

TODAY:

TREES 2

DON'T
FORGET



"MOTIVATION MAY BE
WHAT STARTS YOU ON,
BUT IT'S HABIT THAT
KEEPS YOU GOING
BACK FOR MORE."

- MIYA YAMANOUCHI

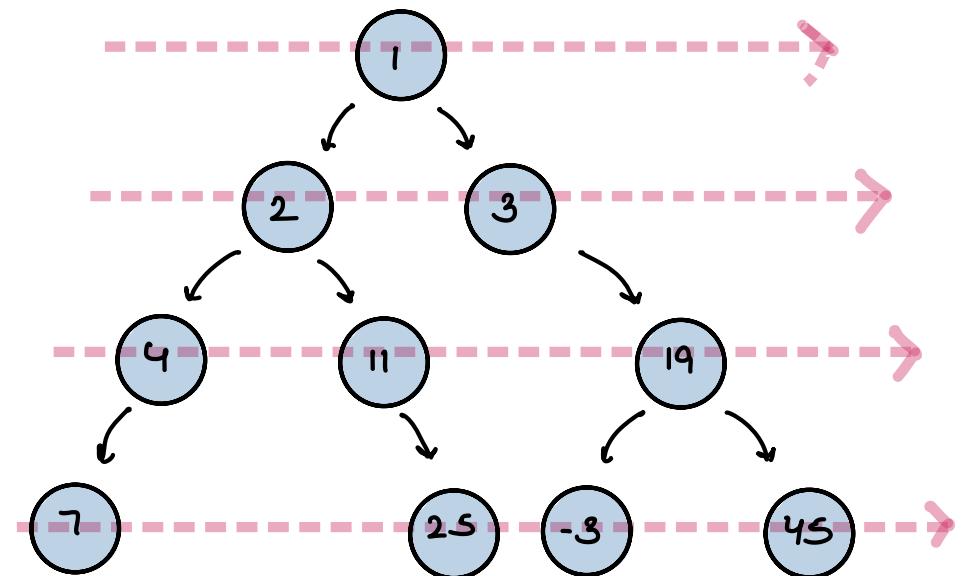
Good
Evening

AGENDA:

Topics for Today

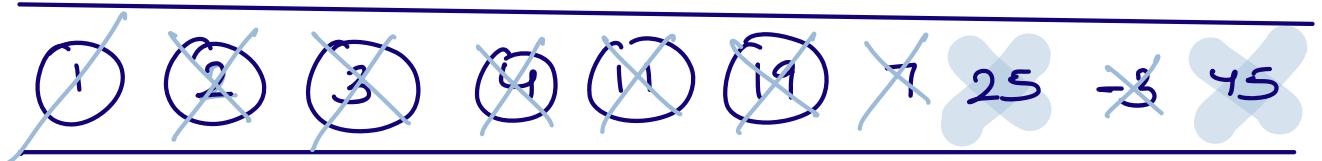
01. Level Order Traversal
02. Left view & Right view
03. Vertical level order Traversal
04. Top view & Bottom view
05. Types of B.T
06. Check height balanced

Level Order Traversal (BFS)



Ans = 1 2 3 4 11 19 7 25 -3 45

Queue ?



{ remove
work
push the child =

Queue < Node > q ;

q. enqueue (root);

while (q.size() > 0)

 Node rem = q.dequeue();

 print (rem.data);

 if (rem.left != null) q.enqueue (rem.left);

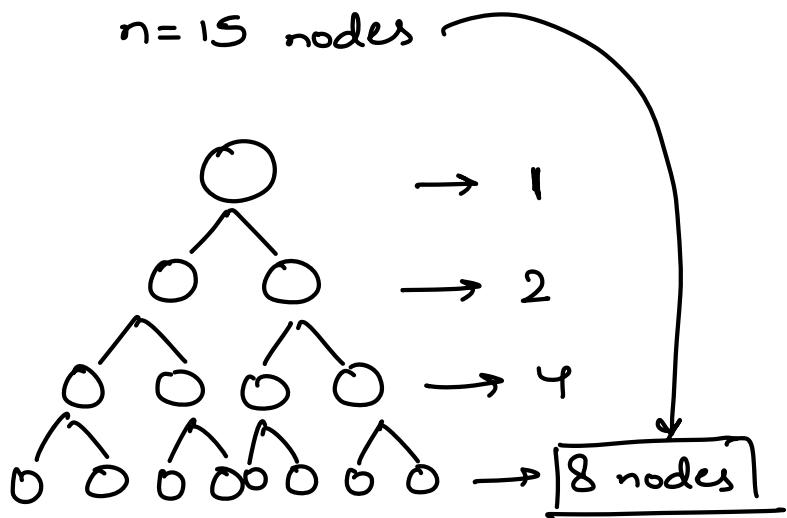
 if (rem.right != null) q.enqueue (rem.right);

