Introduction to Reversing with Z3

RPISEC

Avi Weinstock (aweinstock)

December 3, 2019

Overview

- What are SAT/SMT/Z3?
- How do the z3 python bindings work?
- Solving some small, isolated examples with Z3
- Solving a Cyberseed RE challenge with Z3
- Solving MBE lab1A with Z3

Background - What are SAT/SMT/Z3?

```
static int probable_prime(BIGNUM *rnd, int bits) {
  int i;
  uint16_t mods[NUMPRIMES];
  BN_ULONG delta;
  BN_ULONG maxdelta = BN_MASK2 - primes[NUMPRIMES - 1];
  char is_single_word = bits <= BN_BITS2;

again:
  if (!BN_rand(rnd, bits, BN_RAND_TOP_TWO, BN_RAND_BOTTOM_ODD)) {
    return 0;
  }

  /* we now have a random number 'rnd' to test. */
  for (i = 1; i < NUMPRIMES; i++) {
    BN_ULONG mod = BN_mod_word(rnd, (BN_ULONG)primes[i]);
    if (mod == (BN_ULONG)-1) {
      return 0;
    }
    mods[i] = (uint16_t)mod;
  }
  /* If bits is so small that it fits into a single word then we
   * additionally don't want to exceed that many bits. */
  if (is_single_word) {
    BN_ULONG size_limit;
    if (bits == BN_BITS2) {
      /* Avoid undefined behavior. */
      size_limit = ~((BN_ULONG)0) - get_word(rnd);
    } else {
      size_limit = (((BN_ULONG)1) << bits) - get_word(rnd) - 1;
    }
    if (size_limit < maxdelta) {
      maxdelta = size_limit;
    }
  }
  delta = 0;

loop:
  if (is_single_word) {
    BN_ULONG rnd_word = get_word(rnd);

    /* In the case that the candidate prime is a single word then
     * we check that:
     *   1) It's greater than primes[i] because we shouldn't reject
     *      3 as being a prime number because it's a multiple of
     *      three.
     *   2) That it's not a multiple of a known prime. We don't
     *      check that rnd-1 is also coprime to all the known
     *      primes because there aren't many small primes where
     *      that's true. */
    for (i = 1; i < NUMPRIMES && primes[i] < rnd_word; i++) {
      if ((mods[i] + delta) % primes[i] == 0) {
        delta += 2;
        if (delta > maxdelta) {
          goto again;
        }
        goto loop;
      }
    }
  } else {
    for (i = 1; i < NUMPRIMES; i++) {
      /* check that rnd is not a prime and also
       * that gcd(rnd-1, primes) == 1 (except for 2) */
      if (((mods[i] + delta) % primes[i]) <= 1) {
        delta += 2;
```

What is Z3?

- SAT & SMT solver developed and maintained by Microsoft Research
- Libre and Open Source (MIT Licensed)
- C++, with python bindings (`pip install z3-solver`)
- Based on the CDCL algorithm

What is SAT?

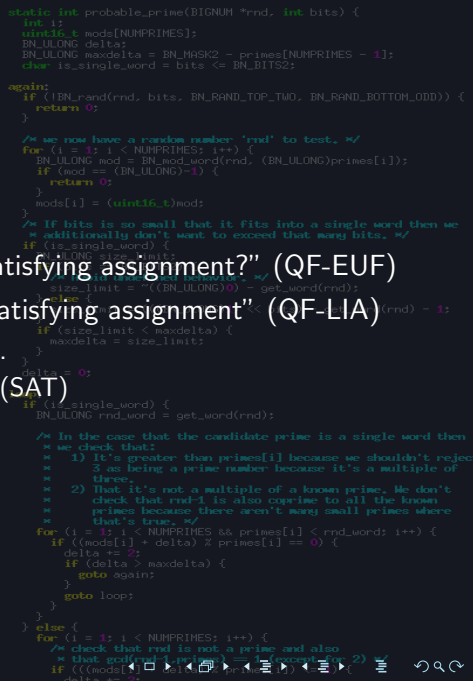
- SAT is the boolean SATisfiability problem
- e.g. "Does the formula $(x \vee \neg y \vee z) \wedge (\neg x \vee y)$ have a satisfying assignment?"
  - $(\neg, \wedge, \vee)$ mean (NOT, AND, OR)
- Brute forceable in $O(2^n)$ by trying all combinations of $\{0, 1\}$ for all variables
- NP-Complete
  - Con: Impossible[1] to solve quickly[2]
  - Pro: Many problems can be expressed as SAT instances, so heuristics for SAT can help solve many problems

