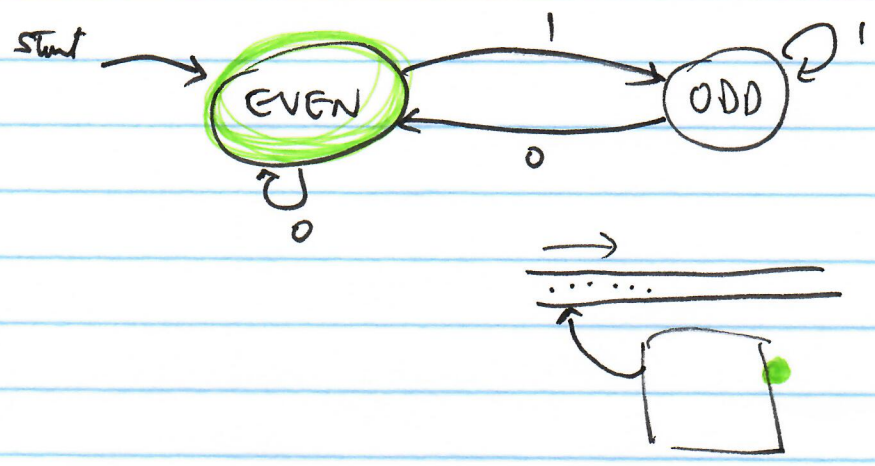
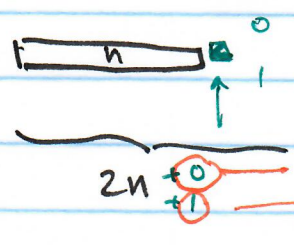


FSM Even integers in binary (L & R)

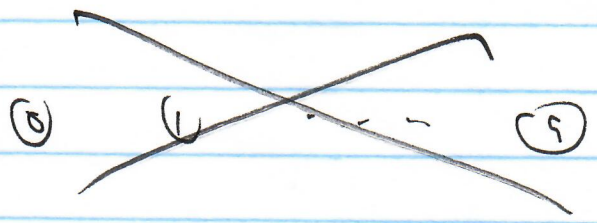
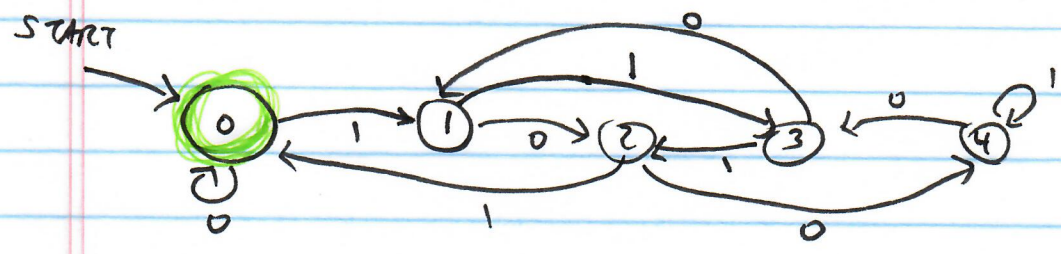


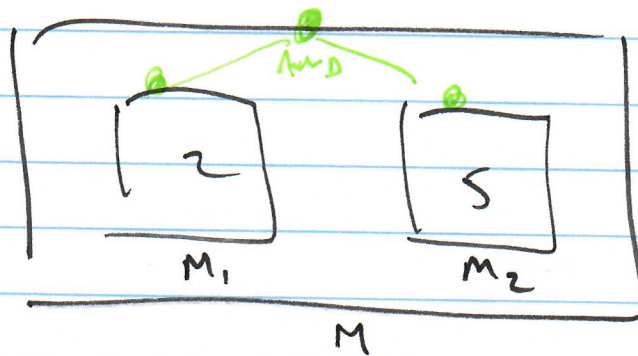
Multiplies of 5 in binary



$n \bmod 5 =$ remainder when n is divided by 5

remainder doubles
doubles, +1





F.S.M. $(\Sigma', S, \lambda_0, F, \delta)$

$$\begin{cases} M_1 = (\{0,1\}, S, \lambda_0, F_1, \delta_1) \\ M_2 = (\{0,1\}, T, t_0, F_2, \delta_2) \end{cases}$$

