```
atic int probable_prime(BIGNUM *rnd, int bits) {
int 1;
in
```

Symbolic Execution with Anground the RPISEC RPISEC

/* If bits is so small that it fits into a single word then a widthionally don't wont to exceed that many bits, */
if (is_single_word) {
 B.U.D.W. size_limits) {
 W | howeld madefined behavior. */
 isze_limits = "(CBR_ULDNRO) - get_word(rnd);

Avi Weinstock (aweinstock), Luke Biery (tiecoon)

December 6, 2019

BRULUONG rnd_word = get_word(rnd);
// In the case that the candidate prime is a single word
// we we check that:
// I. I's greater than primeo[i] because we shouldn't
// S as being a prime number because it's a subtipl
// S as being a prime number because it's a subtipl
// S as being a prime number because it's a subtipl
// S as being a prime number because it's a subtipl
// S as being a prime number because it's as subtipl
// S in the transport of a more prime, Me do
// S in the transport of a more rime, and it's
// S in the transport of a more than a main prime who
// S in the transport of a more than a main prime who
// S in the transport of a more rime for a more rime.
// S in the transport of a more rime in a single word
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the transport of a more rime.
// S in the t

Overview

- ▶ What is Symbolic Execution? What techniques does it compete with?
- How symbolic execution works (theory)
- How symbolic execution works (Angr commands)
- Solving MBE lab1A with Angr

```
Background - What it is and what is the problem space? - at word (red) - 12
```

What is Symbolic Execution?

```
unition to mode(NUMPRIMES):
BN_ULONG delte:
BN_ULONG delte:
BN_ULONG description to be a subject to be a subje
```

- Executes a program with symbolic data (usually input)
- Instead of having concrete data in each variable/address,
 variables/addresses store trees of what to do with the input

```
### Decorate requirements of the considerable prime is a single word then we check that:

### It's greater than primes[i] because we shouldn't reject the sering a prime immber because it's a multiple of the sering a prime immber because it's a multiple of the sering a prime immorphise. We don't check that run't is also conside to all the known is that's true, where aren't many small primes where the series of the serie
```

What problems does Symbolic Execution solve? **RILLION GOOD (NAMPRINES); ** PRINCE OF THE PROBLEM OF THE PROBLE

```
(TBL:mand(mnd, bits, EN_RAND_TOP_TWO, EN_RAND_BOTTON_ODD))

return 0;

As we now have a random number 'nnd' to test. */

for (i = i; i < NUMPRINES: i +) (
BR_ULUNG sod = BR_ucod_sord(nnd, (BR_ULUNG)primes[i]);

if (mod == (BR_ULUNG)-i) (
    return 0;

mods[i] = (uinti6_t)mod;

/* If bits is so small that it fits into a single word then we abilitionally don't work to exceed that many bits, */
```

- What input to provide to reach/avoid a specific line of code?
- ► How is a value deep in the program affected by some specific input? < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () < () <
- Do any inputs lead to any crash?
- On a crashing input, what registers are controlled by the input?

IN_ULUNG rnd_word = get_word(rnd);

```
** In the case that the cardinate prime is a single word then
** we check that:

** 1) It's greater than primes[i] because we shouldn't rejec

** 3 as being a prime number because it's a multiple of

** 2) That it's not a multiple of a known prime, be don't

** check that rud-1 is also coprime to all the known

** primes because there aren't many small primes where

** that's true, */

For (i = 1; i < NUMPRIMES St primes[i] < rnd_word; i++) (

if (delta > madelta) {

goto spain;

} clase {

For (i = 1; i < NUMPRIMES; i++) {

/* check that rud is not a prime and also

** that soft(primes); i++) {

/* check that rud is not a prime and also

** that soft(primes); i++) {

/* check that rud is not a prime and also

** that soft(primes); i++) {

/* check that rud is not a prime and also

** that soft(primes); i++) {
```

Symbolic Execution vs Fuzzing

Symbolic Execution	Fuzzing (uintic t) mod:
+ Explores all inputs	- Only explores table to the second data way bits. */
+ Very detailed output	- Only learn crash vist non-crash - 6-t_word(rnd);
	+ Uses around as much memory time as target program