3

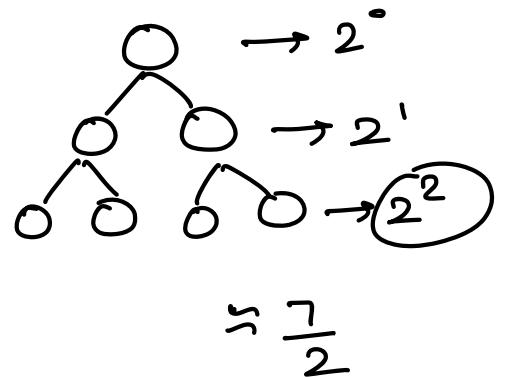
TC : O(n)

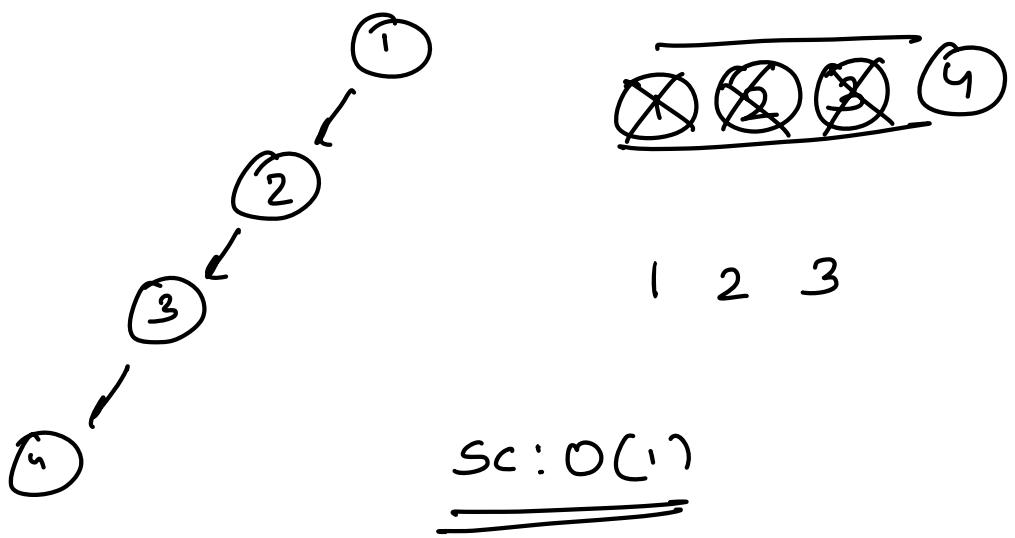
SC : O(n)