[1] Unless P=NP

[2] In polynomial time

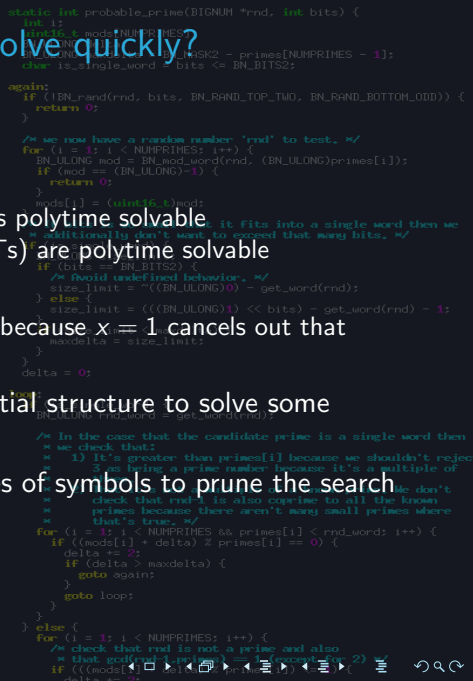
What is SMT?

- SMT is Satisfiability Modulo Theories
- "Does $(f(x,y) \vee z) \wedge (\neg g(x) = f(x,x))$ have a satisfying assignment?" (QF-EUF)
- "Does $(2 * x + y \leq z) \wedge (x + 3 * y \geq z)$ have a satisfying assignment" (QF-LIA)
- Allows more compact translation of problems, e.g.
  - $x = 1 \vee x = 2 \vee x = 3 \vee \ldots \vee x = 99 \vee x = 100$ (SAT)
  - $1 \leq x \wedge x \leq 100$ (SMT)
- Also NP-Complete



Why are SAT/SMT useful if they're hard to solve quickly?

- Not all problems are as hard as the hardest ones
  - 2-SAT (each clause having at most 2 variables) is polytime solvable
  - Monotone circuits (only ANDs and ORs, no NOTs) are polytime solvable
- It's often possible to prune the search space
  - e.g. $x \vee \varphi(a, b, c, \ldots)$ is solvable regardless of φ because $x = 1$ cancels out that subterm
- Algorithms like DPLL and CDCL make use of partial structure to solve some instances faster than others
- SMT can make use of the rules for the extra types of symbols to prune the search space at a higher level

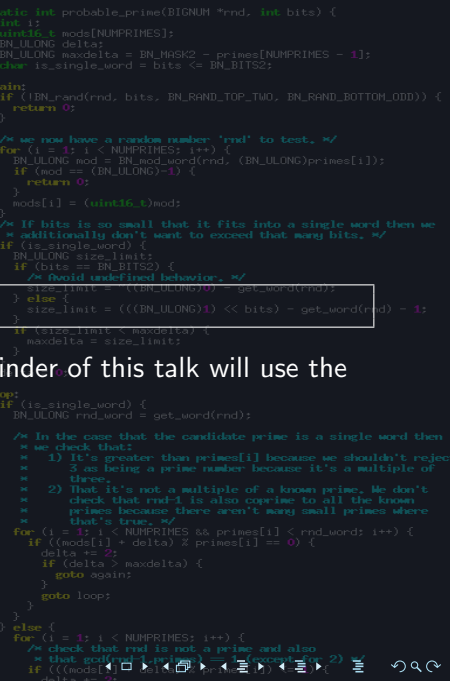


Any questions so far?

The z3 python bindings

Installing and importing z3

- `pip install z3-solver`
- `import z3`
- Some people do `from z3 import *`, but the remainder of this talk will use the qualified import version



Variables and Values

```
import z3
x = z3.Int('x')
y = z3.BitVec('y', 32)
z = z3.BitVec('z', 16)
w = z3.Real('w')
```

- Variables are symbolic
- Operator overloading allows creation of constraints

```
import z3
a = z3.IntVal(42)
b = z3.BitVecVal(0xdeadbeef, 32)
c = z3.BitVecVal(0xbeef, 16)
d = z3.RealVal(__import__('math').pi)
```

- Values are concrete
- Dynamic type checks on them are stricter than Python's defaults

Solver objects

```
import z3
solver = z3.Solver()
```

- The `Solver` class collects constraints/equations to solve
- `solver.add` adds constraints to the current collection
- `solver.check()` checks if the current constraints are solvable, returning one of `{z3.sat, z3.unsat, z3.unknown}`
- After `solver.check()` returns `z3.sat`, `z3.model()` will give you the values that make the constraints true

Any questions so far?

