### **Algorithm Analysis 2**

#### Recursion

Expanding recurrence form results in a simulation of the function.

$$T(n) = 1 + T(n-1) = 1 + 1 + T(n-2) = \dots$$

Which leads to...

$$T(n) = n - 1 + T(1) = O(n)$$
  $T(n) > 2^{n-1/n} * T(1), T(n) = \Omega(2^{n-1/n})$ 

### **Binary Search**

## **Euclid's Algorithm**

- Find the GCD (greatest common divisor) between n and m
  - Given M >= N
- Why is it O(logN)

# **Exponential**

- Calculate  $x^0$
- Complexity O(logN)

```
long pow(long x, int n){
    if(n == 0){
        return 1;
    }
    if(n == 1){
        return x;
    }
    if (isEven(n)){
        return pow(x * x, n/2);
    }
    else
        return pow(x * x, n/2) * x;
}
// O(logN)
```