

No one changes the world who isn't obsessed.

Billie Jean King

QuoteMaster.org



Good
Evening

Today's Agenda

- * Fractional Knapsack
- * 0/1 Knapsack
- * Unbounded Knapsack

Fractional KnapSack

Given N cakes with their happiness & weight.

Find max total happiness that can be kept in a bag with capacity = W .

Note → Cakes can be divided into pastries

$$N = 5$$

$$\text{cap} = 40$$

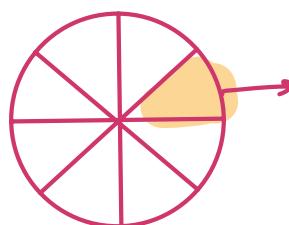
$$\text{happiness}[] = \{3 \quad 8 \quad 10 \quad 2 \quad 5\}$$

$$\text{weight}[] = \{10 \quad 4 \quad 20 \quad 8 \quad 15\}$$

Ideal → Pick the cake with max happiness

$$\sum h = 10 + 8 + 5 + 0.3 \Rightarrow 23.3$$

$$\sum w = 20 + 4 + 15 + 1$$



10 kg → 3 happiness

1 kg → $\frac{3}{10}$ unit = 0.3

$N=5$

$\text{Cap}=40$

$$h[] = \{ 4, 8, 10, 2, 5 \}$$

$$wt[] = \{ 4, 4, 20, 8, 16 \}$$

$$\sum h = 10 + 8 + 5 \Rightarrow \underline{\underline{23}}$$

$$\sum wt = 20 + 4 + 16$$

* Greedy parameter \rightarrow pick max happiness is not working

* Idea 2 \rightarrow Pick the ele with maximum

$$\frac{\text{happiness}}{\text{weight}}$$

$$h[] = \{ 4, 8, 10, 2, 5 \}$$

$$wt[] = \{ 4, 4, 20, 8, 16 \}$$

$$hp/wt = 1, 2, 0.5, 0.25, 0.3125$$

↑ ↑
 4 pastries 4 pastries

we need to pick 40 pastries

$$\Sigma h = 8 + 4 + 10 + 3.75 \rightarrow \underline{\underline{25.75}}$$

$$\Sigma wt = 4 + 4 + 20 + \underline{12} \rightarrow 40 \text{ kg}$$

* Code

01. Sort the cakes ^{in descending} on the basis of $\frac{\text{happiness}}{\text{weight}}$ ratio

02. `for (i=0; i<n; i++) {`

`if (wt[i] ≤ capacity)`

`ans += h[i];`

TC: $O(n \log n)$

`capacity -= wt[i];`

SC: $O(n)$

`else {`

`ans += $\left(\frac{h[i]}{wt[i]} * capacity \right)$`

`break;`

`}`

`3`

`return ans;`

O/I knapsack

Given N items each with a weight & value, find max value which can be obtained by picking items such that total weight of all items $\leq K$

Note 1 :- Every item can be picked at max 1 time

Note 2 :- We cannot take a part of item

Eg:-

N = 4 items

K = 50

Items : 0 1 2 3

Value [] : 100 60 120 150

Weight [] : 20 10 30 40

$v/w = 5 \ 6 \ 4 \ 3.75$

Greedy → Pick item with max v/w ratio 

$$\sum v = 60 + 100 = \underline{160}$$

$$\sum w = 10 + 20$$

Greedy - Pick item with max value

$$\sum v = 150 + 60 = 210$$

$$\sum w = 40 + 10 = 50$$

* Correct ans

$$\Sigma v = 100 + 120 = \underline{\underline{220}}$$

$$\Sigma w = 20 + 30 = 50$$

Idea 2 → Generate all combinations.

Every ele has two choices



& return the ans which is providing the max value.

