

PostgreSQL Performance



PostgreSQL Performance

- Configuration
- Explain
- Planner Statistics
- Indexes
- Understanding the Optimizer
- Tuning Example



Performance Configuration



Configuration

• Several configuration parameters have an affect on performance

Most are tuned via formulas in the default parameter groups

Checkpoint values should be set based on business rules



shared_buffers

Sets the primary cache for the server

Corresponds to the number of shared memory buffers used by the database server.

Each buffer is 8K bytes.

Minimum value must be 16 and at least 2 x max_connections.



work_mem

 Amount of memory in KB to be used by internal sorts and hash tables before switching to temporary disk files.

Minimum allowed value is 64 KB.

• It is set in KB and the default is 1024 KB (1 MB).



maintenance_work_mem

• Maximum memory in KB to be used in maintenance operations such as VACUUM, CREATE INDEX, and ALTER TABLE ADD FOREIGN KEY. Minimum allowed value is 1024 KB.

It is set in KB and the default is 16384 KB (16 MB).

 Performance for vacuuming and restoring database dumps can be improved by increasing this value.



wal_buffers

• Number of disk-page buffers allocated in shared memory for WAL data.

Each buffer is 8K bytes

Need only be large enough to hold the amount of WAL data created by a typical transaction

Minimum allowed value is 32K.

• Default setting is 16MB.



Explain



EXPLAIN

• There are a number of different ways to execute a query

These ways are called a plan

• More complex queries have more possible plans

EXPLAIN returns the plan that PostgreSQL chose



Explain EXPLAIN

```
EXPLAIN SELECT * FROM cities;

QUERY PLAN

Seq Scan on cities (cost=0.00..1230.11 rows=63211 width=43)

(1 row)
```

- Cost
 - First number is the cost for the first row
 - Second number is the cost for all rows
- Rows
 - Number of rows the planner thinks it will be returning
- Width
 - Each row is 43 bytes wide
- The planner has decided that it will be doing a sequential scan



Explain

- Estimated start-up cost (time expended before output scan can start (e.g., time to do sorting in a sort node)
- Estimated total cost if all rows were to be retrieved (e.g., disregards effects of a LIMIT clause)
- Estimated # of rows output if executed to completion
- Estimated average width in bytes of rows output
- (Costs are in units of disk page fetches CPU efforts are converted into disk-page units)



EXPLAIN ANALYZE

```
EXPLAIN ANALYZE SELECT * FROM cities;

QUERY PLAN

Seq Scan on cities (cost=0.00..1230.11 rows=63211 width=43)

(actual time=0.008..5.494 rows=63211 loops=1)

Planning Time: 0.041 ms

Execution Time: 8.504 ms

(3 rows)
```

- The actual runtime of this query.
- Note: the estimate seems to have no correlation to the actual time
- The estimate is in an arbitrary unit which is "How long it takes sequentially read a single page from disk"



Explain Analyze

Actual time in milliseconds of real time (per loop if loops > 1)

Number or rows output (per loop if loops > 1)

Planning Time includes parsing, rewriting, or planning time

Execution Time is executor start-up, shutdown, & processing time



Explain with more than one step

```
EXPLAIN SELECT * FROM cities ORDER BY name;

QUERY PLAN

Sort (cost=8216.52..8374.55 rows=63211 width=43)

Sort Key: name

-> Seq Scan on cities (cost=0.00..1230.11 rows=63211 width=43)
```

- Sort and Sequential scan
- Data flows from the lower steps to the higher steps
- Output of the seq scan is fed into the sort
- The key which is used for the sort is "name"
- First row of the sort is high (a sort can not return rows until it is done)
- Actual cost of the sort is (8216.52 1230.11 = 6986.41)
 - It needs all of the rows before it can sort



Multiple levels

- Indentation is used to show what steps feed into the next step above
 - The Hash is fed from the Seq scan
- The hash join does not wait for the scan, it can start returning immediately