Solving some small examples with Z3

Using Z3 on small examples: 1/3

- $(x \vee \neg y \vee z) \wedge (\neg x \vee y)$
- ```
  import z3
  solver = z3.Solver()
  x, y, z = z3.Bools('x y z')
  solver.add(z3.And(z3.Or(x, z3.Not(y), z), z3.Or(z3.Not(x), y)))
  if solver.check().r == 1:
      print(solver.model())
  ```
- ```
  [z = False, y = False, x = False]
  ```

Using Z3 on small examples: 2/3

- $(2 * x + y \leq z) \wedge (x + 3 * y \geq z) \wedge (z > 1)$
- ```
  import z3
  solver = z3.Solver()
  x, y, z = z3.Ints('x y z')
  solver.add(2*x+y <= z)
  solver.add(x+3*y >= z)
  solver.add(z > 1)
  if solver.check().r == 1:
      print(solver.model())
  ```
- ```
  [z = 5, y = 1, x = 2]
  ```

Using Z3 on small examples: 3/3

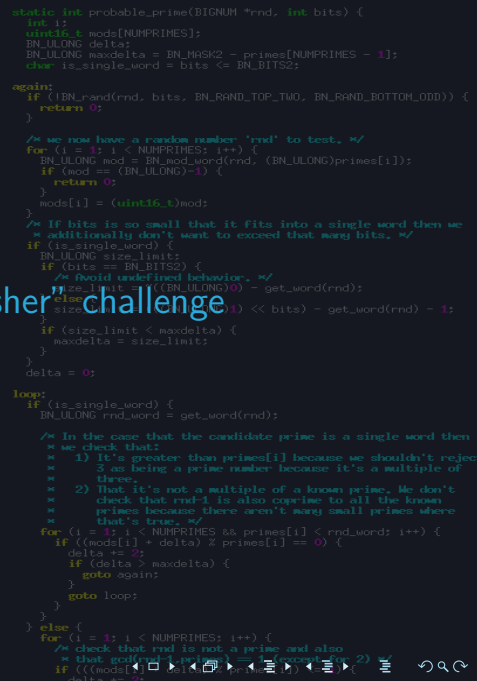
- $$\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix} \begin{bmatrix} x \\ y \\ z \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix}$$

```
import z3
x, y, z = z3.Reals('x y z')
solver = z3.Solver()
solver.add(1*x + 2*y + 3*z == 1)
solver.add(4*x + 5*y + 6*z == 2)
solver.add(7*x + 8*y + 9*z == 3)
status = solver.check()
assert solver.check().r == 1
print(solver.model())
solver.add(x != 0, y != 0, z != 0)
assert solver.check().r == 1
print(solver.model())
```

```
[z = 1/3, y = 0, x = 0]
[z = -1, y = 8/3, x = -4/3]
```

Any questions so far?

Solving Cyberseed 2019's "Hasher" challenge



Cyberseed Hasher - Problem description

- This year's Cyberseed, one of the challenges was a Java crackme
- We were only given a class file, no source
- DDG'ing "java decompiler online" found `https://devtoolzone.com/decompiler/java`
- Now we had source :)

Cyberseed Hasher - Decompiled source

```
public class Hasher {
    private static boolean hash(final String s) {
        int n = 7;
        final int n2 = 593779930;
        for (int i = 0; i < s.length(); ++i) {
            n = n * 31 + s.charAt(i);
        }
        return n == n2;
    }

    public static void main(final String[] array) {
        if (array.length != 1) {
            System.out.println("Usage: java Hasher <password>");
            System.exit(1);
        }
        if (hash(array[0])) {
            System.out.println("Correct");
        }
        else {
            System.out.println("Incorrect");
        }
    }
}
```

Cyberseed Hasher - Z3 script: 1/2

```
import z3
wanted_length = 6

names = ['x{i}'.format(i=i) for i in range(wanted_length)]
vars = [z3.Int(n) for n in names]

expr = 7
for i in range(wanted_length):
    expr *= 31
    expr += vars[i]

prob = z3.Solver()
prob.add((expr % (2**32)) == 593779930)
```

- Concrete input length: `z3.Array` exists, but is more expensive
- Only use `z3.Array` if you need symbolic indexing
- `z3.Int(n)` is faster than `z3.BitVec(n, 32)` here, needed trial and error for this

Cyberseed Hasher - Z3 script: 2/2

```
for v in vars:
    prob.add(ord('0') <= v)
    prob.add(v <= ord('z'))

res = prob.check()
print('z3 check: %r' % (res,))
if res.r == 1:
    soln = prob.model()
    print('z3 solution: %r' % (soln,))
    print(repr(''.join(chr(soln[v].as_long()) for v in vars)))
```

```
z3 check: sat
z3 solution: [x3 = 74, x2 = 98, x1 = 114, x5 = 48, x4 = 51, x0 = 100]
'drbJ30'
```

