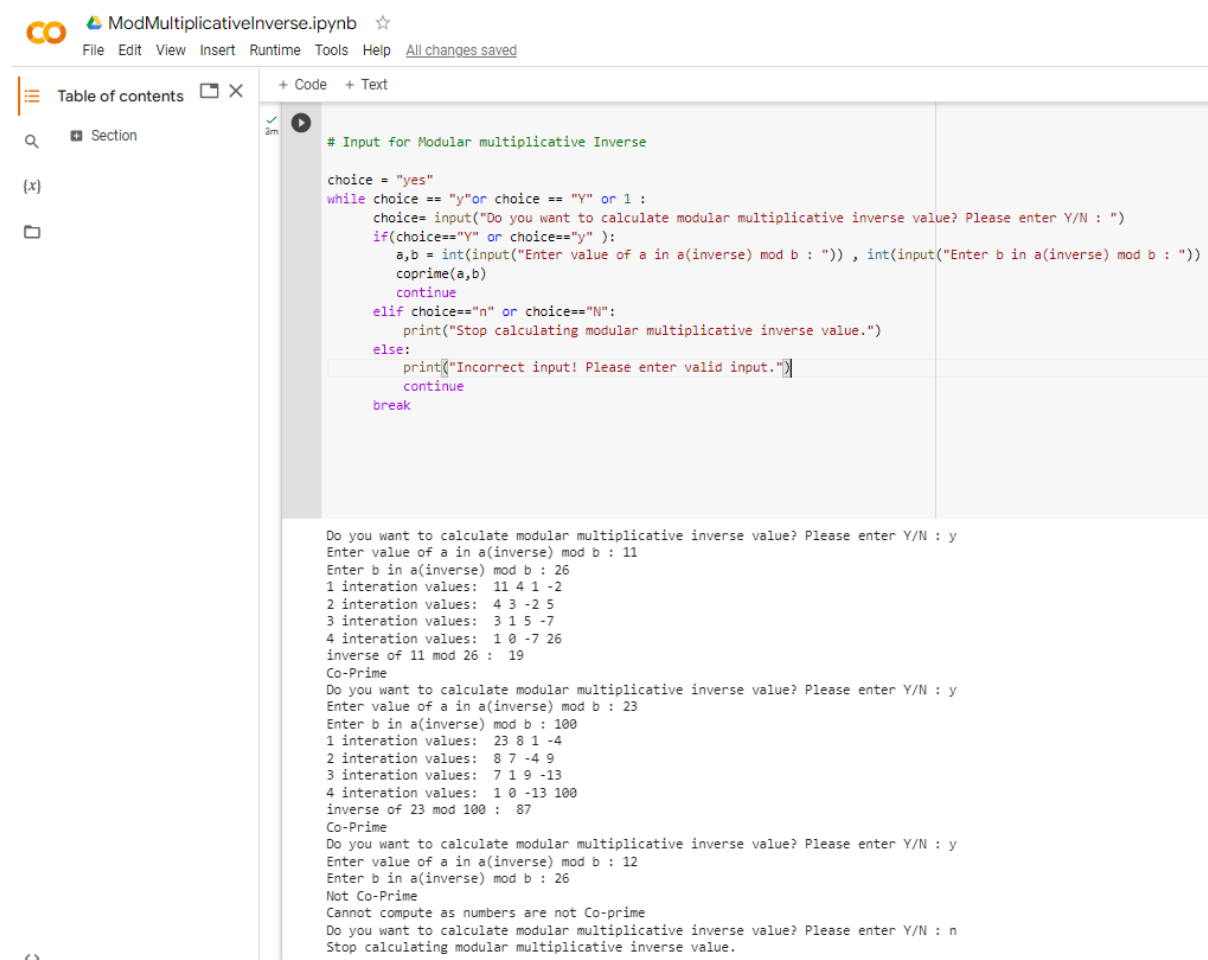


1) Modular multiplicative inverse of a number using extended Euclidian Algorithm

Verify following using program:

- i. $11^{-1} \bmod 26 = 19$
- ii. $23^{-1} \bmod 100 = 87$
- iii. $12^{-1} \bmod 26 =$ The gcd $(26, 12) = 2 \neq 1$, which means there is no multiplicative inverse for 12 in \mathbb{Z}_{26} .

