

Introduction to Categorical Logic

[DRAFT: FEBRUARY 2, 2022]

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Chapter 1

Algebraic Theories

Algebraic theories are descriptions of structures that are given entirely in terms of operations and equations. All such algebraic notions have in common some quite deep and general properties, from the existence of free algebras to Lawvere’s duality theory. The most basic of these are presented in this chapter. The development also serves as a first example and template for the general scheme of *functorial semantics*, to be applied to other logical notions in later chapters.

1.1 Syntax and semantics

We begin with a general approach to algebraic structures such as groups, rings, and lattices. These are characterized by axiomatizations which involve only a single sort of variables and constants, operations, and equations. It is important that the operations are defined everywhere, which excludes two important examples: fields, because the inverse of 0 is undefined, and categories because composition is defined only for certain pairs of morphisms.

Let us start with the quintessential algebraic theory: the theory of groups. In first-order logic, a group can be described as a set G with a binary operation $\cdot : G \times G \rightarrow G$, satisfying the two first-order axioms:

$$\begin{aligned} &\forall x, y, z \in G. (x \cdot y) \cdot z = x \cdot (y \cdot z) \\ &\exists e \in G. \forall x \in G. \exists y \in G. (e \cdot x = x \cdot e = x \wedge x \cdot y = y \cdot x = e) \end{aligned}$$

Taking a closer look at the logical form of these axioms, we see that the second one, which expresses the existence of a unit and inverse elements, is somewhat unsatisfactory because it involves nested quantifiers. Not only does this complicate the interpretation, but it is not really necessary, since the unit element and inverse operation in a group are uniquely determined. Thus we can add them to the structure and reformulate as follows. The unit is to be represented by a distinguished *constant* $e \in G$, and the inverse is to be a unary *operation* $^{-1} : G \rightarrow G$. We then obtain an equivalent formulation in which all axioms can

be expressed as (universally quantified) *equations*:

$$\begin{aligned} x \cdot (y \cdot z) &= (x \cdot y) \cdot z \\ x \cdot e &= x & e \cdot x &= x \\ x \cdot x^{-1} &= e & x^{-1} \cdot x &= e \end{aligned}$$

The universal quantifiers $\forall x \in G, \forall y \in G$, etc. are no longer needed in stating the axioms, since we can interpret all variables as ranging over all elements of G (because of our restriction to totally defined operations). Nor do we really need to explicitly mention the particular set G in the specification. Finally, since the constant e can be regarded as a nullary operation, i.e., a function $e : 1 \rightarrow G$, the specification of the group concept consists solely of operations and equations. This leads to the following general definition of an algebraic theory.

Definition 1.1.1. A *signature* Σ for an algebraic theory consists of a family of sets $\{\Sigma_k\}_{k \in \mathbb{N}}$. The elements of Σ_k are called the *k-ary operations*. In particular, the elements of Σ_0 are the *nullary operations* or *constants*.

The *terms* of a signature Σ are the expressions constructed inductively by the following rules:

1. variables x, y, z, \dots , are terms,
2. if t_1, \dots, t_k are terms and $f \in \Sigma_k$ is a *k-ary operation* then $f(t_1, \dots, t_k)$ is a term.

Definition 1.1.2 (cf. Definition 1.2.10). An *algebraic theory* $\mathbb{T} = (\Sigma_{\mathbb{T}}, A_{\mathbb{T}})$ is given by a signature $\Sigma_{\mathbb{T}}$ and a set $A_{\mathbb{T}}$ of *axioms*, which are equations between terms (formally, pairs of terms).

Algebraic theories are also called *equational theories*. We do not assume that the sets Σ_k or $A_{\mathbb{T}}$ are finite, but the individual terms and equations always involve only finitely many variables.

Example 1.1.3. The theory of a commutative ring with unit is an algebraic theory. There are two nullary operations (constants) 0 and 1, a unary operation $-$, and two binary operations $+$ and \cdot . The equations are:

$$\begin{aligned} (x + y) + z &= x + (y + z) & (x \cdot y) \cdot z &= x \cdot (y \cdot z) \\ x + 0 &= x & x \cdot 1 &= x \\ 0 + x &= x & 1 \cdot x &= x \\ x + (-x) &= 0 & (x + y) \cdot z &= x \cdot z + y \cdot z \\ (-x) + x &= 0 & z \cdot (x + y) &= z \cdot x + z \cdot y \\ x + y &= y + x & x \cdot y &= y \cdot x \end{aligned}$$

Example 1.1.4. The “empty” theory with no operations and no equations is the theory of a set.

Example 1.1.5. The theory with one constant and no equations is the theory of a *pointed set*, cf. Example ??.

Example 1.1.6. Let R be a ring. There is an algebraic theory of left R -modules. It has one constant 0 , a unary operation $-$, a binary operation $+$, and for each $a \in R$ a unary operation \bar{a} , called *scalar multiplication by a* . The following equations hold:

$$\begin{aligned} (x + y) + z &= x + (y + z) , & x + y &= y + x , \\ x + 0 &= x , & 0 + x &= x , \\ x + (-x) &= 0 , & (-x) + x &= 0 . \end{aligned}$$

For every $a, b \in R$ we also have the equations

$$\bar{a}(x + y) = \bar{a}x + \bar{a}y , \quad \bar{a}(\bar{b}x) = \overline{(ab)}x , \quad \overline{(a + b)}x = \bar{a}x + \bar{b}x .$$

Scalar multiplication by a is usually written as $a \cdot x$ instead of $\bar{a}x$. If we replace the ring R by a field \mathbb{F} we obtain the algebraic theory of a vector space over \mathbb{F} (even though the theory of fields is not algebraic!).

Example 1.1.7. In computer science, inductive datatypes are examples of algebraic theories. For example, the datatype of binary trees with leaves labeled by integers might be defined as follows in a programming language:

```
type tree = Leaf of int | Node of tree * tree
```

This corresponds to the algebraic theory with a constant `Leaf n` for each integer n and a binary operation `Node`. There are no equations. Actually, when computer scientists define a datatype like this, they have in mind a particular model of the theory, namely the *free* one.

Example 1.1.8. An obvious non-example is the theory of posets, formulated with a binary relation symbol $x \leq y$ and the usual axioms of reflexivity, transitivity and anti-symmetry, namely:

$$\begin{aligned} x &\leq x \\ x \leq y , y \leq z &\Rightarrow x \leq z \\ x \leq y , y \leq x &\Rightarrow x = y \end{aligned}$$

On the other hand, using an operation of greatest lower bound or “meet” $x \wedge y$, one can make the equational theory of “ \wedge -semilattices”:

$$\begin{aligned} x \wedge x &= x \\ x \wedge y &= y \wedge x \\ x \wedge (y \wedge z) &= (x \wedge y) \wedge z \end{aligned}$$

Then, defining a partial ordering by $x \leq y \iff x \wedge y = x$ we arrive at the notion of a “poset with meets”, which *is* equational (of course, the same can be done with joins $x \vee y$ as well). We will show later (in section ??) that there is no reformulation of the general theory of posets into an equivalent equational one.

Exercise 1.1.9. Let G be a group. Formulate the notion of a (left) G -set (i.e. a functor $G \rightarrow \mathbf{Set}$) as an algebraic theory.

1.1.1 Models of algebraic theories

Let us now consider *models* of an algebraic theory, i.e. *algebras*. Classically, a group can be given by a set G , an element $e \in G$, a function $m : G \times G \rightarrow G$ and a function $i : G \rightarrow G$, satisfying the group axioms:

$$\begin{aligned} m(x, m(y, z)) &= m(m(x, y), z) \\ m(x, i x) &= m(i x, x) = e \\ m(x, e) &= m(e, x) = x \end{aligned}$$

for any $x, y, z \in G$. Observe, however, that this notion can easily be generalized so that we can speak of models of group theory in categories other than \mathbf{Set} . This is accomplished simply by translating the equations between arbitrary elements into equations between the operations themselves: thus a group is given, first, by an object $G \in \mathbf{Set}$ and three morphisms

$$e : 1 \rightarrow G, \quad m : G \times G \rightarrow G, \quad i : G \rightarrow G.$$

The associativity axiom is then expressed by the commutativity of the following diagram:

$$\begin{array}{ccc} G \times G \times G & \xrightarrow{m \times \pi_2} & G \times G \\ \pi_0 \times m \downarrow & & \downarrow m \\ G \times G & \xrightarrow{m} & G \end{array} \quad (1.1)$$

Note that we have omitted the canonical associativity function $G \times (G \times G) \cong (G \times G) \times G$, which should be inserted into the top left corner of the diagram. The equations for the unit and the inverse are similarly expressed by commutativity of the following diagrams:

$$\begin{array}{ccc} G \times 1 & \xrightarrow{1_G \times e} & G \times G \xleftarrow{e \times 1_G} 1 \times G \\ \pi_0 \searrow & & \downarrow m \swarrow \pi_1 \\ & G & \end{array} \quad \begin{array}{ccccc} G & \xrightarrow{\langle 1_G, i \rangle} & G \times G & \xleftarrow{\langle i, 1_G \rangle} & G \\ !_G \downarrow & & \downarrow m & & \downarrow !_G \\ 1 & \xrightarrow{e} & G & \xleftarrow{e} & 1 \end{array} \quad (1.2)$$

This formulation makes sense in any category \mathcal{C} with finite products.

Definition 1.1.10. Let \mathcal{C} be a category with finite products. A *group in \mathcal{C}* consists of an object G equipped with arrows:

$$\begin{array}{ccccc} G \times G & \xrightarrow{m} & G & \xleftarrow{i} & G \\ & & \uparrow e & & \\ & & 1 & & \end{array}$$

such that the above diagrams (1.1) and (1.2) expressing the group equations commute.

There is also an obvious corresponding generalization of a group homomorphism in **Set** to homomorphisms of groups in \mathcal{C} . Namely, an arrow in \mathcal{C} between (the underlying objects of) groups, say $h : M \rightarrow N$, is a homomorphism if it commutes with the interpretations of the basic operations m , i , and e ,

$$h \circ m^M = m^N \circ h^2 \quad h \circ i^M = i^N \circ h \quad h \circ e^M = e^N$$

as indicated in:

$$\begin{array}{ccc} \begin{array}{ccc} M^2 & \xrightarrow{h^2} & N^2 \\ m^M \downarrow & & \downarrow m^N \\ M & \xrightarrow{h} & N \end{array} & \begin{array}{ccc} M & \xrightarrow{h} & N \\ i^M \downarrow & & \downarrow i^N \\ M & \xrightarrow{h} & N \end{array} & \begin{array}{ccc} 1 & \xrightarrow{=} & 1 \\ e^M \downarrow & & \downarrow e^N \\ M & \xrightarrow{h} & N \end{array} \end{array}$$

Together with the evident composition and identity arrows inherited from \mathcal{C} , this gives a category of groups in \mathcal{C} , which we denote:

$$\mathbf{Group}(\mathcal{C})$$

In general, we define an *interpretation* I of a theory \mathbb{T} in a category \mathcal{C} with finite products to consist of an object $I \in \mathcal{C}$ and, for each basic operation f of arity k , a morphism $f^I : I^k \rightarrow I$. (More formally, I is the tuple consisting of an underlying object $|I|$ and the interpretations f^I , but we shall write simply I for $|I|$.) In particular, basic constants are interpreted as morphisms $1 \rightarrow I$. The interpretation is then extended to all terms as follows: a general term t will be interpreted together with a *context of variables* x_1, \dots, x_n (a list without repetitions), where the variables appearing in t are among those appearing in the context. We write

$$x_1, \dots, x_n \mid t \tag{1.3}$$

for a term t in context x_1, \dots, x_n . The interpretation of such a term in context (1.3) is a morphism $t^I : I^n \rightarrow I$, determined by the following specification:

1. The interpretation of a variable x_i among the x_1, \dots, x_n is the i -th projection $\pi_i : I^n \rightarrow I$.

2. A term of the form $f(t_1, \dots, t_k)$ is interpreted as the composite:

$$I^n \xrightarrow{(t_1^I, \dots, t_k^I)} I^k \xrightarrow{f^I} I$$

where $t_i^I : I^n \rightarrow I$ is the interpretation of the subterm t_i , for $i = 1, \dots, k$, and f^I is the interpretation of the basic operation f .

It is clear that the interpretation of a term really depends on the context, and when necessary we shall write $t^I = [x_1, \dots, x_n \mid t]^I$. For example, the term $f x_1$ is interpreted as a morphism $f^I : I \rightarrow I$ in context x_1 , and as the morphism $f^I \circ \pi_1 : I^2 \rightarrow I$ in the context x_1, x_2 .

Suppose u and v are terms in context x_1, \dots, x_n . Then we say that the equation in context $x_1, \dots, x_n \mid u = v$ is *satisfied* by the interpretation I if u^I and v^I are the same morphism in \mathcal{C} . In particular, if $u = v$ is an axiom of the theory, and x_1, \dots, x_n are all the variables appearing in either u or v , we say that I *satisfies the axiom* $u = v$, written

$$I \models u = v,$$

if $[x_1, \dots, x_n \mid u]^I$ and $[x_1, \dots, x_n \mid v]^I$ are the same morphism,

$$I^n \xrightarrow{\begin{array}{c} [x_1, \dots, x_n \mid u]^I \\ [x_1, \dots, x_n \mid v]^I \end{array}} I \quad (1.4)$$

We can then define, as expected:

Definition 1.1.11 (cf. Definition 1.2.10). A *model* M of an algebraic theory \mathbb{T} in a category \mathcal{C} with finite products (also called a \mathbb{T} -*algebra*) is an interpretation of the signature $\Sigma_{\mathbb{T}}$,

$$f^I : I^k \longrightarrow I$$

in \mathcal{C} , for all $f \in \Sigma_{\mathbb{T}}$, that satisfies the axioms $A_{\mathbb{T}}$,

$$I \models u = v,$$

for all $(u = v) \in A_{\mathbb{T}}$.

A *homomorphism* of models $h : M \rightarrow N$ is an arrow in \mathcal{C} that commutes with the interpretations of the basic operations,

$$h \circ f^M = f^N \circ h^k$$

for all $f \in \Sigma_{\mathbb{T}}$, as indicated in:

$$\begin{array}{ccc} M^k & \xrightarrow{h^k} & N^k \\ f^M \downarrow & & \downarrow f^N \\ M & \xrightarrow{h} & N \end{array}$$

The category of \mathbb{T} -models in \mathcal{C} is written,

$$\mathbf{Mod}(\mathbb{T}, \mathcal{C}).$$

A model of the empty theory \mathbb{T}_0 in \mathcal{C} is therefore just an object A , and a homomorphism is just a map, so

$$\text{Mod}(\mathbb{T}_0, \mathcal{C}) = \mathcal{C}.$$

A model of the theory $\mathbb{T}_{\text{Group}}$ of groups in \mathcal{C} is a group in \mathcal{C} , in the above sense, and similarly for homomorphisms, so:

$$\text{Mod}(\mathbb{T}_{\text{Group}}, \mathcal{C}) = \text{Group}(\mathcal{C})$$

as defined above. In particular, a model in **Set** is just a group in the usual sense, so we have:

$$\text{Mod}(\mathbb{T}_{\text{Group}}, \text{Set}) = \text{Group}(\text{Set}) = \text{Group}.$$

An example of a new kind is provided the following.

Example 1.1.12. A model of the theory of groups in a functor category $\text{Set}^{\mathbb{C}}$ is the same thing as a functor from \mathbb{C} into the category groups,

$$\text{Group}(\text{Set}^{\mathbb{C}}) \cong \text{Group}^{\mathbb{C}}.$$

Indeed, for each object $C \in \mathbb{C}$ there is an evaluation functor,

$$\text{eval}_C : \text{Set}^{\mathbb{C}} \rightarrow \text{Set}$$

with $\text{eval}_C(F) = F(C)$, and evaluation preserves products since these are computed point-wise in the functor category. Moreover, every arrow $h : C \rightarrow D$ in \mathbb{C} gives rise to an obvious natural transformation $h : \text{eval}_C \rightarrow \text{eval}_D$. Thus for any group G in $\text{Set}^{\mathbb{C}}$, we have groups $\text{eval}_C(G)$ for each $C \in \mathbb{C}$ and group homomorphisms $h_G : C(G) \rightarrow D(G)$, comprising a functor $G : \mathbb{C} \rightarrow \text{Group}$. Conversely, it is clear that a functor $H : \mathbb{C} \rightarrow \text{Group}$ determines a group H in $\text{Set}^{\mathbb{C}}$ with underlying object $|HC|$, where $|-| : \text{Group} \rightarrow \text{Set}$ is the forgetful functor. These constructions are clearly mutually inverse (up to canonical isomorphisms determined by the choice of products). In this way, *a group in the category of variable sets may be regarded as a variable group*.

