

# Heuristics

CS3243: Introduction to Artificial Intelligence – Lecture 4



**1**

# Administrative Matters

# Midterm Assessment

- Week 7 Lecture Slot
  - Date: 2 October (Wednesday)
  - Time: 1005-1135 hrs
  - Venue: LT16 (COM2)
- Format
  - Duration: 1 hour and 30 minutes
  - Total: 30 marks
  - Closed-book + Cheat Sheet (1 × Double-sided A4 Sheet)
- Topics
  - Everything up to Lecture 5a (i.e., up to and including Local Search)

Practice Papers:  
Canvas > CS3243 > Files >  
Past Papers > Midterm

# Upcoming...

- **Deadlines**

- **Tutorial Assignment 2** (released last week)
  - Post-tutorial submission: **Due this Friday!**
- **Tutorial Assignment 3** (released today)
  - Pre-tutorial Submission: **Due this Sunday!**
  - Post-tutorial Submission: **Due next Friday!**
- **Project 1.1** (released last Monday – 26 August)
  - Due next Sunday! **15 September 2024**

Project 1 Consultations:  
Thursdays via Zoom  
[Canvas > Announcements](#)

Remember that there are  
attendance marks for tutorials!

## Late Penalties

- < deadline +24 hours = 80% of score
- < deadline +48 hours = 50% of score
- $\geq$  deadline +48 hours = 0% of score

# Contents

- Reviewing Informed Search
- Heuristics & Dominance
- Relaxing the Problem
- Reviewing the Midterm

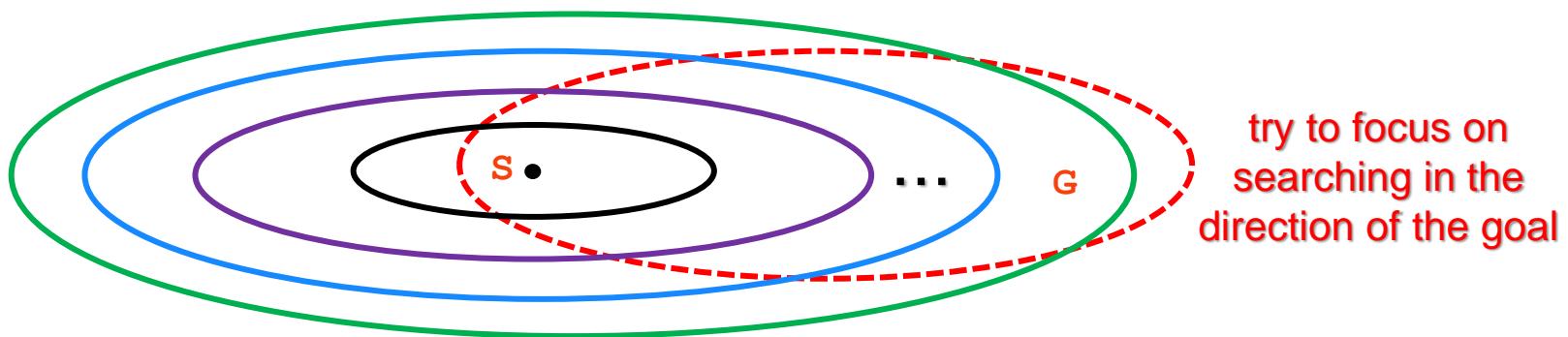
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## Reviewing Informed Search

# Searching Less?

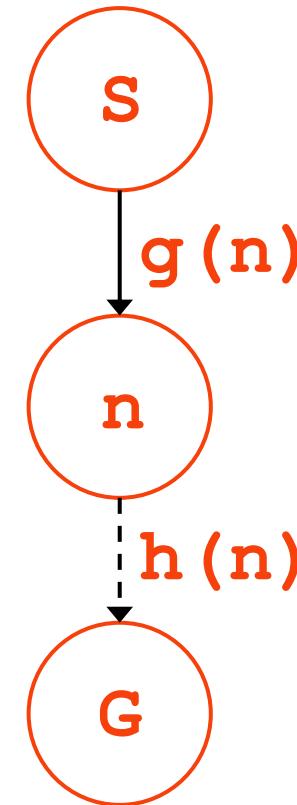
- Uninformed search algorithms are systematic
  - Search outward from the initial state
  - Search in ALL directions
- Search spaces are LARGE
- Can we search more efficiently?

**General Idea**  
Use domain knowledge about the problem environment to determine the cost required to go from a particular state to its nearest goal  
→ informed search



# Approximate Path Costs to Guide Search

- Frontier → Priority Queue
  - Ordered by evaluation function  $f(n)$ 
    - Priority for path corresponding to node  $n$
  - Uniform Cost Search priority:  $f(n) = g(n)$
  - Greedy Best-First Search priority:  $f(n) = h(n)$
  - A\* Search priority:  $f(n) = g(n) + h(n)$
- General idea
  - use domain knowledge design heuristic function  $h$  to estimate the cost from  $n.state$  to  $G$



$g(n)$  : Actual path cost from  $S$  to  $n$  based on the specific path taken by node  $n$

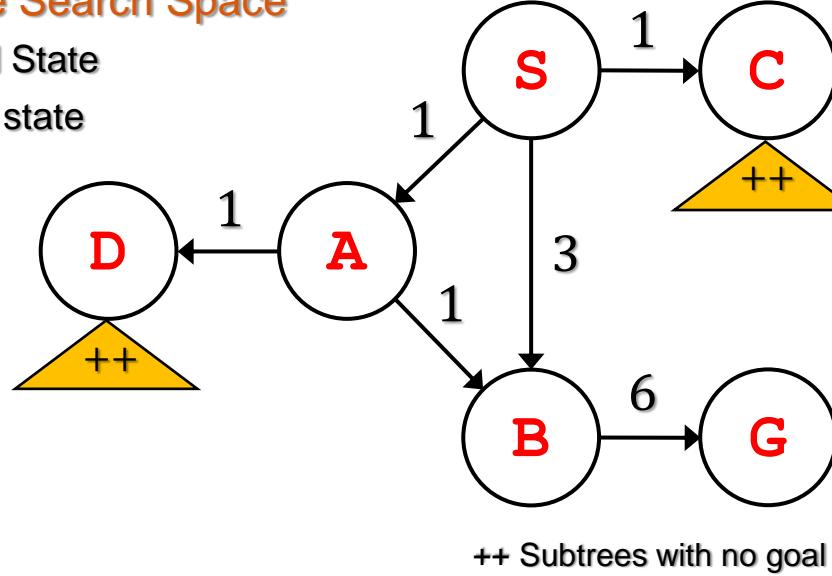
$h(n)$  : Approximate path cost from  $n$  to its nearest goal given by heuristic function  $h$

# Translating a Search Graph to Search Tree

- Example Search Space

**S:** Initial State

**G:** Goal state

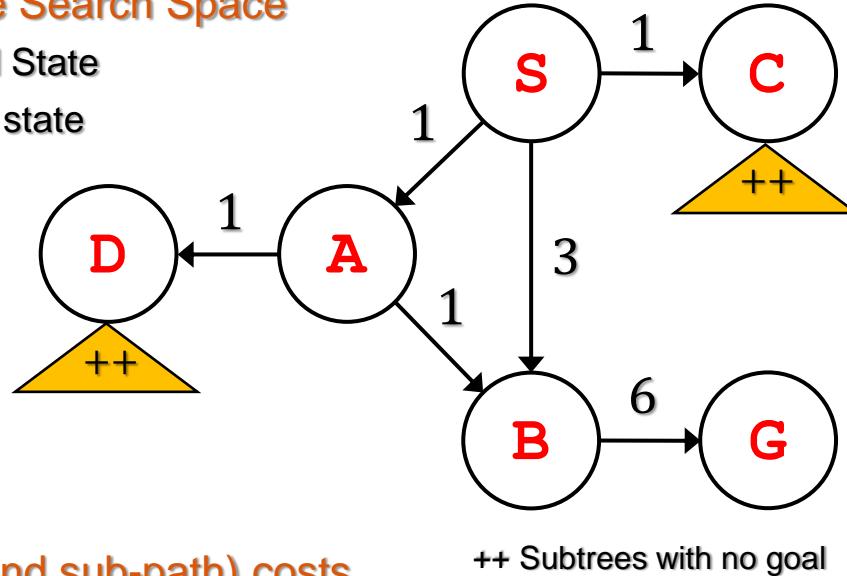


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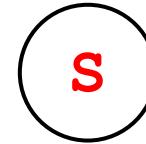
S: Initial State

G: Goal state



- Paths (and sub-path) costs

- Corresponding Search Tree

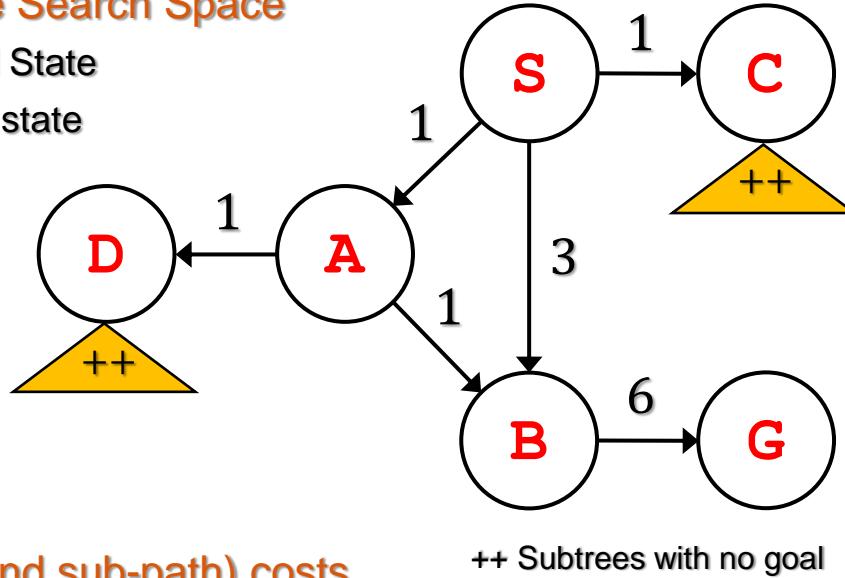


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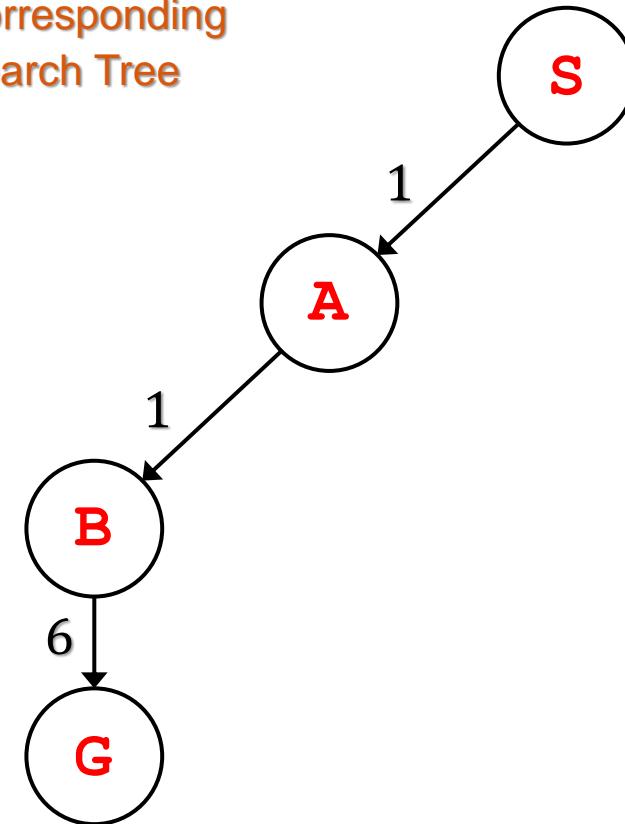
G: Goal state



- Paths (and sub-path) costs

- S-A-B-G : 8 units to goal

- Corresponding Search Tree

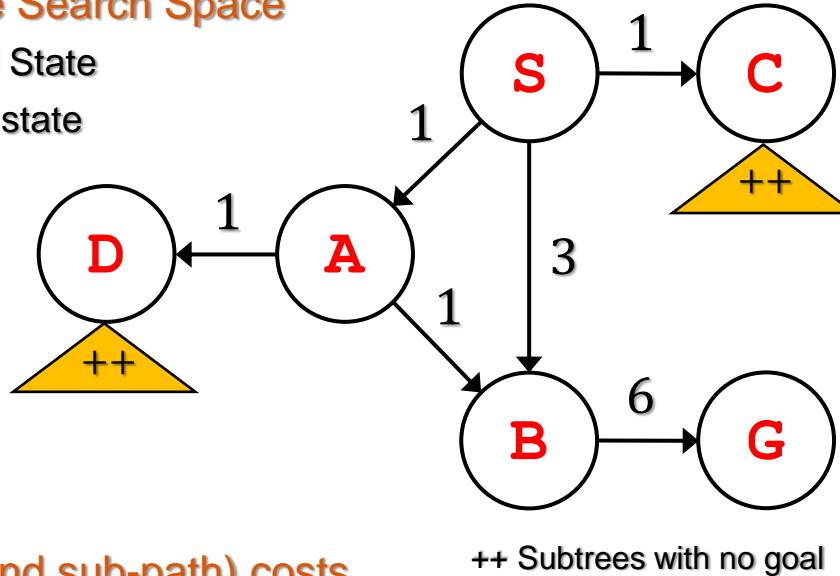


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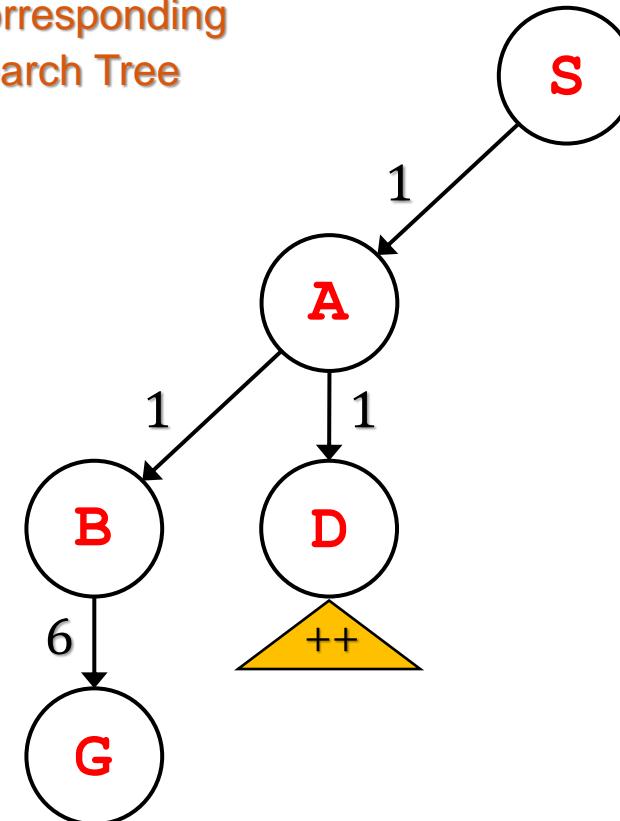
G: Goal state



