

### 4.12.2 Type Argument Inference

Given a signature  $\langle T_1, T_2, \dots, T_n \rangle (p_1: P_1, p_2: P_2, \dots, p_m: P_m)$ , where each parameter type  $P$  references zero or more of the type parameters  $T$ , and an argument list  $(e_1, e_2, \dots, e_m)$ , the task of type argument inference is to find a set of type arguments  $A_1 \dots A_n$  to substitute for  $T_1 \dots T_n$  such that the argument list becomes an applicable signature.

The **inferred type argument** for a particular type parameter is the widened form (section 3.9) of the best common type (section 3.10) of a set of candidate types. In order to compute candidate types, the argument list is processed as follows:

- Initially all inferred type arguments are considered **unfixed** with an empty set of candidate types.
- Proceeding from left to right, each argument expression  $e$  is **inferentially typed** by its corresponding parameter type  $P$ , possibly causing some inferred type arguments to become **fixed**, and candidate type inferences (section 3.8.6) are made for unfixed inferred type arguments from the type computed for  $e$  to  $P$ .

The process of inferentially typing an expression  $e$  by a type  $T$  is the same as that of contextually typing  $e$  by  $T$ , with the following exceptions:

- Where expressions contained within  $e$  would be contextually typed, they are instead inferentially typed.
- Where a contextual type would be included in a candidate set for a best common type (such as when inferentially typing an object or array literal), an inferential type is not.
- When a function expression is inferentially typed (section 4.9.3) and a type assigned to a parameter in that expression references type parameters for which inferences are being made, the corresponding inferred type arguments to become **fixed** and no further candidate inferences are made for them.
- If  $e$  is an expression of a function type that contains exactly one generic call signature and no other members, and  $T$  is a function type with exactly one non-generic call signature and no other members, then any inferences made for type parameters referenced by the parameters of  $T$ 's call signature are **fixed** and, if  $e$ 's call signature can successfully be instantiated in the context of  $T$ 's call signature (section 3.8.5),  $e$ 's type is changed to a function type with that instantiated signature.

In the example

```
function choose<T>(x: T, y: T): T {  
    return Math.random() < 0.5 ? x : y;  
}  
  
var x = choose("Five", 5);
```

inferences for 'T' in the call to 'choose' are made as follows: For the first parameter, an inference is made from type 'string' to 'T'. For the second parameter, an inference is made from type 'number' to 'T'. Since