· PS4: complete version is now on Quercus . TT3... more to follow - but not right now . TTZ: last TA is done, marks out this week! def pal-prefix (s: str) -> int:

"Return length of a longest prefix of s that

is a palindrome."

(len(s))

for p in range (len(s),0,-1): # p=n,n-1,...,1 it is_palind nome (s[o:p]): retam p $WC_{palprefix}(n) = ?$

· Upper bound: Goal: show RT(s) = ?

for all inputs s of size n Approach;
-overestomate, ignore early termination
-simplify as we go
-danger: we could overcount; check with
lower bound, do it more precisely if necessary · loop body takes time = f.p = f.n · loop iterates = n comes from total is $\leq \phi \cdot n \cdot n = \phi \cdot n^2$ $WC_{is-pal}(\rho) \in \Theta(\rho)$ $\longrightarrow WC(n) \in O(n^2)$ simplify to just "p" is okay!

· Lover bound: Goal: Show RT(s) > ! for at least one inputs of size in

input family: concrete input for each size. Approach: - find input family - underestimate, simplify as we go - danger: leave too much out - check against apper bound, redo as necessary with more precision. ROUGH WORK Trick: inputs that make pal-prefix iterate

many times are ones where many prefixes of s are not palindromes, e.g., S=abcd.... all letters are different! but: is-palind nome mus fast (constant) for such stwys ... · inputs that make is pal take a long time (like s = aaa ...a) iterate only once in pal-prefix... Insight: let s = aa...abaa...a

 $\begin{bmatrix} \frac{n}{2} \\ \frac{n}{2} \end{bmatrix} - 1$ e.g. n = 10: $s = \alpha \alpha \alpha \alpha \alpha b \alpha \alpha \alpha \alpha$

aaaa, baaaa,
$$\chi$$
 is-pal takes time $\frac{n}{2}$ is pal takes time $\frac{n}{2}$ -1

 χ is pal takes time \frac

$$\frac{n^2}{8} \implies WC(n) \in \Omega(n^2)$$
Last example

Lef twisty(n):

while n >1:

if n%2 == 0:

$$A(t) = \Omega(n^2)$$

#iterations)

ight: consider maltiple iterations...

even $\frac{n}{2}$ even $\frac{n}{4} = \frac{n}{4}$ upper box as