Diagnosing performance issues

```
=> EXPLAIN ANALYZE SELECT * FROM game changers;
 GroupAggregate (cost=546632.16..549163.19 rows=37497 width=22) (actual time=34264.743..37864.405 rows=2933 loops=1)
   Group Key: p.name
  Filter: (count(*) > 10)
   Rows Removed by Filter: 1180924
   -> Sort (cost=546632.16..546913.39 rows=112490 width=14) (actual time=34263.350..37271.095 rows=2110225 loops=1)
         Sort Key: p.name
        Sort Method: external merge Disk: 50944kB
        -> Nested Loop (cost=177244.72..535269.57 rows=112490 width=14) (actual time=1533.309..25609.351 rows=2110225 loops=1)
               -> Hash Join (cost=177244.29..480421.40 rows=114485 width=10) (actual time=1533.177..6251.774 rows=2218756 loops=1)
                     Hash Cond: ((pt.title id)::text = (r.title id)::text)
                    -> Seq Scan on people_titles pt (cost=0.00..245737.43 rows=15011943 width=20) (actual time=0.007..1398.222 rows=15012000 loops=1)
                     -> Hash (cost=177155.74..177155.74 rows=7084 width=20) (actual time=1533.154..1533.154 rows=35358 loops=1)
                           Buckets: 65536 (originally 8192) Batches: 1 (originally 1) Memory Usage: 2308kB
                           -> Hash Join (cost=21082.40..177155.74 rows=7084 width=20) (actual time=172.764..1522.431 rows=35358 loops=1)
                                Hash Cond: ((t.title id)::text = (r.title id)::text)
                                 -> Hash Join (cost=1.14..148509.24 rows=582567 width=10) (actual time=0.017..1184.387 rows=516421 loops=1)
                                      Hash Cond: (t.kind id = k.kind id)
                                      -> Seg Scan on titles t (cost=0.00..126734.67 rows=5825667 width=14) (actual time=0.005..538.236 rows=5825667
loops=1)
                                       -> Hash (cost=1.12..1.12 rows=1 width=2) (actual time=0.007..0.007 rows=1 loops=1)
                                             Buckets: 1024 Batches: 1 Memory Usage: 9kB
                                             -> Seq Scan on kind k (cost=0.00..1.12 rows=1 width=2) (actual time=0.004..0.005 rows=1 loops=1)
                                                  Filter: ((kind)::text = 'movie'::text)
                                                  Rows Removed by Filter: 9
                                 -> Hash (cost=19849.79..19849.79 rows=70837 width=10) (actual time=172.427..172.427 rows=70794 loops=1)
                                      Buckets: 131072 Batches: 2 Memory Usage: 2474kB
                                      -> Seq Scan on ratings r (cost=0.00..19849.79 rows=70837 width=10) (actual time=0.009..154.582 rows=70794 loops=1)
                                             Filter: ((rating > '5'::numeric) AND (votes > 500))
                                             Rows Removed by Filter: 858059
               -> Index Scan using people pkey on people p (cost=0.43..0.48 rows=1 width=24) (actual time=0.008..0.008 rows=1 loops=2218756)
                     Index Cond: ((name id)::text = (pt.name id)::text)
                    Filter: (death year IS NULL)
                     Rows Removed by Filter: 0
```



Top level Expensive path

This is just the top and the real work is below



Where does the above group get its data?

- This pushes data up to the group nodes
- The sort is spilling to disk, but only accounts for 11.7 seconds of the total 37.8 seconds of the total query time
- Inside the Nested Loop is taking 25.6 seconds



Finally we get to the node taking all the time

- The index scan takes most of the time
- A single iteration takes 0.008ms
- It is performed 2.2 million times taking over 17 seconds



Explain Options

- ANALYZE [boolean]
 - Run the query and return actual values
- COSTS [boolean]
 - show the costs of each node, on by default
- BUFFERS [boolean]
 - show buffer usage
- SUMMARY [boolean]
 - show the timing information after the plan
- TIMING [boolean]
 - show actual timings, on by default
- VERBOSE [boolean]
 - show additional details like the output columns
- FORMAT { TEXT | XML | JSON | YAML }
 - output format, TEXT by default



Explain Options

```
=> EXPLAIN (ANALYZE, BUFFERS true, VERBOSE true) SELECT * FROM cities
ORDER BY name;
                                   OUERY PLAN
 Sort (cost=8216.52...8374.55 rows=63211 width=43)
       (actual time=140.296..176.616 rows=63211 loops=1)
   Output: id, name, state, state name, county, alias
   Sort Key: cities.name
   Sort Method: external merge Disk: 3448kB
   Buffers: shared hit=598, temp read=431 written=433
   -> Seq Scan on public.cities (cost=0.00..1230.11 rows=63211
width=43)
                                  (actual time=0.011..5.731 rows=63211
loops=1)
         Output: id, name, state, state name, county, alias
         Buffers: shared hit=598
Planning Time: 0.053 ms
 Execution Time: 180.949 ms
```