Example 1, 2 and 3: outputs



```

# Input for Modular multiplicative Inverse

choice = "yes"
while choice == "y" or choice == "Y" or 1 :
    choice= input("Do you want to calculate modular multiplicative inverse value? Please enter Y/N : ")
    if(choice=="Y" or choice=="y" ):
        a,b = int(input("Enter value of a in a(inverse) mod b : ")), int(input("Enter b in a(inverse) mod b : "))
        coprime(a,b)
        continue
    elif choice=="n" or choice=="N":
        print("Stop calculating modular multiplicative inverse value.")
    else:
        print("Incorrect input! Please enter valid input.")
        continue
    break

Do you want to calculate modular multiplicative inverse value? Please enter Y/N : y
Enter value of a in a(inverse) mod b : 11
Enter b in a(inverse) mod b : 26
1 iteration values: 11 4 1 -2
2 iteration values: 4 3 -2 5
3 iteration values: 3 1 5 -7
4 iteration values: 1 0 -7 26
inverse of 11 mod 26 : 19
Co-Prime
Do you want to calculate modular multiplicative inverse value? Please enter Y/N : y
Enter value of a in a(inverse) mod b : 23
Enter b in a(inverse) mod b : 100
1 iteration values: 23 8 1 -4
2 iteration values: 8 7 -4 9
3 iteration values: 7 1 9 -13
4 iteration values: 1 0 -13 100
inverse of 23 mod 100 : 87
Co-Prime
Do you want to calculate modular multiplicative inverse value? Please enter Y/N : y
Enter value of a in a(inverse) mod b : 12
Enter b in a(inverse) mod b : 26
Not Co-Prime
Cannot compute as numbers are not Co-prime
Do you want to calculate modular multiplicative inverse value? Please enter Y/N : n
Stop calculating modular multiplicative inverse value.

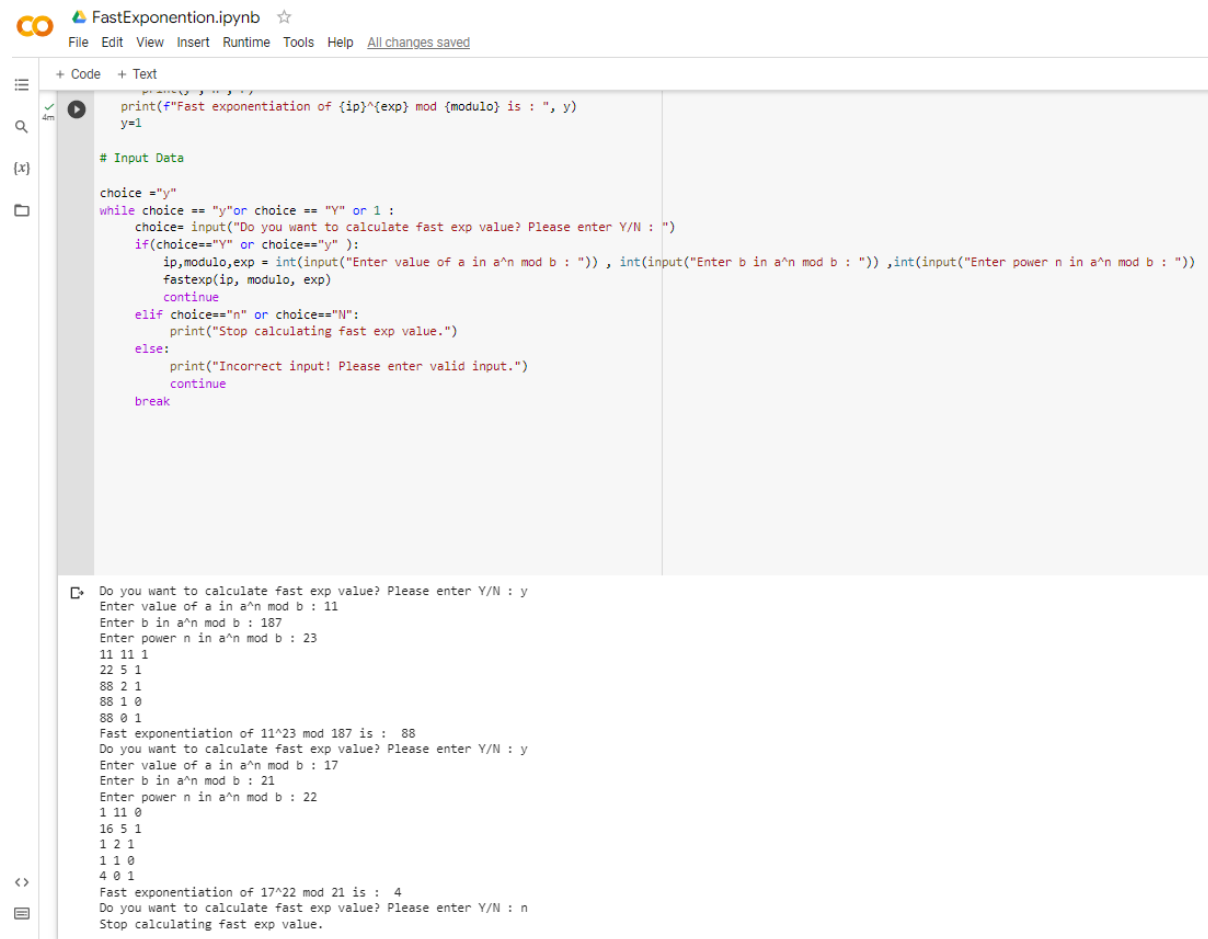
```