Exercise 1.1.13. Verify the details of the isomorphism of categories

$$\text{Mod}(\mathbb{T}, \text{Set}^{\mathbb{C}}) \cong \text{Mod}(\mathbb{T}, \text{Set})^{\mathbb{C}}$$

discussed in example 1.1.12 for an arbitrary algebraic theory \mathbb{T} .

Exercise 1.1.14. Determine what a group is in the following categories: the category of graphs **Graph**, the category of topological spaces **Top**, and the category of groups **Group**. (Hint: Only the last case is tricky. Before thinking too hard about it, prove the following lemma [Bor94, Lemma 3.11.6], known as the Eckmann-Hilton argument. Let G be a set provided with two binary operations \cdot and \star and a common unit e , so that $x \cdot e = e \cdot x = x \star e = e \star x = x$. Suppose the two operations commute, i.e., $(x \star y) \cdot (z \star w) = (x \cdot z) \star (y \cdot w)$. Then they coincide, are *commutative* and associative.)

1.1.2 Theories as categories

The syntactically presented notion of an algebraic theory is a practical convenience, but as a specification of a mathematical concept, say that of a group, it has some defects. We would prefer a *presentation-free* notion that captures the group concept without tying it to a specific syntactic presentation (the example below indicates why). One such notion can be given by a category with a certain universal property, which determines it uniquely, up to equivalence of categories.

Let us consider group theory again. The algebraic axiomatization in terms of unit, multiplication and inverse is not the only possible one. For example, an alternative formulation uses the unit e and a binary operation \odot , called *double division*, along with a single axiom [McC93]:

$$(x \odot (((x \odot y) \odot z) \odot (y \odot e))) \odot (e \odot e) = z .$$

The usual group operations are related to double division as follows:

$$x \odot y = x^{-1} \cdot y^{-1} , \quad x^{-1} = x \odot e , \quad x \cdot y = (x \odot e) \odot (y \odot e) .$$

There may be practical reasons for preferring one formulation of group theory over another, but this should not determine what the general concept of a group is. For example, we would like to avoid particular choices of basic constants, operations, and axioms. This is akin to the situation where an algebra is presented by generators and relations: the algebra itself is regarded as independent of any particular such presentation. Similarly, one usually prefers a basis-free theory of vector spaces: it is better to formulate the general idea of a vector space without referring explicitly to a basis, even though every vector space has one.

As a first step, one could simply take *all* operations built from unit, multiplication, and inverse as basic, and *all* valid equations of group theory as axioms. But we can go a step further and collect all the operations into a category, thus forgetting about which ones were “basic” and which ones “derived”, and which equalities were “axioms”. We first describe this construction of a category $\mathcal{C}_{\mathbb{T}}$ for a general algebraic theory \mathbb{T} , and then determine another characterization of it.

As objects of $\mathcal{C}_{\mathbb{T}}$ we take *contexts*, i.e. sequences of distinct variables,

$$[x_1, \dots, x_n] . \quad (n \geq 0)$$

Actually, it will be convenient to take equivalence classes under renaming of variables, so that $[x_1, x_3] = [x_2, x_1]$. That is to say, the objects are just natural numbers.

A morphism from $[x_1, \dots, x_m]$ to $[x_1, \dots, x_n]$ is then an n -tuple (t_1, \dots, t_n) , where each t_k is a term in the context x_1, \dots, x_m , possibly after renaming the variables. Two such morphisms (t_1, \dots, t_n) and (s_1, \dots, s_n) are equal if, and only if, the axioms of the theory imply that $t_k = s_k$ for every $k = 1, \dots, n$,

$$\mathbb{T} \vdash t_k = s_k$$

Strictly speaking, morphisms are thus (tuples of) *equivalence classes* of terms in context

$$[x_1, \dots, x_m \mid t_1, \dots, t_n] : [x_1, \dots, x_m] \longrightarrow [x_1, \dots, x_n],$$

where two terms are equivalent when the theory proves them to be equal (after renaming the variables). Since it is rather cumbersome to work with such equivalence classes, we shall work with the terms directly, but keeping in mind that equality between them is this equivalence. Note also that the context of the morphism agrees with its domain, so we can omit it from the notation when that domain is clear. The composition of morphisms

$$\begin{aligned} (t_1, \dots, t_m) &: [x_1, \dots, x_k] \rightarrow [x_1, \dots, x_m] \\ (s_1, \dots, s_n) &: [x_1, \dots, x_m] \rightarrow [x_1, \dots, x_n] \end{aligned}$$

is the morphism (r_1, \dots, r_n) whose i -th component is obtained by simultaneously substituting in s_i the terms t_1, \dots, t_m for the variables x_1, \dots, x_m :

$$r_i = s_i[t_1, \dots, t_m/x_1, \dots, x_m] \quad (1 \leq i \leq n)$$

The identity morphism on $[x_1, \dots, x_n]$ is (x_1, \dots, x_n) . Using the usual rules of deduction for equational logic (see Section A.5), it is easy to verify that these specifications are well-defined on equivalence classes, and therefore make $\mathcal{C}_{\mathbb{T}}$ a category.

Definition 1.1.15. The category $\mathcal{C}_{\mathbb{T}}$ just defined is called the *syntactic category* of the theory \mathbb{T} .

The syntactic category $\mathcal{C}_{\mathbb{T}}$ (which may be thought of as the “Lindenbaum-Tarski category” of \mathbb{T} , see ??) contains the same “algebraic” information as the theory \mathbb{T} from which it was built, but in a syntax-invariant way. Two different syntactic presentations of \mathbb{T} — like the ones for groups mentioned above — will give rise to essentially the same category $\mathcal{C}_{\mathbb{T}}$ (i.e. up to isomorphism). In this sense, the category $\mathcal{C}_{\mathbb{T}}$ is the abstract, algebraic object presented by the “generators and relations” (the operations and equations) of \mathbb{T} . But there is another, still more important, sense in which $\mathcal{C}_{\mathbb{T}}$ represents \mathbb{T} , as we next show.

Exercise 1.1.16. Show that the syntactic category $\mathcal{C}_{\mathbb{T}}$ has all finite products.

1.1.3 Models as functors

Having represented an algebraic theory \mathbb{T} by the syntactic category $\mathcal{C}_{\mathbb{T}}$ constructed from it, we next show that $\mathcal{C}_{\mathbb{T}}$ has the universal property that models of \mathbb{T} correspond uniquely to certain functors from $\mathcal{C}_{\mathbb{T}}$. More precisely, given any category with finite products \mathcal{C} (which we shall call an *FP-category*), there is a natural (in \mathcal{C}) equivalence,

$$\frac{\mathcal{M} \in \text{Mod}(\mathbb{T}, \mathcal{C})}{M : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}} \quad (1.5)$$

between models \mathcal{M} of \mathbb{T} in \mathcal{C} and finite product preserving functors (“*FP-functors*”) $M : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$. The equivalence is mediated by a “universal model” \mathcal{U} in $\mathcal{C}_{\mathbb{T}}$, corresponding to the identity functor $1_{\mathcal{C}_{\mathbb{T}}} : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}_{\mathbb{T}}$ in the above displayed correspondence, so that every model \mathcal{M} arises as the functorial image $M(\mathcal{U}) \cong \mathcal{M}$ of \mathcal{U} under an essentially unique FP-functor

$M : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$. The possibility of such universal models is an advantage of the generalized notion of a model in a category other than **Set**.

To give the details of the correspondence (1.5), let \mathbb{T} be an arbitrary algebraic theory and $\mathcal{C}_{\mathbb{T}}$ the syntactic category constructed from \mathbb{T} as in Definition 1.1.15. It is easy to show that the product in $\mathcal{C}_{\mathbb{T}}$ of two objects $[x_1, \dots, x_n]$ and $[x_1, \dots, x_m]$ is the object $[x_1, \dots, x_{n+m}]$, and that $\mathcal{C}_{\mathbb{T}}$ has all finite products (including $1 = [-]$, the empty context). Moreover, there is a distinguished \mathbb{T} -model \mathcal{U} in $\mathcal{C}_{\mathbb{T}}$ consisting of “the language itself”, which we call the *syntactic model*: the underlying object is the context $U = [x_1]$ of length one, and each operation symbol f of, say, arity k is interpreted as itself,

$$f^{\mathcal{U}} = [x_1, \dots, x_k \mid f(x_1, \dots, x_k)] : U^k = [x_1, \dots, x_k] \longrightarrow [x_1] = U.$$

The axioms are all satisfied, because the equivalence relation on terms is determined by \mathbb{T} -provability. Explicitly, for terms s, t , we have:

$$\mathcal{U} \models s = t \iff s^{\mathcal{U}} = t^{\mathcal{U}} \iff \mathbb{T} \vdash s = t. \quad (1.6)$$

We record this fact as the following.

Proposition 1.1.17. *The syntactic model \mathcal{U} in $\mathcal{C}_{\mathbb{T}}$ is “logically generic” in the sense that it satisfies all and only the \mathbb{T} -provable equations.*

Even more importantly, though, the syntactic model \mathcal{U} in $\mathcal{C}_{\mathbb{T}}$ has the following *universal property*:

Proposition 1.1.18. *Any model \mathcal{M} in any finite product category \mathcal{C} is the image of \mathcal{U} under an essentially unique, finite product preserving functor $\mathcal{M}^{\#} : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$.*

In this sense, the syntactic category $\mathcal{C}_{\mathbb{T}}$ is the “free finite product category with a model of \mathbb{T} ”, in the sense given by the following proof. First, observe that any FP-functor $F : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$ takes the syntactic model \mathcal{U} in $\mathcal{C}_{\mathbb{T}}$ (or indeed any model) to a model $F\mathcal{U}$ in \mathcal{C} , with interpretations

$$f^{F\mathcal{U}} = F f^{\mathcal{U}} : F U^k \rightarrow F U \quad \text{for each } f \in \Sigma_k.$$

Moreover, any natural transformation $\vartheta : F \rightarrow G$ between FP-functors determines a homomorphism of models $h = \vartheta_{\mathcal{U}} : F\mathcal{U} \rightarrow G\mathcal{U}$. In more detail, suppose $f : U \times U \rightarrow U$ is a basic operation, then there is a commutative diagram,

$$\begin{array}{ccc}
FU \times FU & \xrightarrow{h \times h} & GU \times GU \\
\downarrow \cong & & \downarrow \cong \\
F(U \times U) & \xrightarrow{\vartheta_{U \times U}} & G(U \times U) \\
\downarrow Ff & & \downarrow Gf \\
FU & \xrightarrow{h = \vartheta_U} & GU
\end{array}
\begin{array}{l}
\text{Left curved arrow: } f^{FU} \\
\text{Right curved arrow: } f^{GU}
\end{array}$$

where the upper square commutes by preservation of products, and the lower one by naturality. Thus the operation “evaluation at \mathcal{U} ” determines a functor,

$$\text{eval}_{\mathcal{U}} : \text{Hom}_{\text{FP}}(\mathcal{C}_{\mathbb{T}}, \mathcal{C}) \longrightarrow \text{Mod}(\mathbb{T}, \mathcal{C}) \quad (1.7)$$

from the category of finite product preserving functors $\mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$, with natural transformations as arrows, into the category of \mathbb{T} -models in \mathcal{C} .

Proposition 1.1.19. *The functor (1.7) is an equivalence of categories, natural in \mathcal{C} .*

Proof. Let \mathcal{M} be any model in an FP-category \mathcal{C} . Then the assignment $f \mapsto f^{\mathcal{M}}$ given by the interpretation part of \mathcal{M} determines a functor $\mathcal{M}^{\sharp} : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$, defined on objects by

$$\mathcal{M}^{\sharp}[x_1, \dots, x_k] = M^k$$

and on morphisms by

$$\mathcal{M}^{\sharp}(t_1, \dots, t_n) = (t_1^{\mathcal{M}}, \dots, t_n^{\mathcal{M}}).$$

In more detail, \mathcal{M}^{\sharp} is defined on morphisms

$$[x_1, \dots, x_k \mid t] : [x_1, \dots, x_k] \rightarrow [x_1, \dots, x_n]$$

in $\mathcal{C}_{\mathbb{T}}$ by the following rules:

1. The morphism

$$(x_i) : [x_1, \dots, x_k] \rightarrow [x_1]$$

is mapped to the i -th projection

$$\pi_i : M^k \rightarrow M.$$

2. The morphism

$$f(t_1, \dots, t_m) : [x_1, \dots, x_k] \rightarrow [x_1]$$

is mapped to the composite

$$M^k \xrightarrow{(\mathcal{M}^\sharp_{t_1}, \dots, \mathcal{M}^\sharp_{t_m})} M^m \xrightarrow{f^\mathcal{M}} M$$

where $\mathcal{M}^\sharp_{t_i} : M^k \rightarrow M$ is the value of \mathcal{M}^\sharp on the morphisms $(t_i) : [x_1, \dots, x_k] \rightarrow [x_1]$, for $i = 1, \dots, m$, and $f^\mathcal{M}$ is the interpretation of the basic operation f .

3. The morphism

$$(t_1, \dots, t_n) : [x_1, \dots, x_k] \rightarrow [x_1, \dots, x_n]$$

is mapped to the morphism $(\mathcal{M}^\sharp_{t_1}, \dots, \mathcal{M}^\sharp_{t_n})$ where $\mathcal{M}^\sharp_{t_i}$ is the value of \mathcal{M}^\sharp on the morphism $(t_i) : [x_1, \dots, x_k] \rightarrow [x_1]$, and

$$(\mathcal{M}^\sharp_{t_1}, \dots, \mathcal{M}^\sharp_{t_n}) : M^k \longrightarrow M^n$$

is the evident n -tuple in the FP-category \mathcal{C} .

That $\mathcal{M}^\sharp : \mathcal{C}_\mathbb{T} \rightarrow \mathcal{C}$ really is a functor now follows from the assumption that the interpretation M is a model, meaning that all the equations of the theory are satisfied by it, so that the above specification is well-defined on equivalence classes. Observe that the functor \mathcal{M}^\sharp is defined in such a way that it obviously preserves finite products, and that, moreover, there is an isomorphism of models,

$$\mathcal{M}^\sharp(\mathcal{U}) \cong \mathcal{M}.$$

Thus we have shown that the functor “evaluation at \mathcal{U} ”,

$$\text{eval}_\mathcal{U} : \text{Hom}_{\text{FP}}(\mathcal{C}_\mathbb{T}, \mathcal{C}) \longrightarrow \text{Mod}(\mathbb{T}, \mathcal{C}) \quad (1.8)$$

is essentially surjective on objects, since $\text{eval}_\mathcal{U}(\mathcal{M}^\sharp) = \mathcal{M}^\sharp(\mathcal{U}) \cong \mathcal{M}$.

We leave the verification that it is full and faithful as an easy exercise.

Exercise 1.1.20. Verify this.

