Learning Objectives

By the end of this worksheet, you will:

- Prove statements using the definition of Big-O and its negation.
- Represent constant functions in Big-O expressions.
- Understand and use the definition of Omega and Theta to compare functions.

For your reference, here is the formal definition of Big-O:

$$g \in \mathcal{O}(f): \exists c, n_0 \in \mathbb{R}^+, \ \forall n \in \mathbb{N}, \ n \ge n_0 \Rightarrow g(n) \le cf(n)$$
 where $f, g: \mathbb{N} \to \mathbb{R}^{\ge 0}$

- 1. Constant functions. As we discussed in class, constant functions, like f(n) = 100, will play an important role in our analysis of running time next week. For now let's get comfortable with the notation.
 - (a) Let $g: \mathbb{N} \to \mathbb{R}^{\geq 0}$. Show how to express the statement $g \in \mathcal{O}(1)$ by expanding the definition of Big-O.

(b) Prove that $100 + \frac{77}{n+1} \in \mathcal{O}(1)$.

Note: this proof isn't too mathematically complex; treat this as another exercise in making sure you understand the definition of Big-O.

Hint: one algebraic property of inequalities is that $\forall x, y \in \mathbb{R}^+, \ x \geq y \Rightarrow \frac{1}{x} \leq \frac{1}{y}$.

¹Remember that we often abbreviate Big-O expressions to just show the function bodies. " $\mathcal{O}(1)$ " is really shorthand for " $\mathcal{O}(f)$, where f is the constant function f(n) = 1."

2. **Omega**. Recall that we can think of Big-O notation as describing an *upper bound* on the rate of growth of a function: saying " $g \in \mathcal{O}(f)$ " is like saying "g grows at most as fast as f." Sometimes we care just as much about a *lower bound* on the rate of growth and for this, we have the symbol Ω (the Greek letter Omega), which is defined analogously to Big-O:

$$g \in \Omega(f): \exists c, n_0 \in \mathbb{R}^+, \ \forall n \in \mathbb{N}, \ n \ge n_0 \Rightarrow g(n) \ge cf(n)$$
 where $f, g : \mathbb{N} \to \mathbb{R}^{\ge 0}$

Using this definition, prove that for all $f, g : \mathbb{N} \to \mathbb{R}^{\geq 0}$, if $g \in \mathcal{O}(f)$, then $f \in \Omega(g)$.

3. **Theta**. Both Big-O and Omega are limited in the same way as inequalities on numbers. " $2 \le 10^{10}$ " is a true statement, but not very insightful; similarly, " $n+1 \in \mathcal{O}(n^{10})$ " and " $2^n+n^2 \in \Omega(n)$ " are both true, but not very precise.

Our final piece of asymptotic notation is Θ (the Greek letter Theta), which we define as:

$$g \in \Theta(f): g \in \mathcal{O}(f) \land g \in \Omega(f)$$
 where $f, g : \mathbb{N} \to \mathbb{R}^{\geq 0}$

Or equivalently,

$$g \in \Theta(f): \exists c_1, c_2, n_0 \in \mathbb{R}^+, \ \forall n \in \mathbb{N}, \ n \ge n_0 \Rightarrow c_1 f(n) \le g(n) \le c_2 f(n)$$
 where $f, g : \mathbb{N} \to \mathbb{R}^{\ge 0}$

When we write $g \in \Theta(f)$, what we mean is "g grows at most as quickly as f and g grows at least as quickly as f"—in other words, that f and g have the same rate of growth. In this case, we call f a **tight bound** on g, since g is essentially squeezed between constant multiples of f.

Prove that for all functions $g: \mathbb{N} \to \mathbb{R}^{\geq 0}$, and all numbers $a \in \mathbb{R}^{\geq 0}$, if $g \in \Omega(1)$, then $a + g \in \Theta(g)$.² (Or in other words, for such functions g, shifting them by a constant amount does not change their "Theta" bound.)

²Here we use a + g to denote the function g_1 defined as $g_1(n) = a + g(n)$ for all $n \in \mathbb{N}$.

- 4. **Negating Big-O**. So far, we have only looked at proving that a function *is* Big-O of another function. In this question, we'll investigate what it means to show that a function *isn't* Big-O of another.
 - (a) Express the statement $g \notin \mathcal{O}(f)$ in predicate logic, using the expanded definition of Big-O. (As usual, simplify so that all negations are pushed as far "inside" as possible.)

(b) Prove that for all positive real numbers a and b, if a > b then $n^a \notin \mathcal{O}(n^b)$.