Explain Options

```
=> EXPLAIN (FORMAT json) SELECT * FROM game changers;
                    QUERY PLAN
     "Plan": {
       "Node Type": "Aggregate",
       "Strategy": "Sorted",
       "Partial Mode": "Simple",
       "Parallel Aware": false,
       "Startup Cost": 546632.16,
       "Total Cost": 549163.19,
       "Plan Rows": 37497,
       "Plan Width": 22,
       "Group Key": ["p.name"],
       "Filter": "(count(*) > 10)",
       "Plans": [
```

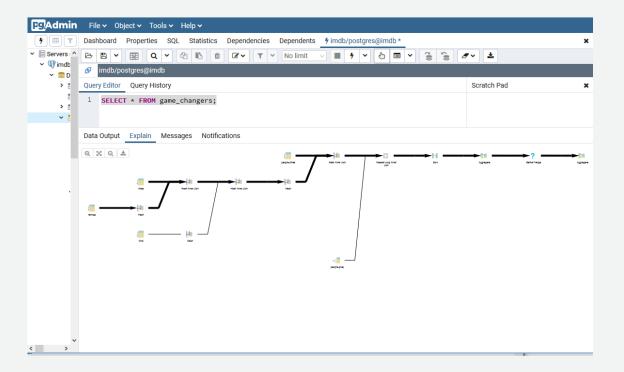
aws

pgAdmin Visual EXPLAIN

 pgAdmin has a visual EXPLAIN feature built into the query tool

 Produces a graphical output of the explain plan

Simpler to see the flow of an execution

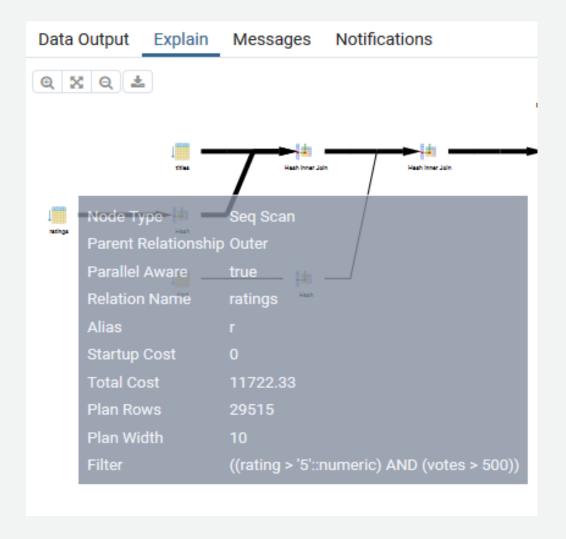




pgAdmin Visual EXPLAIN

 Diagnosing the exact place in the plan can be more difficult

Does not return all of the details necessary



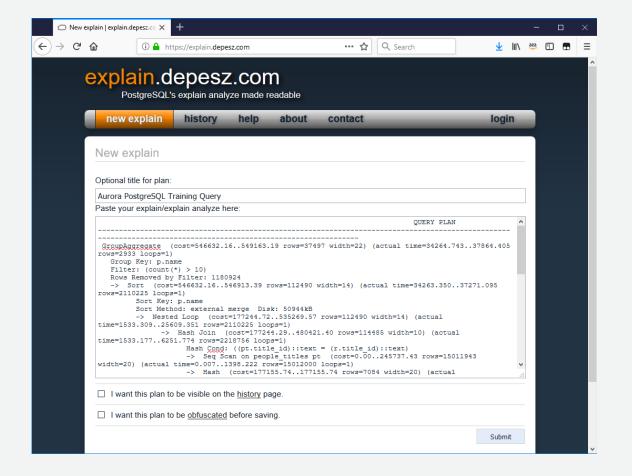


explain.depesz.com

 Simple online tool to interpet EXPLAIN plans

Paste and existing plan into the site

Fully open source so can be run in secure environments

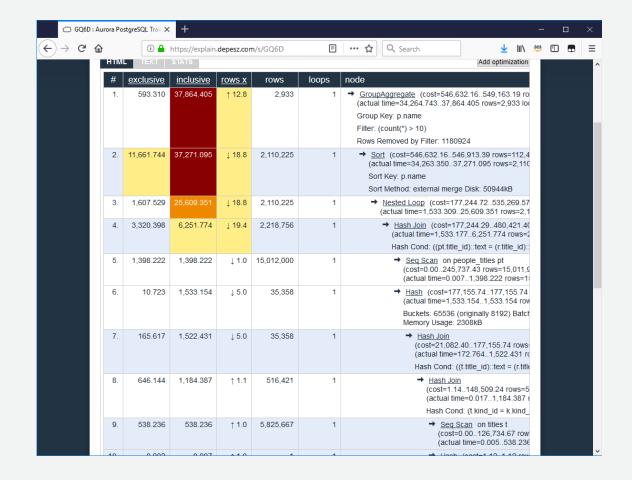




explain.depesz.com

 Highlights areas of the plan that are most expensive

Red and Orange and places to focus

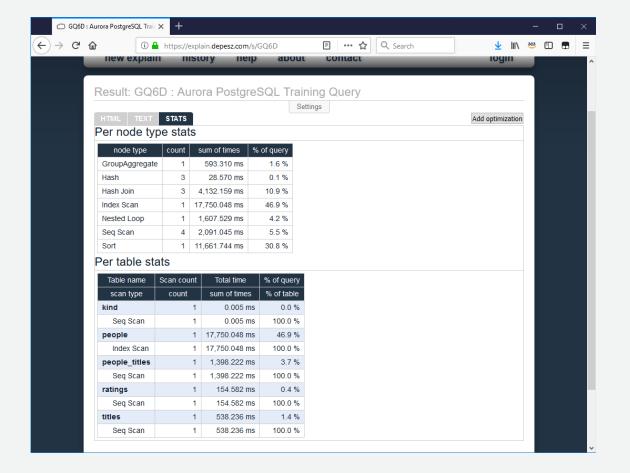




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 Stats show a breakdown by node and table

Easy to focus on problem areas





Summary

• There is a learning curve to understanding EXPLAIN output

EXPLAIN does not give the tuning answer, only where to start looking

Remember that session state affects EXPLAIN results



Planner Statistics



Planner Statistics

- Statistics are used by the planner to to estimate the cost of a query plan
- Collected by
 - ANALYZE
 - VACUUM
 - CREATE INDEX
 - REINDEX
 - CLUSTER
- Stored in pg_class and pg_statistic



Analyze

- Analyze is used to collect statistics about a table
- Results are stored in pg_statistic
- Query planner uses these statistics to increase efficiency
- Analyze should be rerun against tables with many changes