$$cap = 15$$

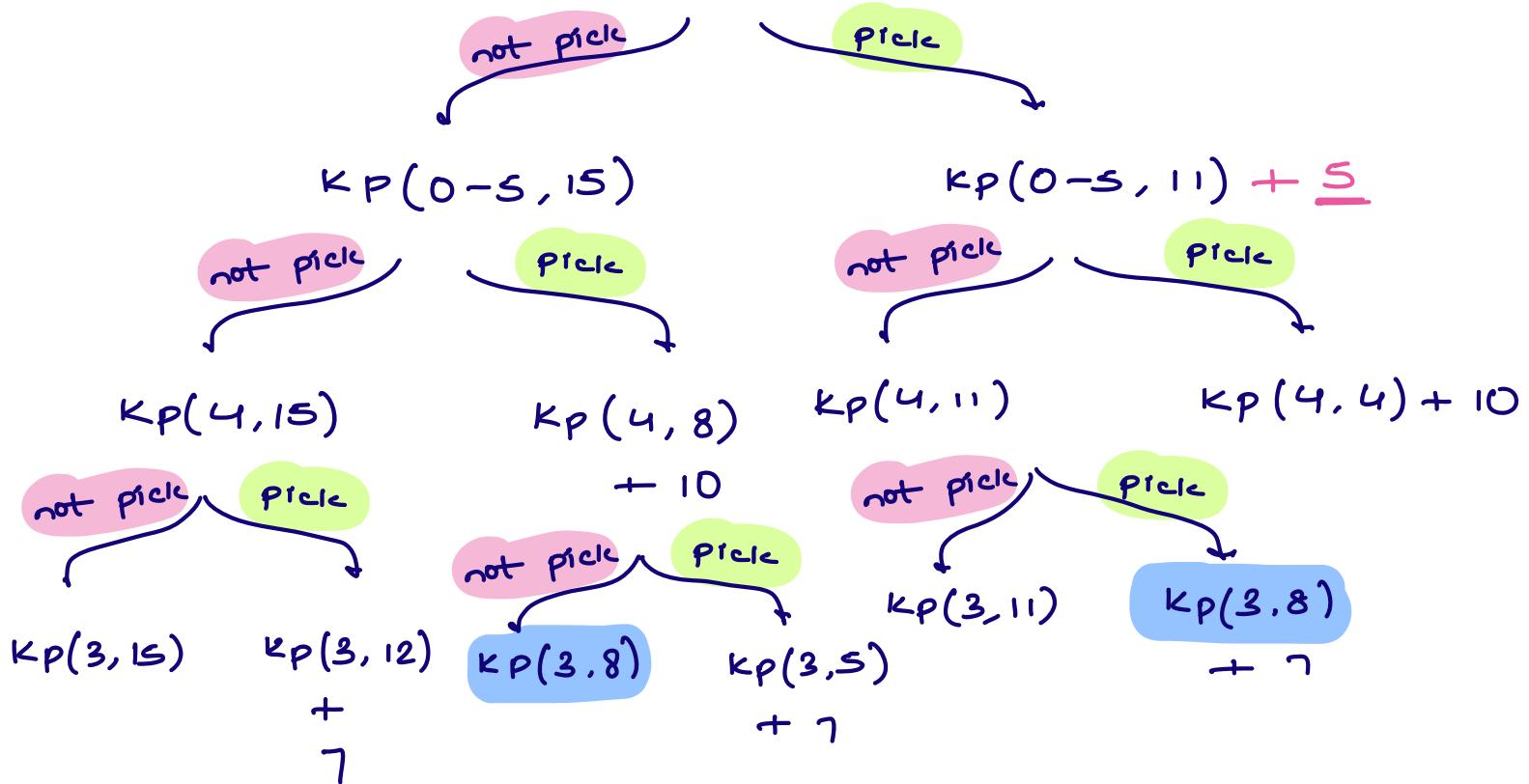
$$N = 7$$

$$wt[] = \{ 4, 1, 5, 4, 3, 7, 4 \}$$

$$wt[] = \{ 0, 1, 2, 3, 4, 5, 6 \}$$

$$val[] = \{ 3, 2, 8, 3, 7, 10, 5 \}$$

KP(0-6, 15)



Code →



```
int knapsack ( int i, int j, int []w, int []val )  
{  
    if ( i < 0 || j < 0 ) return -∞;  
    int not pick = knapsack ( i-1, j, w, val );  
    int pick = 0;  
    if ( j ≥ w[i] )  
    {  
        int pick = knapsack ( i-1, j-w[i], w, val ) + val[i];  
    }  
    return max ( not pick, pick );  
}
```

Top down

int [][] dp = new int[n][k+1]; $\xrightarrow{-1}$

int knapsack (i , j , wt , val)

if (i < 0 || j ≤ 0) return 0;

if (dp[i][j] != -1) return dp[i][j];

int val = knapsack (i-1, j , wt , val)

if (j ≥ wt)

val = max (val , knapsack(i-1, j-wt[i] , wt , val) + val[i])

return dp[i][j] = val ;

3

int [] wt = {3 6 5 2 4} cap=7
0 1 2 3 4

int [] val = {12 20 15 6 10}

int dp[n+1][k+1]