n = 15 nodes

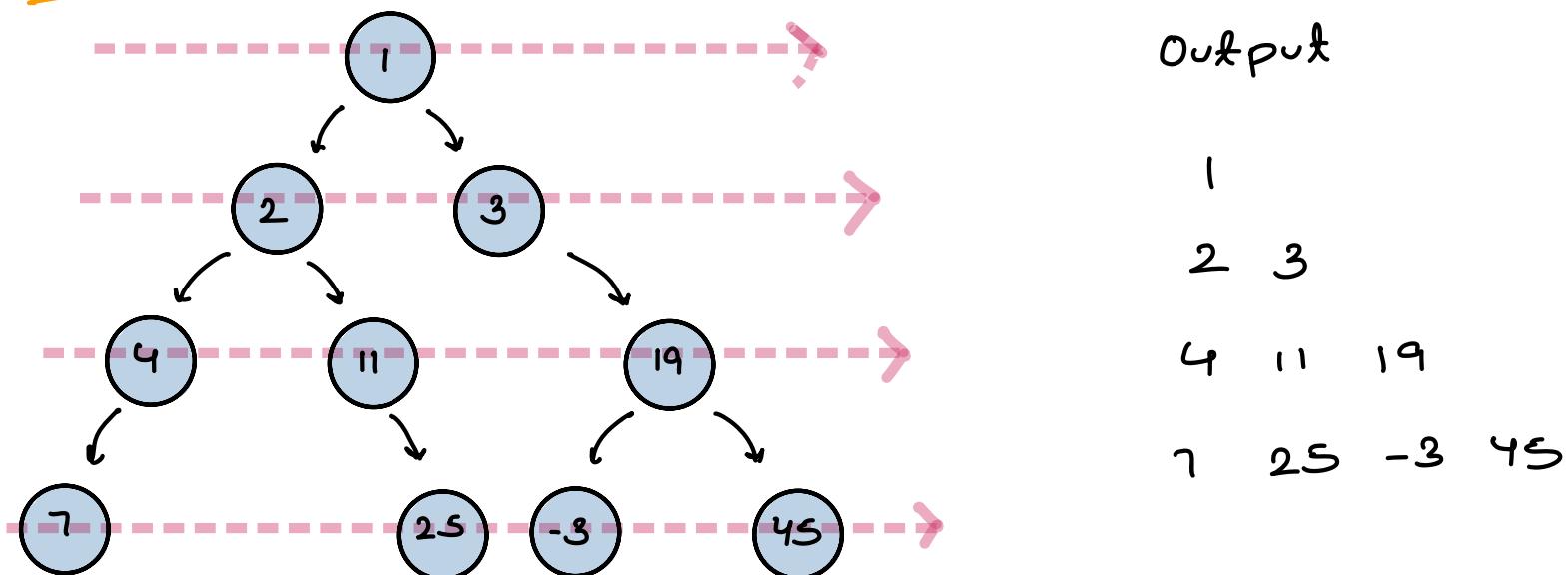


n = 7 nodes



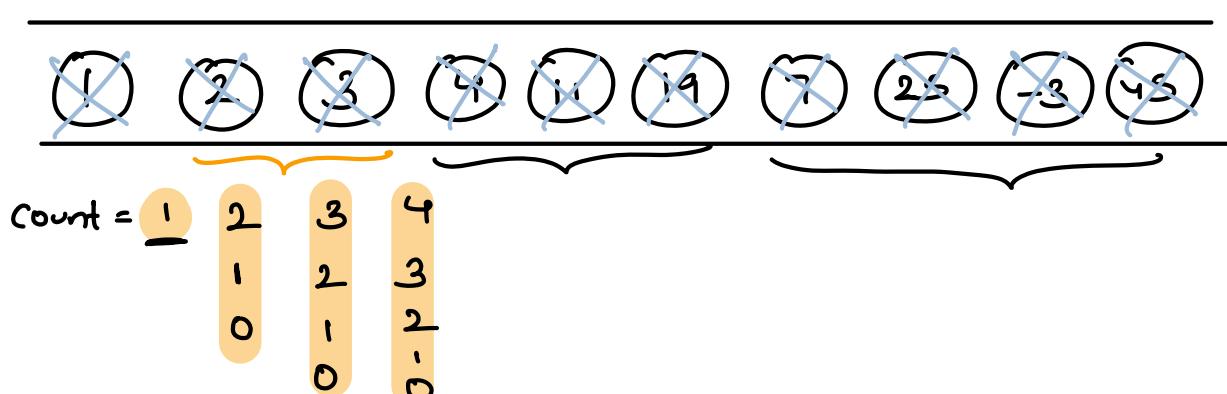


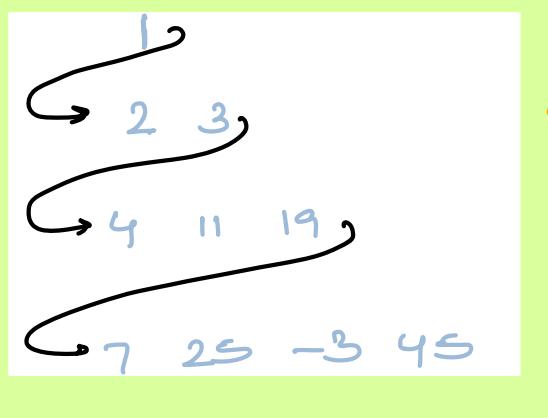
Q2



counter for nodes at a particular level

Queue ↴





Queue <Node> q ;

q. enqueue (root);

while (q.size() > 0)

int sz = q.size();

for (i=1; i ≤ sz; i++) {

Node rem = q.dequeue();

print (rem.data);

if (rem.left != null) q.enqueue (rem.left);