$L(M)$ = set of strings "accepted" by FSM M

$\Rightarrow \underline{M} = (\{0,1\},$

states $S \times T$ set of pairs (a,b) $a \in S$
 $b \in T$

start (λ_0, t_0)

final $F_1 \times F_2$ (a,b) $a \in F_1$
 $b \in F_2$

transition $\delta: S \times T \times \Sigma' \rightarrow S \times T$

$$\delta \begin{matrix} \in S & \in T \\ (a, b, x) \end{matrix} = (\delta_1(a, x), \delta_2(b, x))$$

$L(M) = L(M_1) \cap L(M_2)$

CLOSURE

$\cup (F_1 \times S \cup T \times F_2)$
 OOPS!

A language recognized by a FSM is called REGULAR

Then Even integers in binary are regular

Then Multiple of 5 in binary are regular

Can Multiples of 10 in binary are regular

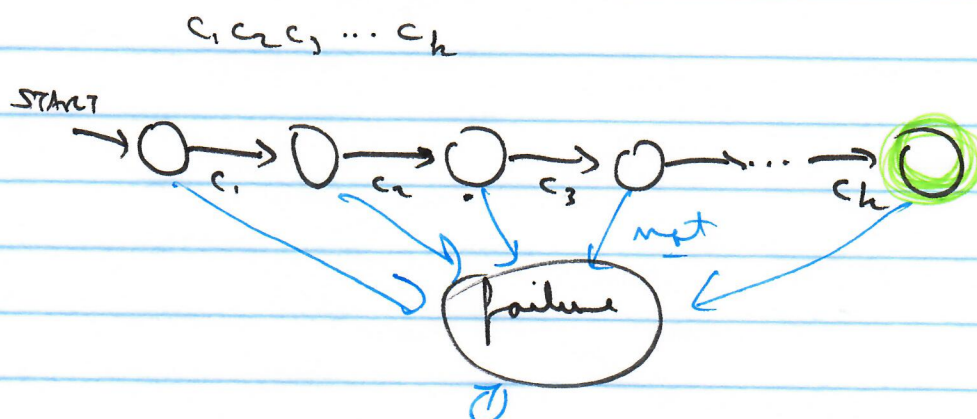
Then If L_1 and L_2 are regular, then
 $L_1 \cap L_2$ is regular

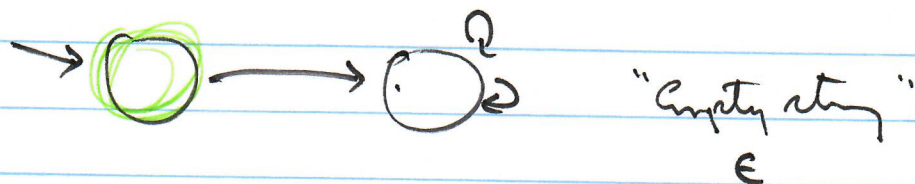
$L_1 \cup L_2$ is regular

closure

Then Any finite set is regular

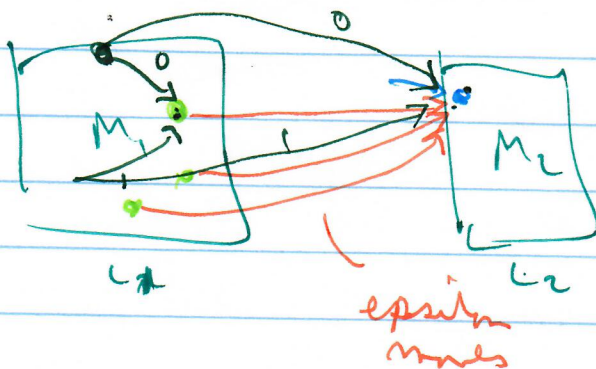
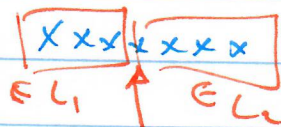
Then If L contains one string, L is regular





L_1, L_2 are regular

$L_1 L_2$ is regular
 $\{xy \mid x \in L_1, y \in L_2\}$



$$P = NP$$