· PS4: complete version is on Quercus .TT3: more on that later... we are open to discussion but today it would take too much of the time we have! Last time... · is-palindrome (s): WC mutime is O(n) · def pal_prefix(s:str) -> int:
"Return length of a longest prefix of s
that is a palindrome." tor p in range (len(s), 0, -1): #p=n, n-1,, 1

if is_palindrome (s[0:p]):
return p

· WC pal-pretix (n) =? for all inputs s of site in · Upper bound: Goal: show RT(s) &O(_) Approach: - overestimate ignore early termination
- simplify as we go
- danger: overcount-upper bound
- danger: overcount-check with lower
is not tight - check with lower
bound -loop body takes time $\leq p \leq n = len(s)$ - # iterations $\leq n$ $-60 + al \leq n^2 \Rightarrow WC(n) \in O(n^2)$ · Lover bound: Goal: show RT(s) € S2(_) for some inpats of site in Approach: -simplify as we go

-underestimate -danger: leave out too much...
try to match appear bound - 's is not a palindrome — and s's prefixes are not palindromes: for loop makes many iterations Here, challenge: -> then, is-palindume nun faster - strings bad for is-pal (s is a palmd nome) are good for pal-prefix: for loop stops early... Insight: want input s that is almost a paladon $S = \underbrace{\alpha \alpha \cdots \alpha}_{\begin{bmatrix} \frac{1}{2} \end{bmatrix} - 1} b \underbrace{\alpha \alpha \cdots \alpha}_{\begin{bmatrix} \frac{1}{2} \end{bmatrix} - 1}$

$$\frac{a a a a a b a a a a}{z \cdot z \cdot pa} = \frac{r}{z \cdot pa} = \frac{r}{z} - \frac{r}{z}$$

$$\frac{n}{z \cdot pa} = \frac{r}{z \cdot pa} = \frac{r}{z} - \frac{r}{z}$$

$$\frac{n}{z \cdot pa} = \frac{r}{z} + \frac{r}{z} + \frac{r}{z}$$

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$$\frac{n}{z} = \frac{r}{z} + \frac{r}{z}$$

prefix that
$$\leq j$$
does not $j=1$
 $w(n) \in \Omega(n^2)$
 $j=p-[\frac{n}{2}]$

Last example... def twisty (n: out) -> None: runtime & O(# iterations) insight: trace multiple iterations