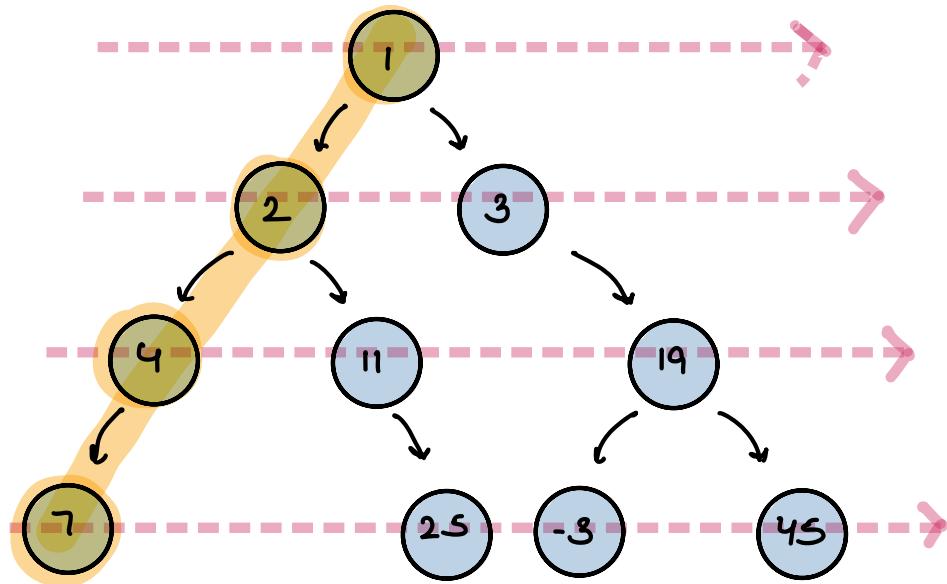
if (rem.right != null) q.enqueue (rem.right);

3
println();

TC : O(n)

SC : O(n)

* left view of tree



Queue < Node > q ;

q. enqueue (root);

while (q.size () > 0)

int sz = q.size();

for (i=1 ; i ≤ sz ; i++) {

Node rem = q.dequeue();

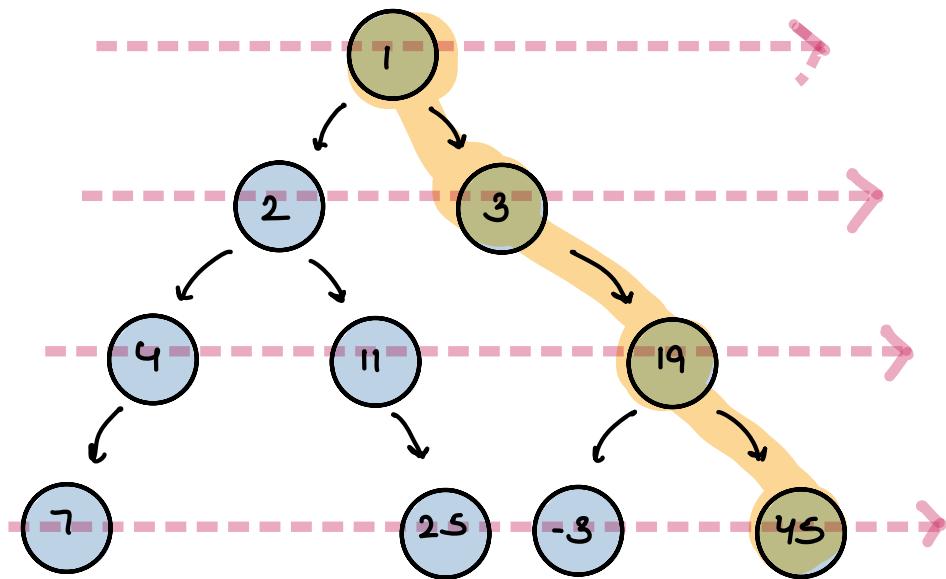
if (i == 1) print (rem.data);

if (rem.left != null) q.enqueue (rem.left);

if (rem.right != null) q.enqueue (rem.right);

3
2
println();

* Right view of binary tree



Queue <Node> q ;

q. enqueue (root);

while (q.size () > 0)

 int sz = q.size();

 for (i=1 ; i ≤ sz ; i++) {

 Node rem = q.dequeue();

 if (i == sz) print (rem.data);