Finally, naturality in \mathcal{C} means the following. Suppose \mathcal{M} is a model of \mathbb{T} in any FP-category \mathcal{C} . Any FP-functor $F : \mathcal{C} \rightarrow \mathcal{D}$ to another FP-category \mathcal{D} then takes \mathcal{M} to a model $F(\mathcal{M})$ in \mathcal{D} . Just as for the special case of \mathcal{U} , the interpretation is given by setting $f^{F(\mathcal{M})} = F(f^\mathcal{M})$ for the basic operations f (and composing with the canonical isos coming from preservation of products, $F(M) \times F(M) \cong F(M \times M)$, etc.). Since equations are described by commuting diagrams, F takes a model to a model, and the same is true for homomorphisms. Thus $F : \mathcal{C} \rightarrow \mathcal{D}$ induces a functor on \mathbb{T} -models,

$$\text{Mod}(\mathbb{T}, F) : \text{Mod}(\mathbb{T}, \mathcal{C}) \longrightarrow \text{Mod}(\mathbb{T}, \mathcal{D}).$$

By naturality of (1.7), we mean that the following square commutes up to natural isomorphism:

$$\begin{array}{ccc}
 \mathrm{Hom}_{\mathrm{FP}}(\mathcal{C}_{\mathbb{T}}, \mathcal{C}) & \xrightarrow{\mathrm{eval}_{\mathcal{U}}} & \mathrm{Mod}(\mathbb{T}, \mathcal{C}) \\
 \downarrow \mathrm{Hom}_{\mathrm{FP}}(\mathcal{C}_{\mathbb{T}}, F) & & \downarrow \mathrm{Mod}(\mathbb{T}, F) \\
 \mathrm{Hom}_{\mathrm{FP}}(\mathcal{C}_{\mathbb{T}}, \mathcal{D}) & \xrightarrow{\mathrm{eval}_{\mathcal{U}}} & \mathrm{Mod}(\mathbb{T}, \mathcal{D})
 \end{array} \tag{1.9}$$

But this is clear, since for any FP-functor $M : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$ we have:

$$\begin{aligned}
 \mathrm{eval}_{\mathcal{U}} \circ \mathrm{Hom}_{\mathrm{FP}}(\mathcal{C}_{\mathbb{T}}, F)(M) &= (\mathrm{Hom}_{\mathrm{FP}}(\mathcal{C}_{\mathbb{T}}, F)(M))(\mathcal{U}) \\
 &= (F \circ M)(\mathcal{U}) \\
 &= F(M(\mathcal{U})) \\
 &= F(\mathrm{eval}_{\mathcal{U}}(M)) \\
 &\cong \mathrm{Mod}(\mathbb{T}, F)(\mathrm{eval}_{\mathcal{U}}(M)) \\
 &= \mathrm{Mod}(\mathbb{T}, F) \circ \mathrm{eval}_{\mathcal{U}}(M).
 \end{aligned}$$

□

The equivalence of categories

$$\mathrm{Hom}_{\mathrm{FP}}(\mathcal{C}_{\mathbb{T}}, \mathcal{C}) \simeq \mathrm{Mod}(\mathbb{T}, \mathcal{C}) \tag{1.10}$$

actually determines $\mathcal{C}_{\mathbb{T}}$ and the universal model \mathcal{U} uniquely, up to equivalence of categories and isomorphism of models. Indeed, to recover \mathcal{U} , just put $\mathcal{C}_{\mathbb{T}}$ for \mathcal{C} and the identity functor $1_{\mathcal{C}_{\mathbb{T}}}$ on the left, to get \mathcal{U} in $\mathrm{Mod}(\mathbb{T}, \mathcal{C}_{\mathbb{T}})$ on the right! To see that $\mathcal{C}_{\mathbb{T}}$ itself is also determined, observe that (1.10) says that the functor $\mathrm{Mod}(\mathbb{T}, \mathcal{C})$ is representable, with representing object $\mathcal{C}_{\mathbb{T}}$, in an appropriate (i.e. bicategorical) sense. As usual, this fact can also be formulated in elementary terms as a universal mapping property of $\mathcal{C}_{\mathbb{T}}$, as follows:

Definition 1.1.21. The *classifying category* of an algebraic theory \mathbb{T} is an FP-category $\mathcal{C}_{\mathbb{T}}$ with a distinguished model \mathcal{U} , called the *universal model*, such that:

- (i) for any model \mathcal{M} in any FP-category \mathcal{C} , there is an FP-functor

$$\mathcal{M}^{\sharp} : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$$

and an isomorphism of models $\mathcal{M} \cong \mathcal{M}^{\sharp}(\mathcal{U})$.

- (ii) for any FP-functors $F, G : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$ and model homomorphism $h : F(\mathcal{U}) \rightarrow G(\mathcal{U})$, there is a unique natural transformation $\vartheta : F \rightarrow G$ with

$$\vartheta_{\mathcal{U}} = h.$$

Observe that (i) says that the evaluation functor (1.7) is essentially surjective, and (ii) that it is full and faithful. The category $\mathcal{C}_{\mathbb{T}}$ is then determined, up to equivalence, by this universal mapping property. Specifically, if $(\mathcal{C}, \mathcal{U})$ and $(\mathcal{D}, \mathcal{V})$ are both classifying categories for the same theory, then there are classifying functors,

$$\mathcal{C} \begin{array}{c} \xrightarrow{\mathcal{V}^\#} \\ \xleftarrow{\mathcal{U}^\#} \end{array} \mathcal{D}$$

the composites of which are necessarily isomorphic to the respective identity functors, since e.g. $\mathcal{U}^\#(\mathcal{V}^\#(\mathcal{U})) \cong \mathcal{U}^\#(\mathcal{V}) \cong \mathcal{U}$.

We have now shown not only that every algebraic theory has a classifying category, but also that the syntactic category is essentially determined by that distinguishing property. We record this as the following.

Theorem 1.1.22. *Every algebraic theory \mathbb{T} has a classifying category $\mathcal{C}_{\mathbb{T}}$, which can be constructed as the syntactic category of \mathbb{T} , in the sense of Definition 1.1.15.*

Example 1.1.23. Let us see explicitly what the foregoing definitions give us in the case of the theory of groups $\mathbb{T}_{\text{Group}}$. Let us write $\mathbb{G} = \mathcal{C}_{\mathbb{T}_{\text{Group}}}$ for the classifying category, which consists of contexts $[x_1, \dots, x_n]$ as objects, and terms built from variables and the group operations (modulo renaming of variables and the group laws) as arrows. A finite product preserving functor $M : \mathbb{G} \rightarrow \mathbf{Set}$ is then determined uniquely, up to natural isomorphism, by its action on the context $[x_1]$ and the terms representing the basic operations. If we set

$$\begin{aligned} G_M &= M[x_1] , & u_M &= M(\cdot \mid e) , \\ i_M &= M(x_1 \mid x_1^{-1}) , & m_M &= M(x_1, x_2 \mid x_1 \cdot x_2) , \end{aligned}$$

then $\mathcal{G}_M = (G_M, u_M, i_M, m_M)$ is just a group, with unit u_M , inverse i_M , and multiplication m_M . That \mathcal{G}_M satisfies the axioms for groups follows from the functoriality of M and preservation of finite products, which implies preservation of the corresponding commutative diagrams. Conversely, any group $\mathcal{G} = (G, u, i, m)$ determines a finite product preserving functor $\mathcal{G}^\# : \mathbb{G} \rightarrow \mathbf{Set}$, by setting $\mathcal{G}^\#[x_1] = G$, etc. Thus $\mathbf{Mod}_{\mathbf{Set}}(\mathbb{G})$ will indeed be equivalent to \mathbf{Group} once we show that both categories have the same notion of morphisms. This is shown just as in the general case above.

Example 1.1.24. Recall from 1.1.12 that a group G in the functor category $\mathbf{Set}^{\mathbb{C}}$ is essentially the same thing as a functor $G : \mathbb{C} \rightarrow \mathbf{Group}$. From the point of view of algebras as functors, this amounts to the observation that product-preserving functors $\mathbb{G} \rightarrow \mathbf{Hom}(\mathbb{C}, \mathbf{Set})$ correspond (by exponential transposition) to functors $\mathbb{C} \rightarrow \mathbf{Hom}_{\mathbf{FP}}(\mathbb{G}, \mathbf{Set})$, where the latter \mathbf{Hom} -set consists just of product-preserving functors. Indeed, the correspondence extends to natural transformations to give the previously observed equivalence of categories,

$$\mathbf{Group}(\mathbf{Set}^{\mathbb{C}}) \simeq (\mathbf{Group}(\mathbf{Set}))^{\mathbb{C}} \simeq \mathbf{Group}^{\mathbb{C}}.$$

1.1.4 Soundness and completeness

Consider an algebraic theory \mathbb{T} and an equation $s = t$ between terms of the theory. If the equation can be proved from the axioms of the theory, $\mathbb{T} \vdash s = t$, then every model \mathcal{M} of the theory in any FP-category satisfies the equation, $\mathcal{M} \models s = t$. This is called the *soundness* of the equational calculus with respect to categorical models, and it can be shown by a straightforward induction on the equational proof that establishes $\mathbb{T} \vdash s = t$. The converse statement reads:

$$\mathcal{M} \models s = t, \text{ for all } \mathcal{M} \quad \Rightarrow \quad \mathbb{T} \vdash s = t .$$

This is called *completeness*, and (together with soundness) it says that the equational calculus suffices for proving all (and only) the ones that hold in the categorical semantics. For functorial semantics, this holds in an especially strong way: by Proposition 1.1.17, we already know that the syntactic model \mathcal{U} in $\mathcal{C}_{\mathbb{T}}$ is logically generic, in the sense that satisfaction by \mathcal{U} is equivalent to provability in \mathbb{T} ,

$$\mathcal{U} \models s = t \iff \mathbb{T} \vdash s = t .$$

But since \mathcal{U} is also universal in the sense of Definition 1.1.21, it follows immediately that we also have soundness and completeness:

Theorem 1.1.25 (Soundness and completeness of equational logic). *For any equation $s = t$, we have $\mathbb{T} \vdash s = t$ if and only if every model \mathcal{M} in every FP-category \mathcal{C} satisfies $s = t$.*

Proof. If $\mathbb{T} \vdash s = t$, then by Proposition 1.1.17 (the syntactic construction of $\mathcal{C}_{\mathbb{T}}$) we have $\mathcal{U} \models s = t$, meaning that $s^{\mathcal{U}} = t^{\mathcal{U}}$. But then for any model \mathcal{M} in an FP-category \mathcal{C} , we obtain $\mathcal{M} \models s = t$ by applying the classifying functor $\mathcal{M}^{\sharp} : \mathcal{C}_{\mathbb{T}} \rightarrow \mathcal{C}$, which preserves the interpretations of s and t ,

$$\mathcal{M}^{\sharp}(s^{\mathcal{U}}) = s^{\mathcal{M}^{\sharp}(\mathcal{U})} = s^{\mathcal{M}}$$

and so from $s^{\mathcal{U}} = t^{\mathcal{U}}$ we can infer $s^{\mathcal{M}} = t^{\mathcal{M}}$.

Conversely, if $\mathcal{M} \models s = t$ for every model \mathcal{M} , then in particular $\mathcal{U} \models s = t$, and so $\mathbb{T} \vdash s = t$, since \mathcal{U} is generic. \square

Classically, it is seldom the case that there exists a generic model; instead, we consider the class of all models in **Set**. Completeness with respect to such a restricted class of models is of course a stronger statement than completeness with respect to all models in all categories. Toward the classical result, we first consider completeness with respect to “variable models” in **Set**, i.e. in presheaf categories $\mathbf{Set}^{\mathcal{C}^{\text{op}}}$.

Proposition 1.1.26. *Let \mathbb{T} be an algebraic theory. The Yoneda embedding*

$$y : \mathcal{C}_{\mathbb{T}} \rightarrow \widehat{\mathcal{C}_{\mathbb{T}}}$$

is a generic model for \mathbb{T} .

Proof. The Yoneda embedding $y : \mathcal{C}_{\mathbb{T}} \rightarrow \widehat{\mathcal{C}_{\mathbb{T}}}$ preserves limits, and in particular finite products, hence it determines a model $\mathcal{Y} = y\mathcal{U}$ in the category of presheaves $\widehat{\mathcal{C}_{\mathbb{T}}}$. Like all models, \mathcal{Y} satisfies all the equations that hold in \mathcal{U} , simply because y is an FP functor. But because y is also faithful, any equation that holds in \mathcal{Y} must already hold in \mathcal{U} , and therewith in all models. \square

Example 1.1.27. We consider group theory one last time. As a presheaf on (the classifying category of) the theory of groups \mathbb{G} , the generic group \mathcal{Y} satisfies every equation that is satisfied by all groups, and no others. Let us describe its underlying object Y explicitly as a variable set. The presheaf Y is represented by the context with one variable,

$$Y = y[x_1] = \mathbb{G}(-, [x_1]) .$$

The values of this functor thus comprise a family of sets parametrized by the objects $[n] = [x_1, \dots, x_n]$ of \mathbb{G} ; namely, for every $n \in \mathbb{N}$, we have the set

$$Y_n = \mathbb{G}([x_1, \dots, x_n], [x_1])$$

consisting of all terms in n variables, modulo the equations of group theory; but this is just the set of elements of the *free group* $F(n)$ on n generators! Thus we have

$$Y_n = \mathbb{G}([x_1, \dots, x_n], [x_1]) \cong |F(n)| = \mathbf{Set}(1, |F(n)|) \cong \mathbf{Group}(F(1), F(n)).$$

Moreover, the unit, inverse, and multiplication operations on Y agree at each stage Y_n with the operations on the free group $F(n)$ (as the reader should verify).

To summarize, the presheaf of groups $\mathcal{Y} : \mathbb{G}^{\text{op}} \rightarrow \mathbf{Group}$ on the theory \mathbb{G} of groups is isomorphic to the functor $F : \mathbb{G}^{\text{op}} \rightarrow \mathbf{Group}$ of free groups $F(n)$ on n -generators (at least pointwise). We will see why this is so in more detail in section 1.2.1 .

Finally, we consider the completeness of equational logic with respect to **Set**-valued models $\mathcal{M} : \mathcal{C}_{\mathbb{T}} \rightarrow \mathbf{Set}$, which of course correspond to classical \mathbb{T} -algebras. We need the following:

Lemma 1.1.28. *For any small category \mathbb{C} , there is a jointly faithful family $(E_i)_{i \in I}$ of FP-functors $E_i : \mathbf{Set}^{\mathbb{C}^{\text{op}}} \rightarrow \mathbf{Set}$, with I a set. That is, for any maps $f, g : A \rightarrow B$ in $\mathbf{Set}^{\mathbb{C}^{\text{op}}}$, if $E_i(f) = E_i(g)$ for all $i \in I$, then $f = g$.*

Proof. We can take $I = \mathbb{C}_0$ and the evaluation functors $E_C = \text{eval}_C : \mathbf{Set}^{\mathbb{C}^{\text{op}}} \rightarrow \mathbf{Set}$, for all $C \in \mathbb{C}$. These are clearly jointly faithful, and they preserve all limits and colimits, which are constructed pointwise in presheaves. \square

Proposition 1.1.29. *Suppose \mathbb{T} is an algebraic theory. For every equation $s = t$ between terms of the theory \mathbb{T} ,*

$$\mathcal{M} \models s = t \quad \text{for all models } \mathcal{M} \text{ in } \mathbf{Set} \iff \mathbb{T} \vdash s = t .$$

*Thus the equational logic of algebraic theories is sound and complete with respect to **Set**-valued semantics.*

Proof. Combine the foregoing lemma with the fact, from Proposition 1.1.26, that the Yoneda embedding is a generic model. \square

Exercise 1.1.30. We described the object part of the functor $Y = \mathbf{y}U : \mathbb{G}^{\text{op}} \rightarrow \mathbf{Set}$ represented by the underlying object $U = [x_1]$ of the universal group \mathcal{U} , in terms of the free groups $F(n)$. What is the action of Y on the arrows of \mathbb{G} in these terms? Also describe the group structure on Y in $\widehat{\mathbb{G}}$ explicitly.

Exercise 1.1.31. Let $t = t(x_1, \dots, x_n)$ be a term of group theory in variables x_1, \dots, x_n . On the one hand we can think of t as an element of the free group $F(n)$, and on the other we can consider the interpretation of t with respect to the representable group \mathcal{Y} in $\widehat{\mathbb{G}}$, namely as a natural transformation $t^{\mathcal{Y}} : Y^n \Rightarrow Y$. Suppose $s = s(x_1, \dots, x_n)$ is another such term in the same variables x_1, \dots, x_n . Show that $s^{\mathcal{Y}} = t^{\mathcal{Y}}$ if, and only if, $s = t$ in the free group $F(n)$.

1.1.5 Functorial semantics

Let us summarize our treatment of algebraic theories thus far. We have reformulated certain traditional *logical* notions in terms of *categorical* ones. The traditional approach may be described as involving the following four different parts:

Terms

There is an underlying *type theory* consisting of types and terms. For algebraic theories there is only one type, which is not even explicitly mentioned. The terms are built from variables and some basic operation symbols.

Equations

Algebraic theories have a particularly simple *logic* that involves only equations between terms and equational reasoning, which is basically substitution of equals for equals.

Theories

An *algebraic theory* is given by a set of basic terms and a set of axioms, which are just equations.

Models

Algebraic theories are interpreted as sets equipped with operations. An interpretation is a *model* if it satisfies all the axioms of the theory, which just means that the functions interpreting the terms that occur in the axioms are actually equal.

The alternative approach of *functorial semantics* developed here may be summarized as follows:

Theories are categories

From a given theory we construct a (structured) category, which captures the same information in a way that it does not depend on a particular presentation by basic operations and axioms.

