

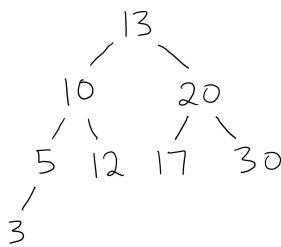
CSC148 - Investigating the Efficiency of Binary Search Trees

1. Recall our implementation of `BinarySearchTree.__contains__`:

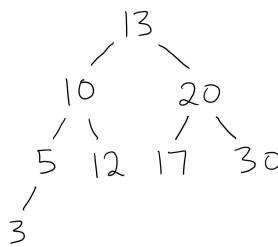
```
def __contains__(self, item: Any) -> bool:
    if self.is_empty():
        return False
    elif item == self._root:
        return True
    elif item < self._root:
        return self._left.__contains__(item)
    else:
        return self._right.__contains__(item)
```

Crucially, after comparing the item against the root of the BST, only *one* subtree is recursed into. To make sure you understand this idea, suppose we have the following BST, and perform three searches, for the items 25, 10, and 4, respectively. Circle all the values that are compared against the target item when we do each search.

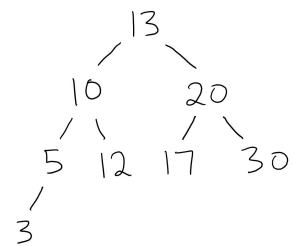
Search for 25:



Search for 10:



Search for 4:



Even though you should have circled a different number of tree values in each of the three diagrams, the number of values is always less than or equal to the *height* of each tree.

Because `BinarySearchTree.__contains__` only ever recurses into one subtree, each recursive call goes down one level in the tree. So at most *one item per level* is ever compared against the `item` being searched for.

So to better understand the running time of this algorithm, we need to investigate how the *height* of a BST can grow as we insert more items into it.

2. Recall that the *BST insertion algorithm* from last week's lab always inserts a new item into an "empty spot" at the bottom of a BST, so that the new item is a leaf of the tree.
- (a) Suppose we start with an empty BST, and then insert the items 1, 2, 3, 4, 5, 6 into the BST, in that order. Draw what the final BST looks like.
- (b) Suppose we start with an empty BST, and then insert the items 4, 2, 3, 5, 1, 6 into the BST, in that order. Draw what the final BST looks like.

3. Write down the *height* of each BST you drew in the previous question. (You can do so beside or underneath the diagram.)

4. Let n be a non-negative integer, and suppose we want a BST that contains the values $1, 2, \dots, n$. What is the *maximum* possible height of the BST?
5. Now suppose we want to find the *minimum* possible height of the BST. We're going to investigate this in an indirect way: for every height h , we're going to investigate what the maximum number of values can be stored in a binary tree of height n . Let's try to find a pattern.
- (a) Suppose we have a BST of height **0**. What's the maximum possible number of values in the BST? (Hint: there's only one kind of BST with height 0.)
- (b) Suppose we have a BST of height **1**. What's the maximum possible number of values in the BST? (Draw an example.)
- (c) Suppose we have a BST of height **2**. What's the maximum possible number of values in the BST? (Draw an example.)
- (d) Repeat for heights 3 and 4.
- (e) Suppose we have a BST of height **148**. What's the maximum possible number of values in the BST? (*Don't* draw an example, but find a pattern from your previous answers.)
- (f) Suppose we have a BST of height h and that contains n values. Write down an inequality of the form $n \leq \dots$ to relate n and h . (This is a generalization of your work so far.)
- (g) Finally, take your inequality from the previous part and isolate h . This produces the answer to our original question: what is the minimum height of a BST with n values?