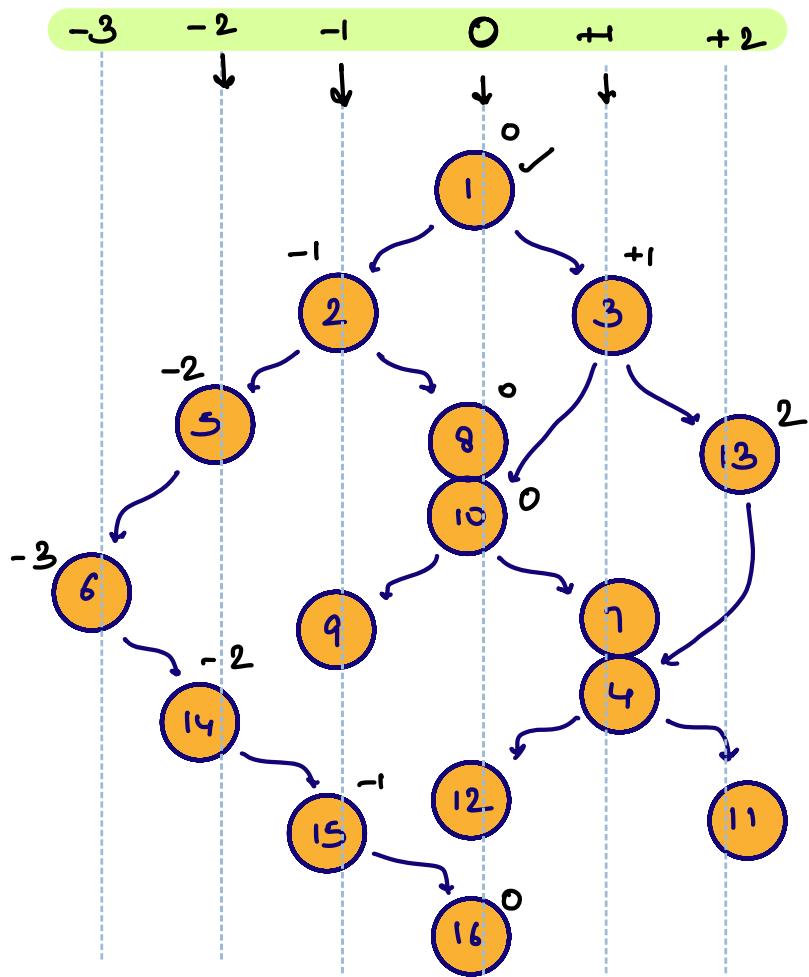
 if (rem.left != null) q.enqueue (rem.left);

 if (rem.right != null) q.enqueue (rem.right);

 }

 println();

* Vertical Order Traversal (Microsoft)



Output

6				
5	14			
2	9	15		
1	8	10	12	16
3	7	4		
13	11			

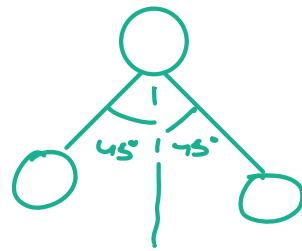
HM

Recursion (DFS)

Iteratively (BFS)

vl Nodes

0	\rightarrow	1, 8, 10, 12, 16
-1	\rightarrow	2, 9
1	\rightarrow	3, 7, 4
-2	\rightarrow	5, 14
2	\rightarrow	13, 11
-3	\rightarrow	6





pair = node, vl

-
10, 0

left = 9, -1

right = 7, +1

7, +1

4, 1

Code

```
class Pair {
```

```
    Node node  
    int vl;
```

```
Pair rootpair = new pair (root, 0)
```

```
Queue<pair> q;
```

```
HashMap<int, list> map;
```

$\min vl = 0$

$\max vl = 0$

q. enqueue (root pair):

TC: O(n)

SC: O(n)

```
while (q.size() > 0)
```

```
    pair rp = q.dequeue();
```

```
    minvl = Min (minvl, rp.vl)
```

```
    maxvl = Max (maxvl, rp.vl);
```

```
    // Add rp.node at rp.vl;
```

```
    hm.get(rp.vl).add(rp.node) // Handle this  
                                carefully
```

```
    if (rp.node.left != null)
```

```
        q.enqueue(new pair(rp.node.left, rp.vl - 1));
```

```
    if (rp.node.right != null)
```

```
        q.enqueue(new pair(rp.node.right, rp.vl + 1));
```

