# Continuous Logic and Learning Bounds

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# Model Theory to Learnability

- Let  $C = \{c_y : y \in Y\}$  be a class of subsets of X indexed by Y.
- C is NIP/stable when there is an NIP/stable formula  $\phi(x; y)$  such that  $x \in c_y \iff \phi(x; y)$ .
- The properties in each row are equivalent:

Model Theory	Combinatorics	Learning Theory
NIP	finite VC dimension	PAC learnable
stable	finite Littlestone dimension	online learnable

# Continuous Logic to Learnability

- Let  $\mathcal{H} = \{h_y : y \in Y\}$  be a class of functions  $X \to [0,1]$  indexed by Y.
- $\mathcal{H}$  is NIP/stable when there is an NIP/stable formula  $\phi(x; y)$  of continuous logic such that  $h_y(x) = \phi(x; y)$ .
- The properties in the table have been generalized to  $\mathcal{H}$ , but the connections are understudied.

### New Learnable Function Classes

### Theorem (A., Benedikt)

A class  $\mathcal{H}$  of functions  $X \to [0,1]$  is stable iff it is online learnable.

### Theorem (A., Benedikt)

The randomization of a PAC/online learnable function class  ${\cal H}$  is also PAC/online learnable.

# Generalizing VC Dimension to Continuous Logic

### Theorem (Ben Yaacov)

A formula  $\phi$  of continuous logic is NIP iff the class of functions it defines has finite  $\gamma$ -fat-shattering dimension for all  $\gamma > 0$ .

#### Definition

Let  $\mathcal{H}$  be a class of functions  $X \to [0,1]$  and let  $\gamma > 0$ . We say  $\mathcal{H}$ has  $\gamma$ -fat-shattering dimension at least n when there are

$$\bullet$$
  $x_1,\ldots,x_n\in X$ 

• 
$$s_1, \ldots, s_n \in [0, 1]$$

$$ullet$$
 For every  $E\subseteq\{1,\ldots,n\}$ , a function  $h_E\in\mathcal{H}$  satisfying

• if 
$$i \in E$$
,  $h_E(x_i) \ge s_i + \gamma$ 

• if 
$$i \notin E$$
,  $h_E(x_i) \leq s_i - \gamma$ .

# Probably Approximately Correct Learning

A class  $\mathcal{H}$  of functions  $X \to [0,1]$  is PAC learnable when for every  $\varepsilon, \delta > 0$ , there is *n* such that when...

- $(x_1, y_1), \dots, (x_n, y_n) \in X \times [0, 1]$  are i.i.d. random,
- we can choose  $h \in \mathcal{H}$  (hoping that  $h(x_{n+1}) \approx y_{n+1}$ ) such that
- with probability at least  $1 \delta$ ,
- $\mathbb{E}[|y_{n+1}-h(x_{n+1})|]$  is within  $\varepsilon$  of the best case for all  $h\in\mathcal{H}$ .

We call  $n = n(\varepsilon, \delta)$  the sample complexity.

# PAC Learning Bound

### Theorem (Bartlett, Long)

The sample complexity  $n(\varepsilon, \delta)$  of PAC-learning  $\mathcal H$  is bounded by

$$O\left(\frac{1}{\epsilon^2} \cdot \left(\operatorname{FatSHDim}_{\frac{\epsilon}{9}}\left(\mathcal{H}\right) \cdot \log^2\left(\frac{1}{\epsilon}\right) + \log\left(\frac{1}{\delta}\right)\right)\right)$$

In fact,  $\mathcal{H}$  is PAC-learnable if and only if the  $\gamma$ -fat-shattering dimension is finite for all  $\gamma > 0$ .

Hu et al. extended this to learning a class of measures on  $\mathcal{H}$ , at the cost of a much worse bound.

We'll use continuous logic to find examples of learnable classes and improve the Hu et al. bound.

