Combining infinite sets of experts

Yoav Freund

January 16, 2020

Dates

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Review

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The Universal prediction machine

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The biased coins set of experts

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Bayes using Jeffrey's prior

Laplace Approximation
Choosing the optimal prior

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Generalization to larger sets of distributions

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- P(M_i) probability of message i
- Arithmetic coding defines a code of length $\lceil -\log_2 P(M_i) \rceil$ for message i
- Conversely: a codebook defines a distribution.

The online Bayes Algorithm

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Prediction of algorithm A

$$\mathbf{p}_{A}^{t} = \frac{\sum_{i=1}^{N} w_{i}^{t} \mathbf{p}_{i}^{t}}{\sum_{i=1}^{N} w_{i}^{t}}$$



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EQUALITY not bound!

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Standardizing online prediction algorithms

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- ▶ $V(\vec{b}, \vec{X}, t)$ is computable (recursively enumerable).

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- ▶ technical details: On iteration t, $|\vec{X}| = t$. Use the predictions of programs \vec{b} such that $|\vec{b}| \le t$ and for which $V(\vec{b}, \vec{X}, 2^t) = 1$.

the unused algorithms predict 1/2 (insuring a loss of 1)

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- Ridiculously bad running time.

Bayes coding is better than two part codes

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- More generally, the regret is smaller if many of the experts perform well.

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- Can we still get a meaningful bound?

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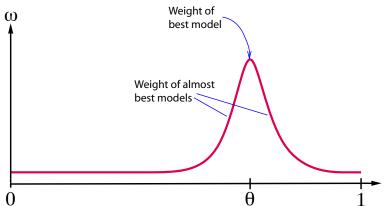
We need a new lower bound on the final total weight

Main Idea

If $\mathbf{w}^t(\theta)$ is large then $\mathbf{w}^t(\theta + \epsilon)$ is also large.

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► Taylor expansion of $g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta})$ around $\theta = \hat{\theta}$.

Laplace Approximation

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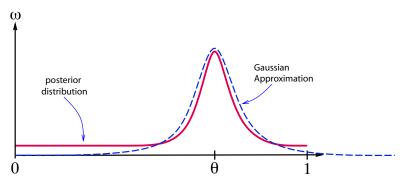
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- ► Third term gives a quadratic expression in the exponent
- → a gaussian approximation of the posterior.



Laplace Approximation, Watson's lemma

$$\int_0^1 w(\theta) e^{T(g(\hat{\theta},\theta)-g(\hat{\theta},\hat{\theta}))} d\theta$$

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$$\int_{0}^{1} w(\theta) e^{T(g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))} d\theta$$

$$= w(\hat{\theta}) \sqrt{\frac{-2\pi}{T \frac{d^{2}}{d\theta^{2}} \Big|_{\theta = \hat{\theta}} (g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))}} + O(T^{-3/2})$$

Choosing the optimal prior

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▶ Choose $w(\theta)$ to maximize the worst-case final total weight

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▶ Make bound equal for all $\hat{\theta} \in [0, 1]$ by choosing

$$w^*(\hat{\theta}) = \frac{1}{Z} \sqrt{\frac{\frac{d^2}{d\theta^2}\Big|_{\theta=\hat{\theta}} (g(\hat{\theta},\theta) - g(\hat{\theta},\hat{\theta}))}{-2\pi}},$$

where Z is the normalization factor:

$$Z = \sqrt{rac{1}{2\pi}} \int_0^1 \left. \sqrt{rac{d^2}{d heta^2}}
ight|_{ heta = \hat{ heta}} (g(\hat{ heta},\hat{ heta}) - g(\hat{ heta}, heta))
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The bound for the optimal prior

Plugging in we get

$$L_{A} - L_{\min} \leq \ln \int_{0}^{1} w^{*}(\theta) e^{T(g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}))} d\theta$$

$$= \ln \left(\sqrt{\frac{2\pi Z}{T}} + O(T^{-3/2}) \right)$$

$$= \frac{1}{2} \ln \frac{T}{2\pi} - \frac{1}{2} \ln Z + O(1/T) .$$

Choosing the optimal prior

Solving for log-loss

The exponent in the integral is

$$g(\hat{\theta}, \theta) - g(\hat{\theta}, \hat{\theta}) = \hat{\theta} \ln \frac{\hat{\theta}}{\theta} + (1 - \hat{\theta}) \ln \frac{1 - \hat{\theta}}{1 - \theta} = D_{KL}(\hat{\theta}||\theta)$$

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$$w^*(\hat{\theta}) = \frac{1}{\pi \sqrt{\hat{\theta}(1-\hat{\theta})}}$$

Known in general as Jeffrey's prior. And, in this case, the Dirichlet-(1/2, 1/2) prior.

Choosing the optimal prior

The cumulative log loss of Bayes using Jeffrey's prior

$$L_A - L_{\min} \le \frac{1}{2} \ln(T+1) + \frac{1}{2} \ln \frac{\pi}{2} + O(1/T)$$

But what is the prediction rule?

As luck would have it the Dirichlet prior is the conjugate prior for the Binomial distribution.

Bayes using Jeffrey's prior

Kritchevski Trofimov Prediction Bule

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This is called the Trichevsky Trofimov prediction rule.

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Suffers larger regret when $\hat{\theta}$ is far from 1/2

Shtarkov Lower bound

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Shtarkov Lower bound

What is the optimal prediction when T is know in advance?

$$L_*^T - \min_{\theta} L_{\theta}^T \geq \frac{1}{2} \ln(T+1) + \frac{1}{2} \ln \frac{\pi}{2} - O(\frac{1}{\sqrt{T}})$$

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► The constant C is optimal.

Exponential Distributions

For any set of distributions from the exponential family defined by k parameters (Some technical conditions on closure of set??) [Rissanen??]

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General Distributions

Characterize distribution family by metric entropy.

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[Haussler and Opper] show that the coefficient in front of In T is optimal for distribution families where the metric entropy is up to

$$N(1/\epsilon) = O(e^{\epsilon^{-\alpha}})$$

For all $\alpha \leq 5/2$.

next Class

Variable-length markov models - a set of distributions with increasing number of parameters.

next Class

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- THe context algorithm: An efficient implementation of the Bayes algorithm which achieves close-to-optimal worst case bounds.