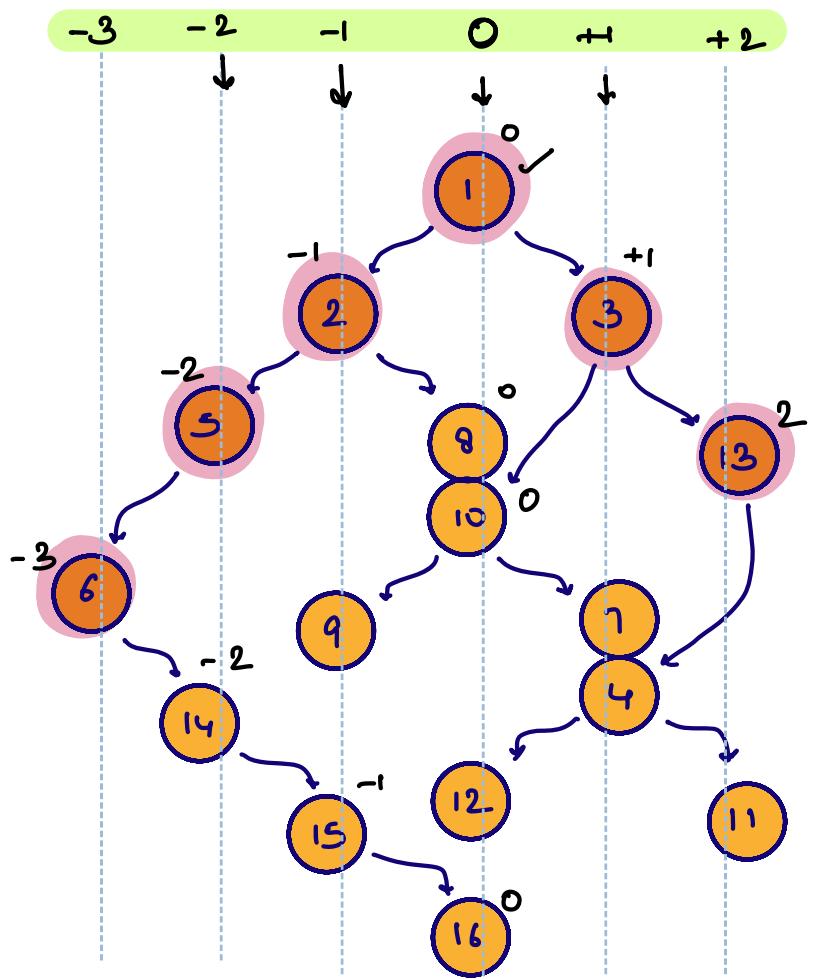
```
}
```

```
for (i = minvl; i ≤ maxvl; i++) {
```

```
    print(map[i]);
```

```
}
```

* Top view



vl	Nodes
0	1, 8, 10, 12, 14
-1	2, 9, 15
1	3, 7, 4
-2	5, 14
2	13, 11
-3	6

Top view → print first ele of every level

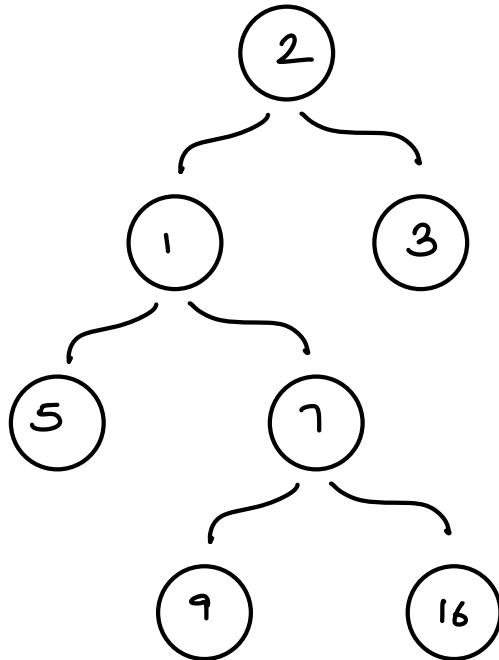
Bottom view → print last ele of every level