- Paths (and sub-path) costs

- S-A-B-G : 8 units to goal
- S-A-D++ :  $\infty$  units to goal

- Corresponding Search Tree

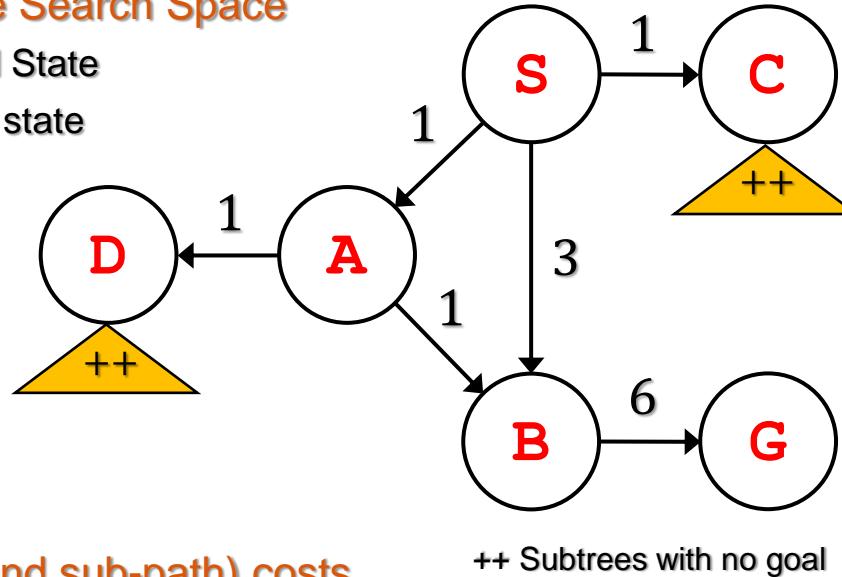


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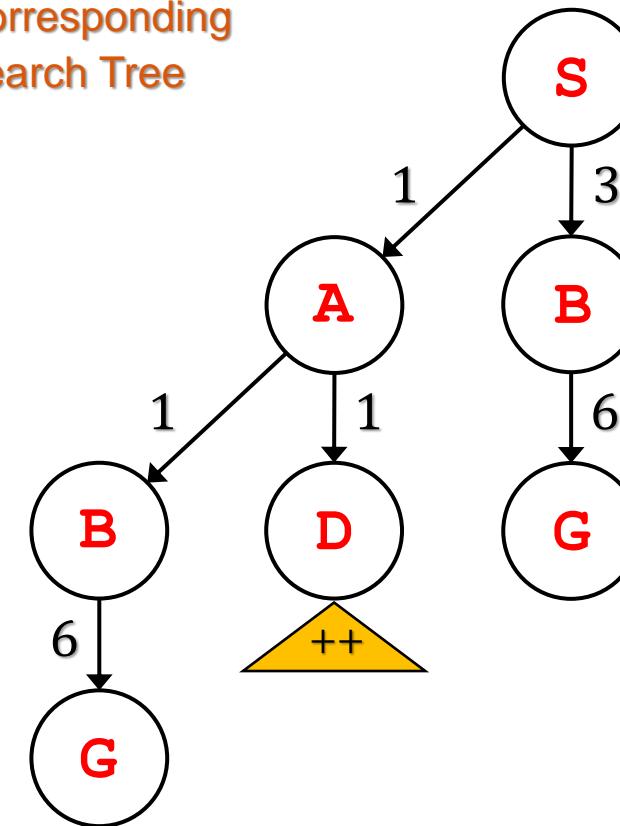
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- Paths (and sub-path) costs

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- Corresponding Search Tree

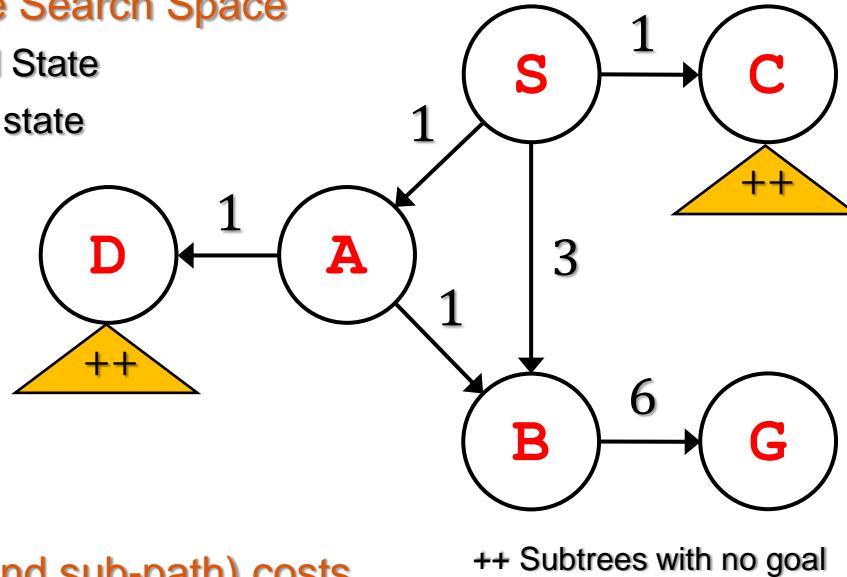


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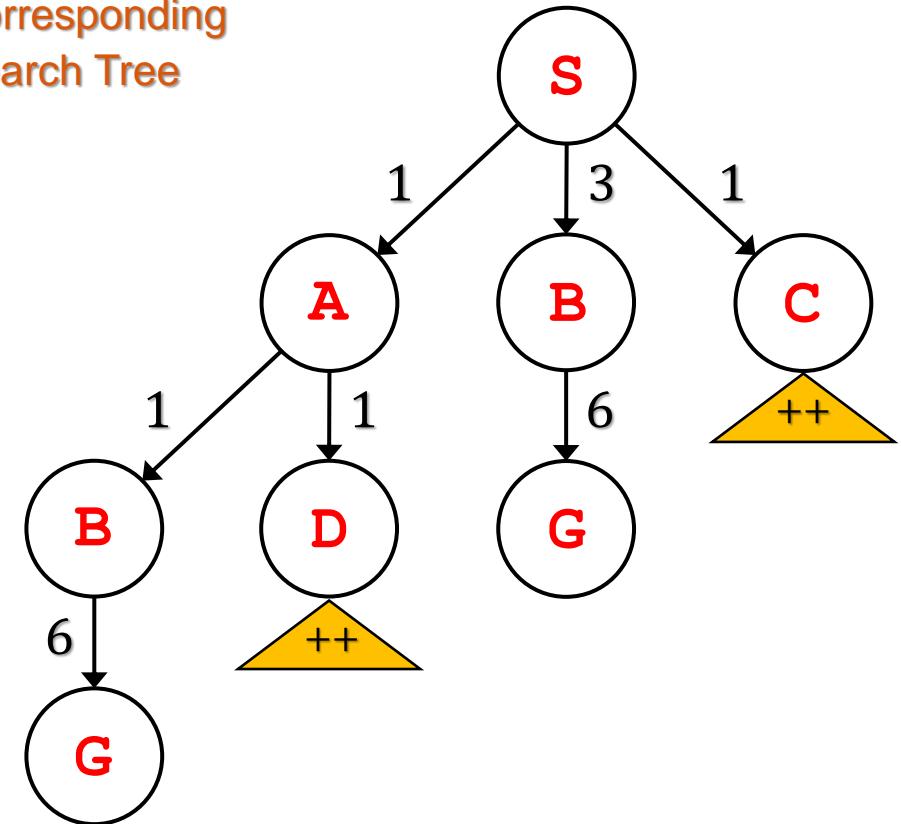
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- Paths (and sub-path) costs

- S-A-B-G : 8 units to goal
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- S-C++ :  $\infty$  units to goal

- Corresponding Search Tree

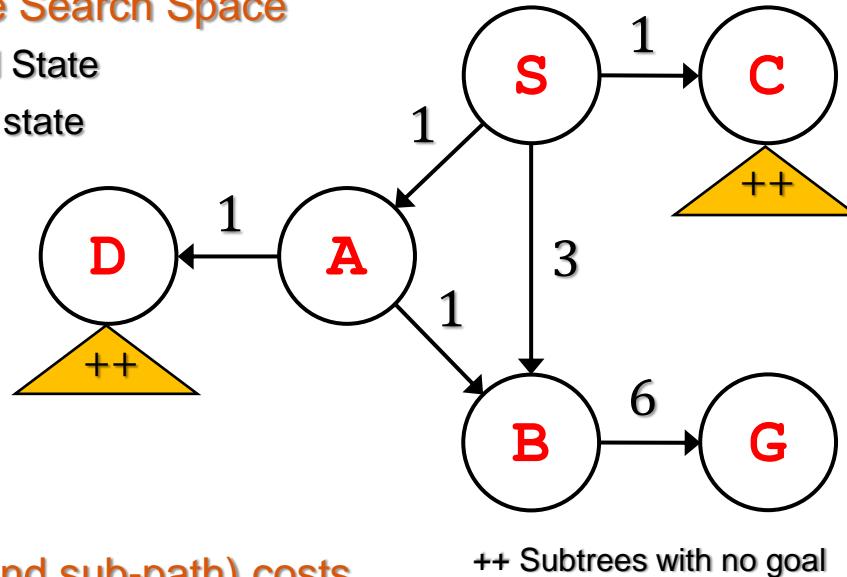


# What are the Paths Discovered by UCS?

- Example Search Space

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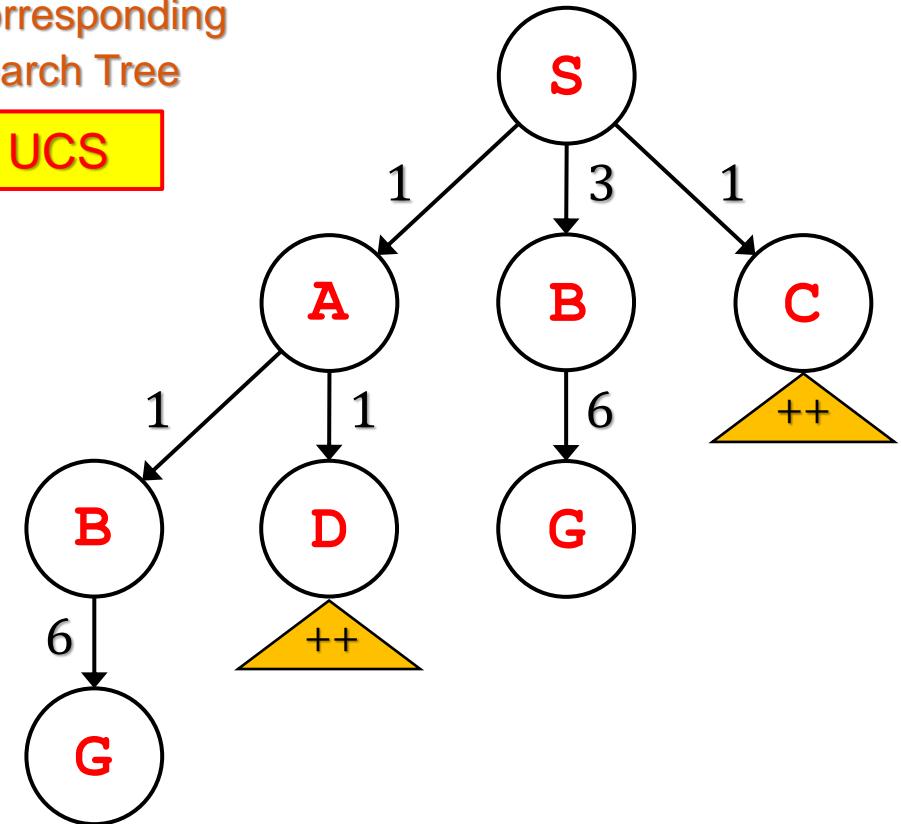


