Trees 2

Binary Trees

- A binary tree is a rooted tree with no vertex
- has more than two children
 - left and right child nodes

```
// Syntax
struct BinaryNode{
    Object element;  // The Data in the node
    BinaryNode *left;  // & of Left Child
    BinaryNode *right; // & of Right Child
}
```

- A binary tree is complete iff every layer except possibly the bottom is fully populated with vertices. All nodes at the bottom level must occupy the leftmost spots consecutively.
- ullet A complete binary tree with n vertices and h height satisfies

```
\begin{array}{l} \bullet \ \ 2^H \leq n < 2^{H+1} \\ \bullet \ \ 2^2 \leq 7 < 2^{2+1}, 2^2 \leq 4 < 2^{2+1} \\ \bullet \ \ 2^H \leq n < 2^{H+1} \end{array}
```

- $\bullet \ \ H \leq log(n) < H+1$
- H = floor(log(n))

• Proof:

- ullet At level $k \leq H-1$, there are 2^K vertices
- ullet At level k=H , there is at least 1 node, and at most 2^H vertices
- Total number of vertices when all levels are fully populated (maximum level k)

```
• n=2^0+2^1+\ldots+2^k
• n=1+2^1+2^2+\ldots 2^k (Geometric Progression)
• n=\frac{1(2^{k+1}-1)}{2-1}
• n=2^{k+1}-1
```

Case 1:

A tree has the maximum number of nodes when all levels are fully populated

- Let k = H
 - $n = 2^{H+1} 1$
 - $n < 2^{H+1}$

Case 2:

The tree has a minimum number of nodes when there is only one node in the bottom level

- Let k = H 1
 - $n' = 2^H 1$
 - $n > n' + 1 = 2^H$

Combining two conditions we have

 $\bullet \ \ 2^H \leq n \leq 2^{H+1}$

Representation of Complete Binary Tree

- All trees can be represented by the generic representation shown in the code above
- Due to the structure of a complete binary tree, it cannot be represented by a vector
 - As long as one can figure out the parent/child relationship
 - Parent/child relationship embedded in the index of parent and child
 - Vector elements carry data
- Tree Structure : Vector
 - Vector indices carry tree structure
 - Index order = levelorder
 - The tree structure is implicit
 - Uses integer arithmetic for tree navigation
 - No need to explicitly store the tree node pointers
- Tree Navigation : Vector
 - The root at v[0]
 - Parent of v[k] = k[(k-1)/2]
 - Left child of v[k] = v[2*k + 1]

Binary Tree Traversal

Inorder traversal

- Definition
 - Left child
 - Vertex
 - Right Child (recursive)
- Algorithm
 - Depth-first search (visit between children)

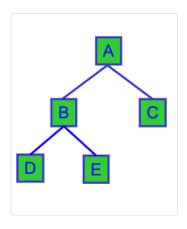
Other Traversals that apply to binary case

- Preorder
 - Trees
- Postorder
 - Trees 🌲
- Level order traversal

A tree can be rebuilt from its inorder and preorder (or postorder) traversal results

Rebuild Tree from Traversal

- Let each node be associated with a letter, traversals print the letters when visiting a node. The results are:
 - Preorder: "ABDEC"Postorder: "DEBCA"
 - Inorder: "DBEAC"



Rebuild tree from preorder + inorder traversal

- Find the root from the preorder result: A
- Decide left and right subtrees
 - Find the letter for the root in the inorder string and decide the inorder string for the two subtrees
- Decide the preorder string for left and right subtrees
 - Inorder for the traversal string length should be equal to another string length, extract preorder strings from the whole preorder string
- Recursively do this to the sub-trees

Build Expression Tree from Postfix Expression

```
stack<T> s;
while(s != /*end of postfix expression*/){
    // Get the next token
    if(token == /* operand */){
        // Create a new node with the operand
        s.push(/* New Node */);
    }
    if(token == /* operator */){
        s.pop(); // corresponding operands from S
        // Create a new node with the operator (and corresponding operands as left/right children)
        s.push(/* New Node */);
    }
} // s.top is the final binary expression
```

Rebuilding with a tree

• Postorder string: FECAHJIGB

• Inorder string: CFEABHGJI

Root: B

Last in the post-order string

Left Subtree: CFEA

• Before root (B) in the inorder string

Right Subtree: HGJI

• After root (B) in the inorder string