Looking at pg_class

• Most columns are structural, but some are statistics

Column	Type
relname	name
relpages	integer
reltuples	real

relname	reltuples	relpages
		T
pgbench_tellers	100	1
pgbench_branches	10	1
pgbench_accounts	1e+06	16394
pgbench history	0	0



Understanding pg_statistic

Contains all of the statistics of all relations

Column	Type	Column	Type
starelid	oid	staop3	oid
staattnum	smallint	staop4	oid
stainherit	boolean	staop5	oid
stanullfrac	real	stanumbers1	real[]
stawidth	integer	stanumbers2	real[]
stadistinct	real	stanumbers3	real[]
stakind1	smallint	stanumbers4	real[]
stakind2	smallint	stanumbers5	real[]
stakind3	smallint	stavalues1	anyarray
stakind4	smallint	stavalues2	anyarray
stakind5	smallint	stavalues3	anyarray
staop1	oid	stavalues4	anyarray
staop2	oid	stavalues5	anyarray

Understanding pg_statistic



Understanding pg_stats

• Contains all of the statistics of all relations

Column	Type
schemaname	name
tablename	name
attname	name
inherited	boolean
null_frac	real
avg_width	integer
n_distinct	real
most_common_vals	anyarray
most_common_freqs	real[]
histogram_bounds	anyarray
correlation	real
most_common_elems	anyarray
most_common_elem_freqs	real[]
elem_count_histogram	real[]

Understanding pg_stats

```
SELECT attname, null frac, avg width, n distinct,
      most common vals, most common freqs, histogram bounds
 FROM pg stats
WHERE schemaname = 'pg catalog' AND tablename = 'pg class';
-[ RECORD 2 ]----+----
attname
                relname
null frac
avg width
                64
n distinct
most common vals (null)
most common freqs (null)
histogram_bounds {_pg_foreign_data_wrappers,_pg_foreign_tables,...
```



Understanding pg_stats



Helping the Optimizer

Understand and react to uneven distributions of data

- default_statistics_target
 - Effects collection of most frequent values and histogram statistics

- ALTER TABLE SET STATISTICS
 - Set low on columns with simple data distributions
 - Set higher on frequently accessed search or join columns with known complex data