- Paths (and sub-path) costs

- S-A-B-G : 8 units to goal
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- Corresponding Search Tree

UCS

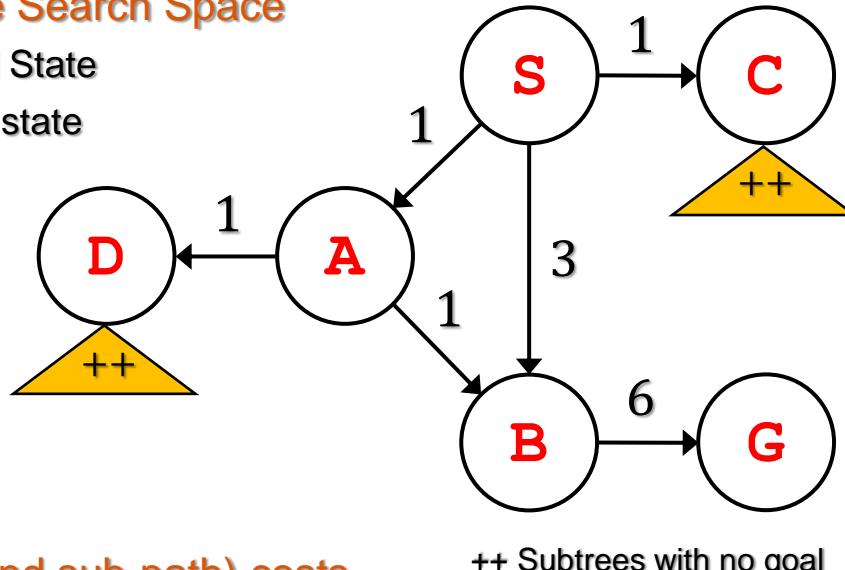


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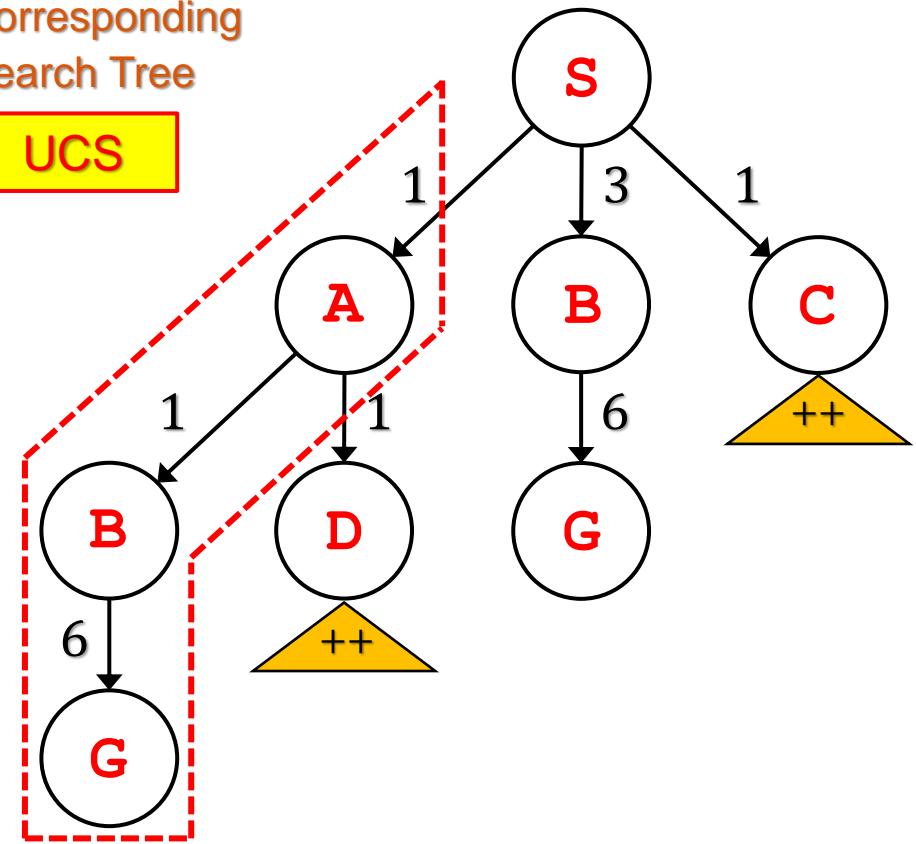


- Paths (and sub-path) costs

- S-A-B-G : 8 units to goal (all explored)
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- Corresponding Search Tree

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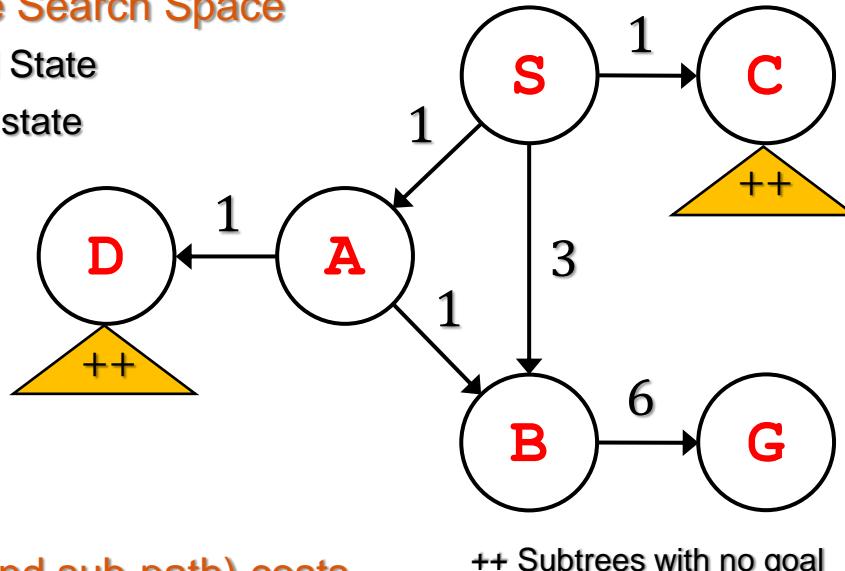


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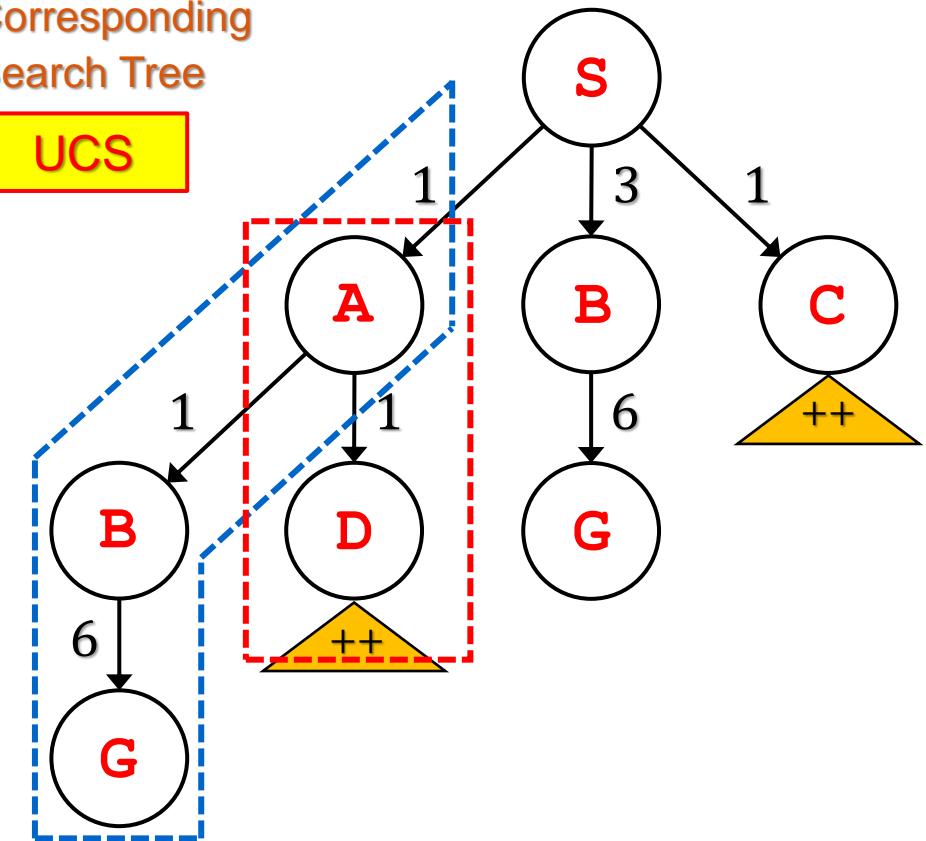
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++ Subtrees with no goal

- Corresponding Search Tree

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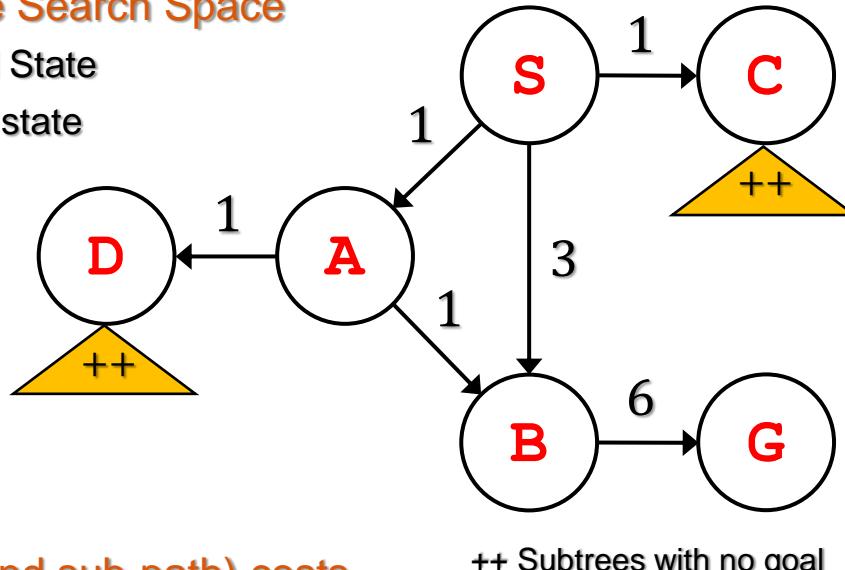


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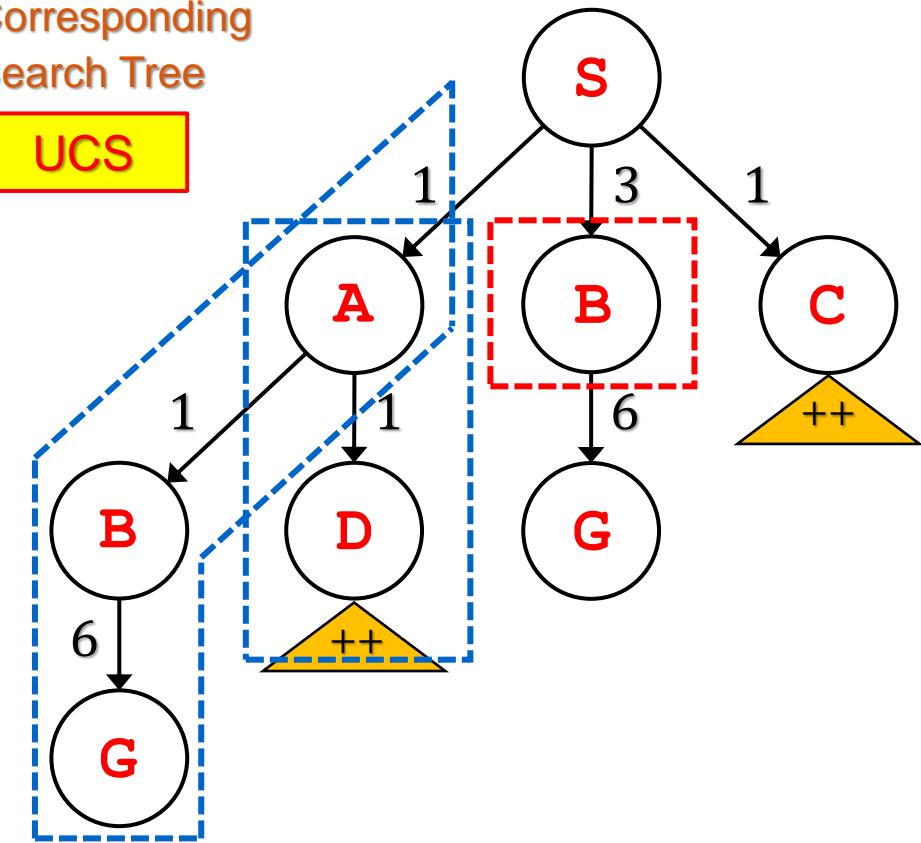
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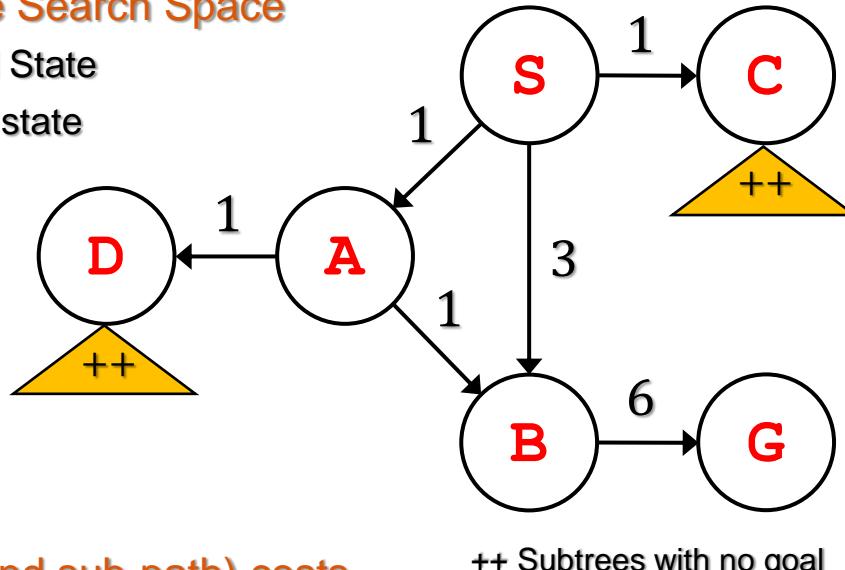