2) Fast exponentiation in congruences

Verify following using program:

- i. $11^{23} \bmod 187 = 88$
- ii. $17^{22} \bmod 21 = 4$

Example 1 and 2 outputs:



```
FastExponention.ipynb ☆
File Edit View Insert Runtime Tools Help All changes saved

+ Code + Text

print(f"Fast exponentiation of {ip}^{exp} mod {modulo} is : ", y)
y=1

# Input Data
choice = "y"
while choice == "y" or choice == "Y" or 1 :
    choice= input("Do you want to calculate fast exp value? Please enter Y/N : ")
    if(choice=="y" or choice=="Y" ):
        ip,module,exp = int(input("Enter value of a in a^n mod b : ")), int(input("Enter b in a^n mod b : ")),int(input("Enter power n in a^n mod b : "))
        fastexp(ip, modulo, exp)
        continue
    elif choice=="n" or choice=="N":
        print("Stop calculating fast exp value.")
    else:
        print("Incorrect input! Please enter valid input.")
        continue
    break

Do you want to calculate fast exp value? Please enter Y/N : y
Enter value of a in a^n mod b : 11
Enter b in a^n mod b : 187
Enter power n in a^n mod b : 23
11 11 1
22 5 1
88 2 1
88 1 0
88 0 1
Fast exponentiation of 11^23 mod 187 is : 88
Do you want to calculate fast exp value? Please enter Y/N : y
Enter value of a in a^n mod b : 17
Enter b in a^n mod b : 21
Enter power n in a^n mod b : 22
1 11 0
16 5 1
1 2 1
1 1 0
4 0 1
Fast exponentiation of 17^22 mod 21 is : 4
Do you want to calculate fast exp value? Please enter Y/N : n
Stop calculating fast exp value.
```