- ▶ Symbolic execution can the path if(input == 0xdeadbeefdeadbeef) { ... }
- Even coverage-guided fuzzing will only find it $\frac{1}{264}$ of the time¹

```
* ue check that:

* 1) It's greater than primes[i] because we shouldn't rejet

* 5 as being a prime number because it's a multiple of

* 5 as being a prime number because it's a multiple of

* 2) Ilbat it's not a multiple of a loosen prime, be don't

check that red-1 is also coprime to all the known

primes became there aren't hangs small primes where

for (i = 1; i < NUMPRIMES as primes[i] < red_word; i++) (

if (mods[i] + delta) × primes[i] == 0) (

delta += 2;

if (delta > maxdelta) (

gotto sgalin'

gotto loop;

}
else {

for (i = 1; i < NUMPRIMES; i++) (

/* debeck that rud is not a prime and also
```

¹Unless the compare is digit-by-digit

Any questions so far?

```
How symbolic execution works in general (CRIC) (CRI
```

Setting up a state for symbolic execution

```
import z3
registers = ['eax', 'ebx', 'ecx', 'edx', 'ebp', 'esp'] # and so on
symstate = {reg: z3.BitVec(reg, 32) for reg in registers} registers}
symstate['memory'] = z3.Array('memory', z3.BitVecSort(32), z3.BitVecSort(8))
```

- ▶ Note that the z3 variable eax in the model will be the starting value of eax
- symstate['eax'] will be mutated throughout the computation, and will contain an expression corresponding to the ending value of eax

```
* S as being a prime number because it's a multiple of a bross prime, be don't be a close that rule! I had a convine to all the bross prime because there are no second there are the many mall primes because there are not many mall primes where the converted to all the bross of the converted to a converted
```

z3. Array vs dict of z3. BitVec for representing memory

- memory = z3.Array('memory', z3.BitVecSort(32), z3.BitVecSort(8))
 symbolically represents an array of 2³² bytes (around 4GB)
- ▶ z3.Store(memory, index, value) represents a modified memory (with value written to index), even with *symbolic* index and value
- memory[index] represents a read from memory, even if index is symbolic
- memory = {i: z3.BitVec('mem[{i}]'.format(i=i), 8) for i in idxs} only allows concrete indices, while still allowing symbolic values, and is more efficient when we know we won't have symbolic-indexed reads/writes

Symbolically executing branch-free code

```
int 1:
int16.t mods[NUMPRIMES];
BN.ULONG delta;
BN.ULONG moxidata = BN.MASK2 - primes[NUMPRIMES - 1];
dwr is_single_word = bits <= BN.DITS2;
graint
if (IRL_rand(rnd, bits, BN_SAND_TOP_TWO, BN_SAND_BOTTOM_ODD)) (
    return 0;
}
/* we now have a random number 'rnd' to test, */
for (1 = 1: i < NUMPRIMES: ;**) (</pre>
```

- ► Translate arithmetic, indexing, etc into SMT constraints
- Angr internally uses VEX for this instead of translating x86 directly

```
mov eax, ebx

symstate['eax'] = symstate['ebx'] symstate['edx'] symstate['edx'] symstate['edx'] symstate['esp'] +0x10

al = z3.Extract(7, 0, symstate['eax']) symstate['memory'] = z3.Store(symstate['memory'], esp_10, al)

star_eax = z3.Select(symstate['memory'], eax) symstate['eax'] symstate['eax'] = z3.SignExt(24, star_eax) symstate['memory'], eax) symstate['eax'] = z3.SignExt(24, star_eax) symstate['eax'] syms
```

Handling symbolic reads with z3.Array vs z3.BitVec

again:
 if (!BN_rand(rnd, bits, BN_RAND_TOP_TWO, BN_RAND_BOTTOM_ODD))
 return 0;

C:

```
tmp = username[i];
tmp ^= serial;
```

Assembly:

```
        0x08048aee
        mov edx, dword [local_14h]

        0x08048af1
        mov eax, dword [arg_8h]

        0x08048af4
        add eax, edx

        0x08048af6
        movzx eax, byte [eax]

        0x08048af9
        movsx eax, al

        0x08048afc
        xor eax, dword [local_10h]
```

List of z3.BitVec:

```
eax = z3.SignExt(24, sym_username[local_14h])
eax ^= local_10h
```

z3.Array:

```
local_14 = symstate['esp']+0x14 # &i
symstate['edx'] = symstate['memory'][local_14]
arg_8 = symstate['ebp']+0x8 # &username
symstate['eax'] = symstate['memory'][arg_8]
symstate['eax'] += symstate['edx']
symstate['eax'] = z3.ZeroExt(24, symstate['eax'])
al = z3.Extract(7, 0, symstate['eax'])
symstate['eax'] = z3.SignExt(24, al)
local_10 = symstate['esp']+0x10 # &serial
symstate['eax'] ^= symstate['memory'][local_10]
```

Any questions so far?