10:13 pm → 10:23 pm

Types of Binary Tree

01. Proper / Full Binary Tree

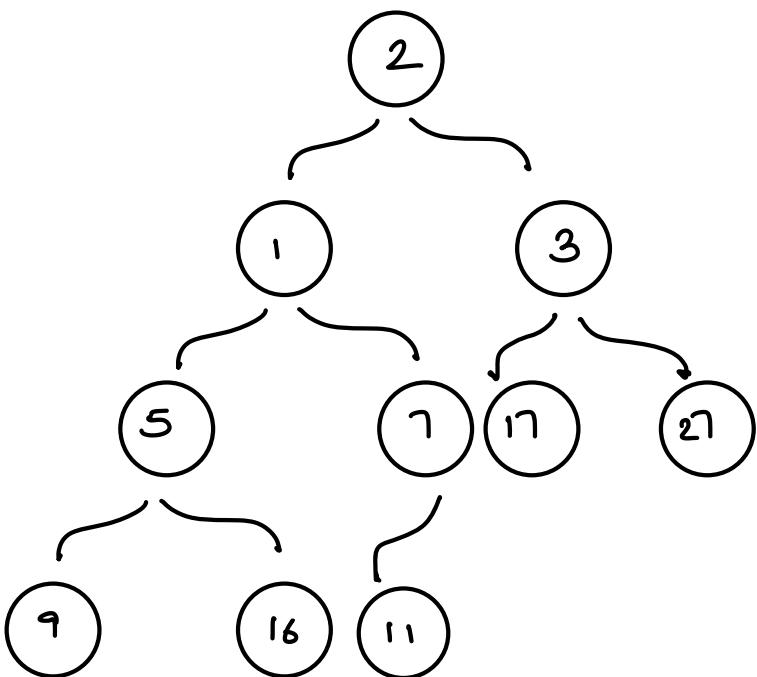
For every node, either 0 or 2 children



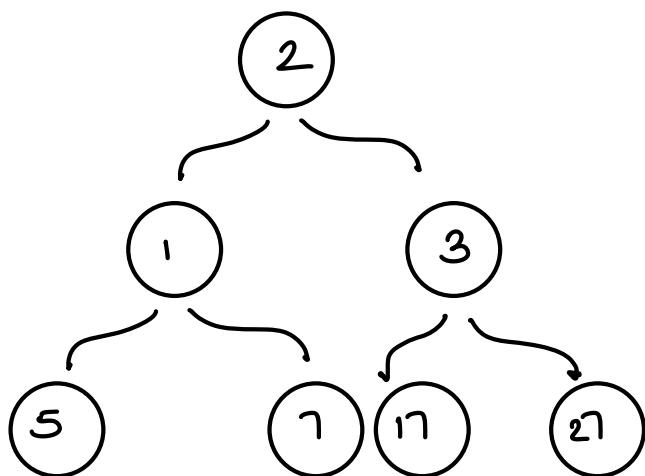
* Complete Binary Tree

→ Every level is completely filled except the last level which may or may not be completely filled

→ In last level, nodes should be filled from LHS to RHS



* Perfect Binary Tree → All levels are completely filled



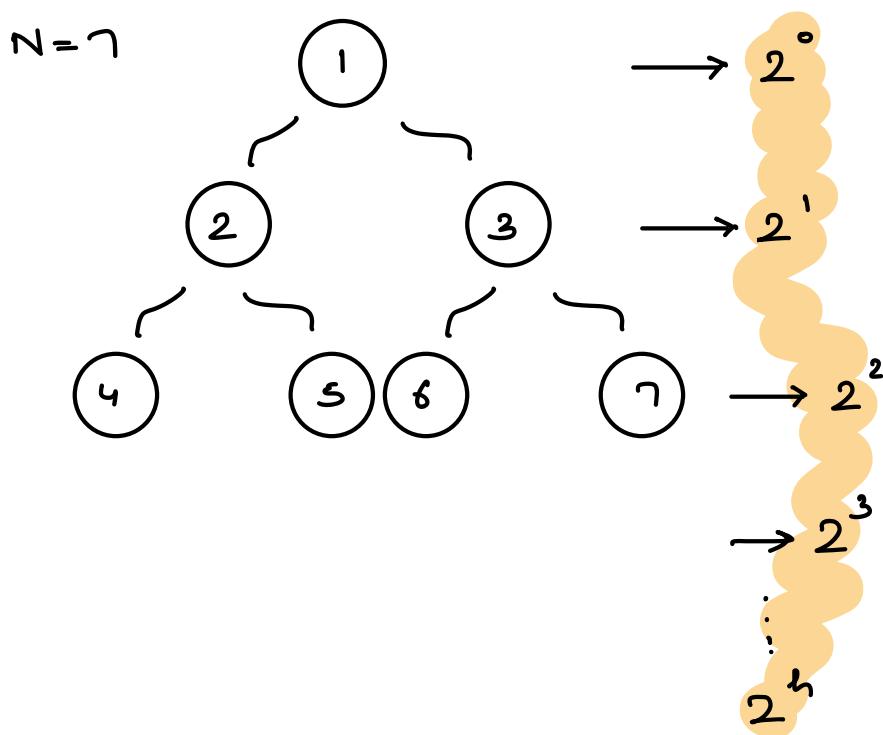
Are all perfect binary Tree also proper? Yes