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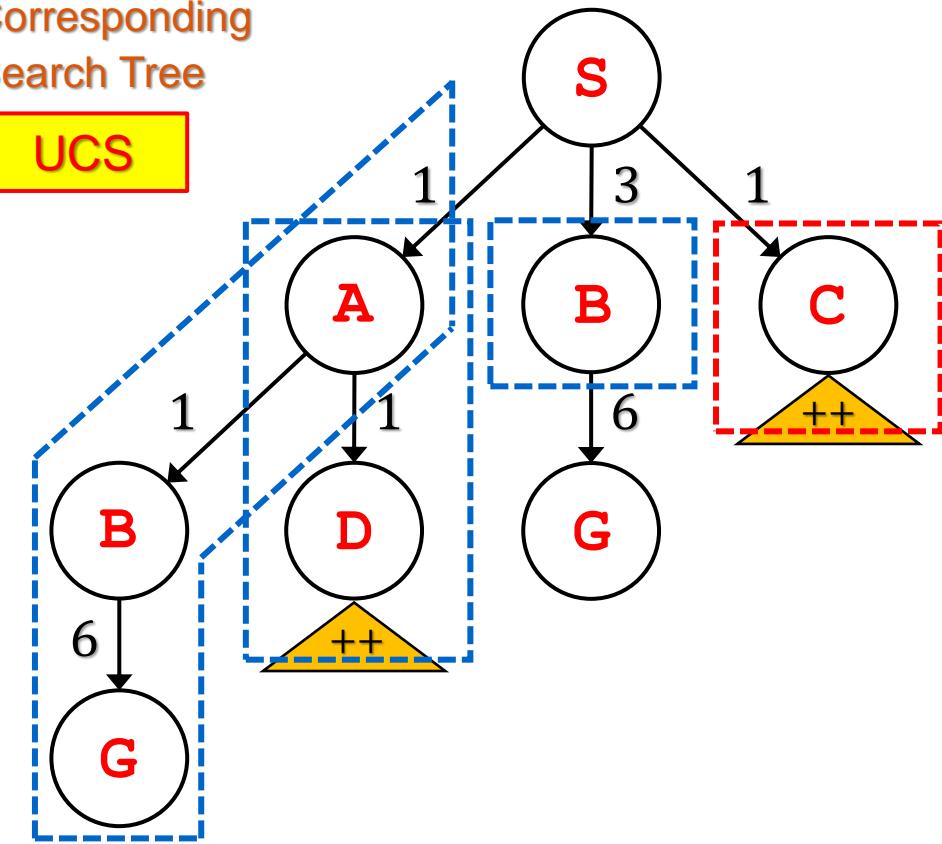
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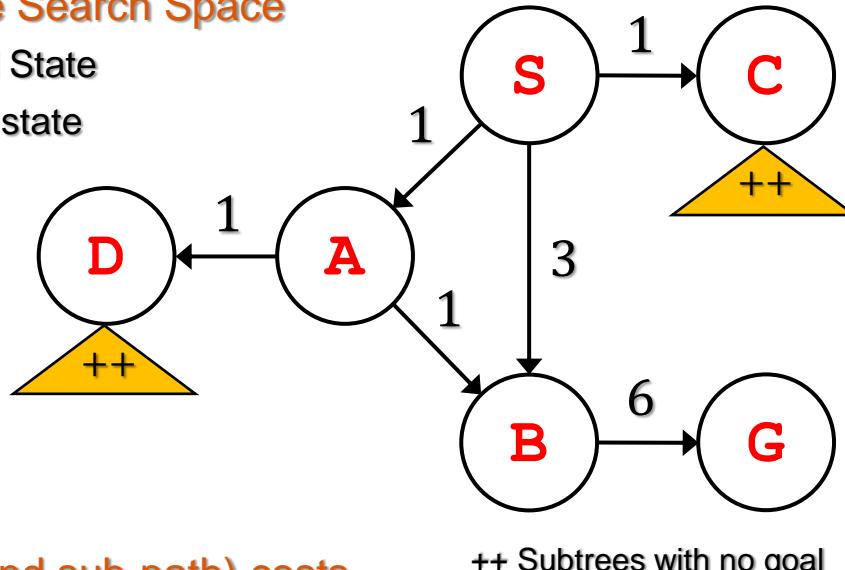


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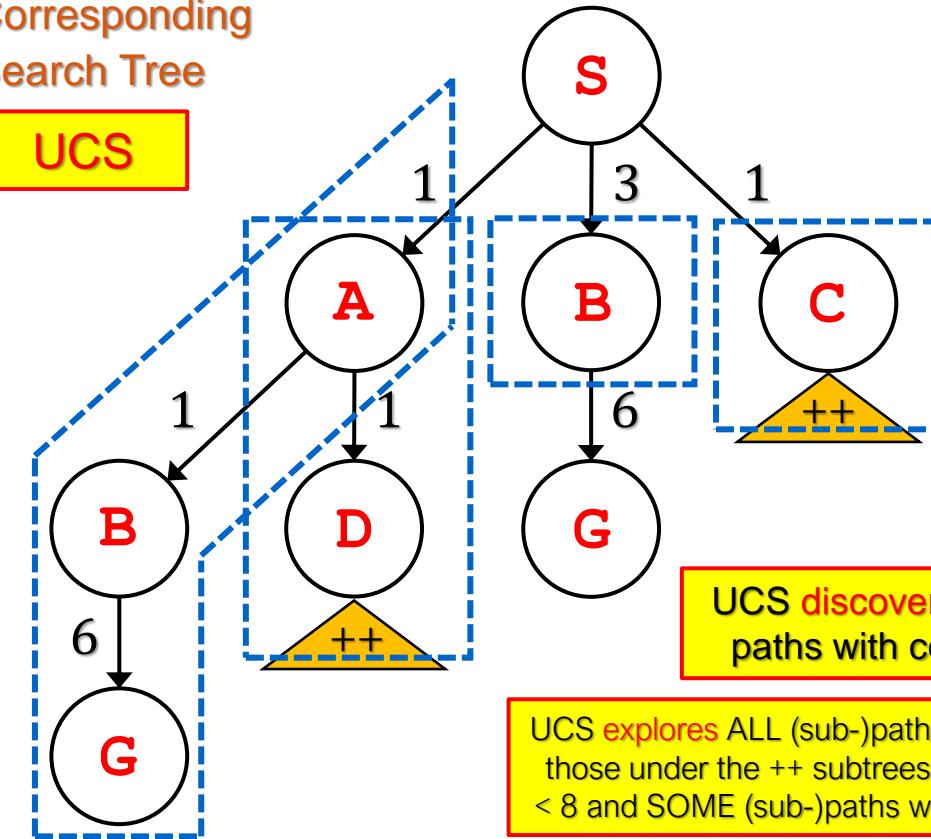
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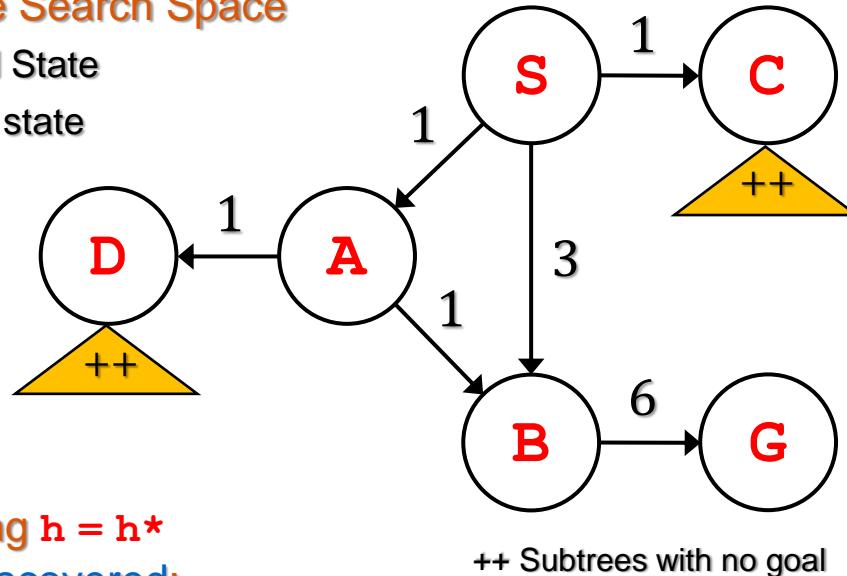


# What are the Paths Discovered by Greedy Best-First Search?

- Example Search Space

S: Initial State

G: Goal state



- Assuming  $h = h^*$

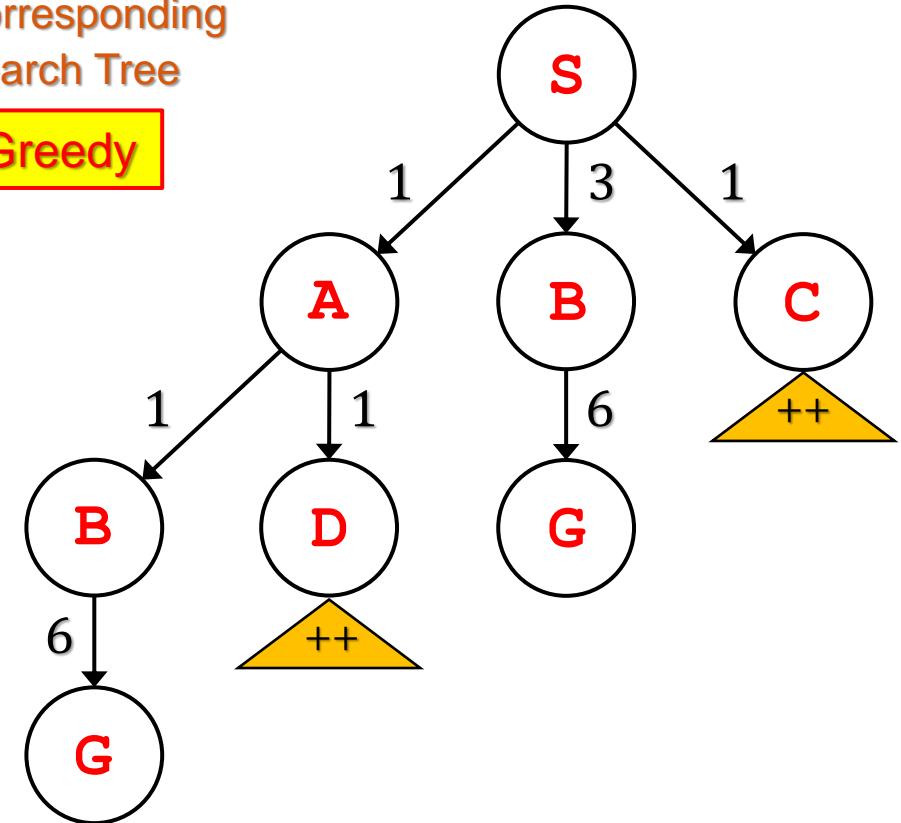
- Paths discovered:

