(4, 4, 4, 8)$	4	$(0.500, 0.0417)$	$(0.492, 0.0387)$
$(4, 4, 6, 8)$	4	$(0.250, 0.0724)$	$(0.241, 0.0694)$

Remarks

- δ^* and p^* are the Shannon thresholds
- $L = 30$, Single system length ≈ 30000

Numerical Results: Small Blocklength



Remarks

- $(L, w) = (30, 3)$, Single system length 1200, Shannon threshold of 0.5
- A total of 10^5 were attempted to encode
- No failures for $\delta < 0.43$

Concluding Remarks

Conclusion

- Spatially-coupled compound codes achieve the capacity of 2-write systems
- Coupling structure is also crucial
  - to achieve optimum thresholds
  - for encoding to succeed

Multi-Write Systems

- Will BPGD work for multi-write systems?

Outline

- SC-LDPC Lattices [C1]
- SC-Compound Codes
  - Side-Information Problems [C2]
  - Coding for WOM [C3]
- Tools from Sparse graph coding for sparse recovery problems
  - Bipartite sparse graph for compressed sensing [C5]
  - Group testing*
  - Pattern matching*
- Uncoordinated multiple access
  - Universal schemes for massive uncoordinated multiple access [C4]
  - Optimal distributions for finite user multiple access*
- Coding for low latency requirements*

C1. A. Vem, Y. C. Huang, K. R. Narayanan and H. D. Pfister, "Multilevel lattices based on spatially-coupled LDPC codes with applications", in *Proc. IEEE. ISIT*, pp. 2336–2340, 2014.
C2. S. Kumar, A. Vem, K. R. Narayanan and H. D. Pfister, "Spatially-coupled codes for side-information problems", in *Proc. IEEE. ISIT*, pp. 516–520, 2014.
C3. S. Kumar, A. Vem, K. R. Narayanan and H. D. Pfister, "Spatially-coupled codes for write-once memories", in *Proc. Allerton. Conf*, pp. 125–131, 2015.
C4. A. Taghavi, A. Vem, J.-F. Chamberland and K. R. Narayanan "On the design of universal schemes for massive uncoordinated multiple access", in *Proc. IEEE. ISIT*, pp. 345–349, 2016.
C5. A. Vem, N. T. Janakiraman, and K. R. Narayanan, "Sub-linear time compressed sensing for support recovery using left and right regular sparse-graph codes", in *Proc. IEEE. ITW*, pp. 429–433, 2016.

*-To be submitted