Indexes



Indexes

Indexes are a common way to enhance performance

- PostgreSQL supports several index types:
 - B-tree (default)
 - Hash Use when a simple comparison using the "=" operator is needed
 - Index on Expressions Use when an often used expression is queried. Inserts and updates will be slower.
 - Partial Index Indexes only rows that satisfy the WHERE clause. A query must include the same WHERE clause.

```
CREATE INDEX <name> on  (<column>);
CREATE INDEX <name> ON  USING HASH (<column>);
CREATE INDEX <name> on (expression(<column(s)>));
CREATE INDEX <name> ON  (<column>) WHERE <where clause>;
```



B-Trees Index

• DELETEs don't remove links – improves concurrency

• INSERTs on Rightmost page most efficient

• Suitable for both unique and highly non-unique data, or very skewed data

Can use index-only scans



B-Trees Index

```
EXPLAIN SELECT original_title FROM titles WHERE tid = 'tt0076759';
                             OUERY PLAN
Index Scan using titles pkey on titles
     (cost=0.43..8.45 rows=1 width=20)
  Index Cond: ((tid)::text = 'tt0076759'::text)
(2 rows)
EXPLAIN SELECT primary title FROM titles WHERE tid = 'tt0076759';
                             OUERY PLAN
Index Only Scan using titles id title idx on titles
     (cost=0.56..4.58 rows=1 width=20)
  Index Cond: (tid = 'tt0076759'::text)
(2 rows)
```



Hash Index

Allows equality search only

Ideal for lookups of large strings like UUIDs

Can not be used as unique constraints



Hash Index

```
CREATE INDEX titles guid hidx ON titles USING HASH (guid);
EXPLAIN ANALYZE SELECT primary title FROM titles
        WHERE guid = 'e6524bb6-761e-11e8-a734-c73f598fdce4';
                              OUERY PLAN
Index Scan using titles guid hidx on titles
  (cost=0.00..8.02 rows=1 width=20)
  (actual time=0.010..0.011 rows=1 loops=1)
  Index Cond: (guid = 'e6524bb6-761e-11e8-a734-c73f598fdce4'::uuid)
Planning time: 0.056 ms
Execution time: 0.024 ms
(4 rows)
```



GIST Index

- Generalized Search Tree
 - Allows application/custom data type specific indexes

- Used by
 - TSEARCH2 Full text indexing
 - POSTGIS R-Trees for spatial searching



GIST Index

```
CREATE INDEX titles title gist idx ON titles
    USING GIST (primary title gist trgm ops);
EXPLAIN ANALYZE SELECT tid, primary title FROM titles
              WHERE primary title ILIKE '%star wars%';
                                       OUERY PLAN
Bitmap Heap Scan on titles (cost=31.94..1746.26 rows=455 width=30)
                     (actual time=789.989..800.880 rows=1668 loops=1)
  Recheck Cond: ((primary title)::text ~~* '%star wars%'::text)
  Rows Removed by Index Recheck: 20
  Heap Blocks: exact=1413
  -> Bitmap Index Scan on titles title gist idx
             (cost=0.00..31.83 \text{ rows}=455 \text{ width}=0)
         (actual time=789.719..789.719 rows=1688 loops=1)
Index Cond: ((primary_title)::text ~~* '%star wars%'::text)
Planning time: 0.568 ms
Execution time: 801.639 ms
(8 rows)
```



GIN Index

- Generalized Inverted Index
 - Allows application/custom data type specific indexes

- Used by
 - TSEARCH2 Full text indexing
 - POSTGIS R-Trees for spatial searching



GIN Index

```
CREATE INDEX titles title gin idx ON titles
     USING GIN (primary_title gin_trgm_ops);
EXPLAIN ANALYZE SELECT tid, primary title FROM titles
                WHERE primary title ILIKE '%star wars%';
                           QUERY PLAN
Bitmap Heap Scan on titles (cost=91.54..1809.51 rows=456 width=30)
                     (actual time=41.174..45.603 rows=1668 loops=1)
  Recheck Cond: ((primary_title)::text ~~* '%star wars%'::text)
  Rows Removed by Index Recheck: 20
  Heap Blocks: exact=1413
  -> Bitmap Index Scan on titles title gin idx
             (cost=0.00..91.42 rows=456 width=0)
             (actual time=40.905..40.905 rows=1688 loops=1)
        Index Cond: ((primary title)::text ~~* '%star wars%'::text)
Planning time: 0.543 ms
Execution time: 45.751 ms
(8 rows)
```



