

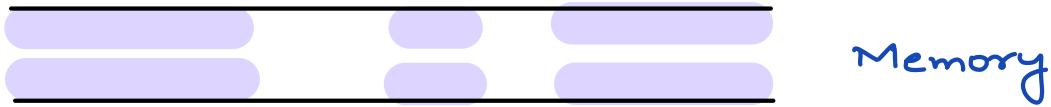


Good  
Evening

## \* Agenda for Today

01. Basics of LinkedList
02. Accessing of element
03. Searching for a value
04. Insertion in LL
05. Deletion in LL
06. Reverse a LinkedList
07. Check if palindrome

## \* Arrays vs Linkedlist



Array → contiguous memory allocations

Linkedlist → It can help in utilising all the free memory available

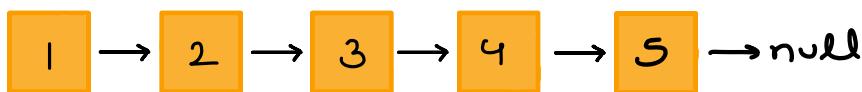
↳ Insertion / Deletion in LL is better



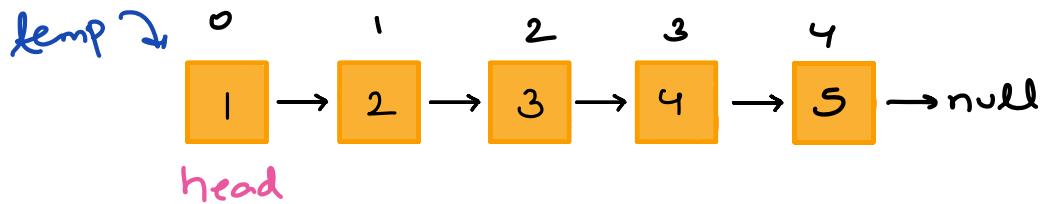
as compared to Arrays

```
class Node {  
    int data;  
    Node next;  
  
    Node (int d){  
        data = d;  
        next = null;  
    }  
}
```

3



$\text{arr} = \{ 1 \ 2 \ 3 \ 4 \ 5 \}$



### O1. Access a particular index

\* Access  $K^{\text{th}}$  idx element in array =  $\text{arr}[k]$  TC:O(1)

Node temp = head;

for ( $i=1$  to  $k$ )

temp = temp.next;

TC: O(k)

}

return temp.data;

O1. Never move the head

O2 Forward =  $\text{temp} = \text{temp}.next;$

O3. Access data =  $\text{temp}.data;$

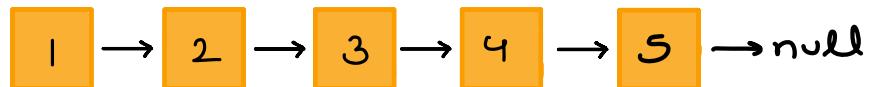
\* Search for an ele

Array → linear search →  $O(n)$

↳ Binary Search (only for sorted) →  $O(\log n)$

## Linkedlist

Search for k=5



head

→ if (head==null) return false; ← O: Empty LL

Node temp = head;

while (temp != null) {

    if (temp.data == k)

TC:  $O(n)$

        return true;

    }

    temp = temp.next;

}

return false;

```
if (head == null) return false;
```

```
Node temp = head;
```

```
while (temp.next != null)
```

```
    if (temp.data == k)
```

```
        return true;
```

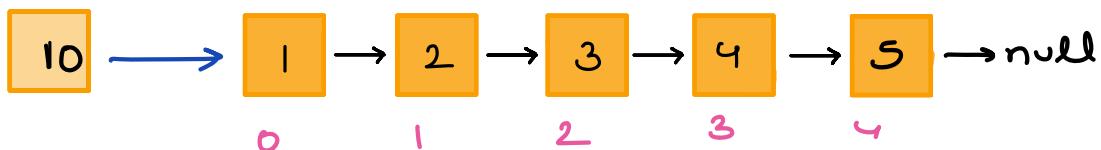
```
    temp = temp.next;
```

```
if (temp.data == k) return T
```

```
return F
```