Symbolically executing branchs - Graphically

```
int f(int x, int y) {
                                                         x = x_0, y = y_0
    if (x > 3) {
         x += 1:
                                                x > 3
    } else {
         y = 2*y+3;
                                       x = x_0 + 1, \ y = y_0
                                                                        x = x_0, v = 2 * v_0 + 3
    if(y != 0) {
         x /= y;
    } else {
                            x = \frac{x_0 + 1}{x_0 + 1}
         x *= 2;
                                             x = 2 * (x_0 + 1)
                                                                                            x = 2 * x_0
                                    y<sub>0</sub>
    return x + v:
```

```
int f(int x, int y) {
    if (x > 3) {
        x += 1;
    } else {
        v = 2*v+3:
    if(y != 0) {
        x /= y;
    } else {
        x *= 2:
    return x + v:
```

```
import z3
x0. v0 = z3.Ints('x0 v0')
states, newstates = [(x0, y0, z3.Solver())], []
for (x, y, s) in states:
  t = s. deepcopy ()
  s.add(x > 3); newstates.append((x+1,\sqrt{x},\sqrt{x}))
  t.add(z3.Not(x > 3)); newstates.append((x, 2*y+3, t))
states, newstates = newstates, []
for (x, y, s) in states:
  t = s._deepcopy_()
  s.add(v != 0); newstates.append((x/v, v, s))
  t.add(z3.Not(y != 0)); newstates.append((2*x, y, t))
for (x, y, s) in newstates:
  print('x: %r; y: %r; s: %r; check: %r' % (x, y, s, s.check(
  if s.check() == z3.sat:
    m = s.model()
    print('m: \( \frac{1}{3}r; \) x: \( \frac{1}{3}r; \) y: \( \frac{1}{3}r' \) \( \frac{1}{3}m.evaluate(x), \) m.evaluate(y)))
    print('-'*5)
```

Symbolically executing loops

```
i \stackrel{\rightarrow}{=} 0
                                                                                                  ret
                                                           i < n
void memcpy(
    char *dest.
                                       mem_1 = Store(mem_0, dst + 0, mem_0[src + 0])
    const char *src,
    size_t n) {
    for(size_t i=0; i<n; i++) {</pre>
        dest[i] = src[i];
                                       mem_2 = Store(mem_1, dst + 1, mem_1[src + 1])
```

Any questions so far?

How to use Angr for symbolic execution (CIRCLELING) (1) - get_word(rind);

Angr

```
intile int probable_prine(BIONUM frnd, int bits) {
    intile is model NUMPRINES];
    BN_ULONG delta;
    RN_ULONG modelta = BN_MASK2 - primes[NUMPRINES - 1];
    disw is_single_word = bits <= BN_BITS2;
    gain;
    if (!BN_rand(rnd, bits, BN_RAND_TOP_TMO, BN_RAND_BOTTOM_ODD)) +
    return 0;
    /* we now have a random number 'rnd' to test. */
    for (i = 1; 1 < NUMPRINES; i++) {
        if (mod = SCR_ULONG) 1);
        if (mod = SCR_ULONG) 2)
        return 0;
    } mode[i] = (uintid.thmod;
    ** will be so small that it fits into a single word then we wightlifted by some to exceed that many bits. */
        if this is so small that it exceed that many bits. */
        if this is so small that it exceed that many bits. */
        if the single_word (if the small that it exceed that many bits. */
        if the single_word then we wish the same to exceed that many bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits. */
        if the single_word then we wish the same bits
```

- ► A very useful tool for RE made by shellphish and now maintained by SEFCOM at Arizona State University as well
- Originally made for DARPAs cyber grand challenge.
- A very strong framework for emulation allowing symbolic values.

/* In the case that the candidate prime is a single word then
** set of the set of

Angr - The basics

```
int is
intile mode[NUMPRIMES];
intile mode[NUMPRIMES];
intile mode[a = BM_HAGK2 - primes[NUMPRIMES - 1];
clust is_single_word = bits <= BM_BITS2;
clust is_single_word = bits <= BM_BITS2;
intile model is_word = bits = bM_BITS2;
intile model is_word = bM_BITS
```

▶ It is written almost entirely in python and as such currently only has python bindings.