Models are functors

A model is a (structure-preserving) functor from the (category representing the) theory to a category with the structure needed to interpret it. The requirement that the axioms of the theory must be satisfied by a model translates to the requirement that the model is a functor, and that it preserves the structure. For models of algebraic theories, we require only that they preserve finite products, which along with functoriality ensures that all valid equations of the theory are preserved, and the axioms are therefore satisfied.

Homomorphisms are natural transformations

We obtain the notion of a homomorphism of models for free: since models are functors, homomorphisms between them are natural transformations. Such homomorphisms between models of algebraic theories turn out to agree with the usual notion of a homomorphism that respects the algebraic structure.

Universal models

By admitting models in categories other than **Set**, functorial semantics admits *universal models*: a model \mathcal{U} in the classifying category $\mathcal{C}_{\mathbb{T}}$, such that any model anywhere is a functorial image of \mathcal{U} by an essentially unique, structure-preserving functor. Such a universal model is then *logically generic*, in the sense that it has all and only those logical properties that are had by all models, since such properties are preserved by the functors in question.

Logical completeness

The construction of the classifying category $\mathcal{C}_{\mathbb{T}}$ from the syntax of the theory \mathbb{T} implies the *soundness and completeness* of the logic with respect to general categorical semantics. Completeness with respect to a restricted class of models, such as the **Set**-valued ones, results from an embedding theorem for the classifying category.

1.2 Lawvere duality

The scheme of functorial semantics that we outlined in the previous section applies to a wide range of logics other than algebraic theories, some of which will be considered in later chapters. A further aspect of this approach is especially transparent in the case of algebraic theories; namely, a deep and fascinating *duality* relating syntax and semantics. We devote the rest of this chapter to its investigation.

1.2.1 Logical duality

There is a remarkable and far-reaching duality in logic of the form

$$\text{Syntax} \simeq \text{Semantics}^{\text{op}}.$$

It was discovered by F.W. Lawvere in the 1960s and presented in some early papers, [Law63a, Law63b, Law65], but it has still hardly been noticed in conventional logic—perhaps because its recognition requires the tools of category theory.

We can see this duality quite clearly in the case of algebraic theories. Let $\mathcal{C}_{\mathbb{T}}$ be the classifying category for an equational theory \mathbb{T} , like the theory of groups, constructed syntactically as in section 1.1.2 above. So the objects of $\mathcal{C}_{\mathbb{T}}$ are contexts of variables $[x_1 \dots, x_n]$, up to renaming, and the arrows $(t_1, \dots, t_n) : [x_1 \dots, x_m] \rightarrow [x_1 \dots, x_n]$ are n -tuples of terms in context $[x_1 \dots, x_m \mid t_i]$, up to \mathbb{T} -provable equality. We will see that this syntactic category $\mathcal{C}_{\mathbb{T}}$ is in fact dual to a certain subcategory of *models* of \mathbb{T} (in **Set**). Specifically, there is a small, full subcategory $\mathbf{mod}(\mathbb{T}) \hookrightarrow \mathbf{Mod}(\mathbb{T})$ and an equivalence of categories,

$$\mathcal{C}_{\mathbb{T}} \simeq \mathbf{mod}(\mathbb{T})^{\mathrm{op}},$$

making the *syntactic* category $\mathcal{C}_{\mathbb{T}}$ dual to a subcategory of the *semantic* category $\mathbf{Mod}(\mathbb{T})$. Thus, in particular, there is an invariant representation of the syntax of the theory \mathbb{T} “hidden” inside the category of models of \mathbb{T} . Indeed, it is quite easy to specify $\mathbf{mod}(\mathbb{T})$: it is the full subcategory on the *finitely generated free models* $F(n)$ of \mathbb{T} ,

$$\mathbf{mod}(\mathbb{T})_0 = \{F(n) \mid n \in \mathbb{N}\}.$$

Theorem 1.2.1. *Let \mathbb{T} be an algebraic theory, and let*

$$\mathbf{mod}(\mathbb{T}) \hookrightarrow \mathbf{Mod}(\mathbb{T})$$

be the full subcategory of finitely generated free models of \mathbb{T} . Then $\mathbf{mod}(\mathbb{T})^{\mathrm{op}}$ classifies \mathbb{T} models. That is to say, for any FP-category \mathcal{C} , there is an equivalence of categories,

$$\mathbf{Hom}_{\mathbf{FP}}(\mathbf{mod}(\mathbb{T})^{\mathrm{op}}, \mathcal{C}) \simeq \mathbf{Mod}(\mathbb{T}, \mathcal{C}), \quad (1.11)$$

which is natural in \mathcal{C} .

Before giving the (somewhat lengthy, but straightforward) proof of the theorem, let us observe that the claimed syntax-semantics duality follows immediately. Indeed, given (1.11), there is then an equivalence,

$$\mathcal{C}_{\mathbb{T}} \simeq \mathbf{mod}(\mathbb{T})^{\mathrm{op}} \quad (1.12)$$

between the syntactic category $\mathcal{C}_{\mathbb{T}}$ and the opposite of the category $\mathbf{mod}(\mathbb{T})$ of finitely generated free models, because by Proposition 1.1.19, both categories $\mathcal{C}_{\mathbb{T}}$ and $\mathbf{mod}(\mathbb{T})^{\mathrm{op}}$ represent the same “semantics” functor $\mathbf{Mod}(\mathbb{T}, \mathcal{C})$.

Proof of Theorem 1.2.1. First, observe that $\mathbf{mod}(\mathbb{T})^{\mathrm{op}}$ has all finite products, since $\mathbf{mod}(\mathbb{T})$ has all finite coproducts. Indeed, we have

$$\begin{aligned} F(n) + F(m) &\cong F(n + m), \\ 0 &\cong F(0), \end{aligned}$$

for the finitely generated free algebras $F(n)$, since the left adjoint F preserves all colimits.

To determine the universal \mathbb{T} -algebra \mathcal{U} in $\mathbf{mod}(\mathbb{T})^{\text{op}}$ let

$$U = F(1),$$

so that every object in $\mathbf{mod}(\mathbb{T})^{\text{op}}$ is a power of U ,

$$F(n) \cong U^n.$$

We next interpret the signature $\Sigma_{\mathbb{T}}$. For each basic operation symbol $f \in \Sigma_{\mathbb{T}}$, with arity k , there is an element of the free algebra $F(k)$ built from the operation $f^{F(k)} : F(k)^k \rightarrow F(k)$ and the k generators $x_1, \dots, x_k \in F(k)$, namely

$$f^{F(k)}(x_1, \dots, x_k).$$

E.g. in the theory of groups, there is the element $x \cdot y$ in the free group on the two generators x, y . By freeness of $F(1)$, each element $t \in F(k)$ determines a unique homomorphism $\bar{t} : F(1) \rightarrow F(k)$ in $\mathbf{mod}(\mathbb{T})$ taking the generator $x \in F(1)$ to $t = \bar{t}(x)$. Thus there is a homomorphism

$$\overline{f^{F(k)}(x_1, \dots, x_k)} : F(1) \rightarrow F(k) \quad \text{in } \mathbf{mod}(\mathbb{T})$$

associated to the element $f^{F(k)}(x_1, \dots, x_k) \in F(k)$. We take this map as the \mathcal{U} -interpretation of the basic operation symbol f ,

$$f^{\mathcal{U}} = \overline{f^{F(k)}(x_1, \dots, x_k)} : U^k \rightarrow U \quad \text{in } \mathbf{mod}(\mathbb{T})^{\text{op}}.$$

It follows that for any term in context $x_1 \dots, x_k \mid t$, the interpretation

$$[x_1 \dots, x_k \mid t]^{\mathcal{U}} : U^k \rightarrow U$$

will be the unique homomorphism $\overline{t^{F(k)}} : F(1) \rightarrow F(k)$ corresponding to the element $t^{F(k)} \in F(k)$ (proof by induction!).

Moreover, it then follows that for every axiom $(s = t)$ of \mathbb{T} , we have $\mathcal{U} \models s = t$. Indeed,

$$[x_1 \dots, x_k \mid s]^{\mathcal{U}} = [x_1 \dots, x_k \mid t]^{\mathcal{U}} : U^k \rightarrow U$$

if, and only if, the corresponding homomorphisms $\bar{s}, \bar{t} : F(1) \rightarrow F(k)$ agree, which they do just if the associated elements of the free algebra $F(k)$ agree, by the freeness of $F(1)$. And the latter holds, in turn, simply because $F(k)$ is a \mathbb{T} -algebra. Indeed, consider the example of the two generators x, y of the free *abelian* group $F(2)$, for which we have $x \cdot y = y \cdot x$ simply because $F(2)$ is abelian. Thus we indeed have a \mathbb{T} -model \mathcal{U} in $\mathbf{mod}(\mathbb{T})^{\text{op}}$, consisting of the free algebras.

We next show that \mathcal{U} has the required universal property, in three steps:

Step 1. Let \mathcal{A} be any \mathbb{T} -algebra in **Set**. Then there is a product-preserving functor,

$$\mathcal{A}^\sharp : \mathbf{mod}(\mathbb{T})^{\mathrm{op}} \rightarrow \mathbf{Set}$$

with $\mathcal{A}^\sharp(\mathcal{U}) \cong \mathcal{A}$ (as \mathbb{T} -models), namely:

$$\mathcal{A}^\sharp(-) = \mathbf{Hom}_{\mathbf{Mod}(\mathbb{T})}(-, \mathcal{A}),$$

where we of course restrict the representable functor $\mathbf{Hom}_{\mathbf{Mod}(\mathbb{T})}(-, \mathcal{A}) : \mathbf{Mod}(\mathbb{T})^{\mathrm{op}} \rightarrow \mathbf{Set}$ along the (full) inclusion

$$\mathbf{mod}(\mathbb{T}) \hookrightarrow \mathbf{Mod}(\mathbb{T})$$

of the finitely generated, free algebras. The functor

$$\mathcal{A}^\sharp : \mathbf{mod}(\mathbb{T})^{\mathrm{op}} \rightarrow \mathbf{Set}$$

clearly preserves products: for each object $U^n \in \mathbf{mod}(\mathbb{T})^{\mathrm{op}}$, we have

$$\mathcal{A}^\sharp(U^n) = \mathbf{Hom}_{\mathbf{Mod}(\mathbb{T})}(F(n), \mathcal{A}) \cong \mathbf{Hom}_{\mathbf{Set}}(n, V(\mathcal{A})) \cong A^n.$$

And in particular $\mathcal{A}^\sharp(U) \cong A$.

Finally, let us show that for any basic operation f , we have $\mathcal{A}^\sharp(f^U) = f^A$, up to isomorphism. Indeed, given any algebra \mathcal{A} and operation $f^A : A^n \rightarrow A$, we have a commutative diagram,

$$\begin{array}{ccc} A^n & \xrightarrow{f^A} & A \\ \cong \downarrow & & \downarrow \cong \\ \mathbf{Hom}(F(n), \mathcal{A}) & \xrightarrow{f^*} & \mathbf{Hom}(F(1), \mathcal{A}) \end{array} \quad (1.13)$$

where f^* is precomposition with the homomorphism

$$F(n) \xleftarrow{\overline{f^{F(n)}(x_1, \dots, x_n)}} F(1)$$

To see that (1.13) commutes, take any $(a_1, \dots, a_n) \in A^n$ with associated homomorphism $\overline{(a_1, \dots, a_n)} : F(n) \rightarrow \mathcal{A}$ and precompose with $\overline{f^{F(n)}(x_1, \dots, x_n)}$ to get a map $F(1) \rightarrow \mathcal{A}$, picking out the element

$$\begin{aligned} \overline{(a_1, \dots, a_n)} \circ \overline{f^{F(n)}(x_1, \dots, x_n)}(x) &= \overline{(a_1, \dots, a_n)}(f^{F(n)}(x_1, \dots, x_n)) \\ &= \overline{(a_1, \dots, a_n)} \circ f^{F(n)}(x_1, \dots, x_n) \\ &= f^A \circ \overline{(a_1, \dots, a_n)}(x_1, \dots, x_n) \\ &= f^A(a_1, \dots, a_n) \end{aligned}$$

where x is the generator of $F(1)$, and using the fact that $\overline{(a_1, \dots, a_n)}$ is a homomorphism and therefore commutes with the respective interpretations of f .

But now note that

$$F(n) \xleftarrow{\overline{f^{F(n)}(x_1, \dots, x_n)}} F(1)$$

in $\mathbf{mod}(\mathbb{T})$ is

$$U^n \xrightarrow{f^U} U$$

in $\mathbf{mod}(\mathbb{T})^{\text{op}}$, and that $\text{Hom}(F(n), \mathcal{A}) = \mathcal{A}^\sharp(U^n)$ and $f^* = \mathcal{A}^\sharp(f^U)$. Thus (1.13) shows that indeed $\mathcal{A}^\sharp(f^U) = f^{\mathcal{A}}$, up to isomorphism. Thus we indeed have $\mathcal{A}^\sharp(\mathcal{U}) \cong \mathcal{A}$ as algebras, as required.

We leave it to the reader to verify that any homomorphism $h : F(\mathcal{U}) \rightarrow G(\mathcal{U})$ of \mathbb{T} -algebras $F(\mathcal{U}), G(\mathcal{U})$ arising from FP-functors $F, G : \mathbf{mod}(\mathbb{T})^{\text{op}} \rightarrow \mathbf{Set}$ is of the form $h = \vartheta_{\mathcal{U}}$ for a unique natural transformation $\vartheta : F \rightarrow G$.

Exercise 1.2.2. Show this.

Step 2. Let \mathbb{C} be any (locally small) category, and \mathcal{A} a \mathbb{T} -algebra in $\mathbf{Set}^{\mathbb{C}}$. Using the isomorphism

$$\mathbf{Mod}(\mathbb{T}, \mathbf{Set}^{\mathbb{C}}) \cong \mathbf{Mod}(\mathbb{T})^{\mathbb{C}},$$

each $\mathcal{A}(C)$ is a \mathbb{T} -algebra (in \mathbf{Set}), which by Step 1 has a classifying functor,

$$\mathcal{A}(C)^\sharp : \mathbf{mod}(\mathbb{T})^{\text{op}} \rightarrow \mathbf{Set}.$$

Together, these determine a single functor $\mathcal{A}^\sharp : \mathbf{mod}(\mathbb{T})^{\text{op}} \rightarrow \mathbf{Set}^{\mathbb{C}}$, defined on any $U^n \in \mathbf{mod}(\mathbb{T})^{\text{op}}$ by

$$(\mathcal{A}^\sharp(U^n))(C) \cong \mathcal{A}(C)^\sharp(U^n) \cong (\mathcal{A}(C))^n.$$

The action on arrows $U^n \rightarrow U^m$ in $\mathbf{mod}(\mathbb{T})^{\text{op}}$ is similarly determined pointwise by the components

$$(\mathcal{A}^\sharp(U^n))(C) \cong \mathcal{A}(C)^\sharp(U^n) \rightarrow \mathcal{A}(C)^\sharp(U^m) = (\mathcal{A}^\sharp(U^m))(C),$$

for all $C \in \mathbb{C}$.

In this way, we have an FP-functor $\mathcal{A}^\sharp : \mathbf{mod}(\mathbb{T})^{\text{op}} \rightarrow \mathbf{Set}^{\mathbb{C}}$, and an isomorphism of models $\mathcal{A}^\sharp(\mathcal{U}) \cong \mathcal{A}$ in $\mathbf{Set}^{\mathbb{C}}$. It is then clear that any natural transformation $\mathcal{A}^\sharp \rightarrow \mathcal{B}^\sharp$ gives rise to a homomorphism $\mathcal{A}^\sharp(\mathcal{U}) \rightarrow \mathcal{B}^\sharp(\mathcal{U})$, and that the resulting functor

$$\text{Hom}_{\text{FP}}(\mathbf{mod}(\mathbb{T})^{\text{op}}, \mathbf{Set}^{\mathbb{C}}) \rightarrow \mathbf{Mod}(\mathbb{T}, \mathbf{Set}^{\mathbb{C}})$$

is an equivalence.