- s (8)

n	$h^*(n)$
S	8
A	7
B	6
C	$\infty$
D	$\infty$
G	0

- Corresponding Search Tree

Greedy

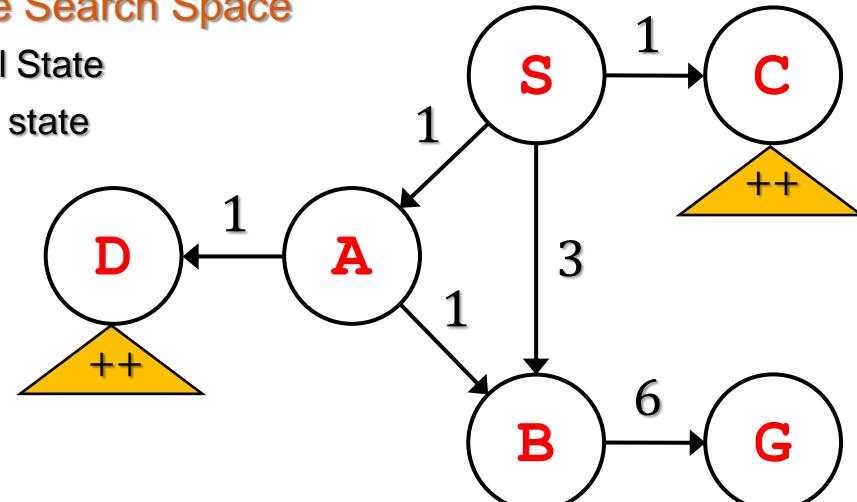


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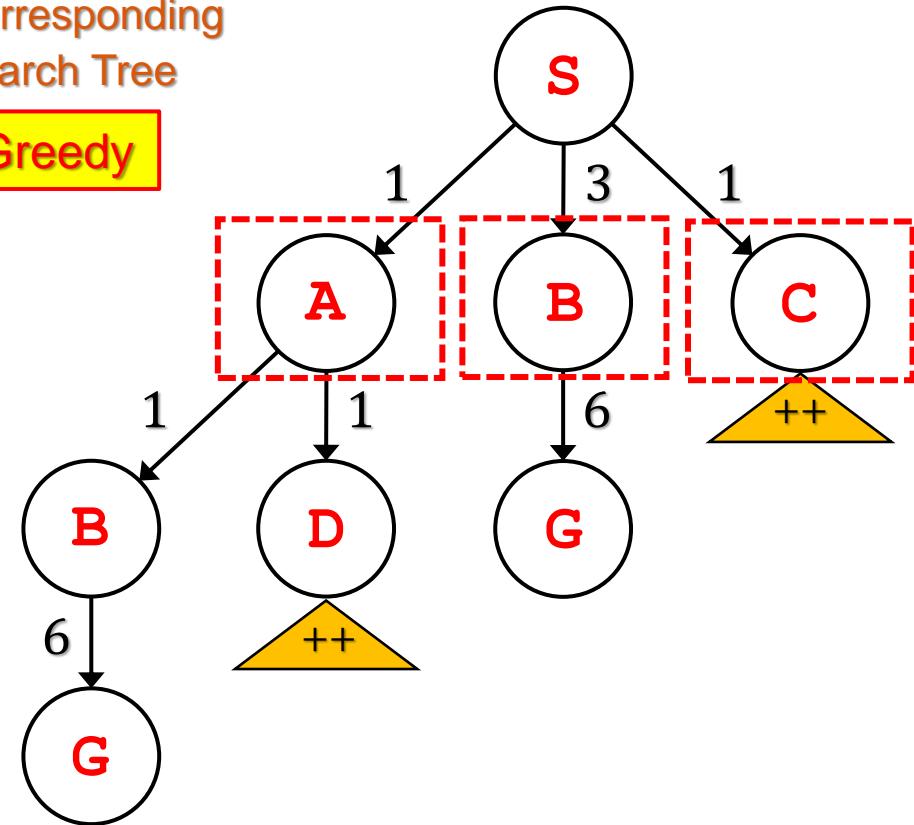
- Paths discovered:

- $s (8)$
- $s-b (6)$ ,  $s-a (7)$ ,  $s-c (\infty)$

n	$h^*(n)$
S	8
A	7
B	6
C	$\infty$
D	$\infty$
G	0

- Corresponding Search Tree

Greedy

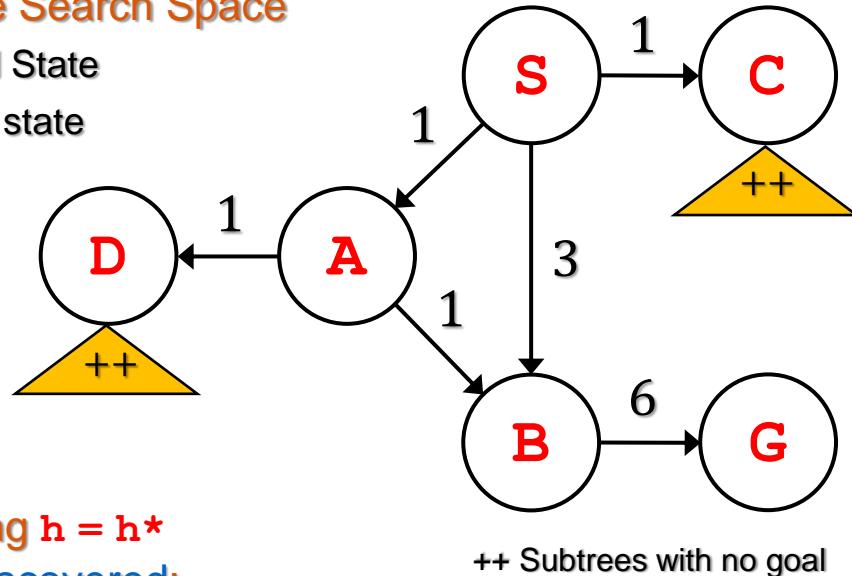


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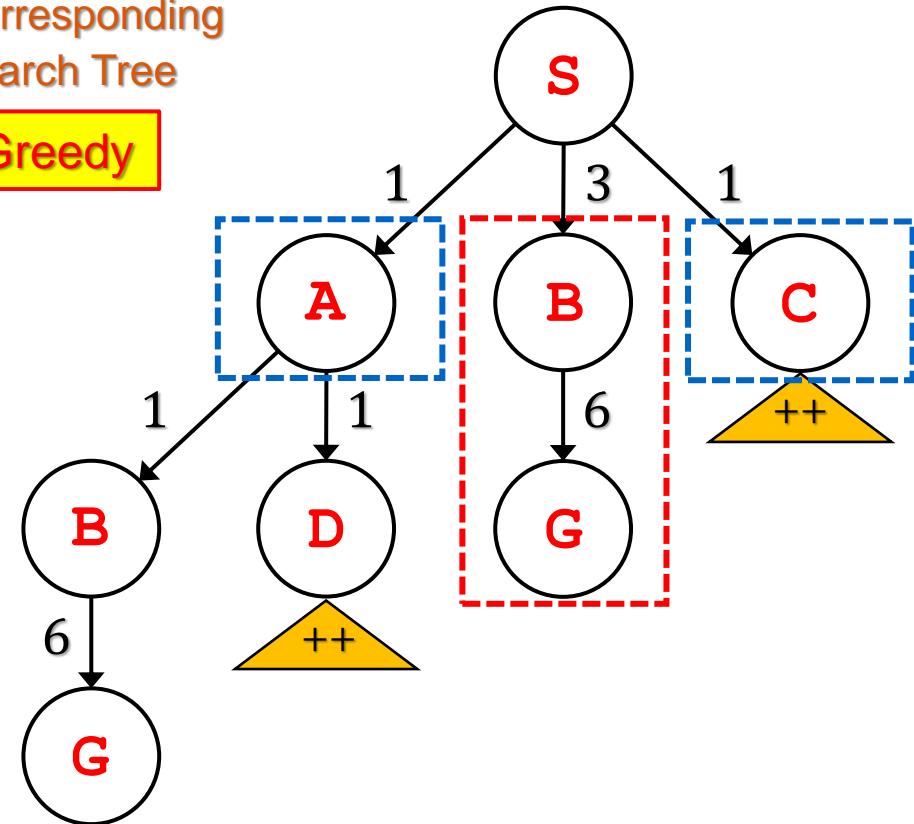
- Paths discovered:

- $s (8)$
- $s-B (6)$ ,  $s-A (7)$ ,  $s-C (\infty)$
- $s-B-G (0) // \text{return}$

n	$h^*(n)$
S	8
A	7
B	6
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G	0

- Corresponding Search Tree

Greedy

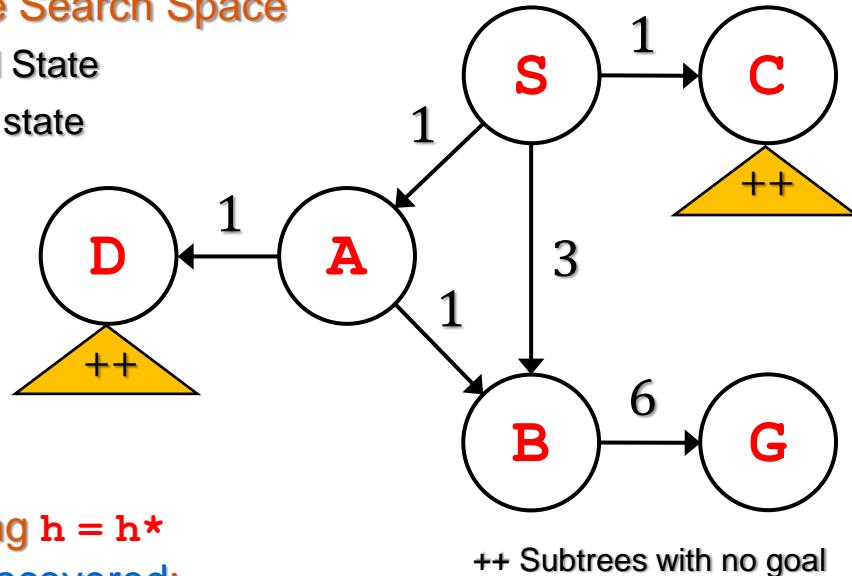


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- Assuming  $h = h^*$

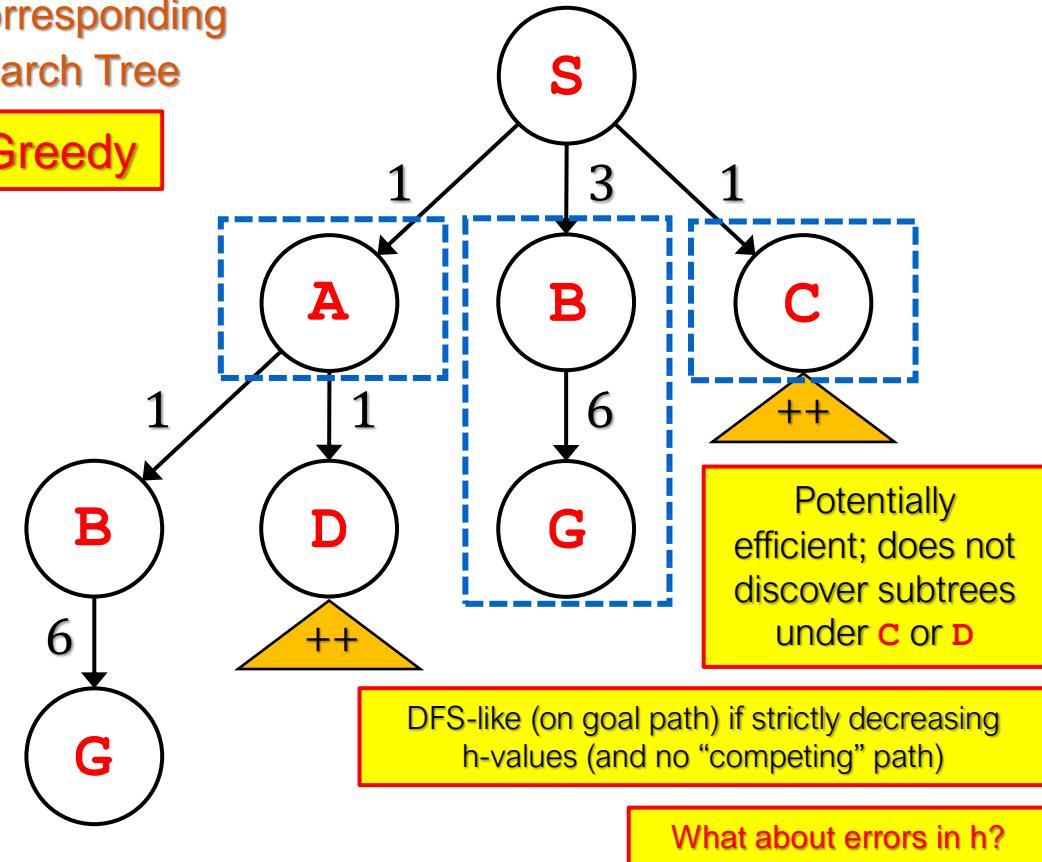
- Paths discovered:

- $s (8)$
- $s-b (6)$ ,  $s-a (7)$ ,  $s-c (\infty)$
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n	$h^*(n)$
S	8
A	7
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- Corresponding Search Tree

Greedy

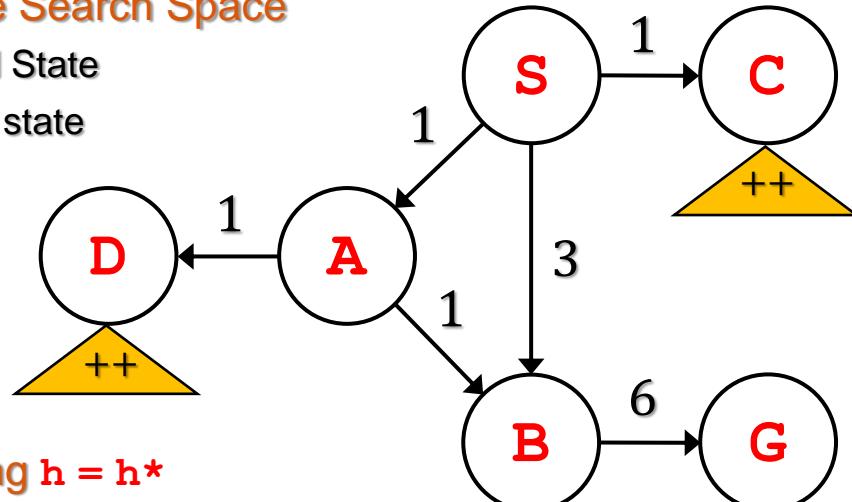


# What are the Paths Discovered by A\* Search?

- Example Search Space

S: Initial State

G: Goal state



- Assuming  $h = h^*$

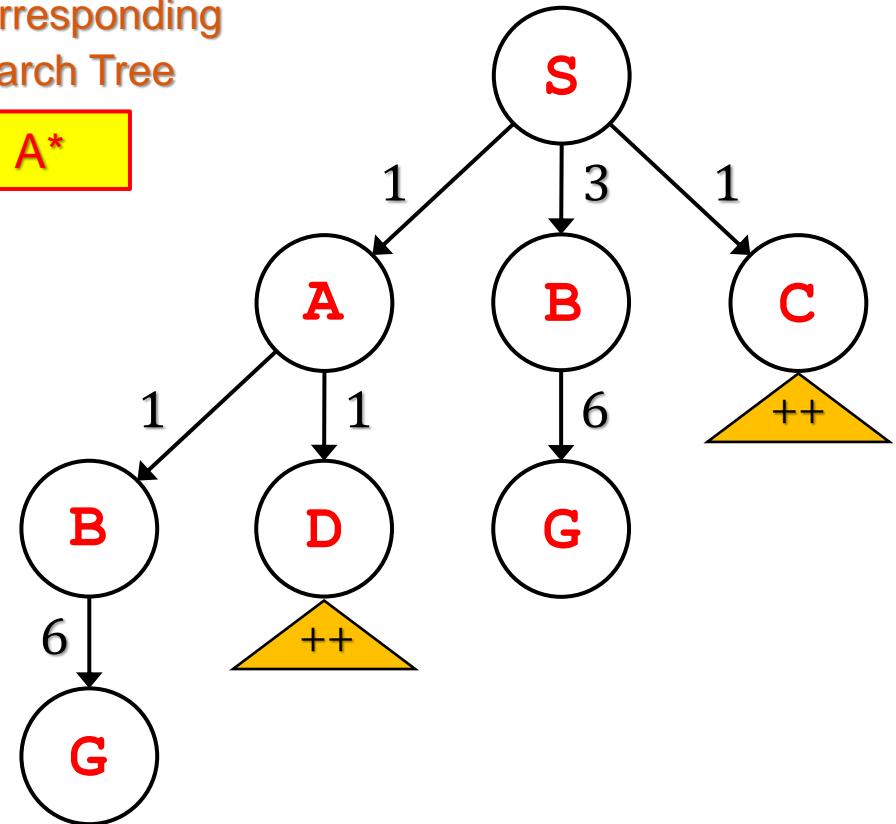
- Paths discovered:

