CPSC 418 / MATH 318 — Introduction to Cryptography
ASSIGNMENT 1

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Problem 1 — Superencipherment for substitution ciphers, 12 marks

1. (a) *Proof.* Encryption using Shift cipher is given by $E_K(M) \equiv (M + K) \mod 26$. Given $\mathcal{M} = \mathcal{C} = \mathcal{K} = \mathbb{Z}_{26}$, $K_1, K_2 \in \mathcal{K}$ and $M \in \mathcal{M}$:
   Let $C_1 \in \mathcal{C}$ be a ciphertext that results from encrypting plaintext M with a key K_1:
   i. $E_{K_1}(M) \equiv (M + K_1) \mod 26$
   ii. $C_1 \equiv E_{K_1}$
   iii. $C_1 \equiv (M + K_1) \mod 26$
   iv. Let $C_2 = E_{K_2}(C_1)$, where $E_{K_2}(C_1) \equiv (C_1 + K_2) \mod 26$
   v. Therefore, by substituting C_1,

   $$\begin{aligned} C_2 &= C_1 + K_2 && \pmod{26} \\ &= M + (K_1 + K_2) && \pmod{26} \\ &= M + K_3 && \pmod{26} \end{aligned}$$

   Where $K_3 \in \mathcal{K}$ and $K_3 = K_1 + K_2$. Therefore, resulting key of multiple encipherment is K_3.

   □

   (b) *Proof.* Based on previous proof, superencipherment using shift cipher can be defined as follows

   $$E_{K_i}(M) = (M + \sum_{\substack{k_i \in \\ i=1}}^{n} k_i) \pmod{26}$$

   Where $M \in \mathcal{M}$, $K_i \in \mathcal{K}$, $n \in \mathbb{Z}$, and $i \geq 1, i \in \mathbb{Z}$.
   i. Base Case: let $n = 1$,

   $$\begin{aligned} C &= M + \sum_{\substack{K_i \in \mathcal{K} \\ i=1}}^{1} K_i && \pmod{26} \\ &= M + K_1 && \pmod{26} \end{aligned} \tag{1}$$

   ii. Induction hypothesis: Assume $n = m$, where $m \in \mathbb{Z}$. Therefore:

   $$C = M + \sum_{\substack{K_i \in \mathcal{K} \\ i=1}}^{m} K_i \pmod{26}$$

Let:

$$
\begin{aligned}
K^{'} &= \sum_{\substack{K_i \in \mathcal{K} \\ i=1}}^{m} K_i \\
C &= M + K^{'} \pmod{26}
\end{aligned}
\tag{2}
$$

Still results in a shift cipher, according to definition.

iii. Inductive case: According to induction hypothesis, we can show that $m+1$ holds true as well:

$$
\begin{aligned}
C &= (M + \sum_{\substack{K_i \in \mathcal{K} \\ i=1}}^{m+1} K_i) && \pmod{26} \\
&= (M + (\sum_{\substack{K_i \in \mathcal{K} \\ i=1}}^{m} K_i + \sum_{\substack{K_i \in \mathcal{K} \\ i=1}}^{1} K_i)) && \pmod{26} \\
&= (M + (\sum_{\substack{K_i \in \mathcal{K} \\ i=1}}^{m} K_i + K_1)) && \pmod{26} \\
&= (M + (K^{'} + K_1)) && \pmod{26} \\
&= (M + K^{"}) && \pmod{26}
\end{aligned}
$$

where K_1 is our base case (1) and $K^{'}$ is our induction hypothesis (2). Since $K_1 \in \mathcal{K}$ and $K^{'} \in \mathcal{K}$, therefore $K^{"} = K_1 + K^{'}, K^{"} \in \mathcal{K}$. Hence, the key of multiple encipherment is sum of all given keys.