GIST vs GIN

GIN is about 3X faster for lookups than GIST

GIN takes about 3X longer to build than GIST

GIN is about 10X slower to update than GIST

• GIN is about 2-3X smaller on disk than GIST



Partial Index

- Allows you to index just some of the rows of a table
 - e.g. Column is unique except when NULL
 - e.g. Fast access required only to Employee.CPRTrained = 'Y'

Defined by WHERE clause on CREATE INDEX

- Must be defined using constants
 - Cannot build index WHERE log_date > current_date 30 since current_date value changes as it is executed



Partial Index

```
CREATE INDEX titles movies id idx ON titles(tid)
  WHERE title type = 'movie';
 EXPLAIN ANALYZE SELECT * FROM titles
 WHERE tid = 'tt0076759' AND title type = 'movie';
                          QUERY PLAN
 Index Scan using titles movies id idx on titles
   (cost=0.42..8.44 rows=1 width=79)
   (actual time=0.047..0.049 rows=1 loops=1)
   Index Cond: ((tid)::text = 'tt0076759'::text)
Planning time: 0.400 ms
Execution time: 0.159 ms
(4 rows)
```



Multiple Indexes

```
EXPLAIN ANALYZE SELECT * FROM people WHERE first name = 'Bruce' AND last name =
'Lee':
                                        OUERY PLAN
Bitmap Heap Scan on people (cost=452.18..456.19 rows=1 width=42) (actual time=4.824..4.844 rows=18 loops=1)
 Recheck Cond: (((first name)::text = 'Bruce') AND ((last name)::text = 'Lee'))
 Heap Blocks: exact=18
 -> BitmapAnd (cost=452.18..452.18 rows=1 width=0)
                 (actual time=4.776..4.776 rows=0 loops=1)
   -> Bitmap Index Scan on people first name idx (cost=0.00..12.88 rows=592
width=0)
                                           (actual time=1.322..1.322 rows=7956
loops=1)
         Index Cond: ((first name)::text = 'Bruce'::text)
   -> Bitmap Index Scan on people last name idx (cost=0.00..439.05 rows=21682
width=0)
                                             (actual time=3.098..3.098 rows=21947
loops=1)
         Index Cond: ((last name)::text = 'Lee'::text)
Planning Time: 0.101 ms
Execution Time: 4.873 ms
```



Compound Indexes

- Multiple columns can be used in a single index
- Order of the columns matter

```
CREATE INDEX people name idx ON people (last name, first name);
EXPLAIN ANALYZE SELECT * FROM people WHERE first name = 'Bruce' AND
last name = 'Lee';
                                      OUERY PLAN
 Index Scan using people name idx on people (cost=0.56..8.58 rows=1
width=42)
                                    (actual time=0.038..0.056 rows=18
loops=1)
   Index Cond: (((last_name)::text = 'Lee') AND ((first_name)::text =
'Bruce'))
Planning Time: 0.304 ms
Execution Time: 0.073 ms
```



• Indexes can help queries perform much better

• The correct type and scope of the index matter

• Indexes come at a cost. Write performance is affected



Understanding the Optimizer



What is the Optimizer?

Creates an optimal execution plan for a query

- Examines a SQL query in a variety of different ways
 - If feasible, it will evaluate each one

- The number of possible plans grows exponentially as the number of joined tables increases
 - Uses the algorithm introduced by the System R database (became DB2)
 - For more than 12 joins, switches to a Genetic Query Optimization

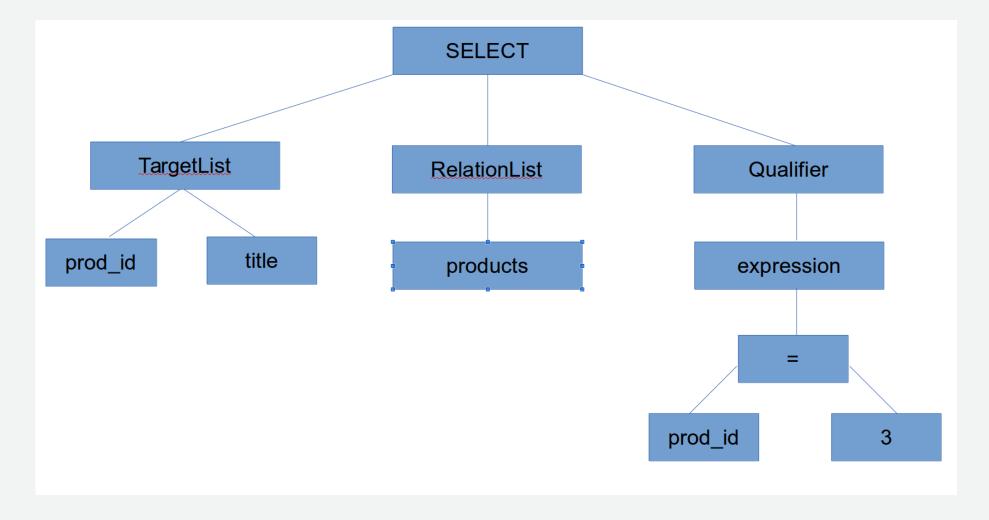
Costs are measured relative to the cost of a sequential page fetch