- ▶ Installation
 - Has a couple dependencies that may conflict with other packages so it is recommended to use python virtual environments.
 - \$ pip install --user angr
 - Now you can just import angr in any python interpreter

Angr - Writing a script

- The basic block in angr is a project
- It is how you tell angr what file to load
- Can be used to obtain a bunch of metadata as well about the file
 - project = angr.Project("./binary")

Angr - Writing a script

- ► Accessing everything else will be done through factory life in so small that it file into a single word then
 - project.factory
- Next you will have to tell angr where it should start
- Most of the time you want this to be entry_state()
 - This will setup everything as it would be at the entry point if you ran the binary
 - ► This is where you would tell it if you wanted a special stdin or args
 - You can use claripy to make it symbolic

```
ecial stuling or args wordered:

/* In the case that the condidate prime is a single word them

** In the case that the condidate prime is a single word them

** In the case that the condidate prime is a single word them

** In the case that the condidate prime is a single word them

** In the case that the case is a single word them

** In the case that the case is a single word them

** In the case that the case is a single word them

** In the case that the case is a single word them

** In the case that the case is a single word that the case is a
```

Angr - Simulation manager

```
initials a modefNUMPRIMES]:
BN_LUNN( delta:
BN_LULONO maddeta = BM_MIGK2 - primes[NUMPRIMES - 1];
claw is_single_word = bits <= BN_BITS2;
again:
If [BN_rand(rnd, bits, BN_RAND_TOP_TMO, BN_RAND_BOTTOM_ODD)) +
    return 0;
}

Z* we now have a random_number 'rynd' to test, */
```

- ▶ This will hold all your various states as they are created and abandoned
 - sm = project.factory.simulation_manager(state)
- Executing
 - ► To execute you will normally want to use .run() or .explore()
 - Explore
 - Allows you to guide execution better and limit computation necessary
 - You can specify find and avoid conditions via addresses or lambda functions
 - Run
 - Will just run every state until exhaustion
 - Theoretically tells you all possible outcomes of the program
 - Prone to path explosion
 - Many more that are more application specific

```
ne program.

" That it's not a multiple of a known prime, be check that rnd-1 is also caprime to all the known prime became there aren't many small primes for (i = 1; i < NUMPRIMES & primes[i] < md_word; i + if (model) + delta) / primes[i] == 0) (word; i + if (model) + delta) / primes[i] == 0) (word; i + if (model) / model) / primes[i] == 0) (word; i + if (model) / model) / primes[i] == 0) (word; i + if (model) / model) / primes[i] == 0) (word; i + if (model) / model) / word / model) / model / model) / model / model) / model / model / model) / model /
```

Angr - States

- States will allow you to introspect, touch or check memory values
- ► For now we will only be touching state.posix
 - This allows access to the various posix style file descriptors on supported systems
 - Alot you can do here by using hooks and moving states between stashes
- Stashes
 - Active
 - Deadended
 - Pruned
 - Unconstrained
 - Unsat

Angr - Symbolic Values

```
idition in the second in
```

- ► Angr uses a library called claripy which is their wrapper for constraint solvers
- Used in place of direct Z3 calls so it is portable between different backends
- Shouldn't need to touch anything besides
 - BV bitvector
 - FP floating point
 - Bool boolean
- ► Then just like in Z3 you can perform math with these, however to combine them one after another you need to use Concat()
- ▶ NOTE: The default python >, <, >=, and <= are unsigned in Claripy. This is different than their behavior in Z3, because it seems more natural in binary analysis.

```
Example: Fairgame RE400 with Angress (Intelligence of the Lord (Intell
```

Fairgame RE400 script and output

Script:

```
import angr
import claripy
p = angr.Project("./re400")
for i in range(0xFF):
    # make a symbolic list of i bytes
    flag_chars = [claripy.BVS("serial_%d" % i, 8) for i in range(j)]
    # combine them all
    flag = claripy.Concat(*flag_chars)
    # tell angr to start at entry
    init = p.factory.entry state(stdin=flag)
    sm = p.factory.simulation_manager(init)
    # find a state with "unlocked" in stdout
    sm.explore(find=lambda s: b"unlocked" in s.posix.dumps(1))
    # try next length if not found
    if len(sm.found) == 0:
        continue
    print('Good length: %d' % (j,))
    print(sm)
    for i in sm.found:
        print(i.posix.dumps(1))
       print(i.posix.dumps(0))
    hreak
```

```
Redacted output:
Good length: 25
 <SimulationManager with 1 active, 1 found>
 b'[+] Welcome to the REvision 400 lock firmware.\n
 [!] Please enter the serial:\n
 [+] REvision 400 lock firmware unlocked.\n'
 b'flag{ if (is_single_word)}\xff'
 python script.py 154.01s user 2.08s system is a single word then
 102% cpu 2:31.85 total
```

Any questions overall?

Resources

- https://github.com/angr/
- https://github.com/Z3Prover/z3/
- ▶ https://github.com/RPISEC/MBE