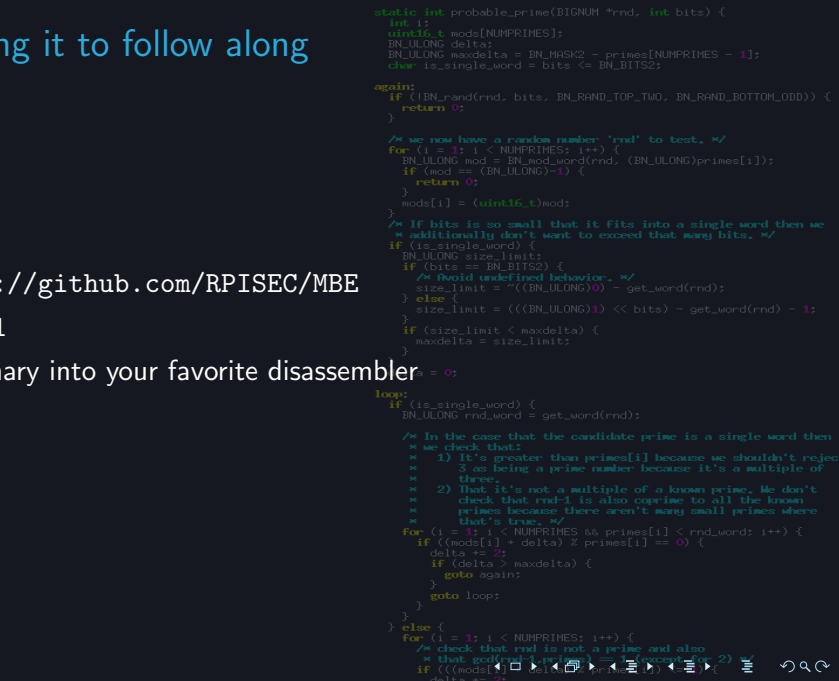
- Requiring that `ord('0') <= v <= ord('z')` mostly-guarantees alphanumeric input (there's a few exceptions in the middle)

Any questions so far?

Solving Modern Binary Exploitation's lab1A

MBE Lab1A - Acquiring it to follow along

- `git clone https://github.com/RPISEC/MBE`
- `cd MBE/src/lab01`
- Load the `lab1A` binary into your favorite disassembler



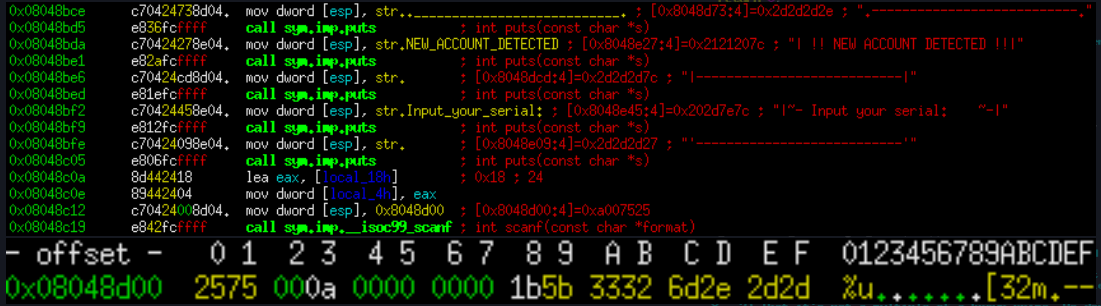
MBE Lab1A - Just Running It

```
avi@aweinstock-debian-ii:~/Documents/cloned-repos/MBE/src/lab01$ ./lab1A
.---------------------------.
|---------  RPISEC  --------|
|+ SECURE LOGIN SYS v. 3.0 +|
|---------------------------|
|~- Enter your Username:  ~-|
'---------------------------'
username
.---------------------------.
| !! NEW ACCOUNT DETECTED !!|
|---------------------------|
|~- Input your serial:    ~-|
'---------------------------'
password
avi@aweinstock-debian-ii:~/Documents/cloned-repos/MBE/src/lab01$ echo $?
1
```