Suppose  $(\Omega, \Sigma, \mu)$  is a probability space and  $\mathcal{F} = (\mathcal{H}_{\omega} : \omega \in \Omega)$  is a family of function classes  $\mathcal{H}_{\omega} = (h_{\omega, y} : y \in Y)$ .

#### Definition

Assuming measurability, define  $\mathbb{E}\mathcal{F}_y:X o[0,1]$  by

$$\mathbb{E}\mathcal{F}_{y}(x) = \mathbb{E}\left[h_{\omega,y}(x)\right].$$

We call the class  $\mathbb{E}\mathcal{F} = \{\mathbb{E}\mathcal{F}_y : y \in Y\}$  the expectation class of  $\mathcal{F}$ .

Think of  $\mathbb{E}\mathcal{F}$  and every  $\mathcal{H}_{\omega}$  as a class of functions  $h:X\to [0,1]$  indexed by Y.

### The Randomization

If the classes  $\mathcal{H} \in \mathcal{F}$  are uniformly definable in some structure M, then  $\mathbb{E}\mathcal{F}$  is definable in a structure  $M^R$  (of continuous logic) called the *randomization* of M, whose elements are M-valued random variables.

### Theorem (Ben Yaacov, Keisler)

- $\circ$  If  $\mathcal F$  is uniformly NIP/stable, then  $\mathbb E\mathcal F$  is NIP/stable.
- If a structure M is NIP/stable, then M<sup>R</sup> is NIP/stable.

# PAC Learning Expectations

### Theorem (A., Benedikt)

If each  $\mathcal{H} \in \mathcal{F}$  has  $\mathrm{FatSHDim}_{\frac{\varepsilon}{50}}(\mathcal{H}) \leq d$ , one can PAC learn the expectation class  $\mathbb{E}\mathcal{F}$  with sample complexity

$$O\left(rac{d}{\epsilon^4} \cdot \log^2 rac{d}{\epsilon} + rac{1}{\epsilon^2} \cdot \log rac{1}{\delta}
ight).$$

- FatSHDim can be used to bound Rademacher mean width
- Rademacher mean width can be used to bound sample complexity
- Adapt Ben Yaacov's proof that Gaussian mean width is preserved under randomization

# Online Learning

- At step i, an adversary chooses  $(x_i, y_i) \in X \times [0, 1]$
- Given  $x_i$ , you guess  $y_i' \approx y_i$  (you can use randomness)
- The adversary tells you  $y_i$ , penalizes you  $|y_i y_i'|$
- After *n* steps, compare to the best strategy  $y'_i = h(x_i)$  for  $h \in \mathcal{H}$ .
- Call the difference in penalty the *regret*.
- $\mathcal{H}$  is online learnable if whatever the adversary does, regret is sublinear in n.

# Online Learning Bounds

To bound regret in online learning, replace our existing notions with *sequential* versions, replacing subsets  $E \subseteq \{1, ..., n\}$  with branches of a binary tree of depth n:

### Theorem (Rakhlin, Sridharan, Tewari)

Finite  $\gamma$ -sequential-fat-shattering dimension is equivalent to online learnability, with bounds given.

Their proof goes through sequential Rademacher mean width.

# Our Online Learning Results

### Theorem (A., Benedikt)

- Stability in continuous logic is equivalent to finite  $\gamma$ -sequential-fat-shattering dimension for all  $\gamma > 0$ .
- Sequential Rademacher mean width, and thus online learnability, is preserved under randomization.

### Theorem (A., Benedikt)

The minimax regret of online learning for the randomization class of  ${\cal H}$  with  $\gamma$ -sequential-fat-shattering dimension at most d on a run of length n is at most

$$4 \cdot \gamma \cdot n + 12 \cdot (1 - \gamma) \cdot \sqrt{d \cdot n \cdot \log\left(\frac{2 \cdot e \cdot n}{\gamma}\right)}.$$

### Thank you!

For downloadable slides, see

https://awainverse.github.io/talks/learningrandom/

For valued fields enthusiasts: Talk to me about RCMVF.