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Question 1

```
{SSN}+ = {SSN, Ename, Bdate, Address, Dnumber, Dname, Dmgr_ssn}  
{Dnumber}+ = {Dnumber, Dname, Dmgr_ssn}
```

Question 2

A set G is not minimal, since dependencies in G on RHS have more than one attribute

Minimal set:

```
G' = {  
  SSN -> Ename,  
  SSN -> Bdate,  
  SSN -> Address,  
  SSN -> Dnumber,  
  DNumber -> Dname,  
  DNumber -> Dmgr_ssn  
}
```

Proof:

```
G+ = {SSN, Ename, Bdate, Address, Dnumber, Dname, Dmgr_ssn}  
G'+ = {{SSN}+, {Dnumber}+}  
      = {SSN, Ename, Bdate, Address, Dnumber, Dname, Dmgr_ssn, {Dnumber}+}  
      = {SSN, Ename, Bdate, Address, Dnumber, Dname, Dmgr_ssn}
```

Therefore $G = G'$

Question 3

1. Primary key is AB :

```
A -> {D, E} => A -> D and A -> E  
A -> D and D -> {I, J} => A -> {I, J}
```

Therefore, $A \rightarrow \{D, E, I, J\}$

$\{A, B\} \rightarrow \{B, D, E, I, J\}$

Given $\{A, B\} \rightarrow C$ and $\{A, B\} \rightarrow \{B, D, E, I, J\}$, $\{A, B\} \rightarrow \{B, C, D, E, I, J\}$

By augmentation, $\{A, B\} \rightarrow \{A, B, C, D, E, I, J\}$

Since $B \rightarrow F$ and $F \rightarrow \{G, H\}$,

we can say $B \rightarrow \{G, H\}$ and therefore $\{A, B\} \rightarrow \{A, G, H\}$

Given $B \rightarrow F$, $\{A, B\} \rightarrow \{A, F\}$ holds too

Therefore, $\{A, B\} \rightarrow \{A, B, C, D, E, F, G, H, I, J\}$

2. Decomposition of R

2NF:

$R1 = \{A, B, C\}$

$R2 = \{A, D, E, I, J\}$

$R3 = \{B, F, G, H\}$

3NF:

$R1 = \{A, B, C\}$

$R2 = \{D, I, J\}$

$R3 = \{A, D, E\}$

$R4 = \{B, F\}$

$R5 = \{F, G, H\}$

Question 4

a

Decomposition has lossless preservation property and dependency preservation property.
Each relation in BCNF

b

Decomposition has lossless preservation property and dependency preservation property. R1 is 1NF, R2 is 2NF, R3 is BCNF

C

Decomposition does not have lossless join preservation property and no dependency preservation property. R1 is in 1NF, R2 is in 1NF, R3, R4, and R5 are in BCNF