MBE Lab1A - Username Entry

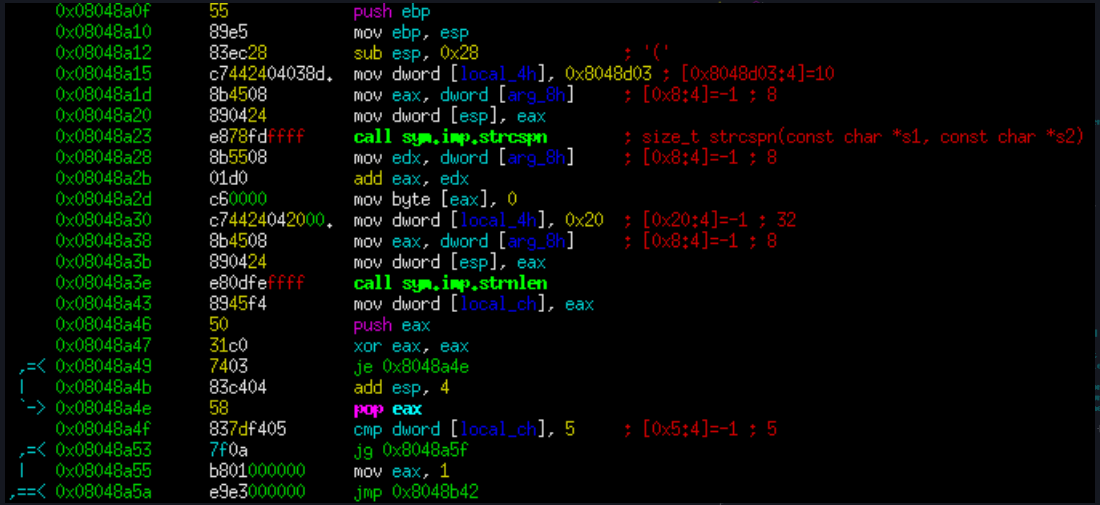
```
0x08048b69      c70424738d04.  mov dword [esp], str.+----------------------------+ ; [0x8048d73:4]=0x2d2d2d2e ; "+----------------------------+"
0x08048b70      e89bfcffff     call sym.imp.puts          ; int puts(const char *s)
0x08048b75      c70424918d04.  mov dword [esp], str.RPISEC ; [0x8048d91:4]=0x2d2d2d7c ; "|---------  RPISEC  --------|"
0x08048b7c      e88ffcffff     call sym.imp.puts          ; int puts(const char *s)
0x08048b81      c70424af8d04.  mov dword [esp], str.SECURE_LOGIN_SYS_v._3.0 ; [0x8048daf:4]=0x53202b7c ; "|+ SECURE LOGIN SYS v. 3.0 +|"
0x08048b88      e883fcffff     call sym.imp.puts          ; int puts(const char *s)
0x08048b8d      c70424cd8d04.  mov dword [esp], str.       ; [0x8048dcd:4]=0x2d2d2d7c ; "|--------------------------|"
0x08048b94      e877fcffff     call sym.imp.puts          ; int puts(const char *s)
0x08048b99      c70424eb8d04.  mov dword [esp], str.Enter_your_Username:_ ; [0x8048deb:4]=0x202d7e7c ; "|~ Enter your Username:  ~-|"
0x08048ba0      e86bfcffff     call sym.imp.puts          ; int puts(const char *s)
0x08048ba5      c70424098e04.  mov dword [esp], str.       ; [0x8048e09:4]=0x2d2d2d27 ; "'--------------------------'"
0x08048bac      e85ffcffff     call sym.imp.puts          ; int puts(const char *s)
0x08048bb1      a160b00408     mov eax, dword [obj.stdin]  ; [0x804b060:4]=0
0x08048bb6      89442408       mov dword [local_8h], eax
0x08048bba      c74424042000.  mov dword [local_4h], 0x20  ; [0x20:4]=-1 ; 32
0x08048bc2      8d44241c       lea eax, [local_1ch]        ; 0x1c ; 28
0x08048bc6      890424         mov dword [esp], eax
0x08048bc9      e802fcffff     call sym.imp.fgets          ; char *fgets(char *s, int size, FILE *stream)
```