Q Insert a value  $x$  at  $k^{th}$  position (0-based) in the given linkedlist ( $0 \leq k \leq n$ )  
length of LL

head



$x = 10$  at  $k = 0$

```
Node nn = new Node(10);
```

```

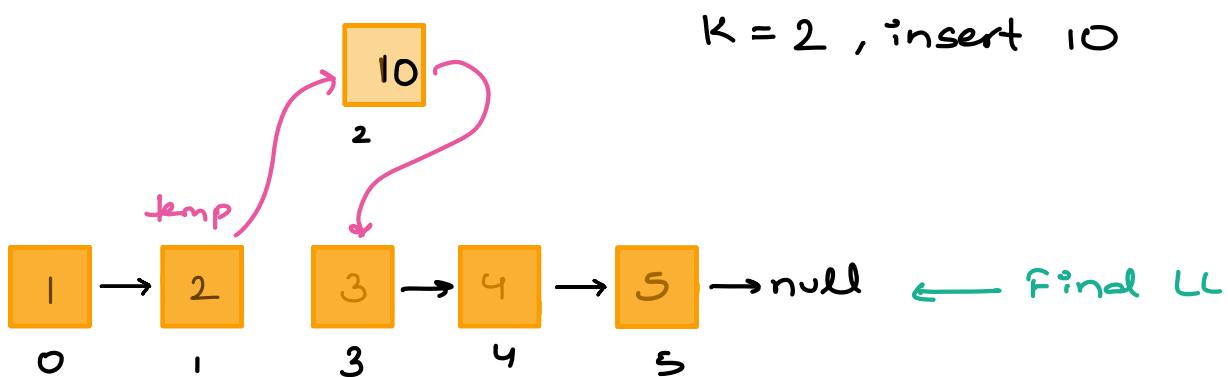
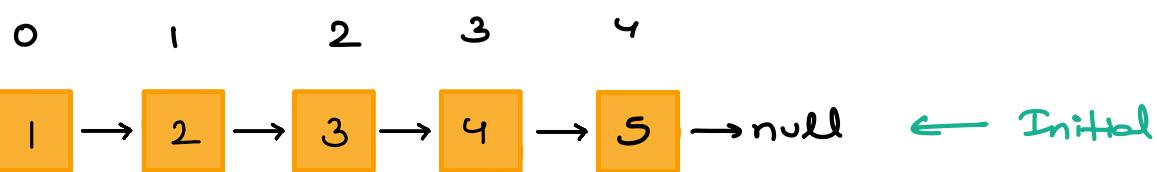
if (k==0){
    nn.next = head
    head = nn;
}

3
else {
    Node temp = head;
    for (i=1; i <= k-1; i++) {
        temp = temp.next
    }
    3
    nn.next = temp.next
    temp.next = nn;
}

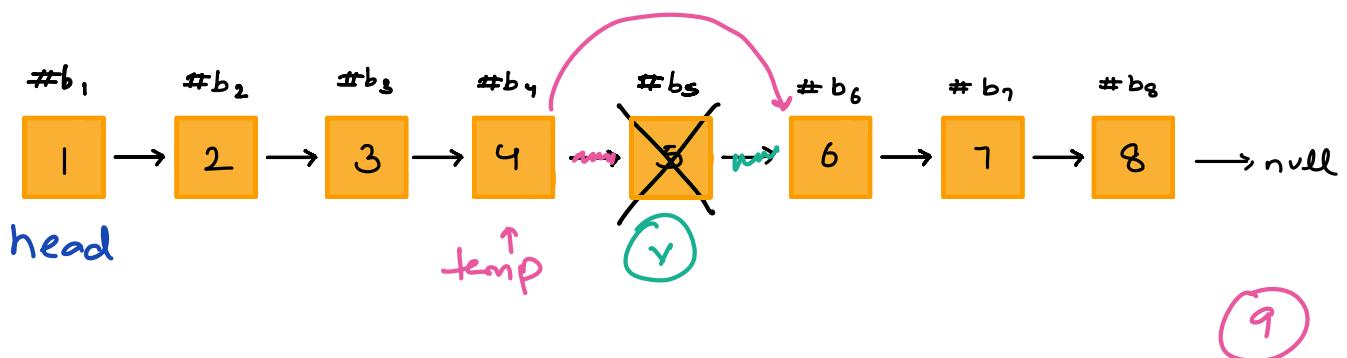
```

$Tc: O(n)$

$Sc: O(1)$



\* Delete the first occurrence of value  $x$  in the given LL.



if ( $\text{head} == \text{null}$ ) return head;

Node temp = head;

Node v = temp.next

if ( $\text{temp.data} == x$ ) {

    head = head.next;

v.next = null

    return head;

}

while ( $\text{temp} != \text{null}$ ) {

    if ( $\text{temp.next} != \text{null}$  &  $\text{temp.next.data} == x$ )

        temp.next = temp.next.next;