- s (0+8)

n	$h^*(n)$
S	8
A	7
B	6
C	$\infty$
D	$\infty$
G	0

- Corresponding Search Tree

A\*

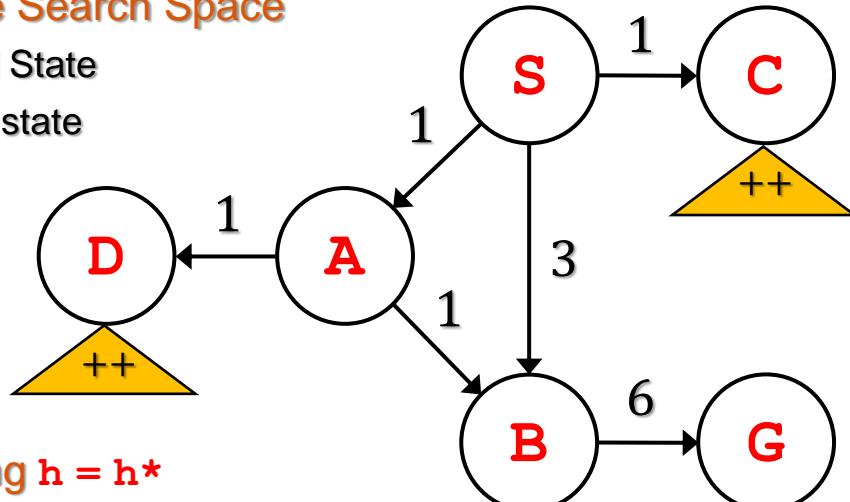


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- Example Search Space

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- Assuming  $h = h^*$

- Paths discovered:

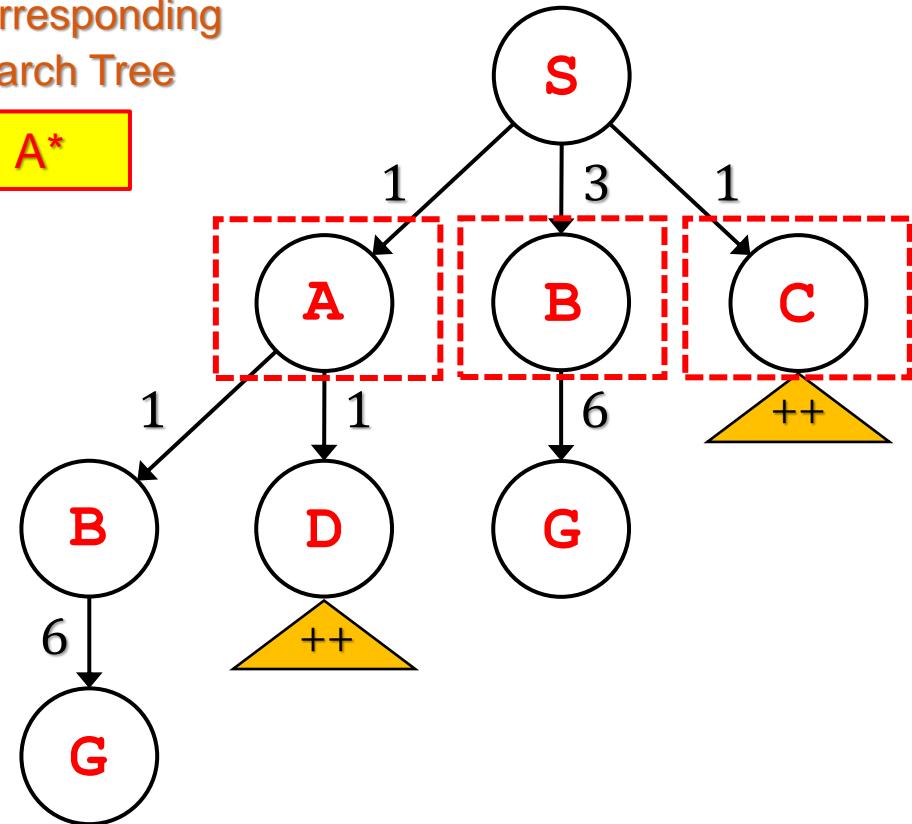
- o  $s (0+8)$

- o  $s-A (1+7)$ ,  $s-B (3+6)$ ,  
 $s-C (1+\infty)$

n	$h^*(n)$
S	8
A	7
B	6
C	$\infty$
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G	0

- Corresponding Search Tree

**A\***

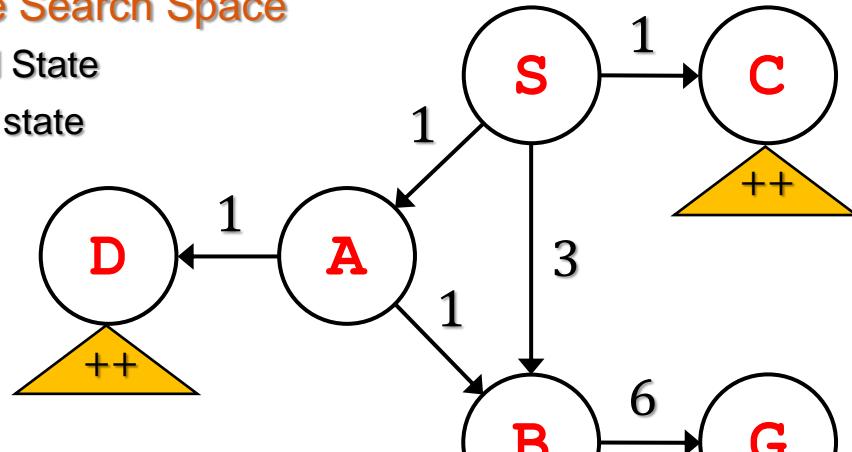


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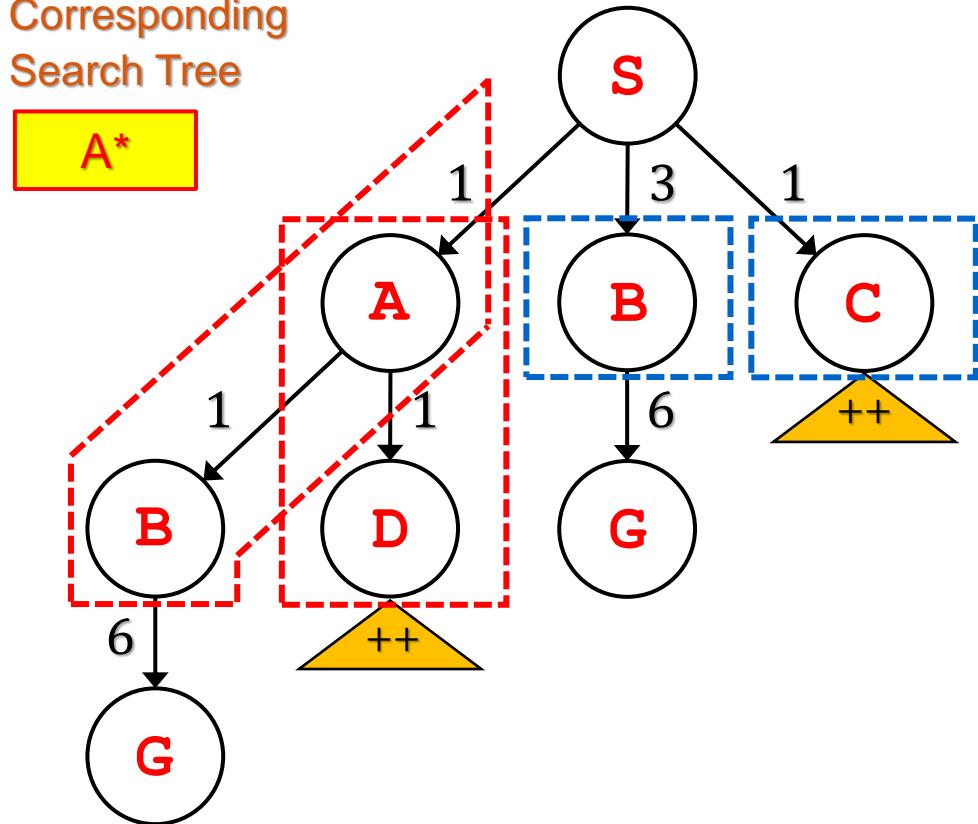
- Paths discovered:

- $s (0+8)$
- $s-A (1+7)$ ,  $s-B (3+6)$ ,  
 $s-C (1+\infty)$
- $s-A-B (2+6)$ ,  $s-A-D (2+\infty)$

++ Subtrees with no goal

n	$h^*(n)$
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- Corresponding Search Tree

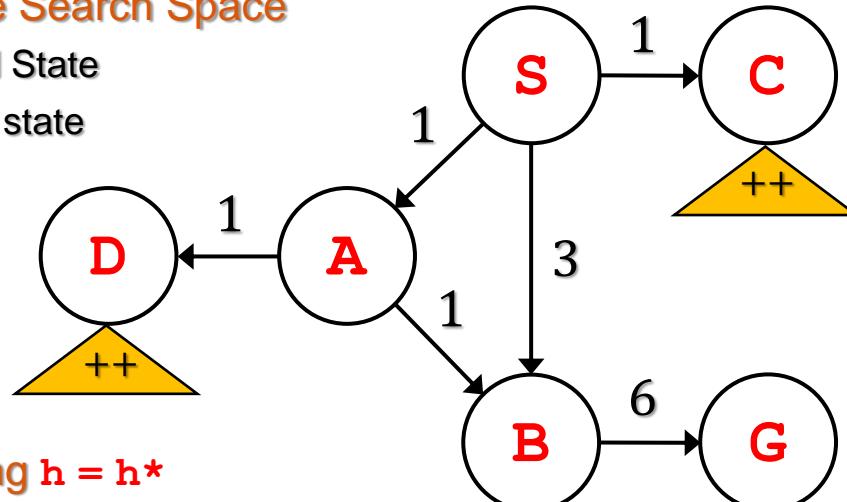


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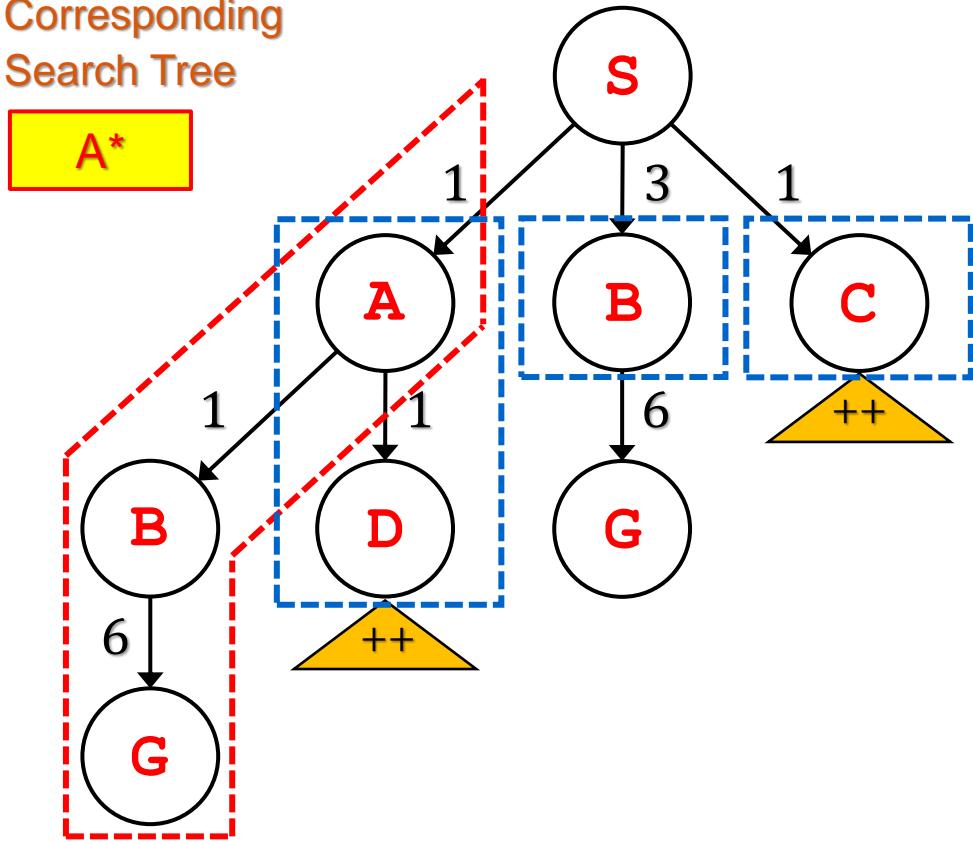
- Assuming  $h = h^*$