MBE Lab1A - Serial Entry



MBE Lab1A - Calling the authentication routine

MBE Lab1A - auth() 1/6: String processing and antidecomp



MBE Lab1A - auth() 2/6: Antidebugging with ptrace

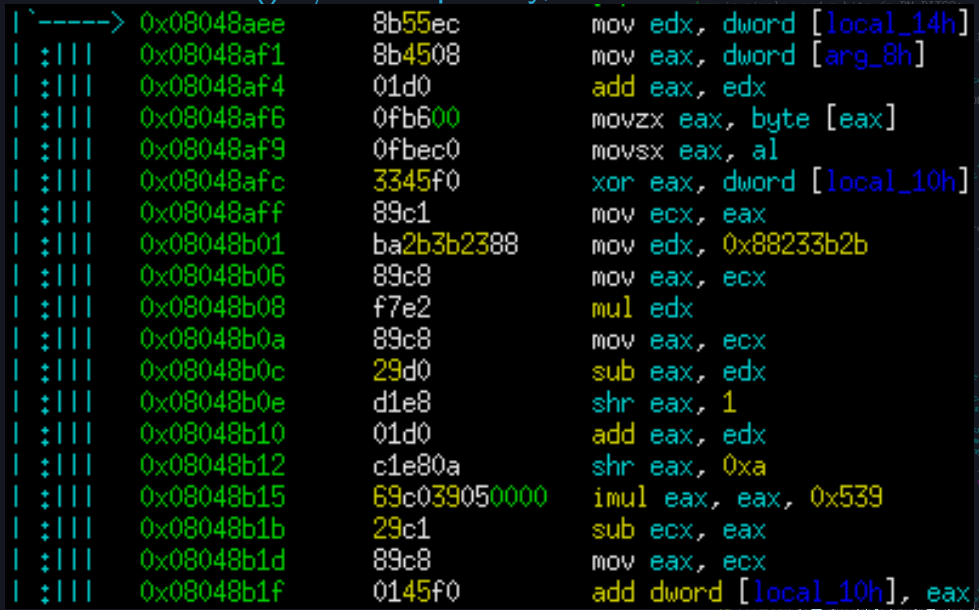
```
0x08048a5f      c744240c0000.   mov dword [local_ch_2], 0
0x08048a67      c74424080100.   mov dword [local_8h], 1
0x08048a6f      c7442404 0000.  mov dword [local_4h], 0
0x08048a77      c70424000000.   mov dword [esp], 0
0x08048a7e      e8edfdffff      call sym.imp.ptrace
0x08048a83      83f8ff          cmp eax, 0xffffffffffffffff
0x08048a86      752e            jne 0x8048ab6
0x08048a88      c70424088d04.   mov dword [esp], str.e_32m._____________________________
0x08048a8f      e87cfdffff      call sym.imp.puts           ; int puts(const char *s)
0x08048a94      c704242c8d04.   mov dword [esp], str.e_31m_____TAMPERING_DETECTED ; [0x
0x08048a9b      e870fdffff      call sym.imp.puts           ; int puts(const char *s)
0x08048aa0      c70424508d04.   mov dword [esp], str.e_32m   ; [0x8048d50:4]=0x32335b1b
0x08048aa7      e864fdffff      call sym.imp.puts           ; int puts(const char *s)
0x08048aac      b801000000      mov eax, 1
0x08048ab1      e98c000000      jmp 0x8048b42
```

MBE Lab1A - auth() 3/6: Pre-loop math

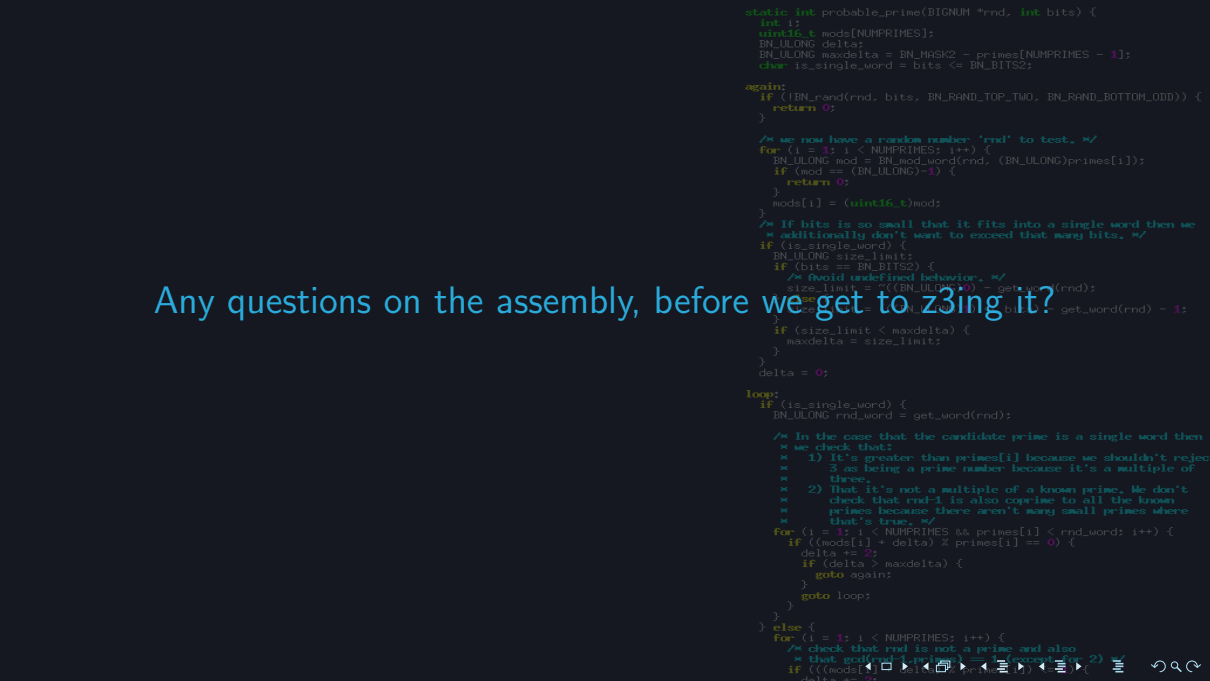
```
0x08048ab6      8b4508          mov eax, dword [arg_8h]
0x08048ab9      83c003          add eax, 3
0x08048abc      0fb600          movzx eax, byte [eax]
0x08048abf      0fbec0          movsx eax, al
0x08048ac2      3537130000      xor eax, 0x1337
0x08048ac7      05eded5e00      add eax, 0x5eeded
0x08048acc      8945f0          mov dword [local_10h], eax
```