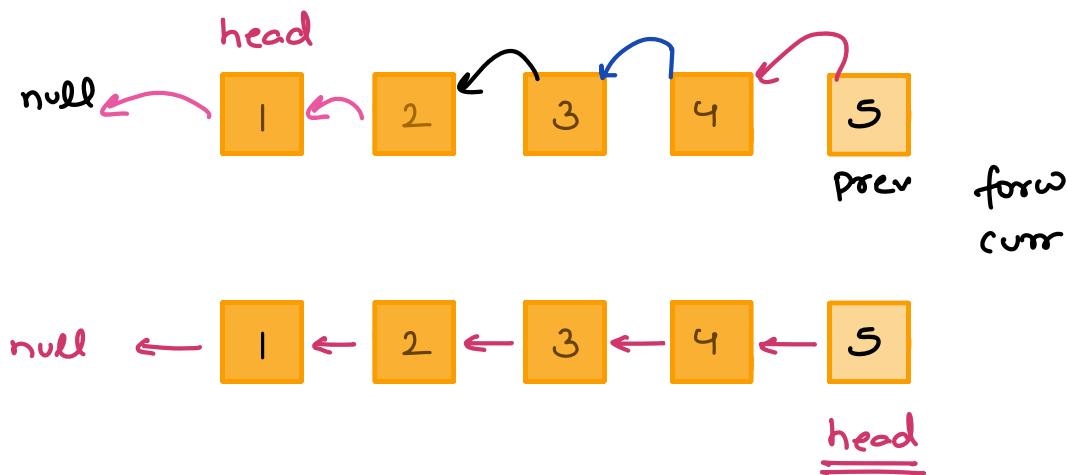
    return head;

    temp = temp.next;

}

return head;

Q Reverse the given linkedlist by updating pointers



Node prev = null

Node curr = head

while ( curr != null ) {

    Node forw = curr.next

    curr.next = prev

    prev = curr

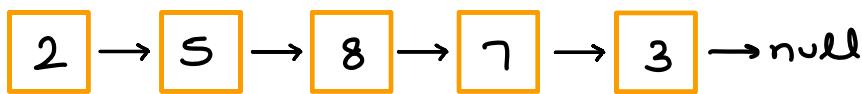
    curr = forw

3

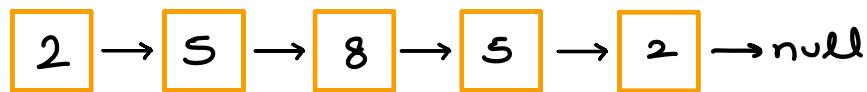
    head = prev;

return head;

Q Check if the given LL is palindrome or not



Ans = false



Ans = true

\* BF approach → Create a copy of this LL &  
reverse it compare if equal or not

Tc : O(n)

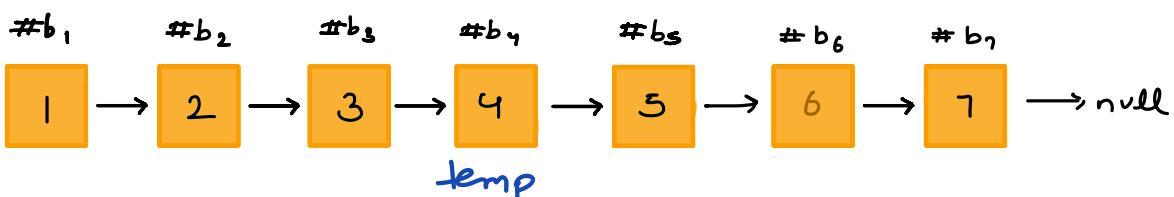
Sc : O(n)

Optimise Sc : O(1)

01. Middle of LL → Length of LL / 2

02. Reverse second half of LL

03. Compare first half & second half



01. n = 0 , temp = head

```

while (temp != null) {
    temp = temp.next;
    n++;
}

```

3

$$n=7 \quad \text{mid} = \frac{7}{2} = 3$$

temp = head;	$i = 1$	$i \leq 3$
for ( $i = 1 \rightarrow \text{mid}$ ) {	2	$2 \leq 3$
temp = temp.next;	3	$3 \leq 3$
}	4	$4 \leq 3 \times$

\* Take temp as head of second LL & reverse it.

\* Compare both of them

abba  $\rightarrow$  ab ba compare = 2

<u>ab c</u> ba	abc ba ✓
↑	<u>ab</u> <u>cba</u> ✓