- Paths discovered:

- $s (0+8)$
- $s-A (1+7)$ ,  $s-B (3+6)$ ,  
 $s-C (1+\infty)$
- $s-A-B (2+6)$ ,  $s-A-D (2+\infty)$
- $s-A-B-G (8+0) // \text{return}$

n	$h^*(n)$
S	8
C	$\infty$
D	$\infty$
B	6
G	0

- Corresponding Search Tree

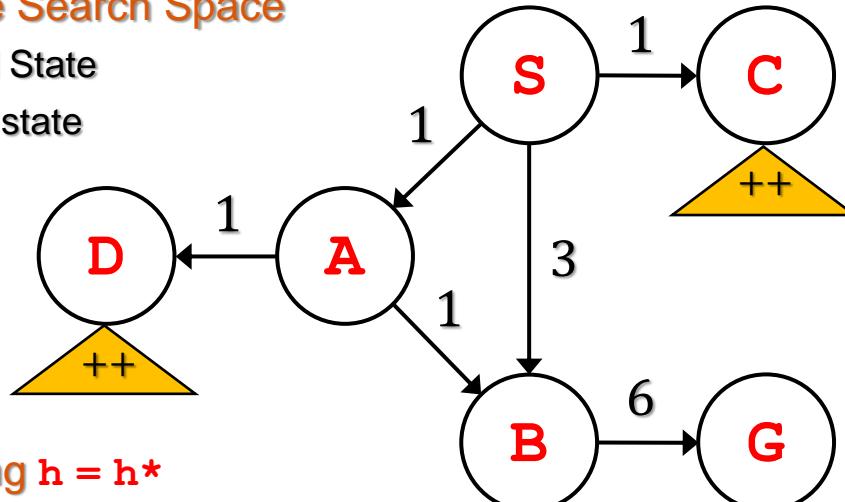


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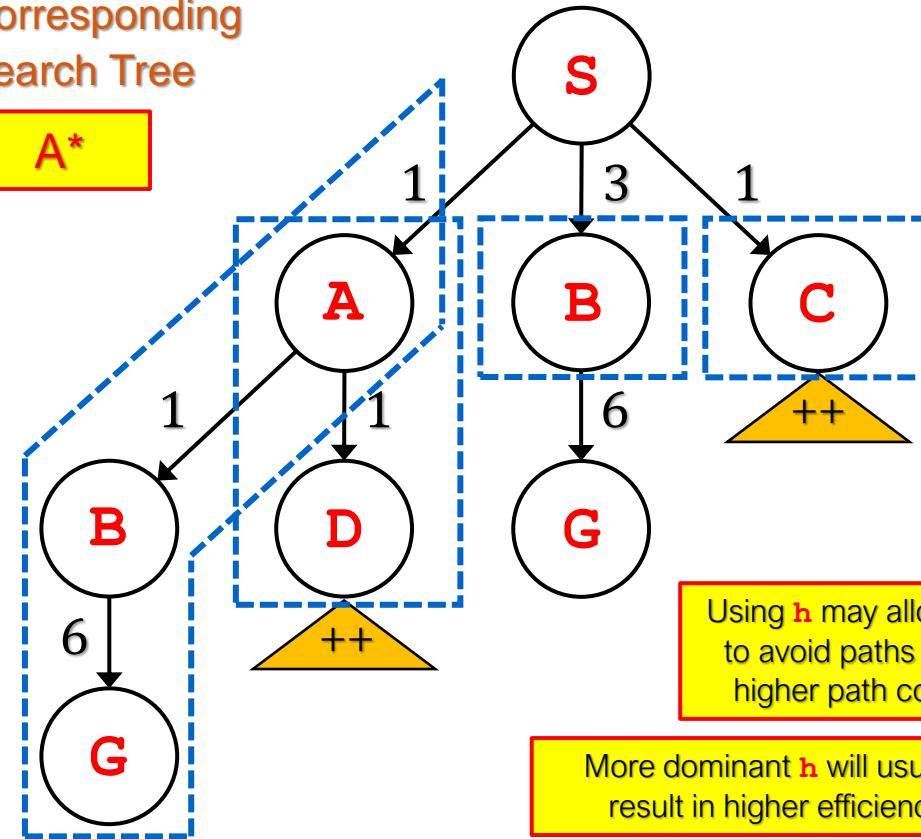
- Paths discovered:

- $s (0+8)$
- $s-A (1+7), s-B (3+6), s-C (1+\infty)$
- $s-A-B (2+6), s-A-D (2+\infty)$
- $s-A-B-G (8+0) // \text{return}$

n	$h^*(n)$
S	8
C	$\infty$
D	$\infty$
B	6
G	0

- Corresponding Search Tree

**A\***



# Admissible & Consistent Heuristics

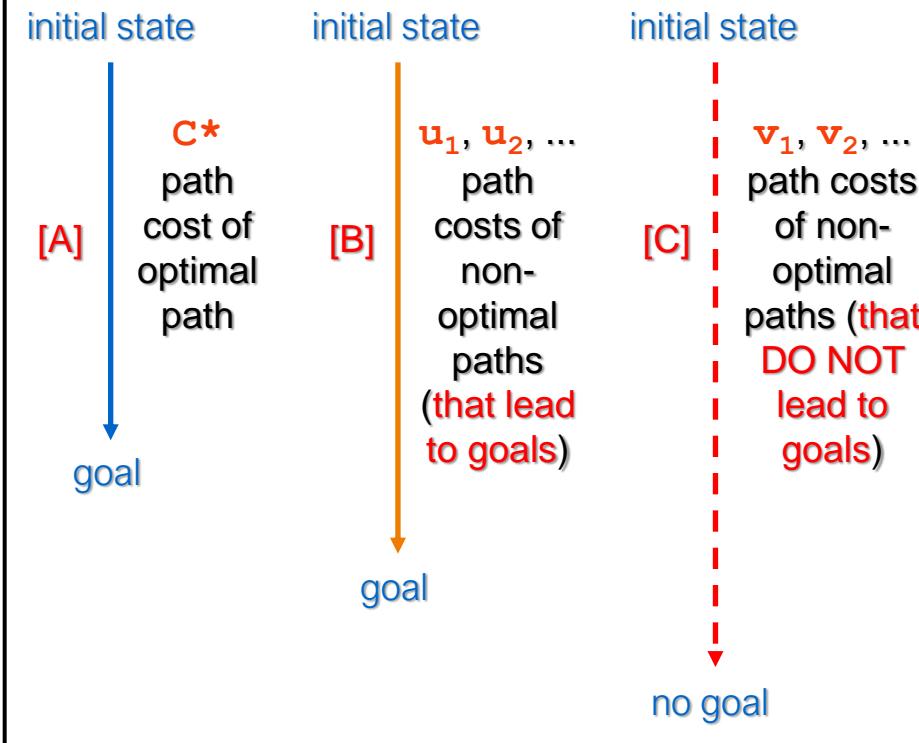
- **Admissible  $h$** 
  - $h(n)$  is **admissible** if  $\forall n: h(n) \leq h^*(n)$
  - $h(n)$  never overestimates the cost
- Theorem: If  $h(n)$  is admissible, then A\* using tree search is optimal
- **Consistent  $h$** 
  - $h(n)$  is **consistent** if  $\forall n$  and each successor of  $n, n'$ :  $h(n) \leq \text{cost}(n, a, n') + h(n')$
- Theorem: If  $h(n)$  is consistent, then A\* using graph search is optimal
  - Under Version 2 & 3 (but not Version 1)!

Main idea: by the time we explore a path to a **goal**, **P**, all paths with actual costs less than **P** must be explored

Main idea: ensures that **f** costs are monotonically increasing along a path

# Taxonomy of Paths Under A\* Search

## Based on True Path Costs



- **Using admissible heuristics (never overestimates)**
  - explores all paths on [A]
  - explores some paths under [B] and [C]
    - Note that severe underestimations are possible (e.g.,  $h(n) = 0$ , which is equivalent to UCS priority)
    - All paths estimated to be less than  $C^*$  will be explored
    - Paths estimated to be equal to  $C^*$  may be explored
- **Using consistent heuristics (never overestimates; monotonically increasing)**
  - only makes a difference under graph search
- **Using neither admissible nor consistent heuristics**
  - depends on the types of errors
  - sequence of paths explored may vary wildly

# Summary

Completeness assumptions:

- $b$  finite AND ( $m$  finite OR has a solution)
- All action costs are  $> \epsilon > 0$

- **UCS**

- **Optimal** under
    - Tree search
    - Graph search (version 2 and 3)
  - Popping node  $n \Rightarrow$  optimal path to  $n$  found

- **Greedy Best-First Search**

- **Not optimal**
  - **Incomplete** under tree search

- **A\* Search**

- Assuming admissible  $h$ 
    - May not have monotonically increasing path cost
    - **Optimal** under
      - Tree search
      - Graph search (version 2)
  - Assuming consistent  $h$ 
    - On popping node  $n$ , optimal path to  $n$  found
    - **Optimal** under
      - Tree search
      - Graph search (version 2 and 3)

Assume graph search version 2 and late goal testing (unless otherwise stated)

3

# Heuristics & Dominance

# Efficiency & Dominance

- Efficiency of A\* depends on the accuracy of its heuristics
  - Higher heuristic accuracy means we need to try fewer paths
  - Specifics not covered in CS3243
- Which heuristics are better?
- If ,  $h_1(n) \geq h_2(n)$  , then  $h_1$  dominates  $h_2$ 
  - If  $h_1$  is also admissible
    - $h_1$  must be closer to  $h^*$  than  $h_2$
    - $h_1$  is usually more efficient than  $h_2$  when used with A\*

Recall that in CS3243, we assume that a dominant heuristic is an admissible one

# Efficiency & Dominance

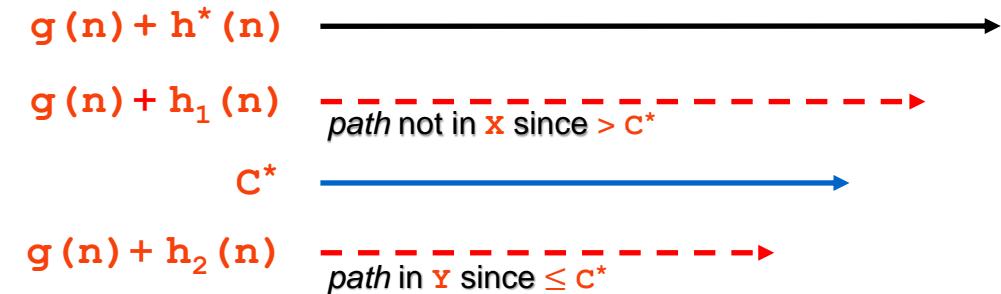
- **Proving a dominant heuristic explores fewer nodes**

- All paths with **approximated path cost  $\leq$  optimal path cost** are **explored**
- A\* explores **fewer nodes** under  $h_1$  since  $|X| \leq |Y|$ , where
  - $X$  denotes the set of paths with costs below  $C^*$  based on  $h_1$
  - $Y$  denotes the set of paths with costs below  $C^*$  based on  $h_2$
- Consider example path, node  $n$ ; assume\*\*
  - $n$  not on optimal path and  $g(n) + h^*(n) > C^*$
  - possible: (1)  $h_2(n) \leq C^* - g(n) < h^*(n) \Rightarrow n \in Y$
  - and that: (2)  $C^* - g(n) < h_1(n) \leq h^*(n) \Rightarrow n \notin X$
- **Consider the reverse of the above:**
  - $h_1(n) \leq C^* - g(n) \leq h^*(n) \& C^* < h_2(n) \leq h^*(n)$
- **The reverse case is impossible!**
  - As we have:  $h_1(n) \leq C^* \& C^* < h_2(n)$
  - Or:  $h_1(n) < h_2(n)$  which **contradicts our dominance assumption**

Assumes tree search with tie-breaking **unnecessary**

Assume  $h_1$  dominates  $h_2$  (i.e., both admissible)

Note: *path* refers to node  $n$



\*\* For completeness, also consider the other possible assumptions over node  $n$  as an exercise

# Gauging Improvement with Effective Branching Factor

- How do we measure the improvement when using dominant  $h$ ?
- Measure effective branching factor,  $b^*$ 
  - Given empirical results
    - $N$  nodes expanded
    - Solution path at depth  $d$
  - Solve for  $b^*$  using
    - $N + 1 = (b^*)^0 + (b^*)^1 + \dots + (b^*)^d$
    - i.e., given  $d$  and  $N$ , solve  $N + 1 = ((b^*)^{d+1} - 1) / ((b^*) - 1)$  to determine effective branching factor  $b^*$

Note that you will NOT be required to calculate  $b^*$  in this manner for CS3243 exams!