MBE Lab1A - auth() 4/6: Loop header, restricting chars

MBE Lab1A - auth() 5/6: Loop body, much math



MBE Lab1A - auth() 6/6: Loop footer, return targets


Any questions on the assembly, before we get to z3ing it?

MBE Lab1A - Z3ing auth() 1/7: Setting up variables

```
import z3
solver = z3.Solver()
wanted_length = 8
assert wanted_length > 5 # checked at 0x08048a4f
sym_username = [z3.BitVec('x{i}'.format(i=i), 8) for i in range(wanted_length)]
sym_serial = z3.BitVec('serial', 32)
```

- 8-bit entries for each character
- 32-bit serial number
- Concrete input length: `z3.Array` is more expensive

MBE Lab1A - Z3ing auth() 2/7: Translating the pre-loop math

```
0x08048ab6      8b4508          mov eax, dword [arg_8h]
0x08048ab9      83c003          add eax, 3 ; eax = (arg_8h + 3)
0x08048abc      0fb600          movzx eax, byte [eax] ; eax = *(uint8_t*)(arg_8h+3)
0x08048abf      0fbec0          movsx eax, al ; eax = (int32_t)*(uint8_t*)(arg_8h+3)
0x08048ac2      3537130000      xor eax, 0x1337
0x08048ac7      05eded5e00      add eax, 0x5eeded
0x08048acc      8945f0          mov dword [local_10h], eax
```

```
eax = z3.SignExt(24, sym_username[3]) # (int32_t)*((uint8_t*)arg_8h + 3)
eax ^= z3.BitVecVal(0x1337, 32)
eax += z3.BitVecVal(0x5eeded, 32)
local_10h = eax
```

- We're wrapping concrete values in `z3.BitVecVal` so that wrapping/truncation happens the x86 way
- If we were using python longs here, we'd have to manually mask them back into range

MBE Lab1A - Z3ing auth() 3/7: Translating the loop header/footer

The following x86:

```
0x08048acf      c745ec000000.   mov dword [local_14h], 0
0x08048ad6      eb4e            jmp 0x8048b26
...
0x08048b22      8345ec01        add dword [local_14h], 1
0x08048b26      8b45ec          mov eax, dword [local_14h]
0x08048b29      3b45f4          cmp eax, dword [local_ch]
0x08048b2c      7caa            jl 0x8048ad8
```

Translates to the following C:

```
for(int local_14h=0; local_14h < local_ch; local_14h++) {
    /* loop body is between 0x08048ad6 and 0x08048b22, exclusive on both ends */
}
```

So we'll gloss that as the following in Python

```
local_ch = len(sym_username) # this is set by the strnlen at 0x08048a3e
for local_14h in range(local_ch):
    pass # we'll translate the loop body here
```

MBE Lab1A - Z3ing auth() 4/7: Translating the loop body: 1/2

x86:

```
0x08048ad8      mov edx, dword [local_14h]
0x08048adb      mov eax, dword [arg_8h]
0x08048ade      add eax, edx
0x08048ae0      movzx eax, byte [eax]
0x08048ae3      cmp al, 0x1f
0x08048ae5      jg 0x8048aee
```

```
0x08048aee      mov edx, dword [local_14h]
0x08048af1      mov eax, dword [arg_8h]
0x08048af4      add eax, edx
0x08048af6      movzx eax, byte [eax]
0x08048af9      movsx eax, al
0x08048afc      xor eax, dword [local_10h]
```

```
0x08048aff      mov ecx, eax
0x08048b01      mov edx, 0x88233b2b
0x08048b06      mov eax, ecx
```

Python:

```
solver.add(sym_username[local_14h] > 0x1f)
```

```
eax = z3.SignExt(24, sym_username[local_14h])
eax ^= local_10h
```

```
ecx = eax
edx = z3.BitVecVal(0x88233b2b, 32)
eax = ecx
```

