

# TODAY!

- average-case analysis
- intro. to graphs

## Average-case analysis

WC(n): worst-case

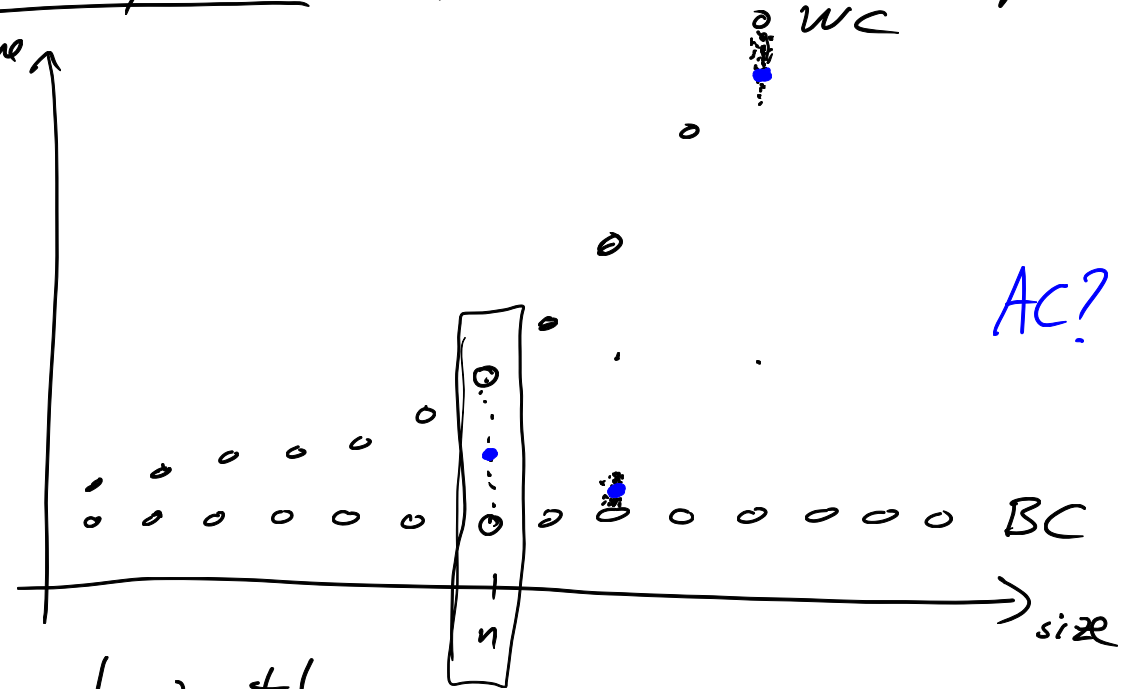
BC(n): best-case

AC(n): average-case

### NOTES:

- NO distinction between WC, BC, AC for algorithms where  $n$  is the only input!

RT(x): running time on input  $x$



• AC(n)  $\neq \frac{BC(n) + WC(n)}{2}$

$AC(n)$  = average of  $RT(x)$  over all inputs  $x$  of size  $n$

Generally, start by defining

$\mathcal{I}_n = \{ \text{all inputs of size } n \}$

input of size  $n$  → another input of size  $n$

imagine  $\mathcal{I}_n = \{x_0, x_1, x_2, \dots, x_m\}$

$$\frac{RT(x_0) + RT(x_1) + RT(x_2) + \dots + RT(x_m)}{m+1}$$

$$AC(n) = \frac{1}{|\mathcal{I}_n|} \sum_{x \in \mathcal{I}_n} RT(x)$$

$m+1$  → runtime on input  $x$   
 → add up over all  $x \in \mathcal{I}_n$   
 → size of  $\mathcal{I}_n$ , i.e., how many inputs of size  $n$  we have

Ex: def search ( $L$ : list,  $x$ : int)  $\rightarrow$  bool:  
for item in  $L$ :  
    if item ==  $x$ :  
        return True  
return False

$\left. \begin{array}{l} \text{for item in } L: \\ \text{if item == } x: \\ \text{return True} \end{array} \right\} \begin{array}{l} WC(n) \in \Theta(n) \\ BC(n) \in \Theta(1) \end{array}$

One input to algo consists of list  $L$  and int  $x$   
 $(L, x)$ .

[For WC lower bound, use input family]  
 $L = [1, 2, \dots, n]$ ,  $x = 0$ .

Q: How to define  $\mathcal{I}_n$ ?

What does "all possible inputs" really mean?

Key insight: need at least one input for each behaviour of the algorithm

For our example, this means inputs where  $x$  appears in  $L$  at each possible location.

• Idea 1:

$$\mathcal{I}_n = \{ (L, x) \mid L \text{ is a permutation of } [1, 2, \dots, n] \text{ and } \underline{x=1} \}$$

Permutation? e.g.  $[1, 2, 3], [1, 3, 2], [2, 1, 3], [2, 3, 1], [3, 1, 2], [3, 2, 1]$  for  $n=3$

Details  $\rightarrow$  see notes...

• Idea 2:

$$\mathcal{I}'_n = \left\{ (L, x) \mid L = [1, 2, \dots, n], x \in \{\underline{0}, 1, 2, \dots, n\} \right\}$$

Note: this includes one possibility  $([1, 2, \dots, n], 0)$  with  $x$  not in  $L$ .

$$\begin{aligned} AC(n) &= \frac{1}{|\mathcal{I}'_n|} \sum_{(L, x) \in \mathcal{I}'_n} RT(L, x) \\ &= \frac{1}{n+1} \left( RT([1, 2, \dots, n], 0) + \sum_{x=1}^n RT([1, \dots, n], x) \right) \\ &= \frac{1}{n+1} \left( (n+1) + \sum_{x=1}^n x \right) \\ &= \dots \text{ arithmetic } \dots \end{aligned}$$

• Idea 3:  $\mathcal{I}_n'' = \{(L, x) \mid L = [1, 2, \dots, n], x \in \{1, 2, \dots, n, \dots, n^2\}\}$

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Graphs — prove a few simple theorems about graphs

[1] Prove  $\forall G = (V, E), |E| \leq \frac{|V|(|V|-1)}{2}$

NOTES:

" $\forall G = (V, E)$ " introduces 3 related variables

- graph  $G$
- vertex set ( $V$ ) of  $G$
- edge set ( $E$ ) of  $G$

$$V = \{a, b, c, d\}$$

$$E = \{\{a, b\}, \{a, c\}, \{b, d\}\}$$

Proof:  $\therefore$  every elem. in  $E$  is a subset of  $V$

of size 2 (by def.)

- from worksheet #10 — there are exactly  $\frac{|V|(|V|-1)}{2}$  many such subsets possible.  $\square$