Are all perfect BT, also complete tree? Yes

Are all complete BT, also proper tree? No

Q Given perfect BT with N nodes, find the height of tree

$$N = 1 \quad \text{h} = 0$$



$$\underbrace{2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^h}_{\text{no. of ele} = h+1} = N$$

$$\frac{a(r^n - 1)}{r - 1}$$

$$\frac{+ [2^{h+1} - 1]}{2 - 1} = N$$

$$2^{h+1} = N + 1$$

$$\log_2 2^{h+1} = \log_2(n+1)$$

$$h+1 = \log_2(n+1)$$

$$h = \log_2(n+1) - 1$$

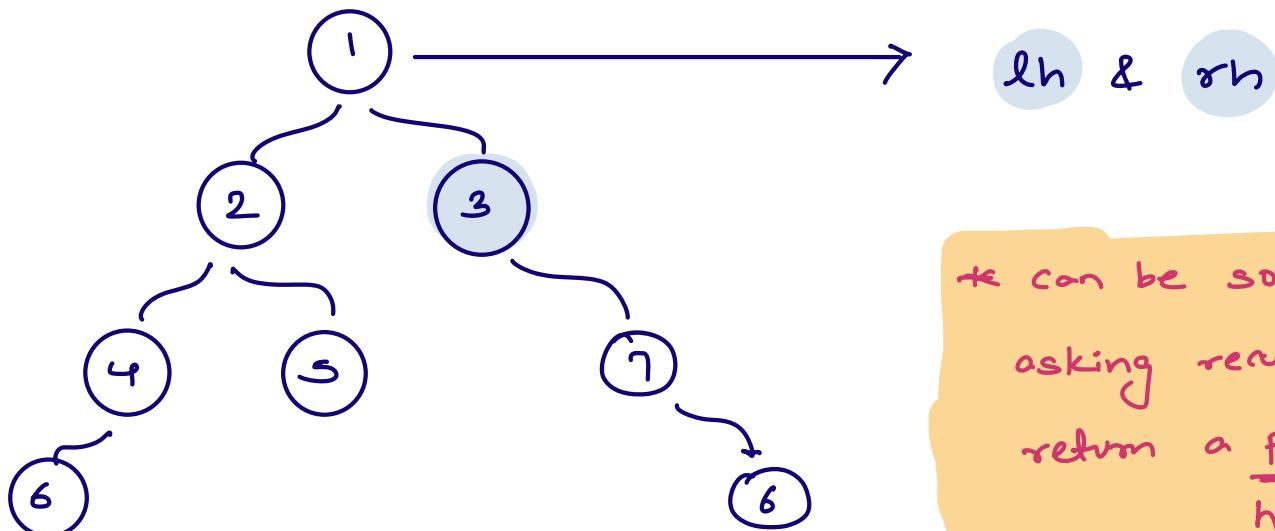
$$O(\log_2 n) \leq h \leq O(n)$$

* Balanced Binary Tree

For all nodes

$$\left| \frac{\text{height of left sub}}{\text{height of right sub}} \right| \leq 1$$

Q Given a BT, check if it is balanced or not



It can be solved by asking recursion to return a pair $\underline{h, \text{isbalanced}}$

```
static boolean isbalanced = true;
```

Travel &
change

```
int height(Node root)
```

```
if (root == null) { return -1; }
```

```
lh = height(root.left);
```

```
rh = height(root.right);
```

```
if (abs(lh - rh) > 1) isbalanced = false;
```

```
return max(lh, rh) + 1;
```