Parse Trees

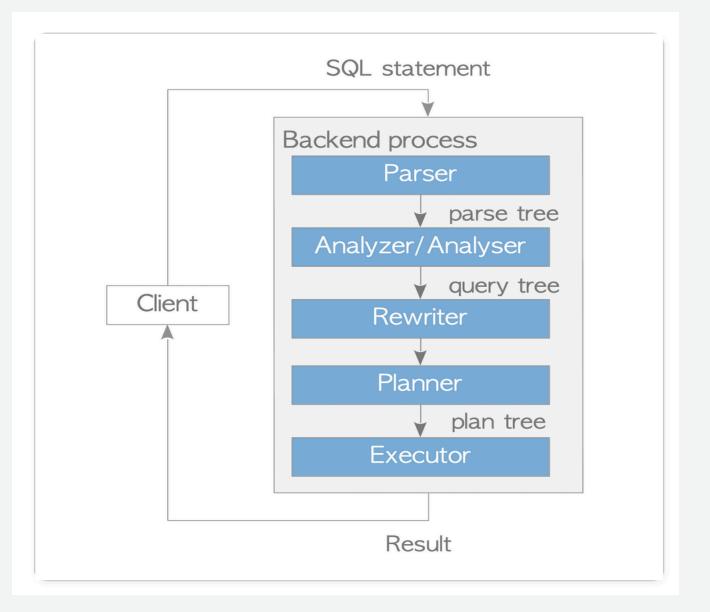
```
select prod_id, title
from products
where prod_id = 3
```







- Parser: Generates a parse tree from the statement in Plain Text
- Analyzer: Carries out a semantic analysis of the parse tree and generates a query tree
- Rewriter: Transforms a query tree using rules in the rules system
- Planner: Planner generates the tree that can be most effectively be executed from the query tree
- Executor: Executes the query by accessing the tables and indexes in the order that was created in the plan tree.





Getting the Parse Tree

Turn on parse debugging

```
SET debug print parse = on;
SET client_min_messages='LOG';
LOG: parse tree:
DETAIL:
            {QUERY
   :commandType 1
   :cteList <>
   :rtable (
      {RTE
      :alias <>
      :eref
         {ALIAS
         :aliasname foo
         :colnames ("row_owner" "col1" "col2")
      :rtekind 0
      :relid 16386
      :relkind r
      :tablesample <>
      :selectedCols (b 10 11)
      :insertedCols (b)
      :updatedCols (b)
      :securityQuals <>
```



Turning a Parse Tree into a Plan

- Uses a series of query operators to return a result set
- The planner generates many execution plans and the optimizer determines the leastexpensive plan
- Uses tables statistics to determine



Table Level Operators

- 6 methods to retrieve data from a table
 - Table Scan (seq scan)
 - Full scan of the table from beginning to end
 - Index Scan
 - Scans an index to find the appropriate row in the table
 - Bitmap Index Scan
 - Scans the full index and finds matching keys
 - Index Only Scan
 - Returns the target list out of the index
 - Remote Scan
 - Scan a foreign server
 - tuple-ID Scan (TID scan)
 - Only used ctid is used as a criteria



Join Operators

- Only 3 possible join strategies
 - Nested Loop Join
 - The right relation scanned once for every relation on the left
 - Merge Sort Join
 - Each relation is sorted and scanned once in parallel
 - Hash Join
 - The right relation is loaded into a hash table and the left relation uses hash keys to find matching rows



Other Operators

- Sort
 - Orders the result set either in-memory or on-disk
- Unique
 - Eliminates duplicate rows from an input set
- Limit
 - Limits the size of a result set
- Aggregate
 - Reads all of the rows in the input set and calculates the aggregate value
- Append
 - Implements a UNION
- Result
 - Used for a query that does not return data



Other Operators (cont.)

