

signature, the result of this process is simply  $A$ . Otherwise, type arguments for  $A$  are inferred from  $B$  producing an instantiation of  $A$  that can be related to  $B$ :

- Using the process described in 3.8.6, inferences for  $A$ 's type parameters are made from each parameter type in  $B$  to the corresponding parameter type in  $A$  for those parameter positions that are present in both signatures, where rest parameters correspond to an unbounded expansion of optional parameters of the rest parameter element type.
- The inferred type argument for each type parameter is the best common type (section 3.10) of the set of inferences made for that type parameter.
- Provided all inferred type arguments satisfy their corresponding type parameter constraints, the result is an instantiation of  $A$  with the inferred type arguments.
- Otherwise,  $A$  cannot be instantiated in the context of  $B$  and the process is unsuccessful.

Given the declarations

```
var f: (x: string) => string[];  
var g: <T>(x: T) => T[];
```

the assignment ' $f = g$ ' causes the signature of ' $g$ ' to be instantiated in the context of the signature of ' $f$ ', which causes inferences for ' $T$ ' to be made from the parameter type ' $string$ ' to the parameter type ' $T$ ', ultimately producing an assignment compatible instantiation of ' $g$ ' with the type argument ' $string$ '.

Changing the example to

```
var f: (x: { a: string; b: number }) => {}[];  
var g: <T>(x: { a: T; b: T }) => T[];
```

inferences for ' $T$ ' are now made from the parameter type ' $\{ a: string; b: number \}$ ' to the parameter type ' $\{ a: T; b: T \}$ ', leading to inferences of ' $string$ ' and ' $number$ ' for ' $T$ '. The best common type of those two types is the empty object type, which becomes the type argument for the resulting assignment compatible instantiation of ' $g$ '.

In the example

```
var f: <T>(x: T, y: T) => { x: T; y: T };  
var g: <U, V>(x: U, y: V) => { x: U; y: V };
```

the assignment ' $f = g$ ' causes ' $T$ ' to be inferred for both ' $U$ ' and ' $V$ ', thus producing an assignment compatible instantiation of ' $g$ '. The reverse assignment ' $g = f$ ' leads to inferences of ' $U$ ' and ' $V$ ' for ' $T$ ', producing an instantiation of ' $f$ ' with the empty object type as an argument, and since an empty object type is not assignable to a type parameter, the assignment is an error.

### 3.8.6 Type Inference

In certain contexts, inferences for a given set of type parameters are made *from* a type  $S$ , in which those type parameters do not occur, *to* another type  $T$ , in which those type parameters do occur. Inferences