□

2. *Proof.* Let $M \in \mathcal{M}$ be a plaintext of length $x \in \mathbb{Z}$. Let $p_0, p_1, p_2, ..., p_{x-1}$ be positions of letters in plaintext M. Given key $w_1 \in \mathcal{K}$ of length $m \in \mathbb{Z}$ and key $w_2 \in \mathcal{K}$ of length $n \in \mathbb{Z}$, let $k_0, k_1, k_2, ..., k_{m-1}$ be positions of letters in key w_1, and $l_0, l_1, l_2, ..., l_{n-1}$ be positions of letters in key w_2. To encrypt plaintext p_i, we use a key k_j, where i is letter position from 0 to $x-1$.

$$j \equiv i \pmod{m}$$

Let ciphertext C_i be ciphertext that corresponds to p_i

$$C_i \equiv p_i + k_j \pmod{26}$$

Where $k_j \equiv i \pmod{m}$.
Let ciphertext C_{2_i} be ciphertext that corresponds to C_i. Therefore, the second round of encryption, using the key w_2, results in following:

$$
\begin{aligned}
C_{2_i} &= C_i + l_j && \pmod{26} \\
&= (p_i + k_j) + l_z && \pmod{26} \\
&= p_i + (k_j + l_z) && \pmod{26}
\end{aligned}
$$

Where $z \equiv i \pmod{n}$ and $l_z \equiv i \pmod{n}$. Therefore that ensures that length of resulting key will be x. $\square$

Problem 2 — Key size versus password size, 21 marks

1. 256^8
2. (a) 94^8
   (b) $\frac{94^8}{256^8} \times 100 \approx 0.033\%$
3. Assuming that all characters are chosen equally likely, then $p(X_i) = \frac{1}{94}$. Therefore, entropy of the key space will be:

$$\begin{aligned} H(X) &= 8 \times \frac{1}{94} \log_2 94 \\ &\approx 0.56 \end{aligned}$$

4.

$$\begin{aligned} H(X) &= 8 \times \frac{1}{26} \log_2 \frac{1}{26} \\ &\approx 1.45 \end{aligned}$$

5. Given $H(X) = 128$, let $n \in \mathbb{Z}$ be a password length.
   (a) $p(X) = \frac{1}{94}$

$$\begin{aligned} n \times \frac{1}{94} \log_2(94) &= 128 \\ n &= \frac{128 \times 94}{\log_2(94)} \\ n &\approx 1836 \end{aligned}$$

   (b) $p(X) = \frac{1}{26}$

$$\begin{aligned} n \times \frac{1}{26} \log_2(26) &= 128 \\ n &= \frac{128 \times 26}{\log_2(26)} \\ n &\approx 709 \end{aligned}$$

Problem 3 — Equiprobability maximizes entropy for two outcomes, 12 marks

1.
$$\begin{aligned} H(X) &= \frac{1}{4}\log_2 4 + \frac{3}{4}\log_2\frac{4}{3} \\ &\approx 0.811 \end{aligned}$$

2. $H(X)$ is maximized if and only if all outcomes are equally likely. For any n, $H(X) = \log_2 n$ is maximal if and only if $p(X_i) = \frac{1}{n}$ for $1 \leq i \leq n$.

   Proof. Let's consider function $H(X) = p\log_2(\frac{1}{p}) + (1-p)\log_2(\frac{1}{(1-p)}) = -p\log_2 p - (1-p)\log_2(1-p)$ as a function of p:

$$\begin{aligned} H(p) &= p\log_2(\frac{1}{p}) + (1-p)\log_2(\frac{1}{(1-p)}) \\ &= -p\log_2 p - (1-p)\log_2(1-p) \end{aligned}$$

   By taking the derivative of $H(p)$, we can determine maximum of the function

$$\begin{aligned} H'(p) &= (-p\log_2 p)' - ((1-p)\log_2(1-p)) \\ &= -\frac{\ln(p)+1}{\ln(2)} + \frac{\ln(1-p)+1}{\ln(2)} \\ &= \frac{\ln(1-p)-\ln(p)}{\ln(2)} \\ &= \log_2(1-p) - \log_2(p) \end{aligned}$$

