



Name : Rasika Santosh Thakur

Class : S. Y. A. Tech

Branch : AI & DS 'B'

Roll No : 71

Subject : Operating Systems

Assignment on Unit III

Q.1 Consider following snapshot of System

Process	Allocation				Max				Available	Need			
	A	B	C	D	A	B	C	D		A	B	C	D
P ₀	3	0	1	4	5	1	1	7	0, 3, 0, 1	2	1	0	3
P ₁	2	2	1	0	3	2	1	1		1	0	0	1
P ₂	3	1	2	1	3	3	2	1		0	2	0	0
P ₃	0	5	1	0	4	6	1	2		4	1	0	2
P ₄	4	2	1	2	6	3	2	5		2	1	1	8

Using Banker's Algorithm determine whether or not each of following state is safe. Illustrate in which order process may completed. Otherwise illustrate why state is unsafe.

1) Available = (0, 3, 0, 1)

2) Available = (1, 0, 0, 2)



→ a) available (0, 3, 0, 1)

work = available = 0, 3, 0, 1

Each time when process is in safe state, work is updated.

Need \leq Work then Work = Work + Allocation

P_0 = Need = 2, 1, 0, 3 & Work = 0, 3, 0, 1
Need > Work

[Unsafe]

P_1 = Need = 1, 0, 0, 1 & Work = 0, 3, 0, 1
Need > Work

[Unsafe]

P_2 = Need = 0, 2, 0, 0 & Work = 0, 3, 0, 1
Need \leq Work

[Safe]

Work = Work + Allocation

= 0, 3, 0, 1 + 3, 1, 2, 1

= 3, 4, 2, 2

Finish [P_2] = true

P_3 = Need = 4, 1, 0, 2 & Work = 3, 4, 2, 2
Need > Work

[Unsafe state]

P_4 = Need = 2, 1, 1, 3 & Work = 3, 4, 2, 2
Need > Work [Unsafe state]

P_5 = Need = 1, 0, 0, 1 & Work = 3, 4, 2, 2
Need \leq Work [Safe state]

Work = 3, 4, 2, 2 + 2, 2, 1, 0

= 5, 6, 3, 2



As, Need of Resource type 'D' of process 'P₁' & 'P₂' is greater than need, it will never complete its execution. So, system is in unsafe state.

h) Available = (1, 0, 0, 2)

Work = available = 1002

Need \leq Work then

Work = Work + Allocation

P_0 : Need = 2002 & Work = 1002

Need $\not\leq$ Work

[Unsafe state]

Work = Work + Allocation

= 1002 + 2210

P_1 : Need = 1001 & Work = 1002

Need \leq Work [Safe state]

Work = Work + Allocation

= 1002 + 2210

= 3212

Finish [P_1] = true



$$P_2 = \text{Need} = 0200 \quad \& \quad \text{Work} = 3212$$

$$\text{Need} < \text{Work}$$

[Safe State]

$$\text{Work} = \text{Work} + \text{Allocation}$$

$$= 3212 + 3121$$

$$= 6333$$

$$\text{Finish}[P_2] = \text{true}$$

$$P_3 = \text{Need} = 4102 \quad \& \quad \text{Work} = 6333$$

$$\text{Need} < \text{Work}$$

[Safe State]

$$\text{Work} = \text{Work} + \text{Allocation}$$

$$= 6333 + 0510$$

$$= 6843$$

$$\text{Finish}[P_3] = \text{true}$$

$$P_4 = \text{Need} = 2113 \quad \& \quad \text{Work} = 6843$$

$$\text{Need} < \text{Work}$$

[Safe State]

$$\text{Work} = \text{Work} + \text{Allocation}$$

$$= 6843 + 4212$$

$$= 11055$$

$$\text{Finish}[P_4] = \text{true}$$

$$P_0 = \text{Need} = 2103 \quad \& \quad \text{Work} = 11055$$

$$\text{Need} < \text{Work}$$

[Safe State]



$$\begin{aligned}\text{Work} &= \text{Work} + \text{Allocation} \\ &= 10 \ 10 \ 5 \ 5 + 30 \ 14 \\ &= 23 \ 20 \ 6 \ 9\end{aligned}$$

$\text{Finish}[P_0] = \text{true}$

$\therefore \text{Safe Sequence} = P_1, P_2, P_3, P_4, P_0$

Q.2 Write a short note on Semaphore

-
- Semaphore is a hardware solution.
 - It is written or given to critical section problem.
 - It is a synchronization tool used in concurrent programming to access control of shared resources among multiple threads & process.
 - Main purpose is to access provide synchronization & mutual exclusion in multi-threaded environment.
 - It prevents race condition & data corruption.
 - It has 2 fundamental operations:
 - 1] Wait &
 - 2] Signal
 - Wait decrements semaphore value to stop until value becomes non-negative where as signal increments semaphore value to unblock process.



- There are two types of Semaphore

i) Binary Semaphore

- Also known as mutex & it has two states
0 & 1 to control critical section

0 - locked

1 - unlocked

ii) Counting Semaphore

- It is used when multiple instances of
resources can execute it's critical section
based on number of instances available.