Step 3. For the general case, let \mathcal{C} be any (locally small) FP-category, and \mathbf{A} a \mathbb{T} -algebra in \mathcal{C} . Use the Yoneda embedding

$$y : \mathcal{C} \hookrightarrow \mathbf{Set}^{\mathcal{C}^{\text{op}}}$$

to send \mathbf{A} to an algebra $\mathcal{A} = y(\mathbf{A})$ in $\mathbf{Set}^{\mathcal{C}^{\text{op}}}$ (since y preserves finite products). Now apply Step 2 to get a classifying functor,

$$\mathcal{A}^\sharp : \mathbf{mod}(\mathbb{T})^{\text{op}} \rightarrow \mathbf{Set}^{\mathcal{C}^{\text{op}}}.$$

We claim that \mathcal{A}^\sharp factors through the Yoneda embedding by an FP-functor \mathbf{A}^\sharp ,

$$\begin{array}{ccc} & & \mathbf{Set}^{\mathcal{C}^{\text{op}}} \\ & \nearrow \mathcal{A}^\sharp & \uparrow y \\ \mathbf{mod}(\mathbb{T})^{\text{op}} & \xrightarrow{\quad \mathbf{A}^\sharp \quad} & \mathcal{C}. \end{array}$$

Indeed, we know that the objects of $\mathbf{mod}(\mathbb{T})^{\text{op}}$ all have the form U^n , and for their images we have

$$\mathcal{A}^\sharp(U^n) \cong \mathcal{A}^\sharp(U)^n \cong y(\mathbf{A})^n \cong y(\mathbf{A}^n).$$

Thus the images of the objects of $\mathbf{mod}(\mathbb{T})^{\text{op}}$ are all representable. Since y is full and faithful, the claim is established, and the resulting functor $\mathbf{A}^\sharp : \mathbf{mod}(\mathbb{T})^{\text{op}} \rightarrow \mathcal{C}$ preserves finite products because \mathcal{A}^\sharp does so, and y creates them. Clearly,

$$\mathbf{A}^\sharp(\mathcal{U}) \cong \mathbf{A},$$

since y reflects isos.

Naturality of the equivalence

$$\mathbf{Hom}_{\mathbf{FP}}(\mathbf{mod}(\mathbb{T})^{\text{op}}, \mathcal{C}) \simeq \mathbf{Mod}(\mathbb{T}, \mathcal{C}),$$

in \mathcal{C} is essentially automatic, using the fact that it is induced by evaluating an FP functor $F : \mathbf{mod}(\mathbb{T})^{\text{op}} \rightarrow \mathcal{C}$ at the universal model \mathcal{U} in $\mathbf{mod}(\mathbb{T})^{\text{op}}$. \square

As already mentioned, since the classifying category is uniquely determined, up to equivalence, by its universal property, combining the foregoing theorem with the syntactic construction of $\mathcal{C}_{\mathbb{T}}$ given in theorem 1.1.19 yields the following:

Corollary 1.2.3 (Logical duality for algebraic theories). *For any algebraic theory \mathbb{T} , there is an equivalence of categories,*

$$\mathcal{C}_{\mathbb{T}} \simeq \mathbf{mod}(\mathbb{T})^{\text{op}} \tag{1.14}$$

between the syntactic category $\mathcal{C}_{\mathbb{T}}$ and the opposite of the category $\mathbf{mod}(\mathbb{T})$ of finitely generated, free models.

Thus the construction of the classifying category $\mathcal{C}_{\mathbb{T}}$ from the syntax of \mathbb{T} , on the one hand, and its semantic construction as $\mathbf{mod}(\mathbb{T})$, taken together, imply that there is an invariant representation of the *syntax* of \mathbb{T} hidden, as it were, in the opposite of the *semantics*, namely the category $\mathbf{Mod}(\mathbb{T})$ of all \mathbb{T} -models. The reader may wish to reflect on the importance of (i) considering the *category* of all models, rather than the mere collection of them, and (ii) generalizing from set-theoretic to categorical models, in arriving at the fundamental logical duality expressed by (1.14).

In section ?? below, we shall consider how to actually *recover* this syntactic category $\mathcal{C}_{\mathbb{T}}$ from the semantic category $\mathbf{Mod}(\mathbb{T})$ by identifying the subcategory $\mathbf{mod}(\mathbb{T})$ intrinsically—indeed, it will be seen to consist of certain continuous functors $\mathbf{Mod}(\mathbb{T}) \rightarrow \mathbf{Set}$. Before doing so, however, let us examine the fundamental equivalence (1.14) explicitly in a few special cases.

Example 1.2.4. Let \mathbb{T}_0 be the “trivial theory”, with no basic operations or equations. A model of \mathbb{T}_0 in \mathbf{Set} is just a set X , equipped with no operations, and satisfying no further conditions (and similarly in any other FP category). Thus \mathbb{T}_0 is the pure theory of equality on an object.

All \mathbb{T}_0 -algebras are free, and the finitely generated ones are just the finite sets, thus

$$\mathbf{mod}(\mathbb{T}_0) = \mathbf{Set}_{\text{fin}}$$

is the category of finite sets (to be specific, let us take a skeleton, with one n -element set $[n]$ for each $n \in \mathbb{N}$). Our theorem 1.2.1 then tells us that, for any FP category \mathcal{C} , there is an equivalence

$$\mathbf{Hom}_{\text{FP}}(\mathbf{Set}_{\text{fin}}^{\text{op}}, \mathcal{C}) \simeq \mathbf{Mod}(\mathbb{T}_0, \mathcal{C}) \simeq \mathcal{C}.$$

This simply says that $\mathbf{Set}_{\text{fin}}^{\text{op}}$ is the free FP category on one object. Equivalently, $\mathbf{Set}_{\text{fin}}$ is the free finite *coproduct* category on one object. And this is indeed the case, as one can easily see directly (the objects are $0, 1, 1 + 1, 1 + 1 + 1, \dots$).

The logical duality of corollary 1.2.3 now tells us that the dual of the category of finite sets is the *syntactic category* of the theory of equality \mathbb{T}_0 . The terms of this theory are simply tuples of variables (x_1, \dots, x_n) , and the valid equations are those that are true of them *as terms*. Our corollary tells us that this is the category of finite sets, which we can see by reading the contexts $[x_1, \dots, x_n]$ as coproducts $1 + \dots + 1$ and a tuple such as $(x_2, x_5) : [x_1, \dots, x_5] \rightarrow [x_1, x_2]$ as a *cotuple* like $[i_2, i_5] : 1 + 1 \rightarrow 1_1 + \dots + 1_5$.

Example 1.2.5. For a less trivial example, consider the theory \mathbb{T}_{Ab} of abelian groups. Duality tells us that the syntactic category $\mathcal{C}_{\mathbb{T}_{\text{Ab}}}$ is dual to the category of finitely generated, free abelian groups \mathbf{Ab}_{fg} ,

$$\mathcal{C}_{\mathbb{T}_{\text{Ab}}} \simeq \mathbf{Ab}_{\text{fg}}^{\text{op}}.$$

This gives us a representation of the syntax of (abelian) group theory in the category of abelian groups, which can be described concretely as follows:

- the basic types of variables $[-] = 1$, $[x_1] = U$, $[x_1, x_2] = U \times U, \dots$ are represented by the abelian groups $\{0\}$, \mathbb{Z} , $\mathbb{Z} + \mathbb{Z} = \mathbb{Z}^2, \dots$,

- the group unit $u : 1 \rightarrow U$ is the zero homomorphism $0 : \mathbb{Z} \rightarrow \{0\}$,
- the inverse $i : U \rightarrow U$ is the homomorphism $\mathbb{Z} \rightarrow \mathbb{Z}$ taking 1 to -1 ,
- the group operation $m : U \times U \rightarrow U$ is the homomorphism $+$: $\mathbb{Z} \rightarrow \mathbb{Z} + \mathbb{Z}$ taking 1 to $\langle 1, 1 \rangle = \langle 1, 0 \rangle + \langle 0, 1 \rangle$ (using $\mathbb{Z} + \mathbb{Z} \cong \mathbb{Z} \times \mathbb{Z}$),
- the laws of abelian groups (and no further ones!) hold under this interpretation, because the group structure on any abelian group A is induced by precomposing with these “co-operations”, by (1.13), as indicated in the following diagram for the sum $a + b$ of elements $a, b \in A$.

$$\begin{array}{ccc}
 \mathbb{Z} + \mathbb{Z} & \xrightarrow{(a, b)} & A \\
 \uparrow + & \nearrow a + b & \\
 \mathbb{Z} & &
 \end{array}$$

Example 1.2.6. The category of *affine schemes* is, by definition, the dual of the category of commutative rings with unit,

$$\mathbf{Scheme}_{\text{aff}} = \mathbf{Ring}^{\text{op}}$$

There is therefore a ring object in affine schemes – called the *affine line* – based on the finitely generated free algebra $F(1) = \mathbb{Z}[x]$, the ring of polynomials in one variable x with integer coefficients. The “co-operations” of $+$ and \cdot are given in rings by the homomorphisms $\mathbb{Z}[x] \rightarrow \mathbb{Z}[x, y]$ taking the generator x to the elements $x + y$ and $x \cdot y$.

Exercise 1.2.7. Prove directly that $\mathbf{Set}_{\text{fin}}$ is the free finite coproduct category on one object.

Exercise 1.2.8. Show that for any algebraic theory \mathbb{T} , the forgetful functor $V : \mathbf{Mod}(\mathbb{T}) \rightarrow \mathbf{Set}$ underlies an algebra in the functor category $\mathbf{Set}^{\mathbf{Mod}(\mathbb{T})}$. In more detail, each n -ary operation symbol f determines a natural transformation $f^V : V^n \rightarrow V$, since the homomorphisms in $\mathbf{Mod}(\mathbb{T})$ commute with the various operations interpreting f . Indeed, given any algebra \mathcal{A} we have the underlying set $V(\mathcal{A}) = A$ and an operation $f^{\mathcal{A}} : A^n \rightarrow A$, and for every homomorphism $h : \mathcal{A} \rightarrow \mathcal{B}$ to another algebra \mathcal{B} , there is a commutative square,

$$\begin{array}{ccc}
 A^n & \xrightarrow{h^n} & B^n \\
 f^{\mathcal{A}} \downarrow & & \downarrow f^{\mathcal{B}} \\
 A & \xrightarrow{h} & B.
 \end{array} \tag{1.15}$$

So we can set $(f^V)_{\mathcal{A}} = f^{\mathcal{A}}$ to get a natural transformation $f^V : V^n \rightarrow V$. Now check that this really is an algebra in $\mathbf{Set}^{\mathbf{Mod}(\mathbb{T})}$.

Exercise 1.2.9. * Show that the algebra described in the previous exercise is represented by the universal one \mathcal{U} in $\mathbf{mod}(\mathbb{T})^{\mathrm{op}} \hookrightarrow \mathbf{Mod}(\mathbb{T})^{\mathrm{op}}$ under the (covariant) Yoneda embedding,

$$y : \mathbf{Mod}(\mathbb{T})^{\mathrm{op}} \longrightarrow \mathbf{Set}^{\mathbf{Mod}(\mathbb{T})}.$$

1.2.2 Lawvere algebraic theories

Nothing in the foregoing account of the functorial semantics for algebraic theories really depended on the primarily syntactic nature of such theories, i.e. their specification in terms of operations and equations. We can thus immediately generalize it to “abstract” algebraic theories, which can be regarded as providing a *presentation-free* notion of an algebraic theory.

Definition 1.2.10 (cf. Definition 1.1.2). A *Lawvere algebraic theory* \mathbb{A} is a small category with finite products whose objects form a sequence A^0, A^1, A^2, \dots such that $A^m \times A^n = A^{m+n}$ for all $m, n \in \mathbb{N}$. In particular, $1 = A^0$ is the terminal object, and every object is a product of finitely many copies of $A = A^1$.

A *model* of a Lawvere algebraic theory \mathbb{A} in any category \mathcal{C} with finite products is a finite-product-preserving functor $M : \mathbb{A} \rightarrow \mathcal{C}$, and a *homomorphism of models* is a natural transformation $\vartheta : M \rightarrow M'$ between such functors.

We could just as well have taken the natural numbers $0, 1, 2, \dots$ themselves as the objects of a Lawvere algebraic theory \mathbb{A} , but the notation A^n is more suggestive. A Lawvere algebraic theory \mathbb{A} in the sense of the above definition determines an algebraic theory in the sense of Definition 1.1.2 as follows. As basic operations with arity k we take all of the morphisms $A^k \rightarrow A$:

$$\Sigma(\mathbb{A})_k = \mathbf{Hom}_{\mathbb{A}}(A^k, A) \tag{1.16}$$

There is a canonical interpretation in \mathbb{A} of terms built from variables and morphisms $A^k \rightarrow A$, namely each morphism is interpreted by itself, and variables are interpreted as product projects, as usual. An equation $u = v$ is taken as an axiom of the theory \mathbb{A} if the canonical interpretations of u and v coincide. Of course, the conventional logical notions of a model 1.1.11 and a homomorphism of models then also correspond to the new, functorial ones in an obvious way.

This more abstract view of algebraic theories immediately suggests some interesting examples.

Example 1.2.11. The algebraic theory \mathcal{C}^∞ of smooth maps is the category whose objects are n -dimensional Euclidean spaces $1, \mathbb{R}, \mathbb{R}^2, \dots$, and whose morphisms are \mathcal{C}^∞ -maps between them. Recall that a \mathcal{C}^∞ -map $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is a function which has all higher partial derivatives, and that a function $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$ is a \mathcal{C}^∞ -map exactly when all of its composites $\pi_k \circ f : \mathbb{R}^n \rightarrow \mathbb{R}$ with the projections $\pi_k : \mathbb{R}^m \rightarrow \mathbb{R}$ are \mathcal{C}^∞ -maps.

A model of this theory in \mathbf{Set} is a finite product preserving functor $A : \mathcal{C}^\infty \rightarrow \mathbf{Set}$. Up to natural isomorphism, such a model can be described as follows. A \mathcal{C}^∞ -model is given

by a set A and for every smooth map $f : \mathbb{R}^n \rightarrow \mathbb{R}$ a function $Af : A^n \rightarrow A$ such that if $f : \mathbb{R}^n \rightarrow \mathbb{R}$ and $g_i : \mathbb{R}^m \rightarrow \mathbb{R}$, $i = 1, \dots, n$, are smooth maps then, for all $a_1, \dots, a_m \in A$,

$$Af((Ag_1)\langle a_1, \dots, a_m \rangle, \dots, (Ag_n)\langle a_1, \dots, a_m \rangle) = A(f \circ \langle g_1, \dots, g_n \rangle)\langle a_1, \dots, a_m \rangle.$$

In particular, since multiplication and addition are smooth maps, A is a commutative ring with unit. Such structures are known as \mathcal{C}^∞ -rings. Therefore, the models in **Set** of the theory of smooth maps are the \mathcal{C}^∞ -rings (cf. [MR91]).

Example 1.2.12. Recall that a (*total*) *recursive function* $f : \mathbb{N}^m \rightarrow \mathbb{N}^n$ is one that can be computed by a Turing machine. This means that there exists a Turing machine which on input $\langle a_1, \dots, a_m \rangle$ outputs the value of $f\langle a_1, \dots, a_m \rangle$. The algebraic theory **Rec** of recursive functions is the category whose objects are finite powers of the natural numbers $1, \mathbb{N}, \mathbb{N}^2, \dots$, and whose morphisms are recursive functions between them. The models of this theory in a category \mathcal{C} with finite products give a notion of computability in \mathcal{C} .

Indeed, let us consider the category of all set-theoretic models $\mathcal{R} = \mathbf{Mod}(\mathbf{Rec})$. First, there is the “identity” model $I \in \mathcal{R}$, defined by $I\mathbb{N}^k = \mathbb{N}^k$ and $If = f$. Given any model $S \in \mathcal{R}$, its object part is determined by $S_1 = S\mathbb{N}$ since $S\mathbb{N}^k = S_1^k$. For every $n \in \mathbb{N}$ there is a morphism $1 \rightarrow \mathbb{N}$ in **Rec** defined by $\star \mapsto n$. Thus we have for each $n \in \mathbb{N}$ an element $s_n = S(\star \mapsto n) : 1 \rightarrow S_1$. This defines a function $s : \mathbb{N} \rightarrow S_1$ which in turn determines a natural transformation $\sigma : I \Rightarrow S$ whose component at \mathbb{N}^k is $s \times \dots \times s : \mathbb{N}^k \rightarrow S_1^k$.

Example 1.2.13. In a category \mathcal{C} with finite products every object $A \in \mathcal{C}$ determines a full subcategory consisting of the finite powers $1, A, A^2, A^3, \dots$ and all morphisms between them. This is the *total theory* $\mathbb{T}(A)$ of the object A in \mathcal{C} .