* Bottom Up Approach

val	wt	ele	cap →							
			0	1	2	3	4	5	6	7
0			0	0	0	0	0	0	0	0
12	3		0	0	0	12	12	12	12	12
20	6		0	0	0	12	12	12	20	20
15	5		0	0	0	12	12	15	20	20
6	2		0	0	6	12	12	18	20	21
10	4		0	0	6	12	12	18	20	22

not pick 4th ele → we have to ask our first 3 ele
 to go & fill the bag of cap = 3
 $\Rightarrow \underline{\underline{12}}$
 & generate max value

pick 4th ele = 6 + rem bag cap = 1
 ask 3 ele to go & fill the
 bag of cap=1 & gen max value

$dp[i][j] = \text{max value that we can generate using } i \text{ elements &} j \text{ as bag capacity}$

$$dp[i][j] = \max \left(dp[i-1][j], val[i-1] + dp[i-1][j - wt[i-1]] \right)$$

↑
not pick

if ($j - wt[i-1] \geq 0$)

* for ($i=0; i < n; i++$)

TC: $O(n * k)$

SC: $O(n * k)$

for ($j=0; j < m; j++$) {

 if ($i == 0 \text{ or } j == 0$) $dp[i][j] = 0$

 else {

$rej = dp[i-1][j]$

$sel = 0$

 if ($j \geq wt[i-1]$)

$select = val[i-1] + dp[i-1][j - wt[i-1]]$;

 3

$dp[i][j] = \max(rej, select)$;

 3

2

return $dp[n][k]$;

Unbounded Knapsack

Given N items each with a weight & value, find max value which can be obtained by picking items such that total weight of all items $\leq K$

Note 1 :- Every item can be picked infinite no. of times

Note 2 :- We cannot take a part of item

Eg:- $N = 4$ items

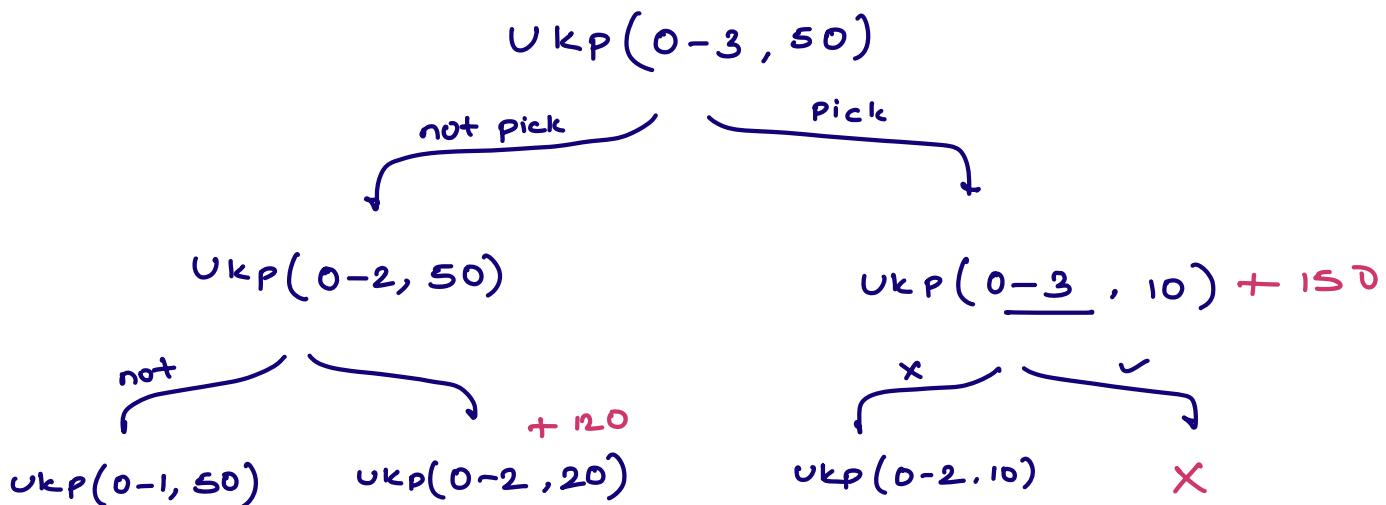
$K = 50$

Items: 0 1 2 3

Weight[]: 20 10 30 40

Value[]: 100 60 120 150

Ans = 300



```

for (i=0; i<n; i++)
    for (j=0; j<m; j++) {
        if (i==0 || j==0) dp[i][j]=0
        else {
            rej = dp[i-1][j]
            sel = 0
            if (j >= wt[i-1])
                select = val[i-1] + dp[i][j-wt[i-1]];
            dp[i][j] = max(rej, select);
        }
    }
return dp[n][k];

```