3) Public key cryptosystems

Diffie Hellman Key exchange Algorithm

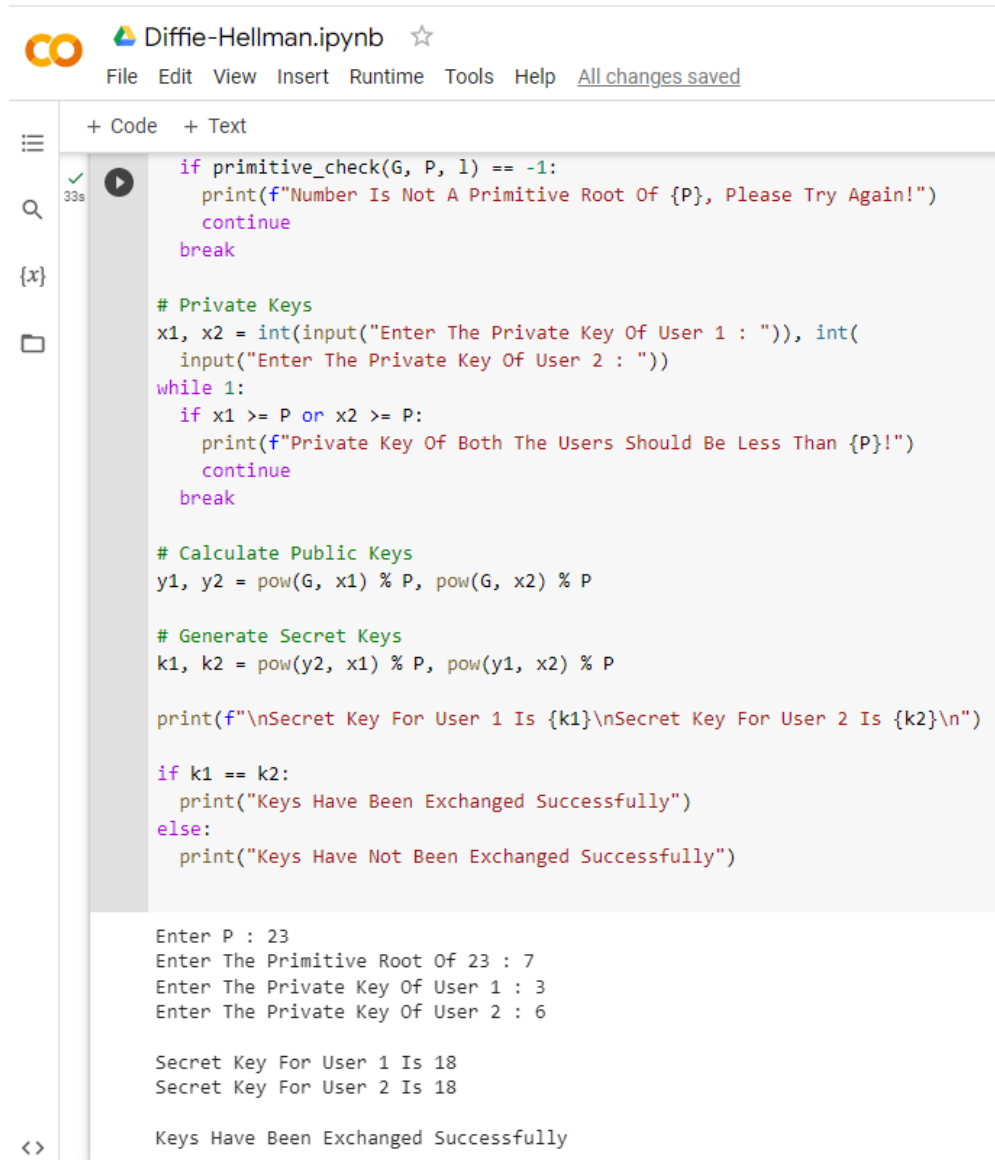
Verify following two examples using program:

Example 1: Let $g = 7$ and $p = 23$.

The steps are as follows:

1. Alice chooses $x = 3$ and calculates $R1 = 7^3 \bmod 23 = 21$.
2. Bob chooses $y = 6$ and calculates $R2 = 7^6 \bmod 23 = 4$.
3. Alice sends the number 21 to Bob.

4. Bob sends the number 4 to Alice.
 5. Alice calculates the symmetric key $K = 4^3 \bmod 23 = 18$.
 6. Bob calculates the symmetric key $K = 21^6 \bmod 23 = 18$.
- The value of K is the same for both Alice and Bob; $g^{xy} \bmod p = 7^{18} \bmod 23 = 18$.



```

CO Diffie-Hellman.ipynb ☆
File Edit View Insert Runtime Tools Help All changes saved

+ Code + Text

if primitive_check(G, P, 1) == -1:
    print(f"Number Is Not A Primitive Root Of {P}, Please Try Again!")
    continue
break

# Private Keys
x1, x2 = int(input("Enter The Private Key Of User 1 : ")), int(
    input("Enter The Private Key Of User 2 : "))
while 1:
    if x1 >= P or x2 >= P:
        print(f"Private Key Of Both The Users Should Be Less Than {P}!")
        continue
    break

# Calculate Public Keys
y1, y2 = pow(G, x1) % P, pow(G, x2) % P

# Generate Secret Keys
k1, k2 = pow(y2, x1) % P, pow(y1, x2) % P

print(f"\nSecret Key For User 1 Is {k1}\nSecret Key For User 2 Is {k2}\n")

if k1 == k2:
    print("Keys Have Been Exchanged Successfully")
else:
    print("Keys Have Not Been Exchanged Successfully")

Enter P : 23
Enter The Primitive Root Of 23 : 7
Enter The Private Key Of User 1 : 3
Enter The Private Key Of User 2 : 6

Secret Key For User 1 Is 18
Secret Key For User 2 Is 18

Keys Have Been Exchanged Successfully

```

Example 2: Alice and Bob have chosen prime value $q = 17$ and primitive root $= 5$. If Alice's secret key is 4 and Bob's secret key is 6.

Both Alice and Bob calculate the value of their public key and exchange with each other.