Free algebras

In order to extend the duality theory of the foregoing section to the abstract case, we will require the notion of a *free model* of an abstract algebraic theory. Of course, we already have the conventional notion of free models determined in terms of the associated conventional algebraic theory given by (1.16). But we can also determine free models directly in terms of the abstract theory, which then of course applies to the conventional case as well.

Let \mathbb{A} be a Lawvere algebraic theory, with objects $1, A, A^2, \dots$. We have the category of models,

$$\mathbf{Mod}(\mathbb{A}) = \mathbf{Hom}_{\mathbf{FP}}(\mathbb{A}, \mathbf{Set}).$$

Let us first define the *forgetful functor* by evaluating at the generating object $A \in \mathbb{A}$,

$$U : \mathbf{Mod}(\mathbb{A}) \rightarrow \mathbf{Set} \tag{1.17}$$

$$(M : \mathbb{A} \rightarrow \mathbf{Set}) \mapsto M(A). \tag{1.18}$$

We shall also write

$$|M| = U(M) = M(A).$$

Now for the (finitary) free functor $F_{\text{fin}} : \mathbf{Set}_{\text{fin}} \rightarrow \mathbf{Mod}(\mathbb{A})$, we set:

$$\begin{aligned} F(0) &= \mathbf{Hom}_{\mathbb{A}}(1, -) \\ F(1) &= \mathbf{Hom}_{\mathbb{A}}(A, -) \\ &\vdots \\ F(n) &= \mathbf{Hom}_{\mathbb{A}}(A^n, -). \end{aligned}$$

Note that this is a composite of the two (contravariant) functors,

$$\mathbf{Set}_{\text{fin}} \rightarrow \mathbb{A}^{\text{op}} \rightarrow \mathbf{Mod}(\mathbb{A}),$$

given by $n \mapsto A^n$ and $X \mapsto \mathbf{Hom}_{\mathbb{A}}(X, -)$, and is therefore (covariantly) functorial. Note also that the representables $\mathbf{Hom}_{\mathbb{A}}(A^n, -)$ do indeed preserve finite products, and are therefore in the full subcategory $\mathbf{Mod}(\mathbb{A}) \hookrightarrow \mathbf{Set}^{\mathbb{A}}$.

For adjointness we need to check that:

$$\mathbf{Hom}_{\mathbf{Mod}(\mathbb{A})}(F(n), M) \cong \mathbf{Hom}_{\mathbf{Set}}(n, |M|) \quad (1.19)$$

(naturally in both arguments, of course). The right-hand side is plainly just

$$\mathbf{Hom}_{\mathbf{Set}}(n, |M|) \cong |M|^n.$$

For the left-hand side we have:

$$\begin{aligned} \mathbf{Hom}_{\mathbf{Mod}(\mathbb{A})}(F(n), M) &= \mathbf{Hom}_{\mathbf{Mod}(\mathbb{A})}(\mathbf{Hom}_{\mathbb{A}}(A^n, -), M) \\ &= \mathbf{Hom}_{\mathbf{Set}^{\mathbb{A}}}(\mathbf{Hom}_{\mathbb{A}}(A^n, -), M) \\ &\cong M(A^n) && \text{(by Yoneda)} \\ &\cong M(A)^n && (M \text{ is FP}) \\ &= |M|^n && (1.17). \end{aligned}$$

The full definition of the free functor

$$F : \mathbf{Set} \rightarrow \mathbf{Mod}(\mathbb{A})$$

is then given by writing an arbitrary set X as a (filtered) colimit of its finite subsets $X_i \subseteq X$, and then setting $F(X) = \text{colim}_i F(X_i)$ in the category $\mathbf{Set}^{\mathbb{A}}$. Since filtered colimits commute with finite products, these colimits taken in $\mathbf{Set}^{\mathbb{A}}$ and will remain in $\mathbf{Mod}(\mathbb{A})$.

Theorem 1.2.14. *For any set X with free algebra $F(X)$ as just defined, there is a natural isomorphism,*

$$\mathbf{Hom}_{\mathbf{Mod}(\mathbb{A})}(F(X), M) \cong \mathbf{Hom}_{\mathbf{Set}}(X, |M|). \quad (1.20)$$

Proof. The proof is an easy exercise. □

By definition, the finitely generated free models $F(n)$ are just the representables $\mathbf{Hom}_{\mathbb{A}}(A^n, -)$; therefore as the “semantic dual” $\mathbf{mod}(\mathbb{A}) \hookrightarrow \mathbf{Mod}(\mathbb{A})$ of the theory \mathbb{A} , in the sense of corollary 1.2.3, we simply have (the full subcategory of $\mathbf{Hom}_{\mathbf{FP}}(\mathbb{A}, \mathbf{Set})$ on) the image of the Yoneda embedding,

$$\begin{array}{ccccc}
 \mathbf{mod}(\mathbb{A}) & \hookrightarrow & \mathbf{Mod}(\mathbb{A}) = \mathbf{Hom}_{\mathbf{FP}}(\mathbb{A}, \mathbf{Set}) & \hookrightarrow & \mathbf{Set}^{\mathbb{A}} \\
 \uparrow \cong & & & & \uparrow \mathbf{y} \\
 \mathbb{A}^{\mathrm{op}} & \xrightarrow{\quad \quad \quad} & \mathbb{A}^{\mathrm{op}} & &
 \end{array}$$

In the abstract case, then, the logical duality

$$\mathbb{A} \simeq \mathbf{mod}(\mathbb{A})^{\mathrm{op}}$$

comes down to the fact that the (contravariant) Yoneda embedding

$$\mathbb{A}^{\mathrm{op}} \hookrightarrow \mathbf{Set}^{\mathbb{A}}$$

presents \mathbb{A} as (the dual of) a full subcategory of (product-preserving!) functors. Thus we have shown:

Theorem 1.2.15. *For any Lawvere algebraic theory \mathbb{A} , there is an equivalence,*

$$\mathbb{A} \simeq \mathbf{mod}(\mathbb{A})^{\mathrm{op}}$$

between \mathbb{A} and the full subcategory of finitely generated free models.

Exercise 1.2.16. Prove theorem 1.2.14.

1.2.3 Algebraic categories

Given an *arbitrary* category \mathcal{A} , we may ask: *When is \mathcal{A} the category of models for some algebraic theory?* Such categories are sometimes called *varieties*, at least in universal algebra, and there are well-known *recognition theorems* such the famous HSP-theorem of Birkhoff, which says that a class of interpretations for some fixed signature are exactly those satisfying a set of equations if the collection is closed under Products, Subalgebras, and Homomorphic images (i.e. quotient algebras). Toward the goal of “recognizing” a *category* of algebras (without even being given the signature!), let us define:

Definition 1.2.17. An *algebraic category* \mathcal{A} is a (locally small) category equivalent to one of the form

$$\mathbf{Hom}_{\mathbf{FP}}(\mathbb{A}, \mathbf{Set}) \hookrightarrow \mathbf{Set}^{\mathbb{A}}$$

where \mathbb{A} is any small finite product category and $\mathbf{Hom}_{\mathbf{FP}}(\mathbb{A}, \mathbf{Set})$ is the full subcategory of finite product preserving functors. If \mathbb{A} is a Lawvere algebraic theory (i.e. the objects are generated under finite products by a single object), then we will say that \mathcal{A} is a *Lawvere algebraic category*.

If \mathcal{A} is the category of models of a Lawvere algebraic theory \mathbb{A} , with generating object A , then in particular there will be a forgetful functor

$$U : \mathcal{A} \rightarrow \mathbf{Set},$$

determined by evaluation at A . It follows immediately that U preserves all limits, and one can show without difficulty that it also preserves all filtered colimits (cf. exercise 1.2.22). We require only one further condition to “recognize” \mathcal{A} as algebraic, namely creation of “ U -absolute coequalizers”.

Definition 1.2.18. In any category \mathcal{C} , a coequalizer $c : Y \rightarrow Z$ of maps $a, b : X \rightrightarrows Y$ is *absolute* if, for every category \mathcal{D} and functor $F : \mathcal{C} \rightarrow \mathcal{D}$, the map $Fc : FY \rightarrow FZ$ is a coequalizer of the maps $Fa, Fb : FX \rightrightarrows FY$. A functor $F : \mathcal{C} \rightarrow \mathcal{D}$ *creates F -absolute coequalizers* if for every parallel pair of maps $a, b : X \rightrightarrows Y$ in \mathcal{C} and absolute coequalizer $q : FY \rightarrow Q$ of $Fa, Fb : FX \rightrightarrows FY$ in \mathcal{D} , there is a unique object Z and map $c : Y \rightarrow Z$ in \mathcal{C} with $FZ = Q$ and $Fc = q$, which, moreover, is a coequalizer in \mathcal{C} .

$$\begin{array}{ccc} \mathcal{C} & X \xrightleftharpoons[b]{a} Y & \xrightarrow{\quad c \quad} Z \\ \downarrow F & & \\ \mathcal{D} & FX \xrightleftharpoons[Fb]{Fa} FY & \xrightarrow{\quad q \quad} Q \end{array} \quad (1.21)$$

Theorem 1.2.19. *Given a category \mathcal{A} equipped with a functor $U : \mathcal{A} \rightarrow \mathbf{Set}$, the following conditions are equivalent.*

1. \mathcal{A} is a Lawvere algebraic category; i.e. there is a Lawvere algebraic theory \mathbb{A} , and an equivalence,

$$\mathcal{A} \simeq \mathbf{Hom}_{\mathbf{FP}}(\mathbb{A}, \mathbf{Set}) \hookrightarrow \mathbf{Set}^{\mathbb{A}}$$

between \mathcal{A} and the full subcategory of finite product preserving functors on \mathbb{A} , associating $U : \mathcal{A} \rightarrow \mathbf{Set}$ to the evaluation at the generating object of \mathbb{A} .

2. $U : \mathcal{A} \rightarrow \mathbf{Set}$ has a left adjoint $F : \mathbf{Set} \rightarrow \mathcal{A}$, preserves all (small) filtered colimits, and creates U -absolute coequalizers.
3. \mathcal{A} is monadic over \mathbf{Set} (via $U : \mathcal{A} \rightarrow \mathbf{Set}$),

$$\mathcal{A} \simeq \mathbf{Set}^T$$

for a finitary monad $T : \mathbf{Set} \rightarrow \mathbf{Set}$.

Proof. (1 \Rightarrow 2) Suppose first that \mathcal{A} is (Lawvere) algebraic, so

$$\mathcal{A} \simeq \mathbf{Mod}(\mathbb{A}) = \mathbf{Hom}_{\mathbf{FP}}(\mathbb{A}, \mathbf{Set}) \hookrightarrow \mathbf{Set}^{\mathbb{A}}$$

for an algebraic theory \mathbb{A} . First, observe that U preserves all limits, since by theorem 1.2.14 it has a left adjoint $F : \mathbf{Set} \rightarrow \mathcal{A}$. It preserves filtered colimits because they commute with finite products, and so a filtered colimit of FP functors, calculated in $\mathbf{Set}^{\mathbb{A}}$, is again an FP functor. This suffices, since colimits are computed pointwise in $\mathbf{Set}^{\mathbb{A}}$ and U is an evaluation functor. Thus U preserves both limits and filtered colimits.

For creation of U -absolute coequalizers, suppose we have maps $f, g : A \rightrightarrows B$ in \mathcal{A} and an absolute coequalizer $c : UB \rightarrow C$ for $Uf, Ug : UA \rightrightarrows UB$ in \mathbf{Set} ; we want to put an algebra structure on C making c a homomorphism $c : B \rightarrow C$ in \mathcal{A} , and a coequalizer of f and g .

$$\begin{array}{ccccc}
 UA^n & \xrightarrow{Uf^n} & UB^n & \xrightarrow{c^n} & C^n \\
 \sigma^A \downarrow & \xrightarrow{Ug^n} & \downarrow \sigma^B & & \downarrow \sigma^C \\
 UA & \xrightarrow{Uf} & UB & \xrightarrow{c} & C \\
 & \xrightarrow{Ug} & & &
 \end{array} \tag{1.22}$$

For each function symbol $\sigma \in \Sigma$ we have commutative squares as on the left in the above diagram, because f and g are homomorphisms. It follows by a simple diagram chase that $c \circ \sigma^B$ coequalizes the pair $Uf^n, Ug^n : UA^n \rightrightarrows UB^n$. Since $c : UB \rightarrow C$ is absolute, it is preserved by the functor $(-)^n$, and therefore $c^n : UB^n \rightarrow C^n$ is a coequalizer of Uf^n, Ug^n . There is therefore a unique map $\sigma^C : C^n \rightarrow C$ as indicated, making the right hand square commute. Doing this for each $\sigma \in \Sigma$ gives an interpretation of Σ on C . This is seen to be an algebra structure because the maps c^n are surjections. Thus $c : B \rightarrow C$ is a homomorphism, which is easily seen to be a coequalizer in \mathcal{A} .

(2 \Rightarrow 3) Taking the standard monad (T, η, μ) on \mathbf{Set} with underlying functor $T = U \circ F$, we want to show that the canonical comparison map

$$\mathcal{A} \rightarrow \mathbf{Set}^T$$

to the category of T -algebras is an isomorphism. This follows from the condition that U creates absolute coequalizers by Beck's theorem; see [Lan71, VI.7]. Moreover, T preserves filtered colimits (i.e. is “finitary”) because each of F and U do so.

(3 \Rightarrow 1) Let (T, η, μ) be a finitary monad on \mathbf{Set} and $U : \mathbf{Set}^T \rightarrow \mathbf{Set}$ the forgetful functor from the category of T -algebras. We want an algebraic theory \mathbb{A} and an equivalence

$$\mathbf{Set}^T \simeq \mathbf{Mod}(\mathbb{A})$$

over U and evaluation at the generator of \mathbb{A} , where recall $\mathbf{Mod}(\mathbb{A}) = \mathbf{Hom}_{\mathbf{FP}}(\mathbb{A}, \mathbf{Set})$. Let

$$\mathbb{A} = \mathbf{FGF}(\mathbf{Set}^T)^{\text{op}} \tag{1.23}$$

be the dual of the full subcategory of finitely generated free T -algebras. The objects of \mathbb{A} are of the form T_0, T_1, T_2, \dots where $T_n = T(n)$, equipped with the multiplication $\mu_n : T^2(n) \rightarrow T(n)$ as algebra structure map. Since, as free algebras, $T(n+m) \cong T(n) + T(m)$ we indeed have $T_n \times T_m \cong T_{n+m}$ as objects of \mathbb{A} , and T_1 as the generating object.

By the first two steps of this proof, we know that the algebraic category $\mathbf{Mod}(\mathbb{A})$ is also (finitary) monadic,

$$\mathbf{Mod}(\mathbb{A}) \simeq \mathbf{Set}^M,$$

with monad $M = U_M \circ F_M$, where $F_M \dashv U_M$ is the free-forgetful adjunction for $\mathbf{Mod}(\mathbb{A}) = \mathbf{Hom}_{\mathbf{FP}}(\mathbb{A}, \mathbf{Set})$, and $U_M \cong \mathbf{eval}_{T_1}$. Thus it will suffice to show that $M \cong T$, as monads, in order to conclude that

$$\mathbf{Mod}(\mathbb{A}) \simeq \mathbf{Set}^M \simeq \mathbf{Set}^T.$$

Moreover, since both M and T are finitary, it suffices to show that their respective restrictions to the dense subcategory $\mathbf{Set}_{\text{fin}} \hookrightarrow \mathbf{Set}$ are isomorphic. By (1.19), we know that the finite free functor $F_M(n)$ has the form

$$F_M(n) = \mathbf{Hom}_{\mathbb{A}}(T_n, -) = \mathbf{Hom}_{\mathbf{FGF}(\mathbf{Set}^T)}(-, \langle T(n), \mu_n \rangle)$$

thus using the fact that $U_M \cong \mathbf{eval}_{T_1}$ we see that

$$\begin{aligned} M(n) &= U_M(F_M(n)) = U_M(\mathbf{Hom}_{\mathbf{FGF}(\mathbf{Set}^T)}(-, \langle T(n), \mu_n \rangle)) \\ &\cong \mathbf{Hom}_{\mathbf{FGF}(\mathbf{Set}^T)}(\langle T(1), \mu_1 \rangle, \langle T(n), \mu_n \rangle) \\ &\cong \mathbf{Hom}_{\mathbf{Set}}(1, T(n)) \cong T(n). \end{aligned}$$

□

Remark 1.2.20. Another “recognition theorem” that can be found in [Bor94] is the following:

Theorem (Borceux II.3.9). *Given a category \mathcal{A} , equipped with a functor $U : \mathcal{A} \rightarrow \mathbf{Set}$, the following conditions are equivalent.*

1. \mathcal{A} is equivalent to the category of models of some Lawvere algebraic theory \mathbb{T} ,

$$\mathcal{A} \simeq \mathbf{Mod}(\mathbb{T})$$

with $U : \mathcal{A} \rightarrow \mathbf{Set}$ the corresponding forgetful functor.