# How To Craft Heuristics?

- Goal is to identify a function that approximates  $h^*(n)$ 
  - Cost from  $n$  to the nearest goal
- Can we implement another search to give us this?
  - E.g., use UCS as  $h$  since it will give us  $h^*(n)$ 
    - $h(n)$  called on each node → need UCS to find the optimal path from the start state to ALL  $n$
    - Defeats the purpose of using  $h$ , which is to improve A\*
    - Better to just use  $h(n) = 0$  and just run UCS once
- We want efficient  $h$ 
  - Ideally,  $h$  is  $O(1)$ , or something reasonably cheap
  - Set our objective to admissible heuristics (since consistent is much more difficult)

# Example Problem: 8-Puzzle

7	2	4
5		6
8	3	1

Example Initial State

	1	2
3	4	5
6	7	8

Goal State

- **Search Problem Specification**
  - State Representation (Initial State):
    - Matrix representing the grid, with each  $(r, c) \in [0, 8]$
    - 0 is the blank cell
  - Actions:
    - Move a chosen cell adjacent  $(r, c)$  to the blank cell  $(r', c')$
    - Same as “moving” blank cell  $(r', c')$  to a chosen cell adjacent  $(r, c)$ 
      - i.e., legal  $(r'-1, c'), (r'+1, c'), (r', c-1), (r', c+1)$  cells → possible actions
  - Goal Test:
    - Current state matrix = goal state matrix
  - Transition Model:
    - Swap the contents of  $(r, c)$  and  $(r', c')$
  - Cost Function:
    - Each action cost is 1 unit

This Puzzle requires the player to shift the numbered squares into the empty cell until the final pattern is obtained

How do we get an admissible heuristic out of this puzzle?

# Questions about the Lecture?

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# 4

# Relaxing the Problem

# 8-Puzzle Example: Defining Admissible Heuristics by Relaxing the Problem (Version A)

7	2	4
5		6
8	3	1

Example Initial State

	1	2
3	4	5
6	7	8

Goal State

- Define an easier problem based on the same context
  - Let heuristic  $h$  be a function that counts actions required to solve the problem
  - The puzzle is constrained by this rule
    - A tile can move from square  $x$  to square  $y$  if  $x$  is adjacent to  $y$  and  $y$  is blank
  - Relaxed 8-Puzzle: Version A ( $h_0$ )
    - Remove all tiles (1 move)
    - Place them in the correct positions (8 moves)
    - $h_0 = 9$ , for any problem
    - $h_0$  overestimates the moves required on some cases

Test admissibility by trying to work backwards from the goal state

We did not properly relax the rules and instead just defined new ones

# 8-Puzzle Example: Defining Admissible Heuristics by Relaxing the Problem (Version B)

7	2	4
5		6
8	3	1

Example Initial State

	1	2
3	4	5
6	7	8

Goal State

- Define an easier problem based on the same context
  - Let heuristic  $h$  be a function that counts actions required to solve the problem
  - The puzzle is constrained by this rule
    - A tile can move from square  $x$  to square  $y$  if  $x$  is adjacent to  $y$  and  $y$  is blank
  - Relaxed 8-Puzzle: Version B ( $h_1$ )
    - A tile can move from square  $x$  to square  $y$  ~~if  $x$  is adjacent to  $y$  and  $y$  is blank~~
    - $h_1$  = number of cells in the wrong position
    - Complexity:  $O(n)$ , where  $n$  is the size of the grid
    - $h_1$  is admissible!

Note: also removes implied constraint that each square can only house one tile – i.e., now we can have more than 1 tile on each square

By properly relaxing the rule, we got an admissible heuristic:  $h_1$  – misplaced tiles

Can we do better? Find an admissible  $h$  that dominates this one

# 8-Puzzle Example: Defining Admissible Heuristics by Relaxing the Problem (Version C)

7	2	4
5		6
8	3	1

Example Initial State

	1	2
3	4	5
6	7	8

Goal State

- Define an easier problem based on the same context
  - Let heuristic  $h$  be a function that counts actions required to solve the problem
  - Relaxed 8-Puzzle: Version B ( $h_1$ )
    - A tile can move from square  $x$  to square  $y$   
~~if  $x$  is adjacent to  $y$  and  $x$  is blank~~
  - Relaxed 8-Puzzle: Version C ( $h_2$ )
    - A tile can move from square  $x$  to square  $y$   
~~if  $x$  is adjacent to  $y$  and  $x$  is blank~~
    - $h_2 = \text{sum over each Manhattan distance}$  between a square and its goal location
    - Complexity:  $O(n)$ , where  $n$  is the size of the grid (given that grid dimensionality here is fixed at 2)
    - $h_2$  is admissible and dominates Version B

Note: this STILL removes the implied constraint that each square can only house one tile – i.e., now we can more than 1 tile on each square

New admissible heuristic:  $h_2 - \text{Manhattan distance}$ 

- $h_2$  dominates  $h_1$  ( $h_1$  is a relaxation of  $h_2$ )
- $h_2$  is admissible ( $h_2$  is a relaxation of the original rule)

# Effects of Dominance Under 8-Puzzle

$d$	Search Cost (nodes generated)			Effective Branching Factor		
	BFS	$A^*(h_1)$	$A^*(h_2)$	BFS	$A^*(h_1)$	$A^*(h_2)$
6	128	24	19	2.01	1.42	1.34
8	368	48	31	1.91	1.40	1.30
10	1033	116	48	1.85	1.43	1.27
12	2672	279	84	1.80	1.45	1.28
14	6783	678	174	1.77	1.47	1.31
16	17270	1683	364	1.74	1.48	1.32
18	41558	4102	751	1.72	1.49	1.34
20	91493	9905	1318	1.69	1.50	1.34
22	175921	22955	2548	1.66	1.50	1.34
24	290082	53039	5733	1.62	1.50	1.36
26	395355	110372	10080	1.58	1.50	1.35
28	463234	202565	22055	1.53	1.49	1.36

Recall: Effective Branching Factor computed by substituting  $d$  and total nodes generated into geometric series sum and solving to find  $b^*$

Data are averaged over 100 puzzles for each solution length  $d$  from 6 to 28

# Rules to Functions

- Will we always be able to define functions  $h_1$  and  $h_2$  to match the relaxed rules (or even define the original rules)?
- Models and approximations
  - Finding functions that model or approximate the quantity you want (efficiently)
  - Constructing models
    - Bottom-up
      - What variables can you efficiently calculate?
      - What can these variables model?
    - Top-down
      - What (dependent) variables do I want to model / approximate?
      - What are the (independent) variables that help to calculate these?

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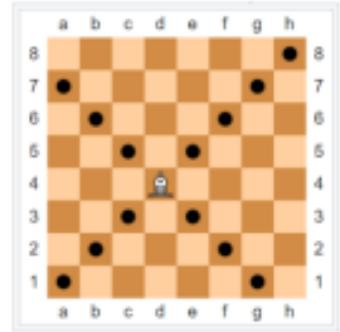


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# Midterm Question 3

- **Heuristic Design**
  - Example: AY22/23 S3 Q3:

3. Consider the Project 1 problem description – i.e., given a Chess board, some obstacles and opposing pieces – find a path from a starting position to a goal position. However, here, instead of the agent corresponding to a Knight piece, the agent will correspond to a Bishop piece. The movement range of the Bishop piece is depicted on the right. As with Project 1, assume that board corresponds to a rectangular grid. For this question, you may further assume that there is only one goal position.

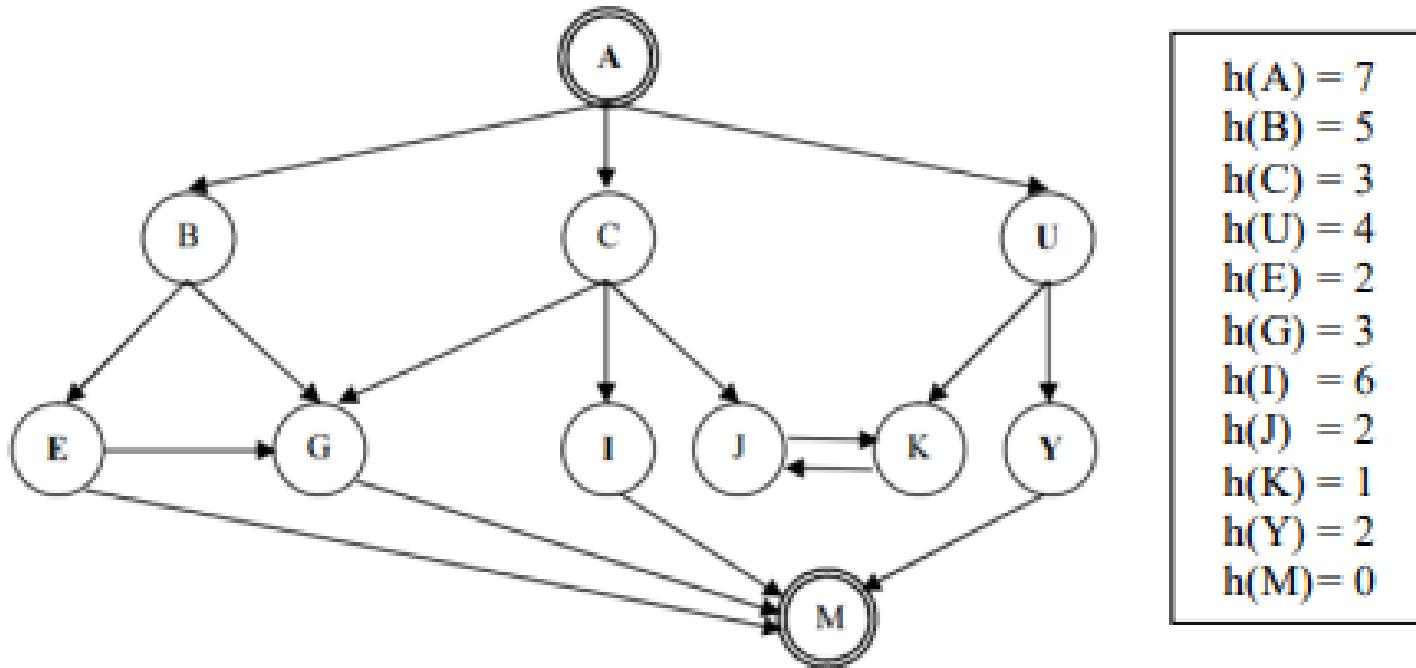


Let a heuristic,  $h_1$ , be defined as follows for any given state  $n$ .

- $h_1(n) = 0$  if  $n = goal$ , where  $n$  and  $goal$  are coordinates on the Chess board
- $h_1(n) = \infty$  if  $(n \neq goal) \wedge (\text{cellColour}(n) \neq \text{cellColour}(goal))$
- $h_1(n) = 0$  if  $(n \neq goal) \wedge (\text{cellColour}(n) = \text{cellColour}(goal)) \wedge (\text{ManhattanDist}(n, goal) \% 2 = 1)$
- $h_1(n) = 1$  if  $(n \neq goal) \wedge (\text{cellColour}(n) = \text{cellColour}(goal)) \wedge (\text{ManhattanDist}(n, goal) \% 2 = 0)$

# Midterm Question 2

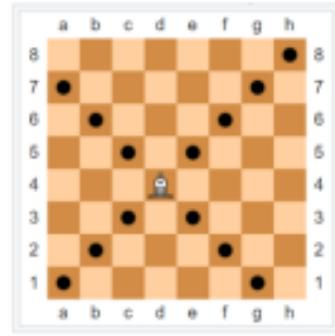
- Search Algorithm Trace (& Proof)
  - Example: AY22/23 S3 Q2b:



# Midterm Question 1

- **Search Problem Modelling**
  - Example: AY21/22 S2 Q1: Cleaning Robot

3. Consider the Project 1 problem description – i.e., given a Chess board, some obstacles and opposing pieces – find a path from a starting position to a goal position. However, here, instead of the agent corresponding to a Knight piece, the agent will correspond to a Bishop piece. The movement range of the Bishop piece is depicted on the right. As with Project 1, assume that board corresponds to a rectangular grid. For this question, you may further assume that there is only one goal position.



Let a heuristic,  $h_1$ , be defined as follows for any given state  $n$ .

- $h_1(n) = 0$  if  $n = goal$ , where  $n$  and  $goal$  are coordinates on the Chess board
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- $h_1(n) = 0$  if  $(n \neq goal) \wedge (\text{cellColour}(n) = \text{cellColour}(goal)) \wedge (\text{ManhattanDist}(n, goal) \% 2 = 1)$
- $h_1(n) = 1$  if  $(n \neq goal) \wedge (\text{cellColour}(n) = \text{cellColour}(goal)) \wedge (\text{ManhattanDist}(n, goal) \% 2 = 0)$

# Midterm Question 3

- **Search Problem Modelling**
  - Example: AY22/23 S1 Q1: Dining Philosophers

1. There are  $n$  philosophers dining together at a circular table (where  $n \geq 2$ ). Each philosopher has their own place at the table. Their philosophical problem in this instance is that the dish served is a kind of spaghetti that must be eaten with two forks.

Each of the  $n$  philosophers has their own plate of spaghetti. However, there is only one fork between each plate (notice therefore that there are also  $n$  forks).

Each philosopher can only alternately think or eat. Moreover, a philosopher can only eat their spaghetti when they have both a left and right fork. Thus, two forks will only be available when their two nearest neighbours are thinking, not eating. After an individual philosopher finishes eating, they may put down both forks.



Image from Wikipedia

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