Public key of Alice

- $= 5^{\text{private key of Alice}} \bmod 17$
- $= 5^4 \bmod 17$
- $= 13$

Public key of Bob

- = 5 private key of Bob mod 17
- = $5^6 \bmod 17$
- = 2

Both the parties calculate the value of secret key at their respective side.

Secret key obtained by Alice

- = 2 private key of Alice mod 7
- = $2^4 \bmod 17$
- = 16

Secret key obtained by Bob

- = 13 private key of Bob mod 7
- = $13^6 \bmod 17$
- = 16

Finally, both the parties obtain the same value of secret key.

The value of common secret key = 16.

```

+ Code + Text
while 1:
    G = int(input(f"Enter The Primitive Root Of {P} : "))
    if primitive_check(G, P, 1) == -1:
        print(f"Number Is Not A Primitive Root Of {P}, Please Try Again!")
        continue
    break

# Private Keys
x1, x2 = int(input("Enter The Private Key Of User 1 : ")), int(
    input("Enter The Private Key Of User 2 : "))
while 1:
    if x1 >= P or x2 >= P:
        print(f"Private Key Of Both The Users Should Be Less Than {P}!")
        continue
    break

# Calculate Public Keys
y1, y2 = pow(G, x1) % P, pow(G, x2) % P

# Generate Secret Keys
k1, k2 = pow(y2, x1) % P, pow(y1, x2) % P

print(f"\nSecret Key For User 1 Is {k1}\nSecret Key For User 2 Is {k2}\n")

if k1 == k2:
    print("Keys Have Been Exchanged Successfully")
else:
    print("Keys Have Not Been Exchanged Successfully")

Enter P : 17
Enter The Primitive Root Of 17 : 5
Enter The Private Key Of User 1 : 4
Enter The Private Key Of User 2 : 6

Secret Key For User 1 Is 16
Secret Key For User 2 Is 16

Keys Have Been Exchanged Successfully

```

RSA Algorithm

Verify following two examples using program:

Example 1: Bob chooses 7 and 11 as p and q and calculates $n = 7 \times 11 = 77$. The value of $\phi(n) = (7 - 1) \times (11 - 1)$ or 60. Now he chooses two exponents, e and d, from Z_{60}^* . If he chooses e to be 13, then d is 37.

Now imagine that Alice wants to send the plaintext 5 to Bob. She uses the public exponent 13 to encrypt 5.

Plaintext: 5	$C = 5^{13} = 26 \text{ mod } 77$	Ciphertext: 26
--------------	-----------------------------------	----------------

Bob receives the ciphertext 26 and uses the private key 37 to decipher the ciphertext:

Ciphertext: 26	$P = 26^{37} = 5 \text{ mod } 77$	Plaintext: 5
----------------	-----------------------------------	--------------

The plaintext 5 sent by Alice is received as plaintext 5 by Bob.

The screenshot shows a Jupyter Notebook interface with the title 'RSA-method.ipynb'. The notebook contains a single code cell with the following Python code:

```

if private_exponent(p,q,e)==-1 :
    print("Please enter correct value of e")
    continue
break

```

Below the code cell, the notebook displays the output of the program, which includes prompts for user input and the results of the RSA calculations:

```

Enter 1st prime number p : 7
Enter 2nd prime number q : 11
Enter public exponent: 13
Co-Prime
1 iteration values: 13 8 1 -4
2 iteration values: 8 5 -4 5
3 iteration values: 5 3 5 -9
4 iteration values: 3 2 -9 14
5 iteration values: 2 1 14 -23
6 iteration values: 1 0 -23 60
inverse of 13 mod 60 : 37
private exponent d : 37
Enter plaintext: 5
5 6 1
5 3 0
45 1 1
26 0 1
ciphertext: 26
26 18 1
26 9 0
45 4 1
45 2 0
45 1 0
5 0 1
plaintext: 5

```

Example 2: Select two prime no's. Suppose $P = 53$ and $Q = 59$. Now first part of the public key: $n = P \cdot Q = 3127$.