2. \mathcal{A} has coequalizers and kernel pairs, and $U : \mathcal{A} \rightarrow \mathbf{Set}$ has a left adjoint $F : \mathbf{Set} \rightarrow \mathcal{A}$, preserves all (small) filtered colimits and regular epimorphisms, and reflects isomorphisms.

Condition (1) is of course equivalent to condition (1) or our theorem 1.2.19, by theorem 1.2.1.

Exercise 1.2.21. A *split coequalizer* for maps $f, g : A \rightrightarrows B$ is a map $e : B \rightarrow C$ together with s and t as indicated below,

$$\begin{array}{ccccc} A & \xrightarrow{f} & B & \xrightarrow{e} & C \\ & \searrow g & \swarrow & & \\ & & & \swarrow s & \\ & & & \searrow t & \end{array} \quad (1.24)$$

satisfying the equations

$$ef = eg, \quad ft = 1_B, \quad gt = se, \quad es = 1_C.$$

Show that a split coequalizer is an absolute coequalizer.

Exercise 1.2.22. A filtered colimit of algebras can be described directly as follows: First consider the case of sets. Let the index category \mathbb{J} be filtered and $D : \mathbb{J} \rightarrow \mathbf{Set}$ a diagram. The colimiting set $\text{colim}_j D_j$ can be described as the quotient of the coproduct $(\coprod_j D_j)/\sim$, where the equivalence relation \sim is defined by:

$$(d_i \in D_i) \sim (d_j \in D_j) \Leftrightarrow t_{ik}(d_i) = t_{jk}(d_j) \text{ for some } t_{ik} : i \rightarrow k \text{ and } t_{jk} : j \rightarrow k \text{ in } \mathbb{J}.$$

1. Show that this is an equivalence relation using the filteredness of \mathbb{J} .
2. Now assume that the D_j all have an algebra structure and that all the transition maps $t_{ik} : D_i \rightarrow D_k$ are homomorphisms. Show that the colimit set $D_\infty = \text{colim}_j D_j$ is also an algebra of the same kind by defining each of the operations $\sigma_\infty : D_\infty \times \dots \times D_\infty \rightarrow D_\infty$ on equivalence classes as

$$\sigma_\infty \langle [d_i], \dots, [d'_j] \rangle = [\sigma_k \langle t_{ik}(d_i), \dots, t_{jk}(d'_j) \rangle]$$

for suitable k . Show that this is well-defined, and that D_∞ , so equipped, also satisfies the equations satisfied by the D_j .

Example 1.2.23. A field is a ring in which every non-zero element has a multiplicative inverse. The theory of fields is (apparently) not algebraic, because the axiom

$$x \neq 0 \Rightarrow \exists y(x \cdot y = 1)$$

is not simply an equation. But in principle there could be an equivalent algebraic formulation of the theory which would somehow circumvent this problem. We can show that this is not the case by proving that the category **Field** of fields and field homomorphisms is not algebraic.

First observe that a category of models $\mathbf{Mod}(\mathbb{A})$ always has a terminal object because **Set** has a terminal object **1**, and the constant functor $\Delta_1 : \mathbb{A} \rightarrow \mathbf{Set}$ which maps everything to **1** is a model. The functor Δ_1 is the terminal object in $\mathbf{Mod}(\mathbb{A})$ because it is the terminal functor in the functor category $\mathbf{Set}^{\mathbb{A}}$. In order to see that **Field** is not algebraic it thus suffices to show that there is no terminal field.

Exercise 1.2.24. Show that the category **Field** does not have a terminal object. (Hint: suppose that T is the terminal field and use the unique homomorphism $\mathbb{Z}_2 \rightarrow T$ to see that $1 + 1 = 0$ in T , then reason similarly using the unique homomorphism $\mathbb{Z}_3 \rightarrow T$.)

1.2.4 Algebraic functors

Classically, a *syntactic translation* of one algebraic theory into another is an assignment of types to types and terms to terms, respecting the tupling operations and substitutions of terms for variables. Such a translation can be described abstractly as a finite product preserving functor,

$$T : \mathbb{A} \rightarrow \mathbb{B}$$

between the associated Lawvere algebraic theories. Every such translation induces a “definable” functor on the semantics, just by precomposition:

$$T^*(M) = M \circ T.$$

$$\mathbf{Mod}(\mathbb{A}) \xleftarrow{T^*} \mathbf{Mod}(\mathbb{B}) \quad (1.25)$$

$$\begin{array}{ccc} \mathbb{A} & \xrightarrow{T} & \mathbb{B} \\ & \searrow T^*(M) & \downarrow M \\ & & \mathbf{Set} \end{array}$$

For instance, let $\mathbb{T}_0 = (\mathbf{Set}_{\text{fin}})^{\text{op}}$, the trivial theory, so that e.g, $\mathbf{Mod}(\mathbb{T}_0) = \mathbf{Set}$. Then for any algebraic theory \mathbb{A} , the generating object $A \in \mathbb{A}$ has a classifying functor

$$A : \mathbb{T}_0 \rightarrow \mathbb{A}$$

which induces the forgetful functor by precomposition,

$$\mathbf{Set} = \mathbf{Mod}(\mathbb{T}_0) \xleftarrow{A^*} \mathbf{Mod}(\mathbb{A})$$

$$\begin{array}{ccc} \mathbb{T}_0 & \xrightarrow{A} & \mathbb{A} \\ & \searrow A^*M & \downarrow M \\ & & \mathbf{Set} \end{array}$$

More generally, by the universal property of \mathbb{A} , every translation corresponds to a “model of \mathbb{A} in the syntax of \mathbb{B} ”:

$$\frac{T : \mathbb{A} \rightarrow \mathbb{B}}{\hat{T} \in \mathbf{Mod}(\mathbb{A}, \mathbb{B})}$$

For instance, since every ring R has an underlying group $\mathbf{Grp}(R)$, the universal ring $U_{\mathbb{R}}$ in the theory of rings \mathbb{R} has one $\mathbf{Grp}(U_{\mathbb{R}})$, which is therefore classified by a unique functor from the theory of groups,

$$\mathbf{Grp}(U_{\mathbb{R}}) : \mathbb{G} \rightarrow \mathbb{R}.$$

This translation induces a functor on the corresponding categories of models,

$$\mathbf{Grp}(U_{\mathbb{R}})^* : \mathbf{Group} \leftarrow \mathbf{Ring}, \quad (1.26)$$

which of course is just the forgetful functor $\mathbf{Grp} : \mathbf{Ring} \rightarrow \mathbf{Group}$.

We can now ask: *Which functors $f : \mathbf{Mod}(\mathbb{B}) \rightarrow \mathbf{Mod}(\mathbb{A})$ between algebraic categories are of the form $f = T^*$ for a translation $T : \mathbb{A} \rightarrow \mathbb{B}$ of theories?* Let us call these *algebraic functors*. We consider first the case where T takes the generator A of \mathbb{A} to the generator B of \mathbb{B} ,

$$T(A) \cong B.$$

Then T^* commutes with the forgetful functors, which, recall, are evaluation at the generators, $U_{\mathbb{A}}(M) = M(A)$:

$$\begin{array}{ccc} \mathbf{Mod}(\mathbb{B}) & \xrightarrow{T^*} & \mathbf{Mod}(\mathbb{A}) \\ & \searrow U_{\mathbb{B}} \quad \swarrow U_{\mathbb{A}} & \\ & \mathbf{Set} & \end{array}$$

This is simply because

$$(U_{\mathbb{A}} \circ T^*)(M) = U_{\mathbb{A}}(M \circ T) = (M \circ T)(A) \cong M(T(A)) \cong M(B) = U_{\mathbb{B}}(M).$$

In fact, this condition is already sufficient!

Proposition 1.2.25. *For Lawvere algebraic theories \mathbb{A} and \mathbb{B} , given any functor $f : \mathbf{Mod}(\mathbb{B}) \rightarrow \mathbf{Mod}(\mathbb{A})$ with*

$$U_{\mathbb{B}} \cong U_{\mathbb{A}} \circ f,$$

there is a unique (up to iso) FP-functor $T : \mathbb{A} \rightarrow \mathbb{B}$ such that

$$f \cong T^*.$$

Proof. (Cf. [Bor94, 3.9.2]) Consider the diagram:

$$\begin{array}{ccc} \mathbf{Mod}(\mathbb{B}) & \xrightarrow{f} & \mathbf{Mod}(\mathbb{A}) \\ & \searrow U_{\mathbb{B}} \quad \swarrow U_{\mathbb{A}} & \\ & \mathbf{Set} & \\ & \nwarrow F_{\mathbb{B}} \quad \nearrow F_{\mathbb{A}} & \end{array}$$

Using the explicit description of \mathbb{A} as the opposite of $\mathbf{mod}_{fg}(\mathbb{A})$, the subcategory of finitely generated free models, (and the same for \mathbb{B}), we therefore seek a finite coproduct preserving functor $f_!$ as indicated in the following

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Now, to define T , take any $\alpha : F_{\mathbb{A}}(n) \rightarrow F_{\mathbb{A}}(m)$ and use

$$\begin{array}{ccc} & & fF_{\mathbb{B}}(m) \\ & \nearrow h_m \alpha & \uparrow h_m \\ F_{\mathbb{A}}(n) & \xrightarrow{\alpha} & F_{\mathbb{A}}(m) \end{array}$$

to get the lift:

$$T(\alpha) := \overline{h_m \alpha} : F_{\mathbb{B}}(n) \longrightarrow F_{\mathbb{B}}(m).$$

This assignment is easily seen to be functorial. Finally, to show that it is induced by pre-composing with $T : \mathbb{A} \rightarrow \mathbb{B}$, we need the following to commute up to natural isomorphism.

$$\begin{array}{ccc} \mathbf{Mod}(\mathbb{B}) & \xrightarrow{f} & \mathbf{Mod}(\mathbb{A}) \\ \cong \downarrow & & \downarrow \cong \\ \mathbf{FP}(\mathbb{B}, \mathbf{Set}) & \xrightarrow{T^*} & \mathbf{FP}(\mathbb{A}, \mathbf{Set}) \end{array}$$

Identifying $\mathbb{B} = \mathbf{Mod}_{\mathbf{fg}}(\mathbb{B})^{\text{op}}$ and $\mathbb{A} = \mathbf{Mod}_{\mathbf{fg}}(\mathbb{A})^{\text{op}}$, going around the Southwest, we have

$$M \mapsto \mathbf{Hom}_{\mathbf{Mod}(\mathbb{B})}(-, M) \mapsto \mathbf{Hom}_{\mathbf{Mod}(\mathbb{A})}(T(-), M).$$

Going around the Northeast gives

$$M \mapsto fM \mapsto \mathbf{Hom}(-, fM).$$

But by (1.27), we have a natural bijection:

$$\frac{F_{\mathbb{A}}(n) \rightarrow fM}{T(F_{\mathbb{A}}(n)) = F_{\mathbb{B}}(n) \rightarrow M}$$

□

Corollary 1.2.26. *For a functor $f : \mathbf{Mod}(\mathbb{B}) \rightarrow \mathbf{Mod}(\mathbb{A})$ between algebraic categories, the following are equivalent.*

1. *f is algebraic: $f = T^*$, for some FP functor $T : \mathbb{A} \rightarrow \mathbb{B}$ which, moreover, preserves the generator, $T(A) \cong B$*
2. *f commutes with the forgetful functors, $U_{\mathbb{A}} \circ f \cong U_{\mathbb{B}}$*

□

Exercise 1.2.27. Show that for any algebraic theory \mathbb{A} , the full inclusion $\mathbf{Mod}(\mathbb{A}) \hookrightarrow \mathbf{Set}^{\mathbb{A}}$ has a left adjoint. (Hint: use the Adjoint Functor Theorem.)

Exercise 1.2.28. Assuming the result of the previous exercise, show that the precomposition functor $T^* : \mathbf{Mod}(\mathbb{B}) \rightarrow \mathbf{Mod}(\mathbb{A})$ induced by *any* translation $T : \mathbb{A} \rightarrow \mathbb{B}$ (not necessarily preserving the generator) always has a left adjoint $T_! : \mathbf{Mod}(\mathbb{A}) \rightarrow \mathbf{Mod}(\mathbb{B})$.

Exercise 1.2.29. Assuming the results of the previous two exercises, show that a functor $f : \mathbf{Mod}(\mathbb{B}) \rightarrow \mathbf{Mod}(\mathbb{A})$ is induced by a translation $T : \mathbb{A} \rightarrow \mathbb{B}$ as $f = T^*$ iff the following conditions hold:

1. f has a left adjoint $f_! : \mathbf{Mod}(\mathbb{A}) \rightarrow \mathbf{Mod}(\mathbb{B})$.
2. $f_!(U_{\mathbb{A}}) \cong (U_{\mathbb{B}})^n$ for some $0 \leq n$, where as usual we identify $\mathbb{A} = \mathbf{mod}_{\mathbf{fg}}(\mathbb{A})^{\mathbf{op}}$ and similarly for \mathbb{B} .

Remark 1.2.30. The foregoing duality can be generalized to many-sorted algebraic theories \mathbb{A} *not* assumed to be generated under finite products by a single object. The corresponding semantic category $\mathbf{Mod}(\mathbb{A})$ still consists of FP-functors $\mathbb{A} \rightarrow \mathbf{Set}$. The question, when is a “semantic functor”

$$f : \mathbf{Mod}(\mathbb{B}) \rightarrow \mathbf{Mod}(\mathbb{A})$$

between such algebraic categories induced by a “syntactic functor” $T : \mathbb{A} \rightarrow \mathbb{B}$ of such algebraic theories, i.e. one that preserves finite products, can also be answered in the abstract, providing a definition of an *algebraic functor*: it is one that preserves all limits, filtered colimits, and regular epimorphisms. This characterization requires a modification in the notion of algebraic theory, however, to include closure under retracts — so-called *Cauchy completeness*. In that setting, there is a neat duality theory between the “syntax category” of all theories and translations and the “semantics category” of all algebraic categories and algebraic functors. This is developed in detail in [ALR03].

Appendix A

Logic

A.1 Concrete and abstract syntax

By *syntax* we generally mean manipulation of finite strings of symbols according to given *grammatical rules*. For instance, the strings “ $7)6 + /(8$ ” and “ $(6 + 8)/7$ ” both consist of the same symbols but you will recognize one as junk and the other as *well formed* because you have (implicitly) applied the grammatical rules for arithmetical expressions.

Grammatical rules are usually quite complicated, as they need to prescribe associativity of operators (does “ $5 + 6 + 7$ ” mean “ $(5 + 6) + 7$ ” or “ $5 + (6 + 7)$ ”?) and their precedence (does “ $6 + 8/7$ ” mean “ $(6 + 8)/7$ ” or “ $6 + (8/7)$ ”?), the role of *white space* (empty space between symbols and line breaks), rules for nesting and balancing parentheses, etc. It is not our intention to dwell on such details, but rather to focus on the mathematical nature of well-formed expressions, namely that they represent inductively generated finite trees.¹ Under this view the string “ $(6 + 8)/7$ ” is just a concrete representation of the tree depicted in Figure A.1.