- Group
 - Used to satisfy a GROUP BY clause
- Subquery Scan
 - Used to satisfy a UNION
- Subplan
 - Used to satisfy a subselect
- Materialize
 - Used to materialize a subselect
- Setop
 - Used for INTERSECT and EXCEPT operators



Affecting the Optimizer

- Configuration settings
 - EFFECTIVE_CACHE_SIZE
 - The estimate of the cache available for an index scan. Higher values favor index scans
 - RANDOM_PAGE_COST
 - The cost of reading a non sequential disk page. Higher values favor sequential scans
 - CPU_INDEX_TUPLE_COST
 - The cost of processing an index row.
 - CPU_OPERATOR_COST
 - The cost of processing each operator in the WHERE clause
 - CPU_TUPLE_COST
 - The cost of processing each row



Operators Costing Algorithms

- Sequential Scan
 - Cost = Number of Pages

- Index Scan
 - Uses the Mackert and Lohman approximation algorithm



Optimizer Hints



Hints (Caution)

• Hints are not magic

• They should be used with extreme care



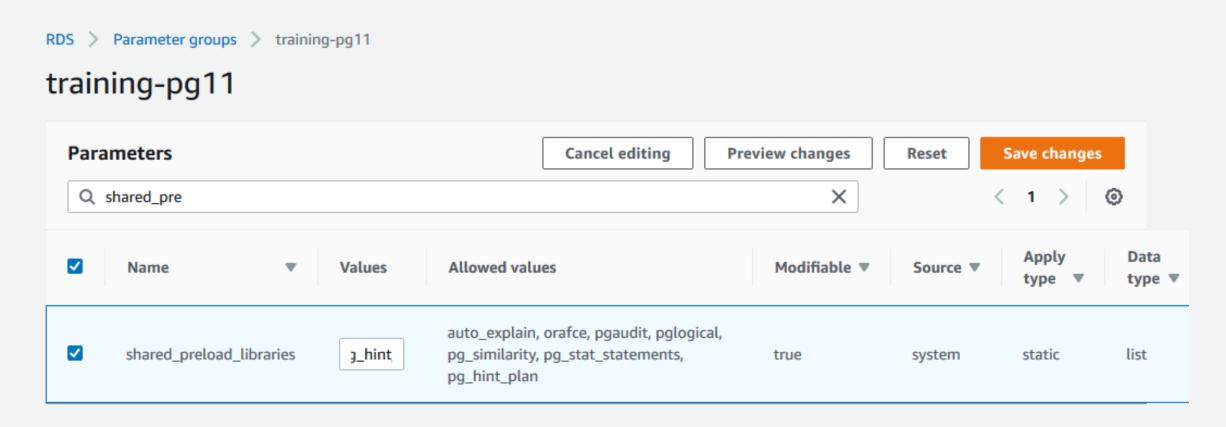
What are hints?

• Allows for the developer or operator to influence the optimizer

• It is a directive to the optimizer to be followed if possible



Using Hints



pg_hint_plan must be added to shared_preload_libraries



Adding Hints

- Hints must be the first element of a statement
 - Will be treated as a comment otherwise

```
/*+
    BitmapScan(titles titles_id_brin_idx)
*/
EXPLAIN ANALYZE SELECT * FROM titles WHERE title_id = 'tt0076759';
```



Hints Types

• Scans - How is data retrieved from a single table

• Joins - How are tables combined

• Environment - Parameters and statistics used by the query



Hints Types - Scans

Format	Description
SeqScan(table)	Forces sequential scan on the table
TidScan(table)	Forces TID scan
IndexScan(table[index])	Forces index scan
IndexOnlyScan(table[index])	Forces index only scan
BitmapScan(table[index])	Forces bitmap scan
NoSeqScan(table)	Do not do sequential scan on the table



Hints Types - Joins

Format	Description
NestLoop(table table[table])	Forces nested loop for the joins
HashJoin(table table[table])	Forces hash join
MergeJoin(table table[table])	Forces merge join
NoNestLoop(table table[table])	Do not do nested loop join
Leading(table table[table])	Forces join order as specified



Hints Types - Environment

Format	Description
Rows(table table[table] correction)	Corrects row number of a result of the joins
Set(param value)	Set the parameter while planner is running



• Hints can be a powerful tool to tune queries

Use with caution





- Use EXPLAIN on all important queries to ensure it is executing the optimal way
- Knowing the data and the data usage helps determine the correct index type
- Just because an index scan is used does not mean it is the best index scan possible
- PostgreSQL does not limit the number of indexes on a table, but each index adds overhead to writes