MBE Lab1A - Z3ing auth() 5/7: Translating the loop body: 2/2

x86:

Python:

0x08048b08	mul edx

```
mul_result = z3.ZeroExt(32, eax) * z3.ZeroExt(32, edx)
edx = z3.Extract(63, 32, mul_result)
eax = z3.Extract(31, 0, mul_result)
```

0x08048b0a	mov eax, ecx
0x08048b0c	sub eax, edx
0x08048b0e	shr eax, 1
0x08048b10	add eax, edx
0x08048b12	shr eax, 0xa

```
eax = ecx
eax -= edx
eax = eax >> 1
eax += edx
eax = eax >> 0xa
```

0x08048b15	imul eax, eax, 0x539

```
eax = z3.Extract(31, 0, z3.SignExt(32, eax) * 0x539)
```

0x08048b1b	sub ecx, eax
0x08048b1d	mov eax, ecx
0x08048b1f	add dword [local_10h], eax

```
ecx -= eax
eax = ecx
local_10h += eax
```

MBE Lab1A - Z3ing auth() 6/7: Solving for a valid serial

```
solver.add(sym_serial == local_10h) # outside the loop
solver.push() # backtracking point for next demo
username = 'username'
for (x, y) in zip(username, sym_username):
    solver.add(ord(x) == y)
assert solver.check().r == 1
model = solver.model()
serial = model.evaluate(sym_serial)
print('serial for name %r is %r' % (username, serial))
```

```
serial for name 'username' is 6234463
```

MBE Lab1A - Z3ing auth() 7/7: Solving for a valid username

```
solver.pop() # remove the constraints on the provided username
solver.add(sym_serial == serial+1)
assert solver.check().r == 1
model = solver.model()
username2 = ''.join(chr(model.evaluate(x).as_long()) for x in sym_username)
serial2 = model.evaluate(sym_serial)
print('serial for name %r is %r' % (username2, serial2))
```

```
serial for name 'sEa2):2-' is 6234464
```

Any questions overall?

```
static int probable_prime(BIGNUM *rnd, int bits) {
  int i;
  uint16_t mods[NUMPRIMES];
  BN_ULONG delta;
  BN_ULONG maxdelta = BN_MASK2 - primes[NUMPRIMES - 1];
  char is_single_word = bits <= BN_BITS2;

again:
  if (!BN_rand(rnd, bits, BN_RAND_TOP_TWO, BN_RAND_BOTTOM_ODD)) {
    return 0;
  }

  /* we now have a random number 'rnd' to test. */
  for (i = 1; i < NUMPRIMES; i++) {
    BN_ULONG mod = BN_mod_word(rnd, (BN_ULONG)primes[i]);
    if (mod == (BN_ULONG)-1) {
      return 0;
    }
    mods[i] = (uint16_t)mod;
  }
  /* If bits is so small that it fits into a single word then we
   * additionally don't want to exceed that many bits. */
  if (is_single_word) {
    BN_ULONG size_limit;
    if (bits == BN_BITS2) {
      /* Avoid undefined behavior. */
      size_limit = ~(BN_ULONG)0 - get_word(rnd);
    } else {
      size_limit = (((BN_ULONG)1) << bits) - get_word(rnd) - 1;
    }
    if (size_limit < maxdelta) {
      maxdelta = size_limit;
    }
  }
  delta = 0;

loop:
  if (is_single_word) {
    BN_ULONG rnd_word = get_word(rnd);

    /* In the case that the candidate prime is a single word then
     * we check that:
     *   1) It's greater than primes[i] because we shouldn't reject
     *      3 as being a prime number because it's a multiple of
     *      three.
     *   2) That it's not a multiple of a known prime. We don't
     *      check that rnd-1 is also coprime to all the known
     *      primes because there aren't many small primes where
     *      that's true. */
    for (i = 1; i < NUMPRIMES && primes[i] < rnd_word; i++) {
      if ((mods[i] + delta) % primes[i] == 0) {
        delta += 2;
        if (delta > maxdelta) {
          goto again;
        }
        goto loop;
      }
    }
  } else {
    for (i = 1; i < NUMPRIMES; i++) {
      /* check that rnd is not a prime and also
       * that gcd(rnd-1, primes) == 1 (except for 2) */
      if (((mods[i] + delta) % primes[i]) <= 1) {
        delta += 2;
```

Resources

- https://github.com/Z3Prover/z3/
- https://pypi.org/project/z3-solver/
- https://rise4fun.com/Z3/tutorialcontent/guide
- https://en.wikipedia.org/wiki/Satisfiability_modulo_theories
- https://github.com/RPISEC/MBE