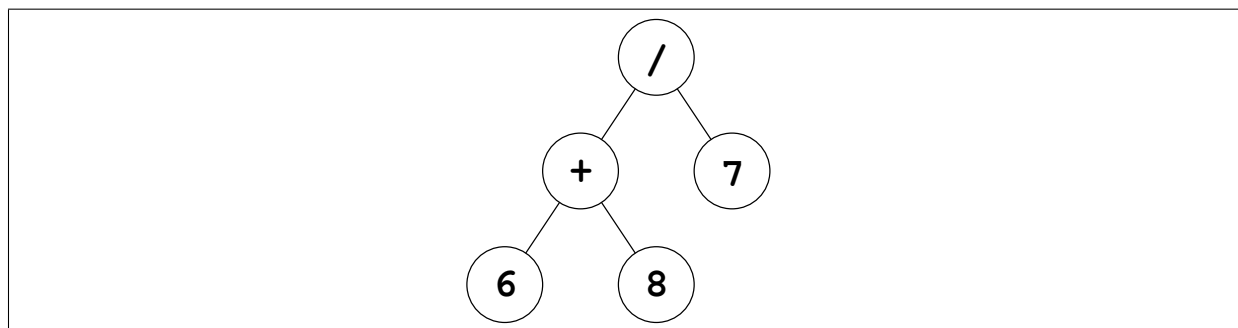


Figure A.1: The tree represented by $(6 + 8)/7$

Concrete representation of expressions as finite strings of symbols is called *concrete syntax*, while in *abstract syntax* we view expressions as finite trees. The passage from the

¹We are limiting attention to the so-called *context-free* grammar, which are sufficient for our purposes. More complicated grammars are rarely used to describe formal languages in logic and computer science.

former to the latter is called *parsing* and is beyond the scope of this book. We will always specify only abstract syntax and assume that the corresponding concrete syntax follows the customary rules for parentheses, associativity and precedence of operators.

As an illustration we give rules for the (abstract) syntax of propositional calculus in *Backus-Naur* form:

Propositional variable $p ::= p_1 \mid p_2 \mid p_3 \mid \dots$

Propositional formula $\phi ::= p \mid \perp \mid \top \mid \phi_1 \wedge \phi_2 \mid \phi_1 \vee \phi_2 \mid \phi_1 \Rightarrow \phi_2 \mid \neg\phi$

The vertical bars should be read as “or”. The first rule says that a propositional variable is the constant p_1 , or the constant p_2 , or the constant p_3 , etc.² The second rule tells us that there are seven inductive rules for building a propositional formula:

- a propositional variable is a formula,
- the constants \perp and \top are formulas,
- if ϕ_1 , ϕ_2 , and ϕ are formulas, then so are $\phi_1 \wedge \phi_2$, $\phi_1 \vee \phi_2$, $\phi_1 \Rightarrow \phi_2$, and $\neg\phi$.

Even though abstract syntax rules say nothing about parentheses or operator associativity and precedence, we shall rely on established conventions for mathematical notation and write down concrete representations of propositional formulas, e.g., $p_4 \wedge (p_1 \vee \neg p_1) \wedge p_4 \vee p_2$.

A word of warning: operator associativity in syntax is not to be confused with the usual notion of associativity in mathematics. We say that an operator \star is *left associative* when an expression $x \star y \star z$ represents the left-hand tree in Figure A.2, and *right associative* when it represents the right-hand tree. Thus the usual operation of subtraction $-$ is left

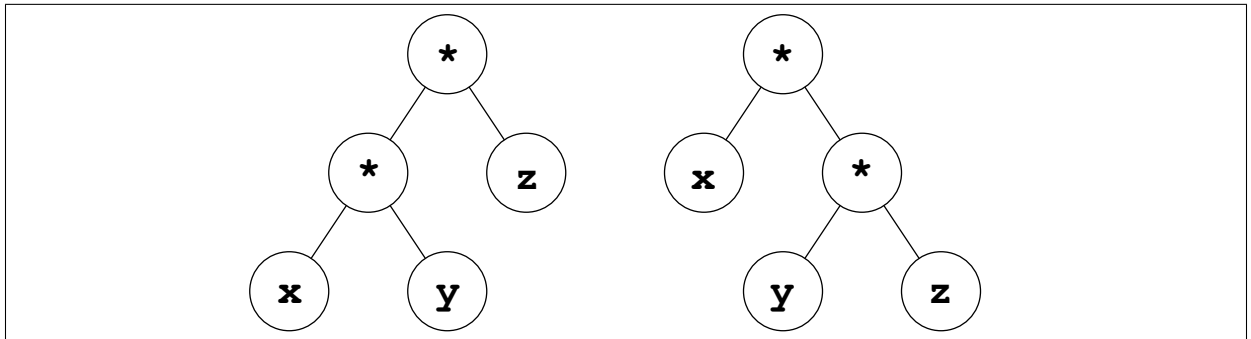


Figure A.2: Left and right associativity of $x \star y \star z$

associative, but is not associative in the usual mathematical sense.

²In an actual computer implementation we would allow arbitrary finite strings of letters as propositional variables. In logic we only care about the fact that we can never run out of fresh variables, i.e., that there are countably infinitely many of them.

A.2 Free and bound variables

Variables appearing in an expression may be *free* or *bound*. For example, in expressions

$$\int_0^1 \sin(a \cdot x) dx, \quad x \mapsto ax^2 + bx + c, \quad \forall x. (x < a \vee x > b)$$

the variables a , b and c are free, while x is bound by the integral operator \int , the function formation \mapsto , and the universal quantifier \forall , respectively. To be quite precise, it is an *occurrence* of a variable that is free or bound. For example, in expression $\phi(x) \vee \exists x. A\psi(x, x)$ the first occurrence of x is free and the remaining ones are bound.

In this book the following operators bind variables:

- quantifiers \exists and \forall , cf. ??,
- λ -abstraction, cf. ??,
- search for others ??.

When a variable is bound we may always rename it, provided the renaming does not confuse it with another variable. In the integral above we could rename x to y , but not to a because the binding operation would *capture* the free variable a to produce the unintended $\int_0^1 \sin(a^2) da$. Renaming of bound variables is called *α -renaming*.

We consider two expressions *equal* if they only differ in the names of bound variables, i.e., if one can be obtained from the other by α -renaming. Furthermore, we adhere to *Barendregt's variable convention* [?, p. 2], which says that bound variables are always chosen so as to differ from free variables. Thus we would never write $\phi(x) \vee \exists x. A\psi(x, x)$ but rather $\phi(x) \vee \exists y. A\psi(y, y)$. By doing so we need not worry about capturing or otherwise confusing free and bound variables.

In logic we need to be more careful about variables than is customary in traditional mathematics. Specifically, we always specify which free variables may appear in an expression.³ We write

$$x_1 : A_1, \dots, x_n : A_n \mid t$$

to indicate that expression t may contain only free variables x_1, \dots, x_n of types A_1, \dots, A_n . The list

$$x_1 : A_1, \dots, x_n : A_n$$

is called a *context* in which t appears. To see why this is important consider the different meaning that the expression $x^2 + y^2 \leq 1$ receives in different contexts:

- $x : \mathbb{Z}, y : \mathbb{Z} \mid x^2 + y^2 \leq 1$ denotes the set of tuples $\{(-1, 0), (0, 1), (1, 0), (0, -1)\}$,
- $x : \mathbb{R}, y : \mathbb{R} \mid x^2 + y^2 \leq 1$ denotes the closed unit disc in the plane, and

³This is akin to one of the guiding principles of good programming language design, namely, that all variables should be *declared* before they are used.

- $x : \mathbb{R}, y : \mathbb{R}, z : \mathbb{R} \mid x^2 + y^2 \leq 1$ denotes the infinite cylinder in space whose base is the closed unit disc.

In single-sorted theories there is only one type or sort A . In this case we abbreviate a context by listing just the variables, x_1, \dots, x_n .

A.3 Substitution

Substitution is a basic syntactic operation which replaces (free occurrences of) distinct variables x_1, \dots, x_n in an expression t with expressions t_1, \dots, t_n , which is written as

$$t[t_1/x_1, \dots, t_n/x_n].$$

We sometimes abbreviate this as $t[\vec{t}/\vec{x}]$ where $\vec{x} = (x_1, \dots, x_n)$ and $\vec{t} = (t_1, \dots, t_n)$. Here are several examples:

$$\begin{aligned} (x^2 + x + y)[(2 + 3)/x] &= (2 + 3)^2 + (2 + 3) + y \\ (x^2 + y)[y/x, x/y] &= y^2 + x \\ (\forall x. (x^2 < y + x^3)) [x + y/y] &= \forall z. (z^2 < (x + y) + z^3). \end{aligned}$$

Notice that in the third example we first renamed the bound variable x to z in order to avoid a capture by \forall .

Substitution is simple to explain in terms of trees. Assuming Barendregt's convention, the substitution $t[u/x]$ means that in the tree t we replace the leaves labeled x by copies of the tree u . Thus a substitution never changes the structure of the tree—it only “grows” new subtrees in places where the substituted variables occur as leaves.

Substitution satisfies the distributive law

$$(t[u/x])[v/y] = (t[v/y])[u[v/y]/x],$$

provided x and y are distinct variables. There is also a corresponding multivariate version which is written the same way with a slight abuse of vector notation:

$$(t[\vec{u}/\vec{x}])[\vec{v}/\vec{y}] = (t[\vec{v}/\vec{y}])[\vec{u}[\vec{v}/\vec{y}]/\vec{x}].$$

A.4 Judgments and deductive systems

A formal system, such as first-order logic or type theory, concerns itself with *judgments*. There are many kinds of judgments, such as:

- The most common judgments are equations and other logical statements. We distinguish a formula ϕ and the judgment “ ϕ holds” by writing the latter as

$$\vdash \phi.$$

The symbol \vdash is generally used to indicate judgments.

- Typing judgments

$$\vdash t : A$$

expressing the fact that a term t has type A . This is not to be confused with the set-theoretic statement $t \in u$ which says that individuals t and u (of type “set”) are in relation “element of” \in .

- Judgments expressing the fact that a certain entity is well formed. A typical example is a judgment

$$\vdash x_1 : A_1, \dots, x_n : A_n \quad \text{ctx}$$

which states that $x_1 : A_1, \dots, x_n : A_n$ is a well-formed context. This means that x_1, \dots, x_n are distinct variables and that A_1, \dots, A_n are well-formed types. This kind of judgement is often omitted and it is tacitly assumed that whatever entities we deal with are in fact well-formed.

A *hypothetical judgement* has the form

$$H_1, \dots, H_n \vdash C$$

and means that hypotheses H_1, \dots, H_n entail consequence C (with respect to a given deductive system). We may also add a typing context to get a general form of judgment

$$x_1 : A_1, \dots, x_n : A_n \mid H_1, \dots, H_m \vdash C.$$

This should be read as: “if x_1, \dots, x_n are variables of types A_1, \dots, A_n , respectively, then hypotheses H_1, \dots, H_m entail conclusion C .” For our purposes such contexts will suffice, but you should not be surprised to see other kinds of judgments in logic.

A *deductive system* is a set of inference rules for deriving judgments. A typical inference rule has the form

$$\frac{J_1 \quad J_2 \quad \dots \quad J_n}{J} C$$

This means that we can infer judgment J if we have already derived judgments J_1, \dots, J_n , provided that the optional side-condition C is satisfied. An *axiom* is an inference rule of the form

$$\overline{J}$$

A *two-way rule*

$$\frac{J_1 \quad J_2 \quad \dots \quad J_n}{K_1 \quad K_2 \quad \dots \quad K_m}$$

is a combination of $n + m$ inference rules stating that we may infer each K_i from J_1, \dots, J_n and each J_i from K_1, \dots, K_m .

A *derivation* of a judgment J is a finite tree whose root is J , the nodes are inference rules, and the leaves are axioms. An example is presented in the next subsection.

The set of all judgments that hold in a given deductive system is generated inductively by starting with the axioms and applying inference rules.

A.5 Example: Equational reasoning

Equational reasoning is so straightforward that one almost doesn't notice it, consisting mainly, as it does, of “substituting equals for equals”. The only judgements are equations between terms, $s = t$, which consist of function symbols, constants, and variables. The inference rules are just the usual ones making $s = t$ a congruence relation on the terms. More formally, we have the following specification of what may be called the *equational calculus*.

$$\begin{aligned} \text{Variable } v &::= x \mid y \mid z \mid \dots \\ \text{Constant symbol } c &::= c_1 \mid c_2 \mid \dots \\ \text{Function symbol } f^k &::= f_1^{k_1} \mid f_2^{k_2} \mid \dots \\ \text{Term } t &::= v \mid c \mid f^k(t_1, \dots, t_k) \end{aligned}$$

The superscript on the function symbol f^k indicates the arity.

The equational calculus has just one form of judgement

$$x_1, \dots, x_n \mid t_1 = t_2$$

where x_1, \dots, x_n is a *context* consisting of distinct variables, and the variables in the equation must occur among the ones listed in the context.

There are four inference rules for the equational calculus. They may be assumed to leave the contexts unchanged, which may therefore be omitted.

$$\begin{array}{cccc} \frac{}{t = t} & \frac{t_1 = t_2}{t_2 = t_1} & \frac{t_1 = t_2, t_2 = t_3}{t_1 = t_3} & \frac{t_1 = t_2, t_3 = t_4}{t_1[t_3/x] = t_2[t_4/x]} \end{array}$$

An *equational theory* \mathbb{T} consists of a set of constant and function symbols (with arities), and a set of equations, called *axioms*.

A.6 Example: Predicate calculus

We spell out the details of single-sorted predicate calculus and first-order theories. This is the most common deductive system taught in classical courses on logic.

The predicate calculus has the following syntax:

$$\begin{aligned} \text{Variable } v &::= x \mid y \mid z \mid \dots \\ \text{Constant symbol } c &::= c_1 \mid c_2 \mid \dots \\ \text{Function symbol}^4 f^k &::= f_1^{k_1} \mid f_2^{k_2} \mid \dots \\ \text{Term } t &::= v \mid c \mid f^k(t_1, \dots, t_k) \\ \text{Relation symbol } R^m &::= R_1^{m_1} \mid R_2^{m_2} \mid \dots \\ \text{Formula } \phi &::= \perp \mid \top \mid R^m(t_1, \dots, t_m) \mid t_1 = t_2 \mid \\ &\quad \phi_1 \wedge \phi_2 \mid \phi_1 \vee \phi_2 \mid \phi_1 \Rightarrow \phi_2 \mid \neg \phi \mid \forall x. \phi \mid \exists x. \phi. \end{aligned}$$

The variable x is bound in $\forall x . \phi$ and $\exists x . \phi$.

The predicate calculus has one form of judgement

$$x_1, \dots, x_n \mid \phi_1, \dots, \phi_m \vdash \phi$$

where x_1, \dots, x_n is a *context* consisting of distinct variables, ϕ_1, \dots, ϕ_m are *hypotheses* and ϕ is the *conclusion*. The free variables in the hypotheses and the conclusion must occur among the ones listed in the context. We abbreviate the context with Γ and Φ with hypotheses. Because most rules leave the context unchanged, we omit the context unless something interesting happens with it.

The following inference rules are given in the form of adjunctions. See Appendix ?? for the more usual formulation in terms of introduction and elimination rules.

$$\begin{array}{c}
 \overline{\phi_1, \dots, \phi_m \vdash \phi_i} \qquad \overline{\Phi \vdash \top} \qquad \overline{\Phi, \perp \vdash \phi} \\
 \\
 \frac{\Phi \vdash \phi_1 \quad \Phi \vdash \phi_2}{\Phi \vdash \phi_1 \wedge \phi_2} \qquad \frac{\Phi, \phi_1 \vdash \psi \quad \Phi, \phi_2 \vdash \psi}{\Phi, \phi_1 \vee \phi_2 \vdash \psi} \qquad \frac{\Phi, \phi_1 \vdash \phi_2}{\Phi \vdash \phi_1 \Rightarrow \phi_2} \\
 \\
 \frac{\Gamma, x, y \mid \Phi, x = y \vdash \phi}{\Gamma, x \mid \Phi \vdash \phi[x/y]} \qquad \frac{\Gamma, x \mid \Phi, \phi \vdash \psi}{\Gamma \mid \Phi, \exists x . \phi \vdash \psi} \qquad \frac{\Gamma, x \mid \Phi \vdash \phi}{\Gamma \mid \Phi \vdash \forall x . \phi}
 \end{array}$$

The equality rule implicitly requires that y does not appear in Φ , and the quantifier rules implicitly require that x does not occur freely in Φ and ψ because the judgments below the lines are supposed to be well formed.

Negation $\neg\phi$ is defined to be $\phi \Rightarrow \perp$. To obtain *classical* logic we also need the law of excluded middle,

$$\overline{\Phi \vdash \phi \vee \neg\phi}$$

Comment on the fact that contraction and weakening are admissible.

Give an example of a derivation.

A *first-order theory* \mathbb{T} consists of a set of constant, function and relation symbols with corresponding arities, and a set of formulas, called *axioms*.

Give examples of a first-order theories.

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