Now we are ready with our – Public Key ($n = 3127$ and $e = 23$) and Private Key ($d = 1967$)

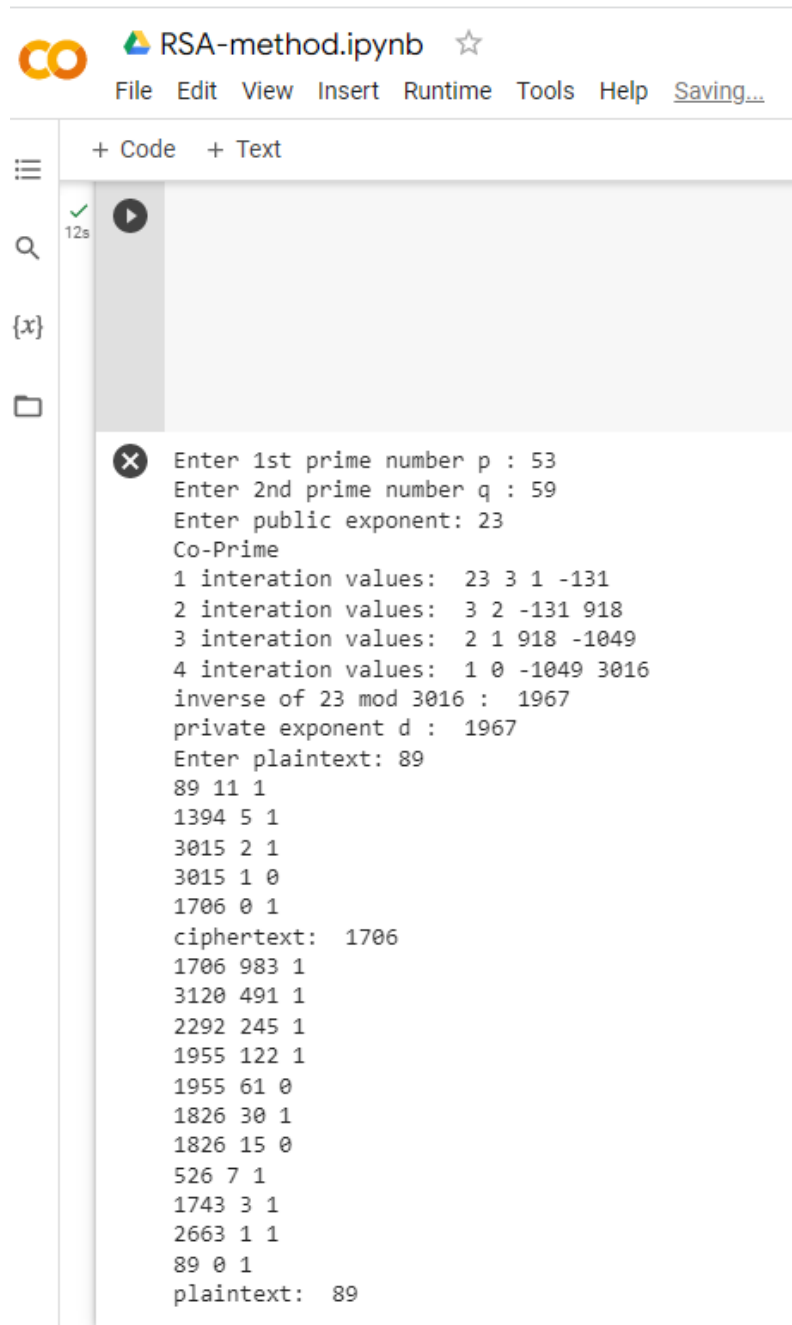
Now we will encrypt “HI”: Convert letters to numbers: $H = 8$ and $I = 9$.

Thus, **Encrypted Data $c = 89^{e \bmod n}$** .

Thus, Encrypted Data comes out to be 1706.

Now we will decrypt **1706: Decrypted Data $= c^{d \bmod n}$** .

Thus, our Encrypted Data comes out to be 89. $8 = H$ and $I = 9$ i.e., “HI”.



```
Enter 1st prime number p : 53
Enter 2nd prime number q : 59
Enter public exponent: 23
Co-Prime
1 iteration values: 23 3 1 -131
2 iteration values: 3 2 -131 918
3 iteration values: 2 1 918 -1049
4 iteration values: 1 0 -1049 3016
inverse of 23 mod 3016 : 1967
private exponent d : 1967
Enter plaintext: 89
89 11 1
1394 5 1
3015 2 1
3015 1 0
1706 0 1
ciphertext: 1706
1706 983 1
3120 491 1
2292 245 1
1955 122 1
1955 61 0
1826 30 1
1826 15 0
526 7 1
1743 3 1
2663 1 1
89 0 1
plaintext: 89
```

Symmetric Key Cryptosystems

4) Simplified DES

Output 1 for plaintext: 654 and Key value = 732. Ciphertext: 1b 71 fe and plaintext received is 654.