   Maximum of $H'(p) = \log_2(1-p) - \log_2(p)$:

$$\begin{aligned} \log_2(1-p) - \log_2(p) &= 0 \\ \log_2(\frac{(1-p)}{p}) &= 0 \\ 1 &= \frac{(1-p)}{p} \\ p &= 1-p \\ p &= \frac{1}{2} \end{aligned}$$

   □

3. Maximal value of $H(X)$, given $p = \frac{1}{2}$:

$$\begin{aligned} H(X) &= -\frac{1}{2}\log_2(\frac{1}{2}) - (1-\frac{1}{2})\log_2(1-\frac{1}{2}) \\ &= \frac{1}{2} + \frac{1}{2} \\ &= 1 \end{aligned}$$

Problem 4 — Conditional entropy, 12 marks

Given conditional entropy

$$H(X|Y) = \sum_{y \in Y} \sum_{x \in X} p(x, y) \log_2(\frac{1}{p(x|y)})$$

1.

2. Assuming that the system provides perfect secrecy $p(M|C) = p(M)$ and $p(M) > 0$ for all $M \in \mathcal{M}$. Given that $H(\mathcal{M}|\mathcal{C}) = \sum_{C \in \mathcal{C}} \sum_{M \in \mathcal{M}} p(M, C) \log_2(\frac{1}{p(M|C)})$ and $H(\mathcal{M}) = \sum_{M \in \mathcal{M}} p(M) \log_2(\frac{1}{p(M)})$, show that $H(\mathcal{M}|\mathcal{C}) = H(\mathcal{M})$. Assume that $H(\mathcal{M}|\mathcal{C}) = H(\mathcal{M})$:

   Proof.

   $$H(\mathcal{M}|\mathcal{C}) = H(\mathcal{M})$$

   $$\sum_{C \in \mathcal{C}} p(C) \sum_{M \in \mathcal{M}} p(M|C) \log_2(\frac{1}{p(M|C)}) = H(\mathcal{M})$$

   Given that the system provides perfect secrecy, we get

   $$\sum_{C \in \mathcal{C}} p(C) \sum_{M \in \mathcal{M}} p(M) \log_2(\frac{1}{p(M)}) = H(\mathcal{M})$$

   $$\sum_{C \in \mathcal{C}} p(C) H(\mathcal{M}) = H(\mathcal{M})$$

   Dividing both sides by $H(\mathcal{M})$, we get

   $$\sum_{C \in \mathcal{C}} p(C) = 1$$

   Therefore, as shown above $\sum_{C \in \mathcal{C}} p(C) = 1$. Using that, we can now show that $H(\mathcal{M}|\mathcal{C}) = H(\mathcal{M})$:

   $$\begin{aligned} H(\mathcal{M}|\mathcal{C}) &= \sum_{C \in \mathcal{C}} p(C) \sum_{M \in \mathcal{M}} p(M|C) \log_2(\frac{1}{p(M|C)}) \\ &= \sum_{M \in \mathcal{M}} p(M|C) \log_2(\frac{1}{p(M|C)}) \\ &= H(\mathcal{M}) \end{aligned}$$

   □

3. *Proof.* Assume the system provides perfect secrecy. Therefore, the following consition should be met $p(M|C) = p(M)$ for $M \in \mathcal{M}$ and $C \in \mathcal{C}$. Let's calculate $p(M_1|C_1)$, if the system provides perfect secrecy, $p(M_1|C_1) = p(M_1) = \frac{1}{4}$

   $$p(M_1|C_1) = p(M_1) + p(M_2) = \frac{1}{4} + \frac{1}{4} = \frac{1}{2}$$

   Therefore, the system does not provide perfect secrecy. □